



**Body Fat DFE MCU**

**BH66R2640**

Revision: V1.00 Date: December 31, 2024

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## Features

### CPU Features

- Operating voltage
  - ♦  $f_{SYS}=8\text{MHz}$ : 2.2V~5.5V
- Up to 0.5 $\mu\text{s}$  instruction cycle with 8MHz system clock at  $V_{DD}=5\text{V}$
- Power down and wake-up functions to reduce power consumption
- Oscillator types
  - ♦ Internal High Speed 8MHz RC – HIRC
  - ♦ Internal Low Speed 32kHz RC – LIRC
- Multi-mode operation: FAST, SLOW, IDLE and SLEEP
- Fully integrated internal oscillators require no external components
- All instructions executed in 1~3 instruction cycles
- Table read instructions
- 115 powerful instructions
- 12-level subroutine nesting
- Bit manipulation instruction

### Peripheral Features

- Flash Program Memory: 4K $\times$ 16
- Data Memory: 512 $\times$ 8
- OTP ROM parameter program – ORPP
- Watchdog Timer function
- 15 bidirectional I/O lines
- Programmable I/O port source current and sink current for LED driving applications
- Two external interrupt lines shared with I/O pins
- One 10-bit CTM for time measure, compare match output or PWM output function
- Serial Interface Module – SIM includes SPI and I<sup>2</sup>C interfaces
- Fully-duplex/Half-duplex Universal Asynchronous Receiver and Transmitter Interface – UART
- Dual Time-Base functions for generation of fixed time interrupt signals
- 2 differential or 4 single-end external channel 24-bit resolution Delta Sigma A/D converter
- Internal LDO with bypass function for PGA, ADC and external sensor power supply
- Body Fat Measurement Circuit
- Low voltage reset function
- Package types: 24-pin SSOP and 32-pin QFN

## General Description

This device is specifically designed for four-electrode AC Body Fat Scale applications. Measuring body fat uses a technique whereby an AC current flowing through the human body is measured and then used to calculate a body fat value. The specialised circuits to do this are a weight measurement circuit and a fat measurement circuit. The weight measurement circuit uses an external load cell to output a signal, which after amplification by an operational amplifier, and then conversion using an A/D converter, reads the corresponding value as the calculated weight. The fat measurement circuit uses an AC signal via an electrode slice to flow through human body. After amplification by an internal operational amplifier, and then conversion by an A/D converter, the measured value is one representing body impedance, which is used to calculate the corresponding body fat value.

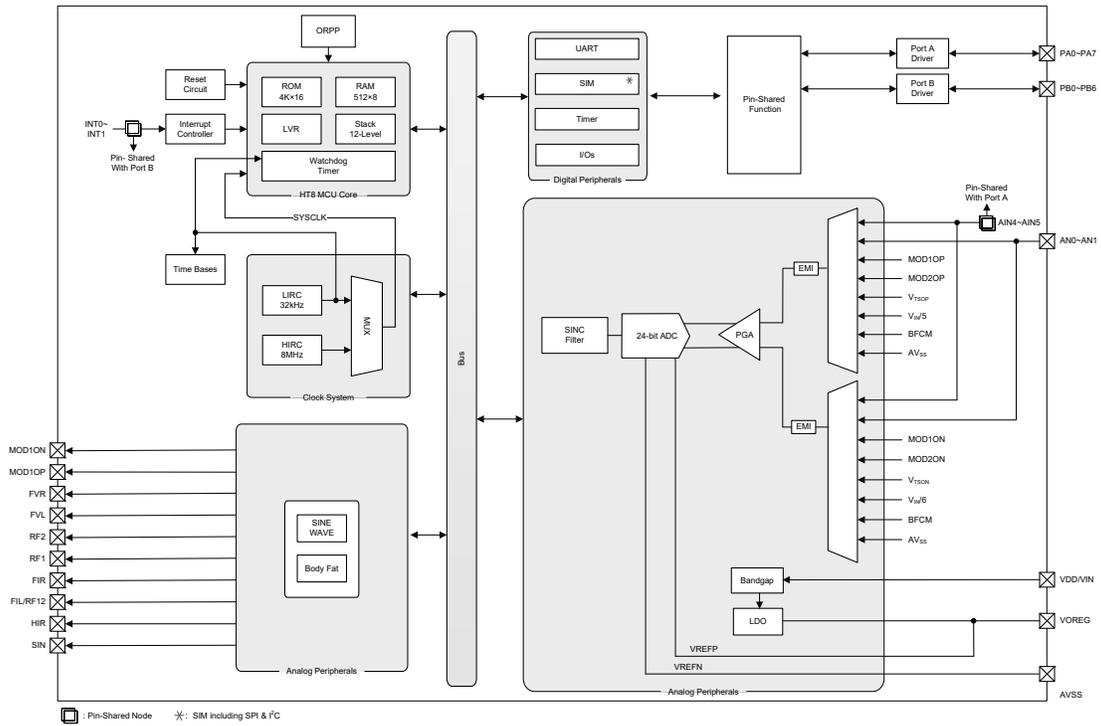
This device is an OTP Program Memory A/D type 8-bit high performance RISC architecture microcontroller which includes a multi-channel 24-bit Delta Sigma A/D converter, designed for applications that interface directly to analog signals and which require a low noise and high accuracy analog-to-digital converter. For memory features, the device is supplied with One-Time Programmable, OTP memory. Other memory includes an area of RAM Data Memory.

Analog features include a multi-channel 24-bit Delta Sigma A/D converter and other circuitry specifically designed for Body Fat Scale applications. Multiple and extremely flexible Timer Module provides timing, compare match output and PWM generation functions. Communication with the outside world is provided by including fully integrated SPI, I<sup>2</sup>C and UART interface functions, these popular interfaces which provide designers with a means of easy communication with external peripheral hardware. Protective features such as an internal Watchdog Timer and Low Voltage Reset coupled with excellent noise immunity and ESD protection ensure that reliable operation is maintained in hostile electrical environments.

A full choice of HIRC and LIRC oscillator functions are provided including a fully integrated system oscillator which requires no external components for its implementation. The ability to operate and switch dynamically between a range of operating modes using different clock sources gives users the ability to optimise microcontroller operation and minimise power consumption.

The inclusion of flexible I/O programming features, Time-Base functions along with many other features ensure that only a minimum of external components is required for application implementation, resulting in reduced component costs and reductions in circuit board areas.

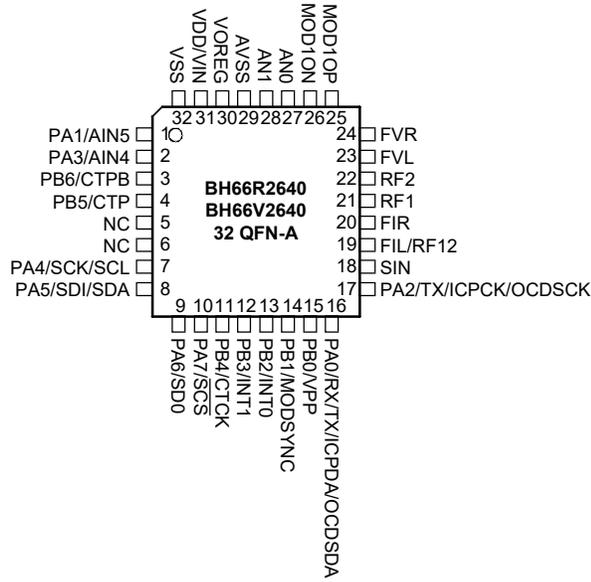
### Block Diagram



### Pin Assignment

NC	1	24	PA2/TX/ICPCK/OCDSCK
NC	2	23	PA0/RX/TX/ICPDA/OCSDA
NC	3	22	PB0/VPP
AN0	4	21	PB1/MODSYNC
AN1	5	20	PB2/INT0
AVSS	6	19	PB3/INT1
VOREG	7	18	PB4/CTCK
VDD/VIN	8	17	PA7/SCS
VSS	9	16	PA6/SD0
PA1/AIN5	10	15	PA5/SDI/SDA
PA3/AIN4	11	14	PA4/SCK/SCL
PB6/CTPB	12	13	PB5/CTP

**BH66R2640/BH66V2640**  
**24 SSOP-A**



- Note:
1. If the pin-shared pin functions have multiple outputs, the desired pin-shared function is determined by the corresponding software control bits.
  2. The OCDSDA and OCDSCK pins are supplied as the OCDS dedicated pins and as such only available for the BH66V2640 device (Flash type) which is the OCDS EV chip for the BH66R2640 device (OTP type).
  3. The VPP pin is the High Voltage input for OTP programming and only available for the BH66R2640 device.
  4. For the less pin count package types there will be unbounded pins which should be properly configured to avoid unwanted power consumption resulting from floating input conditions. Refer to the “Standby Current Considerations” and “Input/Output Ports” sections.

## Pin Description

The function of each pin is listed in the following table, however the details behind how each pin is configured is contained in other sections of the datasheet. As the Pin Description table shows the situation for the package with the most pins, not all pins in the table will be available on smaller package sizes.

Pin Name	Function	OPT	I/T	O/T	Description
PA0/RX/TX/ICPDA/OCDSDA	PA0	PAPU PAWU PAS0	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	RX/TX	PAS0	ST	CMOS	UART serial data input in full-duplex communication or UART serial data input/output in Single Wire Mode communication
	ICPDA	—	ST	CMOS	ICP address/data pin
	OCDSDA	—	ST	CMOS	OCDS address/data pin, for EV chip only
PA1/AIN5	PA1	PAPU PAWU PAS0	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	AIN5	PAS0	AN	—	A/D Converter external input channel

Pin Name	Function	OPT	I/T	O/T	Description
PA2/TX/ICPCK/OCDSCK	PA2	PAPU PAWU PAS0	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	TX	PAS0	—	CMOS	UART TX serial data output
	ICPCK	—	ST	—	ICP clock pin
	OCDSCK	—	ST	—	OCDS clock pin, for EV chip only
PA3/AIN4	PA3	PAPU PAWU PAS0	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	AIN4	PAS0	AN	—	A/D Converter external input channel
PA4/SCK/SCL	PA4	PAPU PAWU PAS1	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	SCK	PAS1	ST	CMOS	SIM SPI serial clock
	SCL	PAS1	ST	CMOS	SIM I <sup>2</sup> C clock line
PA5/SDI/SDA	PA5	PAPU PAWU PAS1	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	SDI	PAS1	ST	—	SIM SPI data input
	SDA	PAS1	ST	CMOS	SIM I <sup>2</sup> C data line
PA6/SDO	PA6	PAPU PAWU PAS1	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	SDO	PAS1	—	CMOS	SPI serial data output
PA7/ $\overline{SCS}$	PA7	PAPU PAWU PAS1	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	$\overline{SCS}$	PAS1	ST	CMOS	SIM SPI slave select pin
PB0/VPP	PB0	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up. This I/O is pin-shared with VPP and could result in increased current consumption if it is set to output high.
	VPP	—	PWR	—	High Voltage input for OTP programming pin, not available for EV chip
PB1/MODSYNC	PB1	PBPU PBS0	ST	CMOS	General purpose I/O. Register enabled pull-up
	MODSYNC	PBS0	ST	—	External IQMOD sync clock input
PB2/INT0	PB2	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up
	INT0	INTEG INTC0	ST	—	External Interrupt 0 input
PB3/INT1	PB3	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up
	INT1	INTEG INTC0	ST	—	External Interrupt 1 input
PB4/CTCK	PB4	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up
	CTCK	—	ST	—	CTM clock input
PB5/CTP	PB5	PBPU PBS1	ST	CMOS	General purpose I/O. Register enabled pull-up
	CTP	PBS1	—	CMOS	CTM output
PB6/CTPB	PB6	PBPU PBS1	ST	CMOS	General purpose I/O. Register enabled pull-up
	CTPB	PBS1	—	CMOS	CTM inverting output
VOREG	VOREG	—	—	PWR	LDO output pin
		—	PWR	—	Positive power supply for ADC and PGA

Pin Name	Function	OPT	I/T	O/T	Description
AVSS	AVSS	—	PWR	—	Positive power supply for ADC and PGA
AN0~AN1	ANn	—	AN	—	A/D Converter external input channel
MOD1ON	MOD1ON	—	—	AN	Demodulator 1 negative output
MOD1OP	MOD1OP	—	—	AN	Demodulator 1 positive output
FVR	FVR	—	AN	AN	Right foot voltage channel
FVL	FVL	—	AN	AN	Left foot voltage channel
RF2	RF2	—	AN	AN	Reference 2 impedance channel
RF1	RF1	—	AN	AN	Reference 1 impedance channel
FIR	FIR	—	AN	AN	Right foot current channel
FIL/RF12	FIL	—	AN	AN	Left foot current channel
	RF12	—	AN	AN	Reference 1/2 impedance channel
SIN	SIN	—	—	AN	Sine wave output
VDD/VIN	VDD	—	PWR	—	Positive power supply
	VIN	—	PWR	—	LDO input pin
VSS	VSS	—	PWR	—	Negative power supply

Legend: I/T: Input type; O/T: Output type;  
 OPT: Optional by register option; PWR: Power;  
 ST: Schmitt Trigger input; CMOS: CMOS output;  
 AN: Analog signal.

## Absolute Maximum Ratings

Supply Voltage .....	$V_{SS}-0.3V$ to $6.0V$
Input Voltage .....	$V_{SS}-0.3V$ to $V_{DD}+0.3V$
Storage Temperature.....	$-60^{\circ}C$ to $150^{\circ}C$
Operating Temperature.....	$-40^{\circ}C$ to $85^{\circ}C$
$I_{OH}$ Total.....	$-80mA$
$I_{OL}$ Total .....	$80mA$
Total Power Dissipation .....	$500mW$

Note: These are stress ratings only. Stresses exceeding the range specified under “Absolute Maximum Ratings” may cause substantial damage to the device. Functional operation of the device at other conditions beyond those listed in the specification is not implied and prolonged exposure to extreme conditions may affect device reliability.

## D.C. Characteristics

For data in the following tables, note that factors such as oscillator type, operating voltage, operating frequency, pin load conditions, temperature and program instruction type, etc., can all exert an influence on the measured values.

### Operating Voltage Characteristics

$T_a=-40^{\circ}C-85^{\circ}C$

Symbol	Parameter	Test Conditions	Min.	Typ.	Max.	Unit
$V_{DD}$	Operating Voltage – HIRC	$f_{SYS}=f_{HIRC}=8MHz$	2.2	—	5.5	V
	Operating Voltage – LIRC	$f_{SYS}=f_{LIRC}=32kHz$	2.2	—	5.5	V

### Standby Current Characteristics

Ta=25°C, unless otherwise specified

Symbol	Standby Mode	Test Conditions		Min.	Typ.	Max.	Max. @85°C	Unit
		V <sub>DD</sub>	Conditions					
I <sub>STB</sub>	SLEEP Mode	2.2V	WDT off	—	0.45	0.80	1.90	μA
		3V		—	0.45	0.90	2.60	
		5V		—	0.5	2.0	5.0	
		2.2V	WDT on	—	1.2	2.4	2.9	μA
		3V		—	1.5	3.0	3.6	
		5V		—	3	5	6	
	IDLE0 Mode – LIRC	2.2V	f <sub>SUB</sub> on	—	2.4	4.0	4.8	μA
		3V		—	3	5	6	
		5V		—	5	10	12	
IDLE1 Mode – HIRC	2.2V	f <sub>SUB</sub> on, f <sub>SYS</sub> =8MHz	—	288	400	480	μA	
	3V		—	360	500	600		
	5V		—	600	800	960		

Note: When using the characteristic table data, the following notes should be taken into consideration:

1. Any digital inputs are setup in a non-floating condition and the I/O pin-shared with VPP is not setup in an output high condition.
2. All measurements are taken under conditions of no load and with all peripherals in an off state.
3. There are no DC current paths.
4. All Standby Current values are taken after a HALT instruction execution thus stopping all instruction execution.

### Operating Current Characteristics

Ta=-40°C~85°C

Symbol	Operating Mode	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
I <sub>DD</sub>	SLOW Mode – LIRC	2.2V	f <sub>SYS</sub> =32kHz	—	8	16	μA
		3V		—	10	20	
		5V		—	30	50	
	FAST Mode – HIRC	2.2V	f <sub>SYS</sub> =8MHz	—	0.6	1.0	mA
		3V		—	0.8	1.2	
		5V		—	1.6	2.4	

Note: When using the characteristic table data, the following notes should be taken into consideration:

1. Any digital inputs are setup in a non-floating condition and the I/O pin-shared with VPP is not setup in an output high condition.
2. All measurements are taken under conditions of no load and with all peripherals in an off state.
3. There are no DC current paths.
4. All Operating Current values are measured using a continuous NOP instruction program loop.

## A.C. Characteristics

For data in the following tables, note that factors such as oscillator type, operating voltage, operating frequency and temperature etc., can all exert an influence on the measured values.

### Internal High Speed Oscillator – HIRC – Frequency Accuracy

During the program writing operation the writer will trim the HIRC oscillator at a user selected HIRC frequency and user selected voltage of either 3V or 5V.

Symbol	Parameter	Test Conditions		Min	Typ	Max	Unit
		V <sub>DD</sub>	Temp.				
f <sub>HIRC</sub>	8MHz Writer Trimmed HIRC Frequency	3V/5V	25°C	-1%	8	+1%	MHz
			-40°C~85°C	-4%	8	+4%	
		2.2V~5.5V	25°C	-2.5%	8	+2.5%	
			-40°C~85°C	-4.5%	8	+4.5%	

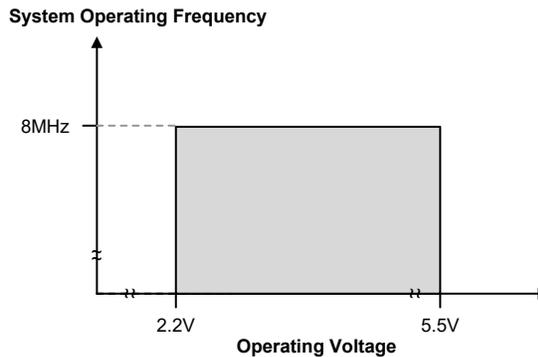
Note: 1. The 3V/5V values for V<sub>DD</sub> are provided as these are the two selectable fixed voltages at which the HIRC frequency is trimmed by the writer.

2. The row below the 3V/5V trim voltage row is provided to show the values for the full V<sub>DD</sub> range operating voltage. It is recommended that the trim voltage is fixed at 3V for application voltage ranges from 2.2V to 3.6V and fixed at 5V for application voltage ranges from 3.3V to 5.5V.

### Internal Low Speed Oscillator – LIRC – Frequency Accuracy

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Temp.				
f <sub>LIRC</sub>	LIRC Frequency	3V	25°C	-20%	32	+20%	kHz
		2.2V~5.5V	-40°C~85°C	-50%	32	+60%	
t <sub>START</sub>	LIRC Start Up Time	—	-40°C~85°C	—	—	500	µs

### Operating Frequency Characteristic Curves



### System Start Up Time Characteristics

Ta=-40°C~85°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
t <sub>SST</sub>	System Start-up Time Wake-up from Condition where f <sub>SYS</sub> is off	—	f <sub>SYS</sub> =f <sub>H</sub> ~ f <sub>H</sub> /64, f <sub>H</sub> =f <sub>HIRC</sub>	—	16	—	t <sub>HIRC</sub>
		—	f <sub>SYS</sub> =f <sub>SUB</sub> =f <sub>LIRC</sub>	—	2	—	t <sub>LIRC</sub>
	System Start-up Time Wake-up from Condition where f <sub>SYS</sub> is on	—	f <sub>SYS</sub> =f <sub>H</sub> ~ f <sub>H</sub> /64, f <sub>H</sub> =f <sub>HIRC</sub>	—	2	—	t <sub>H</sub>
		—	f <sub>SYS</sub> =f <sub>SUB</sub> =f <sub>LIRC</sub>	—	2	—	t <sub>SUB</sub>
	System Speed Switch Time FAST to SLOW Mode or SLOW to FAST Mode	—	f <sub>HIRC</sub> switches from off → on	—	16	—	t <sub>HIRC</sub>
t <sub>RSTD</sub>	System Reset Delay Time Reset Source from Power-on Reset or LVR Hardware Reset	—	RR <sub>POR</sub> =5V/ms	10	16	24	ms
	System Reset Delay Time WDTC/RSTC Software Reset	—	—	—	—	—	—
	System Reset Delay Time Reset Source from WDT Overflow	—	—	10	16	24	ms
t <sub>SRESET</sub>	Minimum Software Reset Width to Reset	—	—	45	90	120	μs

- Note: 1. For the System Start-up time values, whether f<sub>SYS</sub> is on or off depends upon the mode type and the chosen f<sub>SYS</sub> system oscillator. Details are provided in the System Operating Modes section.
2. The time units, shown by the symbols t<sub>HIRC</sub> etc. are the inverse of the corresponding frequency values as provided in the frequency tables. For example t<sub>HIRC</sub>=1/f<sub>HIRC</sub>, t<sub>SYS</sub>=1/f<sub>SYS</sub> etc.
3. If the LIRC is used as the system clock and if it is off when in the SLEEP Mode, then an additional LIRC start up time, t<sub>START</sub>, as provided in the LIRC frequency table, must be added to the t<sub>SST</sub> time in the table above.
4. The System Speed Switch Time is effectively the time taken for the newly activated oscillator to start up.

### Input/Output Characteristics

Ta=-40°C~85°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>IL</sub>	Input Low Voltage for I/O Ports or Input Pins	5V	—	0	—	1.5	V
		—	—	0	—	0.2V <sub>DD</sub>	
V <sub>IH</sub>	Input High Voltage for I/O Ports or Input Pins	5V	—	3.5	—	5	V
		—	—	0.8V <sub>DD</sub>	—	V <sub>DD</sub>	
I <sub>OL</sub>	Sink Current for I/O Pins	3V	V <sub>OL</sub> =0.1V <sub>DD</sub>	16	32	—	mA
				32	65	—	
I <sub>OH</sub>	Source Current for I/O Pins	3V	V <sub>OH</sub> =0.9V <sub>DD</sub> , SLEDC[m+1:m]=00B (m=0, 2, 4, 6)	-0.7	-1.5	—	mA
				-1.5	-2.9	—	
		5V	V <sub>OH</sub> =0.9V <sub>DD</sub> , SLEDC[m+1:m]=01B (m=0, 2, 4, 6)	-1.3	-2.5	—	mA
				-2.5	-5.1	—	
		3V	V <sub>OH</sub> =0.9V <sub>DD</sub> , SLEDC[m+1:m]=10B (m=0, 2, 4, 6)	-1.8	-3.6	—	mA
				-3.6	-7.3	—	
		5V	V <sub>OH</sub> =0.9V <sub>DD</sub> , SLEDC[m+1:m]=11B (m=0, 2, 4, 6)	-4	-8	—	mA
				-8	-16	—	

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
R <sub>PH</sub>	Pull-high Resistance for I/O Ports	3V	—	20	60	100	kΩ
		5V	—	10	30	50	
I <sub>LEAK</sub>	Input Leakage Current	5V	V <sub>IN</sub> =V <sub>DD</sub> or V <sub>IN</sub> =V <sub>SS</sub>	—	—	±1	μA
t <sub>CK</sub>	CTCK Clock Input Pin Minimum Pulse Width	—	—	0.3	—	—	μs
t <sub>INT</sub>	External Interrupt Minimum Pulse Width	—	—	10	—	—	μs

Note: The R<sub>PH</sub> internal pull high resistance value is calculated by connecting to ground and enabling the input pin with a pull-high resistor and then measuring the input sink current at the specified supply voltage level. Dividing the voltage by this measured current provides the R<sub>PH</sub> value.

## Memory Characteristics

T<sub>a</sub>=-40°C~85°C, unless otherwise specified

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
<b>OTP Program Memory</b>							
V <sub>DD</sub>	Operating Voltage for Read – ORPP	—	—	2.2	—	5.5	V
	Operating Voltage for Write – ORPP	—	—	3.0	—	5.5	V
V <sub>PP</sub>	V <sub>PP</sub> for Write – ORPP	—	—	8.25	8.50	8.75	V
t <sub>WR</sub>	Write Cycle Time – ORPP	—	—	—	300	450	μs
E <sub>P</sub>	Cell Endurance	—	—	1	—	—	W
t <sub>RETD</sub>	ROM Data Retention time	—	T <sub>a</sub> =25°C	—	40	—	Year
<b>Flash Program Memory – for BH66V2640 only</b>							
V <sub>DD</sub>	Operating Voltage for Read & Write	—	—	2.2	—	5.5	V
t <sub>WR</sub>	Write Time	—	—	—	2.2	2.7	ms
E <sub>P</sub>	Cell Endurance	—	—	1	—	—	W
t <sub>RETD</sub>	Data Retention Time	—	T <sub>a</sub> = 25°C	—	40	—	Year
t <sub>ACTV</sub>	ROM Activation Time	—	Wake-up from Power Down Mode	32	—	64	μs
<b>RAM Data Memory</b>							
V <sub>DR</sub>	RAM Data Retention Voltage	—	—	1.0	—	—	V

Note: 1. The ROM activation time t<sub>ACTV</sub> should be added when calculating the total system start-up time of a wake-up from the power down mode.  
 2. “W” means Write times.

## 24-bit Delta Sigma A/D Converter Electrical Characteristics

V<sub>DD</sub>=V<sub>IN</sub>, T<sub>a</sub>=25°C, unless otherwise specified  
 LDO Test conditions: MCU enters SLEEP mode, other functions disabled

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
<b>LDO</b>							
V <sub>IN</sub>	Supply Voltage for LDO, Bandgap	—	T <sub>a</sub> =0°C~45°C	2.4	—	5.5	V
I <sub>Q</sub>	LDO Quiescent Current	—	LDOVS[1:0]=00B, V <sub>IN</sub> =3.6V, No load	—	35	50	μA
I <sub>BGP</sub>	Bandgap quiescent current	—	BGPEN=1, V <sub>IN</sub> =3.6V, No load	—	150	200	μA

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>OUT_LDO</sub>	LDO Output Voltage	—	LDOVS[1:0]=00B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =0.1mA	-5%	2.4	+5%	V
		—	LDOVS[1:0]=01B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =0.1mA		2.6		
		—	LDOVS[1:0]=10B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =0.1mA		2.9		
		—	LDOVS[1:0]=11B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =0.1mA		3.3		
ΔV <sub>LOAD</sub>	LDO Load Regulation <sup>(1)</sup>	—	LDOVS[1:0]=00B, V <sub>IN</sub> =V <sub>OUT_LDO</sub> + 0.2V, 0mA ≤ I <sub>LOAD</sub> ≤ 10mA	—	0.105	0.210	%/mA
V <sub>DROP_LDO</sub>	LDO Dropout Voltage <sup>(2)</sup>	—	LDOVS[1:0]=00B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =10mA, ΔV <sub>OUT_LDO</sub> =2%	—	—	220	mV
		—	LDOVS[1:0]=01B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =10mA, ΔV <sub>OUT_LDO</sub> =2%	—	—	200	mV
		—	LDOVS[1:0]=10B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =10mA, ΔV <sub>OUT_LDO</sub> =2%	—	—	180	mV
		—	LDOVS[1:0]=11B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =10mA, ΔV <sub>OUT_LDO</sub> =2%	—	—	160	mV
TC <sub>LDO</sub>	LDO Temperature Coefficient	—	T <sub>a</sub> =-40°C~85°C, LDOVS[1:0]=00B, V <sub>IN</sub> =3.6V, I <sub>LOAD</sub> =100μA	—	—	200	ppm/°C
ΔV <sub>LINE_LDO</sub>	LDO Line Regulation	—	LDOVS[1:0]=00B, 2.6V ≤ V <sub>IN</sub> ≤ 5.5V, I <sub>LOAD</sub> =100μA	—	—	0.7	%/V
		—	LDOVS[1:0]=00B, 2.6V ≤ V <sub>IN</sub> ≤ 3.6V, I <sub>LOAD</sub> =100μA	—	—	0.2	%/V
<b>ADC &amp; ADC internal reference voltage (Delta Sigma A/D Converter)</b>							
V <sub>OREG</sub>	Supply Voltage for ADC and PGA	—	LDOEN=0	2.4	—	3.3	V
		—	LDOEN=1	2.4	—	3.3	V
I <sub>ADC</sub>	Additional Current for ADC Enable	—	BGPEN=1	—	400	550	μA
I <sub>ADSTB</sub>	Standby Current	—	MCU enters SLEEP mode, no load	—	—	1	μA
N <sub>R</sub>	Resolution	—	—	—	—	24	Bit
INL	Integral Nonlinearity	—	V <sub>OREG</sub> =3.3V, V <sub>REF</sub> =V <sub>OREG</sub> /2, ΔSI=±450mV, PGA gain=1	—	±50	—	ppm
PSRR	Power Supply Rejection Ratio	—	PGA gain=128, ΔV=0.1V	—	146	—	dB
CMRR	Common Mode Rejection Ratio	5V	PGA gain=128, ΔV=0.1V	—	101	—	dB
I <sub>bias</sub>	Input Bias Current	—	—	—	0.8	—	nA
NFB	Noise Free Bits	—	f <sub>MCLK</sub> =4MHz, FLMS[1:0]=00B, V <sub>OREG</sub> =3.3V, V <sub>REF</sub> =V <sub>OREG</sub> /2 PGA gain=128, OSR=16384	—	15.4	—	Bit
ENOB	Effective Number of Bits	—	f <sub>MCLK</sub> =4MHz, FLMS[1:0]=00B, V <sub>OREG</sub> =3.3V, V <sub>REF</sub> =V <sub>OREG</sub> /2 PGA gain=128, OSR=16384	—	18.1	—	Bit
f <sub>ADCK</sub>	ADC Clock Frequency	—	—	40.0	409.6	440.0	kHz
f <sub>ADO</sub>	ADC Output Data Rate	—	f <sub>MCLK</sub> =4MHz, FLMS[1:0]=00B, SINCS=0	4	—	1042	Hz
		—	f <sub>MCLK</sub> =4MHz, FLMS[1:0]=10B, SINCS=0	10	—	2604	Hz

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>REFP</sub>	External Reference Input Voltage	—	—	V <sub>REFN</sub> +0.8	—	V <sub>OREG</sub>	V
V <sub>REFN</sub>		—		0	—	V <sub>REFP</sub> -0.8	V
V <sub>REF</sub>		—		V <sub>REF</sub> =V <sub>REFP</sub> - V <sub>REFN</sub>	0.80	—	1.75
<b>PGA</b>							
V <sub>CM_PGA</sub>	Common Mode Voltage Range	—	—	0.4	—	V <sub>OREG</sub> -0.95	V
ΔD <sub>I</sub>	Differential Input Voltage Range	—	Gain=PGAGN×ADGN	-V <sub>REF</sub> /Gain	—	+V <sub>REF</sub> /Gain	V
<b>Temperature sensor</b>							
TC <sub>TS</sub>	Temperature Sensor Temperature Coefficient	—	T <sub>a</sub> =-40°C~85°C	—	175	—	μV/°C

Note: 1. Load regulation is measured at a constant junction temperature, using pulse testing with a low ON time and is guaranteed up to the maximum power dissipation. Power dissipation is determined by the input/output differential voltage and the output current. Guaranteed maximum power dissipation will not be available over the full input/output range. The maximum allowable power dissipation at any ambient temperature is  $P_D=(T_{J(MAX)}-T_a)/\theta_{JA}$ .

2. Dropout voltage is defined as the input voltage minus the output voltage that produces a 2% change in the output voltage from the value at appointed V<sub>IN</sub>.

**Effective Number of Bits (ENOB)**

V<sub>REF</sub>=1.2V, f<sub>ADCK</sub>=133kHz

Data Rate (SPS)	PGA Gain							
	1	2	4	8	16	32	64	128
4	21.2	21.2	21.2	21.0	20.6	20.1	19.2	18.1
8	20.7	20.7	20.6	20.6	20.2	19.5	18.6	17.7
16	20.2	20.1	20.2	20.0	19.7	19.0	18.2	17.2
33	19.5	19.5	19.5	19.5	19.2	18.6	17.7	16.7
65	19.2	19.2	19.1	18.9	18.6	18.0	17.0	16.1
130	18.4	18.4	18.4	18.2	17.9	17.3	16.5	15.5
260	16.9	16.8	16.8	16.8	16.7	16.5	15.9	15.0
521	15.0	15.0	14.9	14.9	14.9	14.9	14.7	14.3

V<sub>REF</sub>=1.2V, f<sub>ADCK</sub>=333kHz

Data Rate (SPS)	PGA Gain							
	1	2	4	8	16	32	64	128
10	21.2	21.1	21.0	20.8	20.2	19.4	18.6	17.6
20	20.7	20.7	20.6	20.3	19.8	19.0	18.0	17.0
41	20.2	20.1	20.0	19.8	19.3	18.5	17.6	16.6
81	19.6	19.5	19.5	19.2	18.7	17.9	17.1	16.0
163	19.2	19.1	19.0	18.8	18.2	17.4	16.5	15.5
326	18.4	18.4	18.3	18.0	17.7	16.8	16.0	15.0
651	16.8	16.8	16.8	16.8	16.6	16.2	15.5	14.6
1302	15.0	14.9	14.9	14.9	14.9	14.8	14.6	14.1

## LVR Electrical Characteristics

Ta=-40°C~85°C

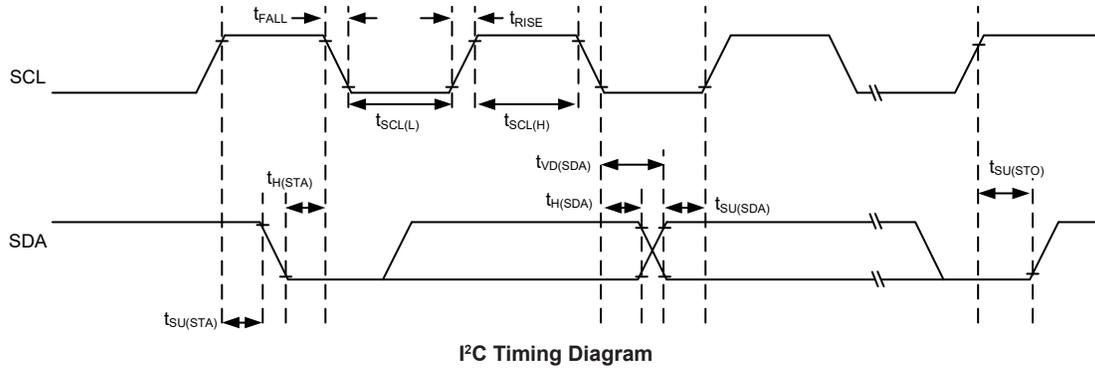
Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>LVR</sub>	Low Voltage Reset Voltage	—	LVR enable, voltage select 2.1V	-5%	2.1	+5%	V
I <sub>LVR</sub>	Operating Current	5V	LVR enable, V <sub>LVR</sub> =2.1V	—	15	25	μA
t <sub>LVR</sub>	Minimum Low Voltage Width to Reset	—	TLVR[1:0]=00B	120	240	480	μs
			TLVR[1:0]=01B	0.5	1.0	2.0	
			TLVR[1:0]=10B	1	2	4	ms
			TLVR[1:0]=11B	2	4	8	

## I<sup>2</sup>C Electrical Characteristics

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
f <sub>I2C</sub>	I <sup>2</sup> C Standard Mode (100kHz) f <sub>sys</sub> Frequency <small>(Note)</small>	—	No clock debounce	2	—	—	MHz
		—	2 system clock debounce	4	—	—	MHz
		—	4 system clock debounce	4	—	—	MHz
	I <sup>2</sup> C Fast Mode (400kHz) f <sub>sys</sub> Frequency <small>(Note)</small>	—	No clock debounce	4	—	—	MHz
		—	2 system clock debounce	8	—	—	MHz
		—	4 system clock debounce	8	—	—	MHz
f <sub>SCL</sub>	SCL Clock Frequency	3V/5V	Standard mode	—	—	100	kHz
			Fast mode	—	—	400	
t <sub>SCL(H)</sub>	SCL Clock High Time	3V/5V	Standard mode	3.5	—	—	μs
			Fast mode	0.9	—	—	
t <sub>SCL(L)</sub>	SCL Clock Low Time	3V/5V	Standard mode	3.5	—	—	μs
			Fast mode	0.9	—	—	
t <sub>FALL</sub>	SCL and SDA Fall Time	3V/5V	Standard mode	—	—	1.3	μs
			Fast mode	—	—	0.34	
t <sub>RISE</sub>	SCL and SDA Rise Time	3V/5V	Standard mode	—	—	1.3	μs
			Fast mode	—	—	0.34	
t <sub>SU(SDA)</sub>	SDA Data Setup Time	3V/5V	Standard mode	0.25	—	—	μs
			Fast mode	0.1	—	—	
t <sub>H(SDA)</sub>	SDA Data Hold Time	3V/5V	—	0.1	—	—	μs
t <sub>VD(SDA)</sub>	SDA Data Valid Time	3V/5V	—	—	—	0.6	μs
t <sub>SU(STA)</sub>	Start Condition Setup Time	3V/5V	Standard mode	3.5	—	—	μs
			Fast mode	0.6	—	—	
t <sub>H(STA)</sub>	Start Condition Hold Time	3V/5V	Standard mode	4.0	—	—	μs
			Fast mode	0.6	—	—	
t <sub>SU(STO)</sub>	Stop Condition Setup Time	3V/5V	Standard mode	3.5	—	—	μs
			Fast mode	0.6	—	—	

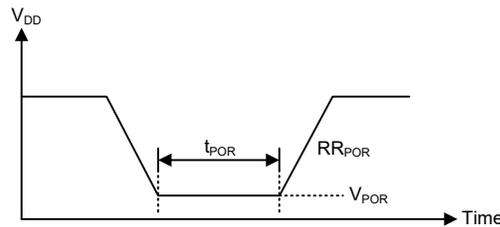
Note: Using the debounce function can make the transmission more stable and reduce the probability of communication failure due to interference.



### Power-on Reset Characteristics

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>POR</sub>	V <sub>DD</sub> Start Voltage to Ensure Power-on Reset	—	—	—	—	100	mV
RR <sub>POR</sub>	V <sub>DD</sub> Rising Rate to Ensure Power-on Reset	—	—	0.035	—	—	V/ms
t <sub>POR</sub>	Minimum Time for V <sub>DD</sub> Stays at V <sub>POR</sub> to Ensure Power-on Reset	—	—	1	—	—	ms

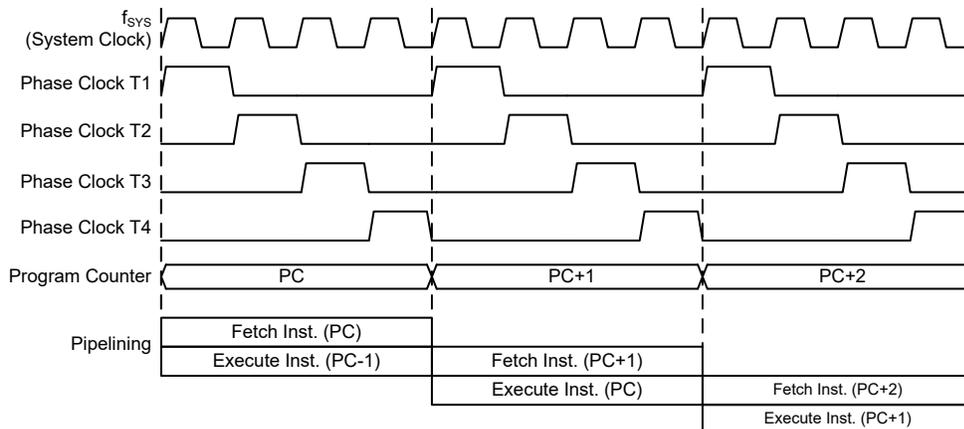


### System Architecture

A key factor in the high-performance features of the Holtek range of microcontrollers is attributed to their internal system architecture. The range of the device take advantage of the usual features found within RISC microcontrollers providing increased speed of operation and enhanced performance. The pipelining scheme is implemented in such a way that instruction fetching and instruction execution are overlapped, hence instructions are effectively executed in one or two cycles for most of the standard or extended instructions respectively, with the exception of branch or call instructions which need one more cycle. An 8-bit wide ALU is used in practically all instruction set operations, which carries out arithmetic operations, logic operations, rotation, increment, decrement, branch decisions, etc. The internal data path is simplified by moving data through the Accumulator and the ALU. Certain internal registers are implemented in the Data Memory and can be directly or indirectly addressed. The simple addressing methods of these registers along with additional architectural features ensure that a minimum of external components is required to provide a functional I/O and A/D control system with maximum reliability and flexibility. This makes the device suitable for affordable, high-volume production for controller applications.

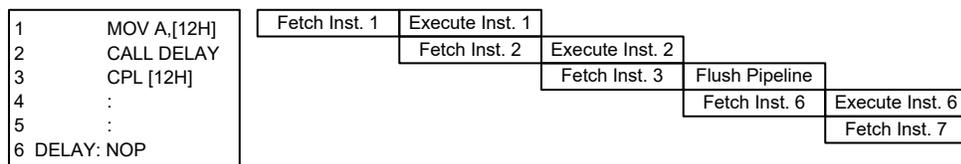
### Clocking and Pipelining

The main system clock, derived from either an HIRC or LIRC oscillator is subdivided into four internally generated non-overlapping clocks, T1~T4. The Program Counter is incremented at the beginning of the T1 clock during which time a new instruction is fetched. The remaining T2~T4 clocks carry out the decoding and execution functions. In this way, one T1~T4 clock cycle forms one instruction cycle. Although the fetching and execution of instructions takes place in consecutive instruction cycles, the pipelining structure of the microcontroller ensures that instructions are effectively executed in one instruction cycle. The exception to this are instructions where the contents of the Program Counter are changed, such as subroutine calls or jumps, in which case the instruction will take one more instruction cycle to execute.



**System Clocking and Pipelining**

For instructions involving branches, such as jump or call instructions, two machine cycles are required to complete instruction execution. An extra cycle is required as the program takes one cycle to first obtain the actual jump or call address and then another cycle to actually execute the branch. The requirement for this extra cycle should be taken into account by programmers in timing sensitive applications.



**Instruction Fetching**

### Program Counter

During program execution, the Program Counter is used to keep track of the address of the next instruction to be executed. It is automatically incremented by one each time an instruction is executed except for instructions, such as “JMP” or “CALL” that demand a jump to a non-consecutive Program Memory address. Only the lower 8 bits, known as the Program Counter Low Register, are directly addressable by the application program.

When executing instructions requiring jumps to non-consecutive addresses such as a jump instruction, a subroutine call, interrupt or reset, etc., the microcontroller manages program control by loading the required address into the Program Counter. For conditional skip instructions, once the condition has been met, the next instruction, which has already been fetched during the present instruction execution, is discarded and a dummy cycle takes its place while the correct instruction is obtained.

Program Counter	
Program Counter High Byte	PCL Register
PC12~PC8	PCL7~PCL0

**Program Counter**

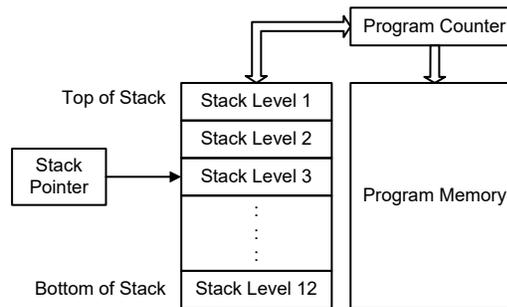
The lower byte of the Program Counter, known as the Program Counter Low register or PCL, is available for program control and is a readable and writeable register. By transferring data directly into this register, a short program jump can be executed directly; however, as only this low byte is available for manipulation, the jumps are limited to the present page of memory that is 256 locations. When such program jumps are executed it should also be noted that a dummy cycle will be inserted. Manipulating the PCL register may cause program branching, so an extra cycle is needed to pre-fetch.

**Stack**

This is a special part of the memory which is used to save the contents of the Program Counter only. The stack is organized into 12 levels and is neither part of the data nor part of the program space, and is neither readable nor writeable. The activated level is indexed by the Stack Pointer, and is neither readable nor writeable. At a subroutine call or interrupt acknowledge signal, the contents of the Program Counter are pushed onto the stack. At the end of a subroutine or an interrupt routine, signaled by a return instruction, RET or RETI, the Program Counter is restored to its previous value from the stack. After a device reset, the Stack Pointer will point to the top of the stack.

If the stack is full and an enabled interrupt takes place, the interrupt request flag will be recorded but the acknowledge signal will be inhibited. When the Stack Pointer is decremented, by RET or RETI, the interrupt will be serviced. This feature prevents stack overflow allowing the programmer to use the structure more easily. However, when the stack is full, a CALL subroutine instruction can still be executed which will result in a stack overflow. Precautions should be taken to avoid such cases which might cause unpredictable program branching.

If the stack is overflow, the first Program Counter save in the stack will be lost.



**Arithmetic and Logic Unit – ALU**

The arithmetic-logic unit or ALU is a critical area of the microcontroller that carries out arithmetic and logic operations of the instruction set. Connected to the main microcontroller data bus, the ALU receives related instruction codes and performs the required arithmetic or logical operations after which the result will be placed in the specified register. As these ALU calculation or operations may result in carry, borrow or other status changes, the status register will be correspondingly updated to reflect these changes. The ALU supports the following functions:

- Arithmetic operations:  
 ADD, ADDM, ADC, ADCM, SUB, SUBM, SBC, SBCM, DAA,

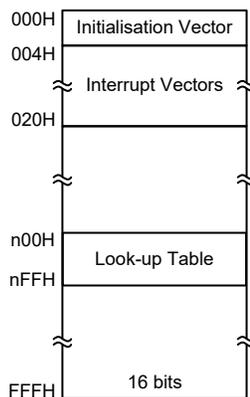
- LADD, LADDM, LADC, LADCM, LSUB, LSUBM, LSBC, LSBCM, LDAA
- Logic operations:
  - AND, OR, XOR, ANDM, ORM, XORM, CPL, CPLA,
  - LAND, LANDM, LOR, LORM, LXOR, LXORM, LCPL, LCPLA
- Rotation:
  - RRA, RR, RRCA, RRC, RLA, RL, RLCA, RLC,
  - LRR, LRRR, LRRCA, LRRC, LRLA, LRL, LRLCA, LRLC
- Increment and Decrement:
  - INCA, INC, DECA, DEC,
  - LINCA, LINC, LDECA, LDEC
- Branch decision:
  - JMP, SZ, SZA, SNZ, SIZ, SDZ, SIZA, SDZA, CALL, RET, RETI,
  - LSNZ, LSZ, LSZA, LSIZ, LSIZA, LSDZ, LSDZA

## OTP Program Memory

The Program Memory is the location where the user code or program is stored. The device is supplied with One-Time Programmable, OTP memory where users can program their application code into the device.

### Structure

The Program Memory has a capacity of 4K×16 bits. The Program Memory is addressed by the Program Counter and also contains data, table information and interrupt entries. Table data, which can be set in any location within the Program Memory, is addressed by a separate table pointer register.



**Program Memory Structure**

### Special Vectors

Within the Program Memory, certain locations are reserved for the reset and interrupts. The location 000H is reserved for use by the device reset for program initialisation. After a device reset is initiated, the program will jump to this location and begin execution.

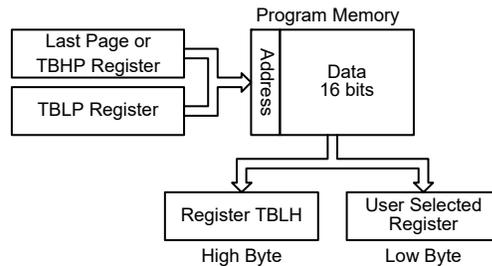
### Look-up Table

Any location within the Program Memory can be defined as a look-up table where programmers can

store fixed data. To use the look-up table, the table pointer must first be setup by placing the address of the look up data to be retrieved in the table pointer registers, TBLP and TBHP. These registers define the total address of the look-up table.

After setting up the table pointer, the table data can be retrieved from the Program Memory using the “TABRD [m]” or “TABRDL [m]” instructions respectively when the memory [m] is located in sector 0. If the memory [m] is located in other sectors except sector 0, the data can be retrieved from the program memory using the corresponding extended table read instruction such as “LTABRD [m]” or “LTABRDL [m]” respectively. When the instruction is executed, the lower order table byte from the Program Memory will be transferred to the user defined Data Memory register [m] as specified in the instruction. The higher order table data byte from the Program Memory will be transferred to the TBLH special register. Any unused bits in this transferred higher order byte will be read as 0.

The accompanying diagram illustrates the addressing data flow of the look-up table.



### Table Program Example

The following example shows how the table pointer and table data is defined and retrieved from the microcontroller. This example uses raw table data located in the Program Memory which is stored there using the ORG statement. The value at this ORG statement is “0F00H” which refers to the start address of the last page within the 4K words Program Memory of the microcontroller. The table pointer low byte register is setup here to have an initial value of “06H”. This will ensure that the first data read from the data table will be at the Program Memory address “0F06H” or 6 locations after the start of the last page. Note that the value for the table pointer is referenced to the specified address pointed by the TBLP and TBHP registers if the “TABRD [m]” or “LTABRD [m]” instruction is being used. The high byte of the table data which in this case is equal to zero will be transferred to the TBLH register automatically when the “TABRD [m]” or “LTABRD [m]” instruction is executed.

Because the TBLH register is a read/write register and can be restored, care should be taken to ensure its protection if both the main routine and Interrupt Service Routine use table read instructions. If using the table read instructions, the Interrupt Service Routines may change the value of the TBLH and subsequently cause errors if used again by the main routine. As a rule it is recommended that simultaneous use of the table read instructions should be avoided. However, in situations where simultaneous use cannot be avoided, the interrupts should be disabled prior to the execution of any main routine table-read instructions. Note that all table related instructions require two instruction cycles to complete their operation.

### Table Read Program Example

```

tempreg1 db? ; temporary register #1
tempreg2 db? ; temporary register #2
:
:
mov a,06h ; initialise low table pointer - note that this address is referenced
mov tblp,a ; to the last page or the page that tbhp pointed
mov a,0Fh ; initialise high table pointer
  
```

```

mov tbhp,a    ; it is not necessary to set tbhp if executing tabrdl
:
:
tabrd tempreg1 ; transfers value in table referenced by table pointer
                ; data at program memory address "0F06H" transferred to tempreg1 and
                ; TBLH
dec tblp      ; reduce value of table pointer by one
tabrd tempreg2 ; transfers value in table referenced by table pointer
                ; data at program memory address "0F05H" transferred to tempreg2 and
                ; TBLH
                ; in this example the data "1AH" is transferred to
                ; tempreg1 and data "0FH" to tempreg2
                ; the value "00H" will be transferred to the high byte register TBLH
:
:
org 0F00h     ; set initial address of last page
dc 00Ah,00Bh,00Ch,00Dh,00Eh,00Fh,01Ah,01Bh

```

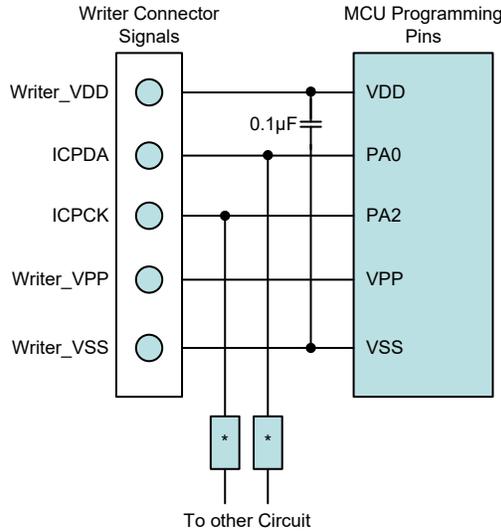
### In Circuit Programming – ICP

The OTP type Program Memory is provided for users to program their application code one time into the device. As an additional convenience, Holtek has provided a means of programming the microcontroller in-circuit using a 5-pin interface. This provides manufacturers with the possibility of manufacturing their circuit boards complete with an un-programmed microcontroller, and then programming the program at a later stage.

Holtek Writer Pins	MCU Programming Pins	Pin Description
ICPDA	PA0	Programming Serial Data/Address
ICPCK	PA2	Programming Clock
VPP	VPP	Programming OTP ROM power supply (8.5V)
VDD	VDD	Power Supply. A 0.1µF capacitor is required to be connected between VDD and VSS for programming.
VSS	VSS	Ground

The Program Memory can be programmed serially in-circuit using this 5-wire interface. Data is downloaded and uploaded serially on a single pin with an additional line for the clock. Three additional lines are required for the power supply. The technical details regarding the in-circuit programming of the device is beyond the scope of this document and will be supplied in supplementary literature.

During the programming process, the user must take care of the ICPDA and ICPCK pins for data and clock programming purposes to ensure that no other outputs are connected to these two pins.



- Note: 1. A 0.1µF capacitor is required to be connected between VDD and VSS for ICP programming, and as close to these pins as possible.
2. \* may be resistor or capacitor. The resistance of \* must be greater than 1kΩ or the capacitance of \* must be less than 1nF.

### On-Chip Debug Support – OCDS

There is an EV chip named BH66V2640 which is used to emulate the BH66R2640 device. This EV chip device also provides an “On-Chip Debug” function to debug the real MCU device during the development process. The EV chip and the real MCU device are almost functionally compatible except for “On-Chip Debug” function. Users can use the EV chip device to emulate the real chip device behavior by connecting the OCSDSA and OCDSCK pins to the Holtek HT-IDE development tools. The OCSDSA pin is the OCDS Data/Address input/output pin while the OCDSCK pin is the OCDS clock input pin. When users use the EV chip for debugging, other functions which are shared with the OCSDSA and OCDSCK pins in the device will have no effect in the EV chip. However, the two OCDS pins which are pin-shared with the ICP programming pins are still used as the Program Memory programming pins for ICP. For more detailed OCDS information, refer to the corresponding document named “Holtek e-Link for 8-bit MCU OCDS User’s Guide”.

Holtek e-Link Pins	EV Chip Pins	Pin Description
OCSDSA	OCSDSA	On-Chip Debug Support Data/Address input/output
OCDSCK	OCDSCK	On-Chip Debug Support Clock input
VDD	VDD	Power Supply
VSS	VSS	Ground

### OTP ROM Parameter Program – ORPP

This device contains an ORPP function. The provision of the ORPP function offers users the convenience of OTP Memory programming features. Note that the Write operation only writes data to the last page of OTP Program Memory, and the data can only be written once and cannot be erased.

Before the write operation is implemented, the VPP pin must be connected to an 8.5V power and after the write operation is completed, the high voltage power should be removed from the VPP pin. If the VPP function is pin-shared with an I/O port, the corresponding I/O port cannot be set as an output when it is used as the VPP function.

## ORPP Registers

Three registers control the overall operation of the internal ORPP function. These are data registers ODL and ODH, and a control register OCR.

Register Name	Bit							
	7	6	5	4	3	2	1	0
OCR	—	—	—	—	WREN	WR	—	—
ODL	D7	D6	D5	D4	D3	D2	D1	D0
ODH	D15	D14	D13	D12	D11	D10	D9	D8

ORPP Register List

### • ODL Register

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: ORPP program memory data bit 7~bit 0

### • ODH Register

Bit	7	6	5	4	3	2	1	0
Name	D15	D14	D13	D12	D11	D10	D9	D8
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D15~D8**: ORPP program memory data bit 15~bit 8

### • OCR Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	WREN	WR	—	—
R/W	—	—	—	—	R/W	R/W	—	—
POR	—	—	—	—	0	0	—	—

Bit 7~4 Unimplemented, read as “0”

Bit 3 **WREN**: ORPP Write Enable  
0: Disable  
1: Enable

This is the ORPP Write Enable Bit which must be set high before write operations are carried out. This bit will be automatically reset to zero by the hardware after the write cycle has finished. Clearing this bit to zero will inhibit ORPP write operations.

Bit 2 **WR**: ORPP Write Control  
0: Write cycle has finished  
1: Activate a write cycle

This is the ORPP Write Control Bit and when set high by the application program will activate a write cycle. This bit will be automatically reset to zero by the hardware after the write cycle has finished. Setting this bit high will have no effect if the WREN has not first been set high.

Bit 1~0 Unimplemented, read as “0”

Note: 1. The WREN and WR cannot be set high at the same time in one instruction.

2. Note that the CPU will be stopped when a write operation is successfully activated.

3. Ensure that the  $f_{SUB}$  clock is stable before executing the write operation.

4. Ensure that the write operation is totally complete before executing other operations.

### **ORPP Writing Data to the OTP Program Memory**

For ORPP write operation the data to be written should be placed in the ODH and ODL registers and the desired write address should first be placed in the TBLP register. To write data to the OTP Program Memory, the write enable bit, WREN, in the OCR register must first be set high to enable the write function. After this, the WR bit in the OCR register must be immediately set high to initiate a write cycle. These two instructions must be executed in two consecutive instruction cycles to activate a write operation successfully. The global interrupt bit EMI should also first be cleared before implementing any write operations, and then set again after a valid write activation procedure has completed. Note that the CPU will be stopped when a write operation is successfully activated. When the write cycle terminates, the CPU will resume executing the application program. And the WR bit will be automatically cleared to zero by the microcontroller, informing the user that the data has been written to the OTP Program Memory.

### **ORPP Reading Data from the OTP Program Memory**

For ORPP read operation the desired address should first be placed in the TBLP register. Then the data can be retrieved from the program memory using the “TABRDL [m]” instruction. When the instruction is executed, the lower order table byte from the Program Memory will be transferred to the user defined Data Memory register [m] as specified in the instruction. The higher order table data byte from the Program Memory will be transferred to the TBLH special register.

### **Programming Considerations**

Care must be taken that data is not inadvertently written to the OTP Program Memory. Protection can be enhanced by ensuring that the Write Enable bit is normally cleared to zero when not writing. Although certainly not necessary, consideration might be given in the application program to the checking of the validity of new write data by a simple read back process.

When writing data the WR bit must be set high immediately after the WREN bit has been set high, to ensure the write cycle executes correctly. The global interrupt bit EMI should also be cleared before a write cycle is executed and then set high again after a write activation procedure has completed. Note that the device should not enter the IDLE or SLEEP mode until the ORPP write operation is totally complete. Otherwise, the ORPP write operation will fail.

### **Programming Examples**

#### **ORPP Reading Data from the OTP Program Memory**

```
Tempreg1 db?      ; temporary register
MOV A, 03H
MOV TBLP, A       ; set read address 03H
TABRDL Tempreg1   ; transfers value in table (last page) referenced by table
                  ; pointer, data at program memory address "0F03H" transferred
                  ; to tempreg1 and TBLH
```

#### **ORPP Writing Data to the OTP Program Memory**

```
MOV A, ORPP_ADRES ; user defined address
MOV TBLP, A
MOV A, ORPP_DATA_L ; user defined data
MOV ODL, A
MOV A, ORPP_DATA_H
MOV ODH, A
MOV A, 00H
MOV OCR, A
CLR EMI
```

```

SET WREN          ; set WREN bit, enable write operation
SET WR           ; start Write Cycle - set WR bit - executed immediately
                 ; after setting WREN bit

SET EMI
BACK:
SZ WR            ; check for write cycle end
JMP BACK
NOP

```

## Data Memory

The Data Memory is a volatile area of 8-bit wide RAM internal memory and is the location where temporary information is stored.

Categorized into two types, the first of these is an area of RAM where special function registers are located. These registers have fixed locations and are necessary for correct operation of the device. Many of these registers can be read from and written to directly under program control, however, some remain protected from user manipulation. The second area of Data Memory is reserved for general purpose use. All locations within this area are read and write accessible under program control.

There is another area of the Data Memory reserved for the Body Fat Sine Wave Pattern. The addresses of the Sine Wave Pattern Memory area overlap those in the Special Purpose Data Memory area.

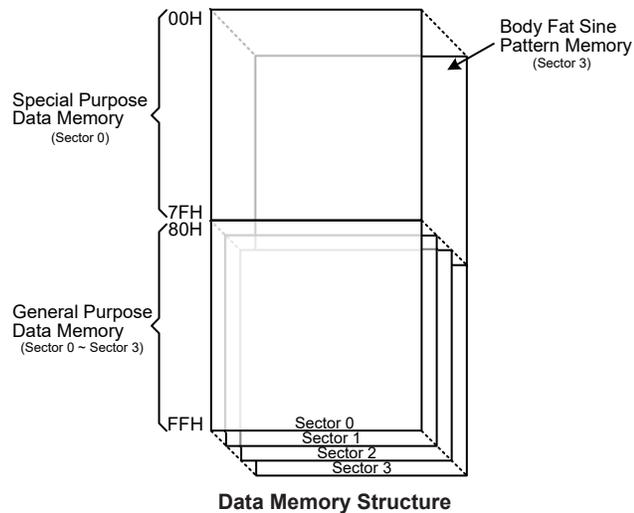
### Structure

The Data Memory is subdivided into several sectors, all of which are implemented in 8-bit wide RAM. Each of the Data Memory Sector is categorised into two types, the special Purpose Data Memory and the General Purpose Data Memory. The address range of the Special Purpose Data Memory for the device is from 00H to 7FH while the General Purpose Data Memory address range is from 80H to FFH. Note that the 00H~7FH of Sector 3 is Sine Pattern Memory.

Switching between the different Data Memory sectors is achieved by properly setting the Memory Pointers to correct value if using the indirect addressing method.

Special Purpose Data Memory	General Purpose Data Memory		Body Fat Sine Pattern Data Memory	
Located Sectors	Capacity	Sector: Address	Capacity	Sector: Address
0	512×8	0: 80H~FFH 1: 80H~FFH 2: 80H~FFH 3: 80H~FFH	128×8	3: 00H~7FH

**Data Memory Summary**



### Data Memory Addressing

For the device that supports the extended instructions, there is no Bank Pointer for Data Memory. For Data Memory the desired Sector is pointed by the MP1H or MP2H register and the certain Data Memory address in the selected sector is specified by the MP1L or MP2L register when using indirect addressing access.

Direct Addressing can be used in all sectors using the extended instructions which can address all available data memory space. For the accessed data memory which is located in any data memory sectors except sector 0, the extended instructions can be used to access the data memory instead of using the indirect addressing access. The main difference between standard instructions and extended instructions is that the data memory address “m” in the extended instructions has 10 valid bits for this device, the high byte indicates a sector and the low byte indicates a specific address.

### General Purpose Data Memory

All microcontroller programs require an area of read/write memory where temporary data can be stored and retrieved for use later. It is this area of RAM memory that is known as General Purpose Data Memory. This area of Data Memory is fully accessible by the user programming for both reading and writing operations. By using the bit operation instructions individual bits can be set or reset under program control giving the user a large range of flexibility for bit manipulation in the Data Memory.

### Special Purpose Data Memory

This area of Data Memory is where registers, necessary for the correct operation of the microcontroller, are stored. Most of the registers are both readable and writeable but some are protected and are readable only, the details of which are located under the relevant Special Function Register section. Note that for locations that are unused, any read instruction to these addresses will return the value “00H”.

Sector 0		Sector 0	
00H	IAR0	40H	
01H	MP0	41H	
02H	IAR1	42H	
03H	MP1L	43H	
04H	MP1H	44H	
05H	ACC	45H	
06H	PCL	46H	
07H	TBLP	47H	
08H	TBLH	48H	
09H	TBHP	49H	
0AH	STATUS	4AH	
0BH		4BH	CTMC0
0CH	IAR2	4CH	CTMC1
0DH	MP2L	4DH	CTMDL
0EH	MP2H	4EH	CTMDH
0FH	RSTFC	4FH	CTMAL
10H	SCC	50H	CTMAH
11H	HIRCC	51H	
12H		52H	
13H		53H	
14H	PA	54H	
15H	PAC	55H	
16H	PAPU	56H	
17H	PAWU	57H	SLEDC
18H	RSTC	58H	
19H		59H	
1AH		5AH	PAS0
1BH		5BH	PAS1
1CH	MFI	5CH	PBS0
1DH		5DH	PBS1
1EH	WDC	5EH	
1FH	INTEG	5FH	
20H	INTC0	60H	OCR
21H	INTC1	61H	ODL
22H	INTC2	62H	ODH
23H	TLVRC	63H	
24H	PB	64H	PWRC
25H	PBC	65H	PGAC0
26H	PBPU	66H	PGAC1
27H		67H	PGACS
28H		68H	ADRL
29H		69H	ADRM
2AH		6AH	ADRH
2BH	PSC0R	6BH	ADCR0
2CH	PSC1R	6CH	ADCR1
2DH	TB0C	6DH	ADCS
2EH	TB1C	6EH	
2FH		6FH	
30H	USR	70H	
31H	UCR1	71H	
32H	UCR2	72H	
33H	UCR3	73H	SGC
34H	BRDH	74H	SGN
35H	BRDL	75H	SGDN
36H	UFCR	76H	OPAC
37H	TXR_RXR	77H	SWC0
38H	RxCNT	78H	SWC1
39H	SIMC0	79H	SWC2
3AH	SIMC1	7AH	SWC3
3BH	SIMD	7BH	CMPC0
3CH	SIMA/SIMC2	7CH	CMPC1
3DH	SIMTOC	7DH	BDA0C
3EH		7EH	BDA1C
3FH		7FH	BDA2C

: Unused, read as 00H      : Reserved, cannot be changed unless otherwise specified

**Special Purpose Data Memory**

## Special Function Register Description

Most of the Special Function Register details will be described in the relevant functional sections, however several registers require a separate description in this section.

### Indirect Addressing Registers – IAR0, IAR1, IAR2

The Indirect Addressing Registers, IAR0, IAR1 and IAR2, although having their locations in normal RAM register space, do not actually physically exist as normal registers. The method of indirect addressing for RAM data manipulation uses these Indirect Addressing Registers and Memory Pointers, in contrast to direct memory addressing, where the actual memory address is specified. Actions on the IAR0, IAR1 and IAR2 registers will result in no actual read or write operation to these registers but rather to the memory location specified by their corresponding Memory Pointers, MP0, MP1L/MP1H or MP2L/MP2H. Acting as a pair, IAR0 and MP0 can together access data only from Sector 0 while the IAR1 register together with the MP1L/MP1H register pair and IAR2 register together with the MP2L/MP2H register pair can access data from any Data Memory Sector. As the Indirect Addressing Registers are not physically implemented, reading the Indirect Addressing Registers will return a result of “00H” and writing to the registers will result in no operation.

### Memory Pointers – MP0, MP1L/MP1H, MP2L/MP2H

Five Memory Pointers, known as MP0, MP1L, MP1H, MP2L, MP2H, are provided. These Memory Pointers are physically implemented in the Data Memory and can be manipulated in the same way as normal registers providing a convenient way with which to address and track data. When any operation to the relevant Indirect Addressing Registers is carried out, the actual address that the microcontroller is directed to is the address specified by the related Memory Pointer. MP0, together with Indirect Addressing Register, IAR0, are used to access data from Sector 0, while MP1L/MP1H together with IAR1 and MP2L/MP2H together with IAR2 are used to access data from all sectors according to the corresponding MP1H or MP2H register. Direct Addressing can be used in all sectors using the extended instruction which can address all available data memory space.

The following example shows how to clear a section of four Data Memory locations already defined as locations adres1 to adres4.

### Indirect Addressing Program Example

#### Example 1

```
data .section 'data'
adres1 db ?
adres2 db ?
adres3 db ?
adres4 db ?
block db ?
code .section at 0 'code'
org 00h
start:
    mov a, 04h                ; setup size of block
    mov block, a
    mov a, offset adres1     ; Accumulator loaded with first RAM address
    mov mp0, a               ; setup memory pointer with first RAM address
loop:
    clr IAR0                 ; clear the data at address defined by MP0
    inc mp0                  ; increment memory pointer
    sdz block                 ; check if last memory location has been cleared
    jmp loop
continue:
```

### Example 2

```

data .section 'data'
adres1 db ?
adres2 db ?
adres3 db ?
adres4 db ?
block db ?
code .section at 0 'code'
org 00h
start:
    mov a, 04h          ; setup size of block
    mov block, a
    mov a, 01h          ; setup the memory sector
    mov mplh, a
    mov a, offset adres1 ; Accumulator loaded with first RAM address
    mov mpll, a         ; setup memory pointer with first RAM address
loop:
    clr IAR1           ; clear the data at address defined by MP1L
    inc mpll           ; increment memory pointer MP1L
    sdz block          ; check if last memory location has been cleared
    jmp loop
continue:

```

The important point to note here is that in the example shown above, no reference is made to specific Data Memory addresses.

### Direct Addressing Program Example using extended instructions

```

data .section 'data'
temp db ?
code .section at 0 'code'
org 00h
start:
    lmov a, [m]         ; move [m] data to acc
    lsub a, [m+1]       ; compare [m] and [m+1] data
    snz c               ; [m]>[m+1]?
    jmp continue       ; no
    lmov a, [m]         ; yes, exchange [m] and [m+1] data
    mov temp, a
    lmov a, [m+1]
    lmov [m], a
    mov a, temp
    lmov [m+1], a
continue:

```

Note: here “m” is a data memory address located in any data memory sectors. For example, m=1F0H, it indicates address 0F0H in Sector 1.

### Accumulator – ACC

The Accumulator is central to the operation of any microcontroller and is closely related with operations carried out by the ALU. The Accumulator is the place where all intermediate results from the ALU are stored. Without the Accumulator it would be necessary to write the result of each calculation or logical operation such as addition, subtraction, shift, etc., to the Data Memory resulting in higher programming and timing overheads. Data transfer operations usually involve the temporary storage function of the Accumulator; for example, when transferring data between one user-defined register and another, it is necessary to do this by passing the data through the Accumulator as no direct transfer between two registers is permitted.

## **Program Counter Low Register – PCL**

To provide additional program control functions, the low byte of the Program Counter is made accessible to programmers by locating it within the Special Purpose area of the Data Memory. By manipulating this register, direct jumps to other program locations are easily implemented. Loading a value directly into this PCL register will cause a jump to the specified Program Memory location, however, as the register is only 8-bit wide, only jumps within the current Program Memory page are permitted. When such operations are used, note that a dummy cycle will be inserted.

## **Look-up Table Registers – TBLP, TBHP, TBLH**

These three special function registers are used to control operation of the look-up table which is stored in the Program Memory. TBLP and TBHP are the table pointers and indicate the location where the table data is located. Their value must be setup before any table read commands are executed. Their value can be changed, for example using the “INC” or “DEC” instructions, allowing for easy table data pointing and reading. TBLH is the location where the high order byte of the table data is stored after a table read data instruction has been executed. Note that the lower order table data byte is transferred to a user defined location.

## **Status Register – STATUS**

This 8-bit register contains the SC flag, CZ flag, zero flag (Z), carry flag (C), auxiliary carry flag (AC), overflow flag (OV), power down flag (PDF), and watchdog time-out flag (TO). These arithmetic/logical operation and system management flags are used to record the status and operation of the microcontroller.

With the exception of the TO and PDF flags, bits in the status register can be altered by instructions like most other registers. Any data written into the status register will not change the TO or PDF flag. In addition, operations related to the status register may give different results due to the different instruction operations. The TO flag can be affected only by a system power-up, a WDT time-out or by executing the “CLR WDT” or “HALT” instruction. The PDF flag is affected only by executing the “HALT” or “CLR WDT” instruction or during a system power-up.

The Z, OV, AC, C, SC and CZ flags generally reflect the status of the latest operations.

- C is set if an operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation; otherwise C is cleared. C is also affected by a rotate through carry instruction.
- AC is set if an operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction; otherwise AC is cleared.
- Z is set if the result of an arithmetic or logical operation is zero; otherwise Z is cleared.
- OV is set if an operation results in a carry into the highest-order bit but not a carry out of the highest-order bit, or vice versa; otherwise OV is cleared.
- PDF is cleared by a system power-up or executing the “CLR WDT” instruction. PDF is set by executing the “HALT” instruction.
- TO is cleared by a system power-up or executing the “CLR WDT” or “HALT” instruction. TO is set by a WDT time-out.
- CZ is the operational result of different flags for different instructions. Refer to register definitions for more details.
- SC is the result of the “XOR” operation which is performed by the OV flag and the MSB of the current instruction operation result.

In addition, on entering an interrupt sequence or executing a subroutine call, the status register will

not be pushed onto the stack automatically. If the contents of the status registers are important and if the subroutine can corrupt the status register, precautions must be taken to correctly save it.

• **STATUS Register**

Bit	7	6	5	4	3	2	1	0
Name	SC	CZ	TO	PDF	OV	Z	AC	C
R/W	R/W	R/W	R	R	R/W	R/W	R/W	R/W
POR	x	x	0	0	x	x	x	x

"x": unknown

- Bit 7      **SC**: The result of the "XOR" operation which is performed by the OV flag and the MSB of the instruction operation result.
- Bit 6      **CZ**: The operational result of different flags for different instructions.  
For SUB/SUBM/LSUB/LSUBM instructions, the CZ flag is equal to the Z flag.  
For SBC/SBCM/LSBC/LSBCM instructions, the CZ flag is the "AND" operation result which is performed by the previous operation CZ flag and current operation zero flag.  
For other instructions, the CZ flag will not be affected.
- Bit 5      **TO**: Watchdog Time-out flag  
0: After power up or executing the "CLR WDT" or "HALT" instruction  
1: A watchdog time-out occurred.
- Bit 4      **PDF**: Power down flag  
0: After power up or executing the "CLR WDT" instruction  
1: By executing the "HALT" instruction
- Bit 3      **OV**: Overflow flag  
0: No overflow  
1: An operation results in a carry into the highest-order bit but not a carry out of the highest-order bit or vice versa.
- Bit 2      **Z**: Zero flag  
0: The result of an arithmetic or logical operation is not zero  
1: The result of an arithmetic or logical operation is zero
- Bit 1      **AC**: Auxiliary flag  
0: No auxiliary carry  
1: An operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction
- Bit 0      **C**: Carry flag  
0: No carry-out  
1: An operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation  
The "C" flag is also affected by a rotate through carry instruction.

## Oscillators

Various oscillator options offer the user a wide range of functions according to their various application requirements. The flexible features of the oscillator functions ensure that the best optimisation can be achieved in terms of speed and power saving. Oscillator selections and operation are selected through the application program by using some control registers.

### Oscillator Overview

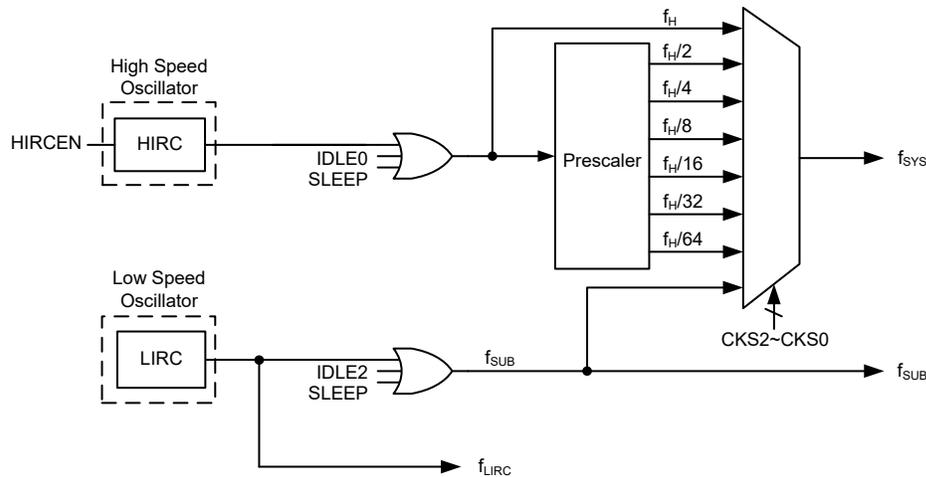
In addition to being the source of the main system clock the oscillators also provide clock sources for the Watchdog Timer and Time Base Interrupts. Two fully integrated internal oscillators, requiring no external components, are provided to form a wide range of both fast and slow system oscillators. The higher frequency oscillator provides higher performance but carry with it the disadvantage of higher power requirements, while the opposite is of course true for the lower frequency oscillator. With the capability of dynamically switching between fast and slow system clock, the device have the flexibility to optimize the performance/power ratio, a feature especially important in power sensitive portable applications.

Type	Name	Frequency
Internal High Speed RC	HIRC	8MHz
Internal Low Speed RC	LIRC	32kHz

Oscillator Types

### System Clock Configurations

There are two methods of generating the system clock, one high speed oscillator and one low speed oscillator. The high speed oscillator is the internal 8MHz RC oscillator, HIRC. The low speed oscillators are the internal 32kHz RC oscillator, LIRC. Selecting whether the low or high speed oscillator is used as the system oscillator is implemented using the CKS2~CKS0 bits in the SCC register and as the system clock can be dynamically selected.



System Clock Configurations

### Internal High Speed RC Oscillator – HIRC

The internal RC oscillator is a fully integrated system oscillator requiring no external components. The internal RC oscillator has a fixed frequency of 8MHz. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised.

### **Internal 32kHz Oscillator – LIRC**

The Internal 32kHz System Oscillator is the low frequency oscillator. It is a fully integrated RC oscillator with a typical frequency of 32kHz, requiring no external components for its implementation.

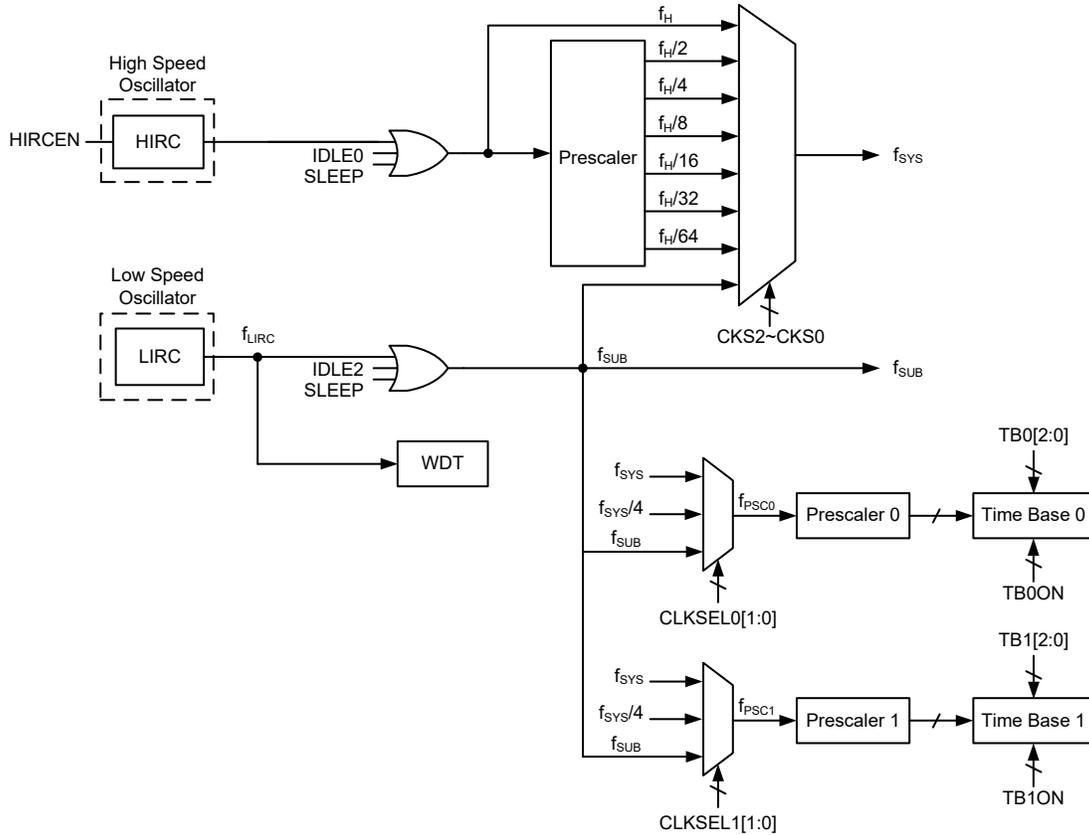
## **Operating Modes and System Clocks**

Present day applications require that their microcontrollers have high performance but often still demand that they consume as little power as possible, conflicting requirements that are especially true in battery powered portable applications. The fast clocks required for high performance will by their nature increase current consumption and of course vice versa, lower speed clocks reduce current consumption. As Holtek has provided the device with both high and low speed clock sources and the means to switch between them dynamically, the user can optimise the operation of their microcontroller to achieve the best performance/power ratio.

### **System Clocks**

The device has many different clock sources for both the CPU and peripheral function operation. By providing the user with a wide range of clock options using register programming, a clock system can be configured to obtain maximum application performance.

The main system clock, can come from either a high frequency,  $f_H$ , or low frequency,  $f_{SUB}$ , source, and is selected using the CKS2~CKS0 bits in the SCC register. The high speed system clock is sourced from the HIRC oscillator. The low speed system clock source is sourced from the LIRC oscillator. The other choice, which is a divided version of the high speed system oscillator has a range of  $f_H/2 \sim f_H/64$ .



### Device Clock Configurations

Note: When the system clock source  $f_{SYS}$  is switched to  $f_{SUB}$  from  $f_H$ , the high speed oscillator will stop to conserve the power or continue to oscillate to provide the clock source,  $f_H \sim f_H/64$ , for peripheral circuit to use, which is determined by configuring the corresponding high speed oscillator enable control bit.

### System Operation Modes

There are six different modes of operation for the microcontroller, each one with its own special characteristics and which can be chosen according to the specific performance and power requirements of the application. There are two modes allowing normal operation of the microcontroller, the FAST Mode and SLOW Mode. The remaining four modes, the SLEEP, IDLE0, IDLE1 and IDLE2 Mode are used when the microcontroller CPU is switched off to conserve power.

Operation Mode	CPU	Register Setting			$f_{SYS}$	$f_H$	$f_{SUB}$	$f_{LIRC}$
		FHIDEN	FSIDEN	CKS2~CKS0				
FAST	On	x	x	000~110	$f_H \sim f_H/64$	On	On	On
SLOW	On	x	x	111	$f_{SUB}$	On/Off <sup>(1)</sup>	On	On
IDLE0	Off	0	1	000~110	Off	Off	On	On
				111	On			
IDLE1	Off	1	1	xxx	On	On	On	On
IDLE2	Off	1	0	000~110	On	On	Off	On
				111	Off			
SLEEP	Off	0	0	xxx	Off	Off	Off	On/Off <sup>(2)</sup>

"x": Don't care

Note: 1. The  $f_H$  clock will be switched on or off by configuring the corresponding oscillator enable bit in the SLOW mode.

- The  $f_{LIRC}$  clock can be switched on or off which is controlled by the WDT function being enabled or disabled in the SLEEP mode.

### FAST Mode

This is one of the main operating modes where the microcontroller has all of its functions operational and where the system clock is provided by one of the high speed oscillators. This mode operates allowing the microcontroller to operate normally with a clock source will come from HIRC oscillator. The high speed oscillator will however first be divided by a ratio ranging from 1 to 64, the actual ratio being selected by the CKS2~CKS0 bits in the SCC register. Although a high speed oscillator is used, running the microcontroller at a divided clock ratio reduces the operating current.

### SLOW Mode

This is also a mode where the microcontroller operates normally although now with a slower speed clock source. The clock source used will be from  $f_{SUB}$ .

### SLEEP Mode

The SLEEP Mode is entered when a HALT instruction is executed and when the FHIDEN and FSIDEN bit are low. In the SLEEP mode the CPU will be stopped. The  $f_{SUB}$  clock provided to the peripheral function will also be stopped, too. However the  $f_{LIRC}$  clock can still continue to operate if the WDT function is enabled by the WDTC register.

### IDLE0 Mode

The IDLE0 Mode is entered when a HALT instruction is executed and when the FHIDEN bit in the SCC register is low and the FSIDEN bit in the SCC register is high. In the IDLE0 Mode the CPU will be switched off but the low speed oscillator will be turned on to drive some peripheral functions.

### IDLE1 Mode

The IDLE1 Mode is entered when a HALT instruction is executed and when the FHIDEN bit in the SCC register is high and the FSIDEN bit in the SCC register is high. In the IDLE1 Mode the CPU will be switched off but both the high and low speed oscillators will be turned on to provide a clock source to keep some peripheral functions operational.

### IDLE2 Mode

The IDLE2 Mode is entered when a HALT instruction is executed and when the FHIDEN bit in the SCC register is high and the FSIDEN bit in the SCC register is low. In the IDLE2 Mode the CPU will be switched off but the high speed oscillator will be turned on to provide a clock source to keep some peripheral functions operational.

## Control Registers

The registers, SCC and HIRCC, are used to control the system clock and the corresponding oscillator configurations.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SCC	CKS2	CKS1	CKS0	—	—	—	FHIDEN	FSIDEN
HIRCC	—	—	—	—	—	—	HIRCF	HIRCEN

**System Operating Mode Control Register List**

• **SCC Register**

Bit	7	6	5	4	3	2	1	0
Name	CKS2	CKS1	CKS0	—	—	—	FHIDEN	FSIDEN
R/W	R/W	R/W	R/W	—	—	—	R/W	R/W
POR	0	0	0	—	—	—	0	0

Bit 7~5 **CKS2~CKS0**: System clock selection

- 000:  $f_H$
- 001:  $f_H/2$
- 010:  $f_H/4$
- 011:  $f_H/8$
- 100:  $f_H/16$
- 101:  $f_H/32$
- 110:  $f_H/64$
- 111:  $f_{SUB}$

These three bits are used to select which clock is used as the system clock source. In addition to the system clock source directly derived from  $f_H$  or  $f_{SUB}$ , a divided version of the high speed system oscillator can also be chosen as the system clock source.

Bit 4~2 Unimplemented, read as “0”

Bit 1 **FHIDEN**: High Frequency oscillator control when CPU is switched off

- 0: Disable
- 1: Enable

This bit is used to control whether the high speed oscillator is activated or stopped when the CPU is switched off by executing a “HALT” instruction.

Bit 0 **FSIDEN**: Low Frequency oscillator control when CPU is switched off

- 0: Disable
- 1: Enable

This bit is used to control whether the low speed oscillator is activated or stopped when the CPU is switched off by executing a “HALT” instruction.

Note: A certain delay is required before the relevant clock is successfully switched to the target clock source after any clock switching setup using the CKS2~CKS0 bits. A proper delay time must be arranged before executing the following operations which require immediate reaction with the target clock source.

Clock switching delay time =  $4 \times t_{SYS} + [0 \sim (1.5 \times t_{curr} + 0.5 \times t_{Tar})]$ , Where  $t_{curr}$  indicates the current clock period,  $t_{Tar}$  indicates the target clock period and the  $t_{SYS}$  indicates the current system clock period.

• **HIRCC Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	HIRCF	HIRCEN
R/W	—	—	—	—	—	—	R	R/W
POR	—	—	—	—	—	—	0	1

Bit 7~2 Unimplemented, read as “0”

Bit 1 **HIRCF**: HIRC oscillator stable flag

- 0: HIRC unstable
- 1: HIRC stable

This bit is used to indicate whether the HIRC oscillator is stable or not. When the HIRCEN bit is set high to enable the HIRC oscillator, the HIRCF bit will first be cleared to zero and then set high after the HIRC oscillator is stable.

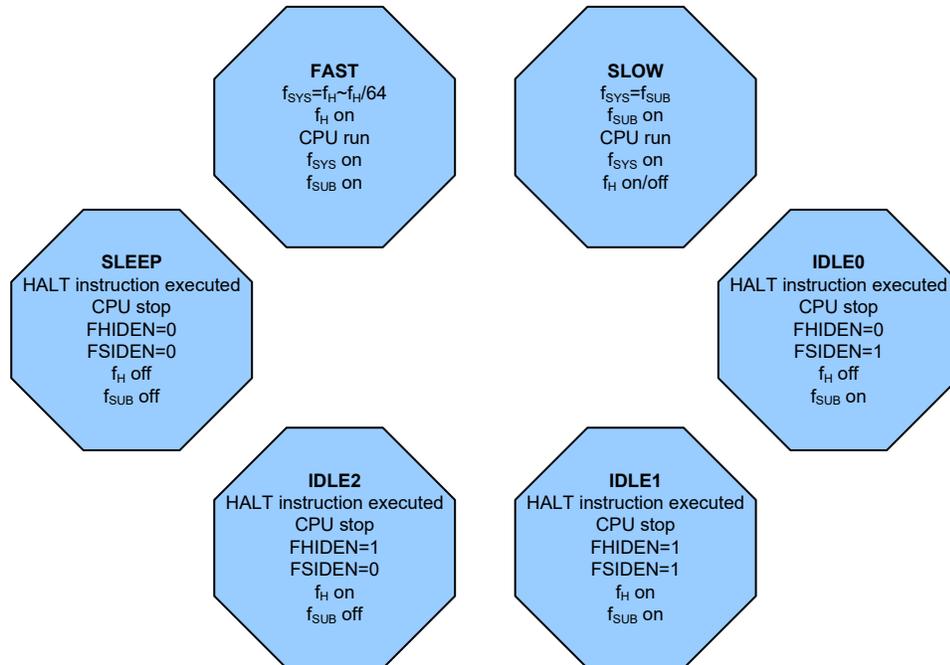
Bit 0 **HIRCEN**: HIRC oscillator enable control

- 0: Disable
- 1: Enable

### Operating Mode Switching

The device can switch between operating modes dynamically allowing the user to select the best performance/power ratio for the present task in hand. In this way microcontroller operations that do not require high performance can be executed using slower clocks thus requiring less operating current and prolonging battery life in portable applications.

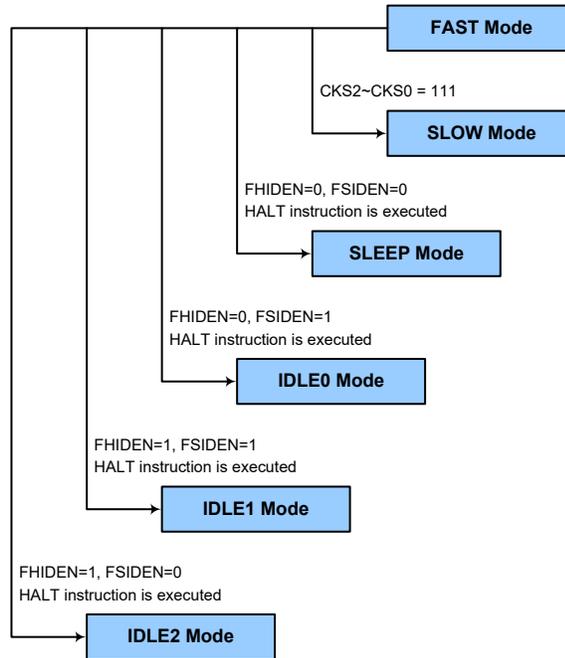
In simple terms, Mode Switching between the FAST Mode and SLOW Mode is executed using the CKS2~CKS0 bits in the SCC register while Mode Switching from the FAST/SLOW Modes to the SLEEP/IDLE Modes is executed via the HALT instruction. When a HALT instruction is executed, whether the device enters the IDLE Mode or the SLEEP Mode is determined by the condition of the FHIDEN and FSIDEN bits in the SCC register.



### FAST Mode to SLOW Mode Switching

When running in the FAST Mode, which uses the high speed system oscillator, and therefore consumes more power, the system clock can switch to run in the SLOW Mode by setting the CKS2~CKS0 bits to “111” in the SCC register. This will then use the low speed system oscillator which will consume less power. Users may decide to do this for certain operations which do not require high performance and can subsequently reduce power consumption.

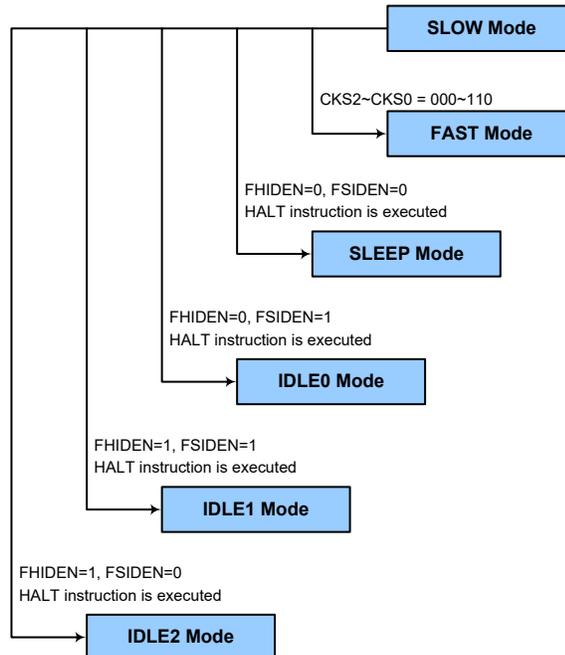
The SLOW Mode is sourced from the LIRC oscillator and therefore requires this oscillator to be stable before full mode switching occurs.



### SLOW Mode to FAST Mode Switching

In SLOW mode the system clock is derived from  $f_{SUB}$ . When system clock is switched back to the FAST mode from  $f_{SUB}$ , the  $CKS2-CKS0$  bits should be set to “000”~“110” and then the system clock will respectively be switched to  $f_H \sim f_H/64$ .

However, if  $f_H$  is not used in SLOW mode and thus switched off, it will take some time to re-oscillate and stabilise when switching to the FAST mode from the SLOW Mode. This is monitored using the HIRCF bit in the HIRCC register. The time duration required for the high speed system oscillator stabilization is specified in the System Start Up Time Characteristics.



### Entering the SLEEP Mode

There is only one way for the device to enter the SLEEP Mode and that is to execute the “HALT” instruction in the application program with both the FHIDEN and FSIDEN bits in the SCC register equal to “0”. In this mode all the clocks and functions will be switched off except the WDT function. When this instruction is executed under the conditions described above, the following will occur:

- The system clock will be stopped and the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.
- The WDT will be cleared and resume counting if the WDT function is enabled. If the WDT function is disabled, the WDT will be cleared and then stopped.

### Entering the IDLE0 Mode

There is only one way for the device to enter the IDLE0 Mode and that is to execute the “HALT” instruction in the application program with the FHIDEN bit in the SCC register equal to “0” and the FSIDEN bit in the SCC register equal to “1”. When this instruction is executed under the conditions described above, the following will occur:

- The  $f_H$  clock will be stopped and the application program will stop at the “HALT” instruction, but the  $f_{SUB}$  clock will be on.
- The Data Memory contents and registers will maintain their present condition.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.
- The WDT will be cleared and resume counting if the WDT function is enabled. If the WDT function is disabled, the WDT will be cleared and then stopped.

### Entering the IDLE1 Mode

There is only one way for the device to enter the IDLE1 Mode and that is to execute the “HALT” instruction in the application program with both the FHIDEN and FSIDEN bits in the SCC register equal to “1”. When this instruction is executed under the conditions described above, the following will occur:

- The  $f_H$  and  $f_{SUB}$  clocks will be on but the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.
- The WDT will be cleared and resume counting if the WDT function is enabled. If the WDT function is disabled, the WDT will be cleared and then stopped.

### Entering the IDLE2 Mode

There is only one way for the device to enter the IDLE2 Mode and that is to execute the “HALT” instruction in the application program with the FHIDEN bit in the SCC register equal to “1” and the FSIDEN bit in the SCC register equal to “0”. When this instruction is executed under the conditions described above, the following will occur:

- The  $f_H$  clock will be on but the  $f_{SUB}$  clock will be off and the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.
- The WDT will be cleared and resume counting if the WDT function is enabled. If the WDT function is disabled, the WDT will be cleared and then stopped.

### **Standby Current Considerations**

As the main reason for entering the SLEEP or IDLE Mode is to keep the current consumption of the device to as low a value as possible, perhaps only in the order of several micro-amps except in the IDLE1 and IDLE2 Mode, there are other considerations which must also be taken into account by the circuit designer if the power consumption is to be minimised. Special attention must be made to the I/O pins on the device. All high-impedance input pins must be connected to either a fixed high or low level as any floating input pins could create internal oscillations and result in increased current consumption. This also applies to the device which has different package types, as there may be unbonded pins. These must either be setup as outputs or if setup as inputs must have pull-high resistors connected. In addition, the I/O pin-shared with VPP must not be set to output high, as this could result in increased current consumption.

Care must also be taken with the loads, which are connected to I/O pins, which are setup as outputs. These should be placed in a condition in which minimum current is drawn or connected only to external circuits that do not draw current, such as other CMOS inputs. Also note that additional standby current will also be required if the LIRC oscillator has enabled.

In the IDLE1 and IDLE2 Mode the high speed oscillator is on, if the peripheral function clock source is derived from the high speed oscillator, the additional standby current will also be perhaps in the order of several hundred micro-amps.

### **Wake-up**

To minimise power consumption the device can enter the SLEEP or any IDLE Mode, where the CPU will be switched off. However, when the device is woken up again, it will take a considerable time for the original system oscillator to restart, stabilise and allow normal operation to resume.

After the system enters the SLEEP or IDLE Mode, it can be woken up from one of various sources listed as follows:

- An external falling edge on Port A
- A system interrupt
- A WDT overflow

When the device executes the “HALT” instruction, it will enter the SLEEP or IDLE mode and the PDF flag will be set to 1. The PDF flag will be cleared to 0 if the device experiences a system power-up or executes the clear Watchdog Timer instruction. If the system is woken up by a WDT overflow, a Watchdog Timer reset will be initiated and the TO flag will be set to 1. The TO flag is set if a WDT time-out occurs and causes a wake-up that only resets the Program Counter and Stack Pointer, other flags remain in their original status.

Each pin on Port A can be setup using the PAWU register to permit a negative transition on the pin to wake-up the system. When a pin wake-up occurs, the program will resume execution at the instruction following the “HALT” instruction. If the system is woken up by an interrupt, then two

possible situations may occur. The first is where the related interrupt is disabled or the interrupt is enabled but the stack is full, in which case the program will resume execution at the instruction following the “HALT” instruction. In this situation, the interrupt which woke-up the device will not be immediately serviced, but will rather be serviced later when the related interrupt is finally enabled or when a stack level becomes free. The other situation is where the related interrupt is enabled and the stack is not full, in which case the regular interrupt response takes place. If an interrupt request flag is set high before entering the SLEEP or IDLE Mode, the wake-up function of the related interrupt will be disabled.

## Watchdog Timer

The Watchdog Timer is provided to prevent program malfunctions or sequences from jumping to unknown locations, due to certain uncontrollable external events such as electrical noise.

### Watchdog Timer Clock Source

The Watchdog Timer clock source is provided by the internal clock,  $f_{LIRC}$  which is sourced from the LIRC oscillator. The LIRC internal oscillator has an approximate frequency of 32kHz and this specified internal clock period can vary with  $V_{DD}$ , temperature and process variations. The Watchdog Timer source clock is then subdivided by a ratio of  $2^8$  to  $2^{18}$  to give longer timeouts, the actual value being chosen using the WS2~WS0 bits in the WDTC register.

### Watchdog Timer Control Register

A single register, WDTC, controls the required time-out period as well as the WDT enable/disable and software reset MCU operation. This register controls the overall operation of the Watchdog Timer.

#### • WDTC Register

Bit	7	6	5	4	3	2	1	0
Name	WE4	WE3	WE2	WE1	WE0	WS2	WS1	WS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	1	0	1	0	0	1	1

Bit 7~3 **WE4~WE0**: WDT function software control

10101: Disable

01010: Enable

Other values: Reset MCU

When these bits are changed to any other values due to environmental noise the microcontroller will be reset; this reset operation will be activated after a delay time,  $t_{SRESET}$  and the WRF bit in the RSTFC register will be set high.

Bit 2~0 **WS2~WS0**: WDT time-out period selection

000:  $2^8/f_{LIRC}$

001:  $2^{10}/f_{LIRC}$

010:  $2^{12}/f_{LIRC}$

011:  $2^{14}/f_{LIRC}$

100:  $2^{15}/f_{LIRC}$

101:  $2^{16}/f_{LIRC}$

110:  $2^{17}/f_{LIRC}$

111:  $2^{18}/f_{LIRC}$

These three bits determine the division ratio of the Watchdog Timer source clock, which in turn determines the timeout period.

• **RSTFC Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	RSTF	LVRF	—	WRF
R/W	—	—	—	—	R/W	R/W	—	R/W
POR	—	—	—	—	0	x	—	0

“x”: unknown

Bit 7~4 Unimplemented, read as “0”

Bit 3 **RSTF**: Reset control register software reset flag  
Refer to Internal Reset Control section.

Bit 2 **LVRF**: LVR function reset flag  
Refer to the Low Voltage Reset section.

Bit 1 Unimplemented, read as “0”

Bit 0 **WRF**: WDT control register software reset flag  
0: Not occurred  
1: Occurred

This bit is set high by the WDT Control register software reset and cleared to zero by the application program. Note that this bit can only be cleared to zero by the application program.

**Watchdog Timer Operation**

The Watchdog Timer operates by providing a device reset when its timer overflows. This means that in the application program and during normal operation the user has to strategically clear the Watchdog Timer before it overflows to prevent the Watchdog Timer from executing a reset. This is done using the clear watchdog instruction. If the program malfunctions for whatever reason, jumps to an unknown location, or enters an endless loop, the clear instruction will not be executed in the correct manner, in which case the Watchdog Timer will overflow and reset the device. There are five bits, WE4~WE0, in the WDTC register to offer the enable/disable control and reset control of the Watchdog Timer. The WDT function will be disabled when the WE4~WE0 bits are set to a value of 10101B while the WDT function will be enabled if the WE4~WE0 bits are equal to 01010B. If the WE4~WE0 bits are set to any other values, other than 01010B and 10101B, it will reset the device after a delay time,  $t_{\text{RESET}}$ . After power on these bits will have a value of 01010B.

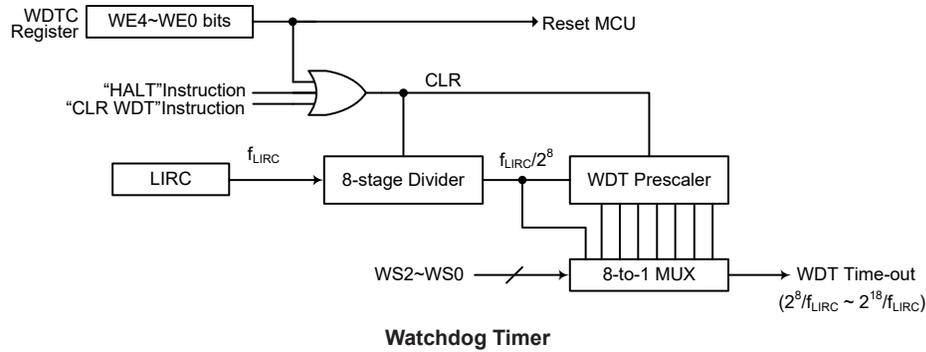
WE4~WE0 Bits	WDT Function
10101B	Disable
01010B	Enable
Any other value	Reset MCU

**Watchdog Timer Enable/Disable Control**

Under normal program operation, a Watchdog Timer time-out will initialise a device reset and set the status bit TO. However, if the system is in the SLEEP or IDLE Mode, when a Watchdog Timer time-out occurs, the TO bit in the status register will be set and only the Program Counter and Stack Pointer will be reset. Three methods can be adopted to clear the contents of the Watchdog Timer. The first is a WDTC software reset, which means a certain value except 01010B and 10101B written into the WE4~WE0 bits, the second is using the Watchdog Timer software clear instruction, the third is via a HALT instruction.

There is only one method of using software instruction to clear the Watchdog Timer. That is to use the single “CLR WDT” instruction to clear the WDT.

The maximum time-out period is when the  $2^{18}$  division ratio is selected. As an example, with a 32kHz LIRC oscillator as its source clock, this will give a maximum watchdog period of around 8s for the  $2^{18}$  division ratio, and a minimum timeout of 8ms for the  $2^8$  division ration.



## Reset and Initialisation

A reset function is a fundamental part of any microcontroller ensuring that the device can be set to some predetermined condition irrespective of outside parameters. The most important reset condition is after power is first applied to the microcontroller. In this case, internal circuitry will ensure that the microcontroller, after a short delay, will be in a well-defined state and ready to execute the first program instruction. After this power-on reset, certain important internal registers will be set to defined states before the program commences. One of these registers is the Program Counter, which will be reset to zero forcing the microcontroller to begin program execution from the lowest Program Memory address.

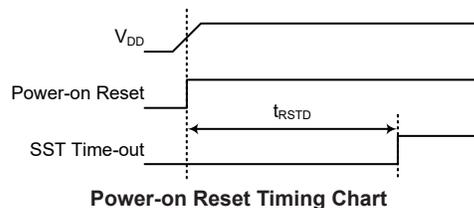
In addition to the power-on reset, another reset exists in the form of a Low Voltage Reset, LVR, where a full reset is implemented in situations where the power supply voltage falls below a certain threshold. Another type of reset is when the Watchdog Timer overflows and resets the microcontroller. All types of reset operations result in different register conditions being setup.

### Reset Functions

There are several ways in which a microcontroller reset can occur, through events occurring internally.

#### Power-on Reset

The most fundamental and unavoidable reset is the one that occurs after power is first applied to the microcontroller. As well as ensuring that the Program Memory begins execution from the first memory address, a power-on reset also ensures that certain other registers are preset to known conditions. All the I/O port and port control registers will power up in a high condition ensuring that all pins will be first set to inputs.



#### Internal Reset Control

There is an internal reset control register, RSTC, which is used to provide a reset when the device operates abnormally due to the environmental noise interference. If the content of the RSTC register is set to any value other than 01010101B or 10101010B, it will reset the device after a delay time,  $t_{SRESET}$ . After power on the register will have a value of 01010101B.

RSTC7~RSTC0 Bits	Reset Function
01010101B	No operation
10101010B	No operation
Any other value	Reset MCU

**Internal Reset Function Control**

• **RSTC Register**

Bit	7	6	5	4	3	2	1	0
Name	RSTC7	RSTC6	RSTC5	RSTC4	RSTC3	RSTC2	RSTC1	RSTC0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	1	0	1	0	1	0	1

Bit 7~0 **RSTC7~RSTC0**: Reset function control  
 01010101: No operation  
 10101010: No operation  
 Other values: Reset MCU

If these bits are changed due to adverse environmental conditions, the microcontroller will be reset. The reset operation will be activated after a delay time,  $t_{SRESET}$  and the RSTF bit in the RSTFC register will be set to 1.

• **RSTFC Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	RSTF	LVRF	—	WRF
R/W	—	—	—	—	R/W	R/W	—	R/W
POR	—	—	—	—	0	x	—	0

“x”: unknown

Bit 7~4 Unimplemented, read as “0”

Bit 3 **RSTF**: Reset control register software reset flag  
 0: Not occurred  
 1: Occurred

This bit is set high by the RSTC control register software reset and cleared to zero by the application program. Note that this bit can only be cleared to 0 by the application program.

Bit 2 **LVRF**: LVR function reset flag  
 Refer to the Low Voltage Reset section.

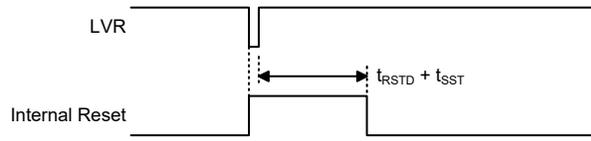
Bit 1 Unimplemented, read as “0”

Bit 0 **WRF**: WDT control register software reset flag  
 Refer to the Watchdog Timer Control Register section

**Low Voltage Reset – LVR**

The microcontroller contains a low voltage reset circuit in order to monitor the supply voltage of the device and provides an MCU reset should the value fall below a certain predefined level.

The LVR function is always enabled in FAST and SLOW Mode with a specific LVR voltage  $V_{LVR}$ . If the supply voltage of the device drops to within a range of  $0.9V \sim V_{LVR}$  such as might occur when changing the battery in battery powered applications, the LVR will automatically reset the device internally and the LVRF bit in the RSTFC register will also be set to 1. For a valid LVR signal, a low supply voltage, i.e., a voltage in the range between  $0.9V \sim V_{LVR}$  must exist for a time greater than that specified by  $t_{LVR}$  in the LVR Electrical Characteristics. If the duration of the low supply voltage state does not exceed this value, the LVR will ignore the low supply voltage and will not perform a reset function. The actual  $V_{LVR}$  is 2.1V, the LVR will reset the device after a delay time,  $t_{SRESET}$ . Note that the LVR function will be automatically disabled when the device enters the SLEEP/IDLE mode.



**Low Voltage Reset Timing Chart**

• **TLVRC Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	TLVR1	TLVR0
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	1

Bit 7~2 Unimplemented, read as “0”

Bit 1~0 **TLVR1~TLVR0**: Minimum low voltage width to reset time ( $t_{LVR}$ ) selection

- 00:  $(7\sim 8) \times t_{LIRC}$
- 01:  $(31\sim 32) \times t_{LIRC}$
- 10:  $(63\sim 64) \times t_{LIRC}$
- 11:  $(127\sim 128) \times t_{LIRC}$

• **RSTFC Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	RSTF	LVRF	—	WRF
R/W	—	—	—	—	R/W	R/W	—	R/W
POR	—	—	—	—	0	x	—	0

“x”: unknown

Bit 7~4 Unimplemented, read as “0”

Bit 3 **RSTF**: Reset control register software reset flag  
Refer to the Internal Reset Control section.

Bit 2 **LVRF**: LVR function reset flag  
0: Not occur  
1: Occurred

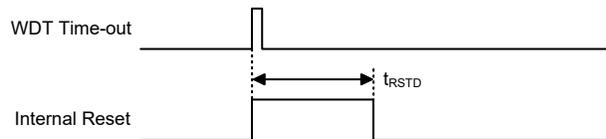
This bit is set high when a specific Low Voltage Reset situation condition occurs. This bit can only be cleared to zero by the application program.

Bit 1 Unimplemented, read as “0”

Bit 0 **WRF**: WDT control register software reset flag  
Refer to the Watchdog Timer Control Register section

**Watchdog Time-out Reset during Normal Operation**

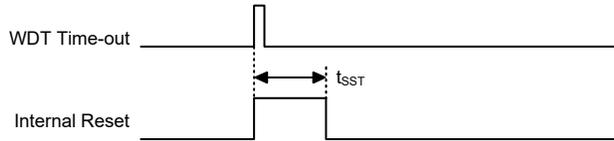
The Watchdog time-out Reset during normal operations in the FAST or SLOW mode is the same as a LVR reset except that the Watchdog time-out flag TO will be set to “1”.



**WDT Time-out Reset during Normal Operation Timing Chart**

**Watchdog Time-out Reset during SLEEP or IDLE Mode**

The Watchdog time-out Reset during SLEEP or IDLE Mode is a little different from other kinds of reset. Most of the conditions remain unchanged except that the Program Counter and the Stack Pointer will be cleared to “0” and the TO and PDF flags will be set to “1”. Refer to the System Start Up Time Characteristics for  $t_{SST}$  details.



**WDT Time-out Reset during SLEEP or IDLE Timing Chart**

**Reset Initial Conditions**

The different types of reset described affect the reset flags in different ways. These flags, known as PDF and TO are located in the status register and are controlled by various microcontroller operations, such as the SLEEP or IDLE Mode function or Watchdog Timer. The reset flags are shown in the table:

TO	PDF	RESET Conditions
0	0	Power-on reset
u	u	LVR reset during FAST or SLOW Mode operation
1	u	WDT time-out reset during FAST or SLOW Mode operation
1	1	WDT time-out reset during IDLE or SLEEP Mode operation

“u” stands for unchanged

The following table indicates the way in which the various components of the microcontroller are affected after a power-on reset occurs.

Item	Condition After RESET
Program Counter	Reset to zero
Interrupts	All interrupts will be disabled
WDT, Time Bases	Clear after reset, WDT begins counting
Timer Module	Timer Module will be turned off
Input/Output Ports	I/O ports will be setup as inputs
Stack Pointer	Stack Pointer will point to the top of the stack

The different kinds of resets all affect the internal registers of the microcontroller in different ways. To ensure reliable continuation of normal program execution after a reset occurs, it is important to know what condition the microcontroller is in after a particular reset occurs. The following table describes how each type of reset affects each of the microcontroller internal registers.

Register	Power On Reset	LVR Reset (Normal Operation)	WDT Time-out (Normal Operation)	WDT Time-out (IDLE/SLEEP)
IAR0	0000 0000	0000 0000	0000 0000	uuuu uuuu
MP0	0000 0000	0000 0000	0000 0000	uuuu uuuu
IAR1	0000 0000	0000 0000	0000 0000	uuuu uuuu
MP1L	0000 0000	0000 0000	0000 0000	uuuu uuuu
MP1H	0000 0000	0000 0000	0000 0000	uuuu uuuu
ACC	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
PCL	0000 0000	0000 0000	0000 0000	0000 0000
TBLP	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
TBLH	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
TBHP	---- xxxx	---- uuuu	---- uuuu	---- uuuu
STATUS	xx00 xxxx	uuuu uuuu	uu1u uuuu	uu11 uuuu
IAR2	0000 0000	0000 0000	0000 0000	uuuu uuuu
MP2L	0000 0000	0000 0000	0000 0000	uuuu uuuu
MP2H	0000 0000	0000 0000	0000 0000	uuuu uuuu
RSTFC	---- 0x-0	---- u1-u	---- uu-u	---- uu-u

Register	Power On Reset	LVR Reset (Normal Operation)	WDT Time-out (Normal Operation)	WDT Time-out (IDLE/SLEEP)
SCC	000- --00	000- --00	000- --00	uuu- --uu
HIRCC	---- --01	---- --01	---- --01	---- --uu
PA	1111 1111	1111 1111	1111 1111	uuuu uuuu
PAC	1111 1111	1111 1111	1111 1111	uuuu uuuu
PAPU	0000 0000	0000 0000	0000 0000	uuuu uuuu
PAWU	0000 0000	0000 0000	0000 0000	uuuu uuuu
RSTC	0101 0101	0101 0101	0101 0101	uuuu uuuu
MFI	--00 --00	--00 --00	--00 --00	--uu --uu
WDTC	0101 0011	0101 0011	0101 0011	uuuu uuuu
INTEG	---- 0000	---- 0000	---- 0000	---- uuuu
INTC0	-000 0000	-000 0000	-000 0000	-uuu uuuu
INTC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
INTC2	---- ---0	---- ---0	---- ---0	---- ---u
TLVRC	---- --01	---- --01	---- --01	---- --uu
PB	-111 1111	-111 1111	-111 1111	-uuu uuuu
PBC	-111 1111	-111 1111	-111 1111	-uuu uuuu
PBPU	-000 0000	-000 0000	-000 0000	-uuu uuuu
PSC0R	---- --00	---- --00	---- --00	---- --uu
PSC1R	---- --00	---- --00	---- --00	---- --uu
TB0C	0--- -000	0--- -000	0--- -000	u--- -uuu
TB1C	0--- -000	0--- -000	0--- -000	u--- -uuu
USR	0000 1011	0000 1011	0000 1011	uuuu uuuu
UCR1	0000 00x0	0000 00x0	0000 00x0	uuuu uuuu
UCR2	0000 0000	0000 0000	0000 0000	uuuu uuuu
UCR3	---- ---0	---- ---0	---- ---0	---- ---u
BRDH	0000 0000	0000 0000	0000 0000	uuuu uuuu
BRDL	0000 0000	0000 0000	0000 0000	uuuu uuuu
UFCR	--00 0000	--00 0000	--00 0000	--uu uuuu
TXR_RXR	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
RxCNT	---- -000	---- -000	---- -000	---- -uuu
SIMC0	111- 0000	111- 0000	111- 0000	uuu- uuuu
SIMC1	1000 0001	1000 0001	1000 0001	uuuu uuuu
SIMD	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
SIMA/SIMC2	0000 0000	0000 0000	0000 0000	uuuu uuuu
SIMTOC	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMC0	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMDL	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMDH	---- --00	---- --00	---- --00	---- --uu
CTMAL	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMAH	---- --00	---- --00	---- --00	---- --uu
SLEDC	0000 0000	0000 0000	0000 0000	uuuu uuuu
PAS0	0000 0000	0000 0000	0000 0000	uuuu uuuu
PAS1	0000 0000	0000 0000	0000 0000	uuuu uuuu
PBS0	---- 00--	---- 00--	---- 00--	---- uu--
PBS1	--00 00--	--00 00--	--00 00--	--uu uu--
OCR	---- 00--	---- 00--	---- 00--	---- uu--

Register	Power On Reset	LVR Reset (Normal Operation)	WDT Time-out (Normal Operation)	WDT Time-out (IDLE/SLEEP)
ODL	0000 0000	0000 0000	0000 0000	uuuu uuuu
ODH	0000 0000	0000 0000	0000 0000	uuuu uuuu
PWRC	00-0 0000	00-0 0000	00-0 0000	uu-u uuuu
PGAC0	0000 0000	0000 0000	0000 0000	uuuu uuuu
PGAC1	0 0 0 - - - - -	0 0 0 - - - - -	0 0 0 - - - - -	uuu- - - - -
PGACS	0000 0000	0000 0000	0000 0000	uuuu uuuu
ADRL	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
ADRM	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
ADRH	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
ADCR0	0000 0000	0000 0000	0000 0000	uuuu uuuu
ADCR1	- - 0 0 0 - -	- - 0 0 0 - -	- - 0 0 0 - -	- - uu u - -
ADCS	- - - 0 0000	- - - 0 0000	- - - 0 0000	- - - u uuuu
SGC	0 0 - - - - 0 0	0 0 - - - - 0 0	0 0 - - - - 0 0	uu - - - - uu
SGN	0000 0000	0000 0000	0000 0000	uuuu uuuu
SGDN	- - 0 0 0000	- - 0 0 0000	- - 0 0 0000	- - uu uuuu
OPAC	0 0 1 0 0000	0 0 1 0 0000	0 0 1 0 0000	uuuu uuuu
SWC0	0000 0000	0000 0000	0000 0000	uuuu uuuu
SWC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
SWC2	0000 0000	0000 0000	0000 0000	uuuu uuuu
SWC3	0000 0000	0000 0000	0000 0000	uuuu uuuu
CMPC0	0-0 0 0-0 0	0-0 0 0-0 0	0-0 0 0-0 0	u-uu u-uu
CMPC1	- - - - 0-0 0	- - - - 0-0 0	- - - - 0-0 0	- - - - u-uu
BDA0C	0-0 0 0000	0-0 0 0000	0-0 0 0000	u-uu uuuu
BDA1C	0-0 0 0000	0-0 0 0000	0-0 0 0000	u-uu uuuu
BDA2C	0-0 0 0000	0-0 0 0000	0-0 0 0000	u-uu uuuu

Note: “u” stands for unchanged  
“x” stands for unknown  
“-” stands for unimplemented

## Input/Output Ports

Holtek microcontrollers offer considerable flexibility on their I/O ports. With the input or output designation of every pin fully under user program control, pull-high selections for all ports and wake-up selections on certain pins, the user is provided with an I/O structure to meet the needs of a wide range of application possibilities.

The device provides bidirectional input/output lines labeled with port names PA~PB. These I/O ports are mapped to the RAM Data Memory with specific addresses as shown in the Special Purpose Data Memory table. All of these I/O ports can be used for input and output operations. For input operation, these ports are non-latching, which means the inputs must be ready at the T2 rising edge of instruction “MOV A, [m]”, where m denotes the port address. For output operation, all the data is latched and remains unchanged until the output latch is rewritten.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PA	PA7	PA6	PA5	PA4	PA3	PA2	PA1	PA0
PAC	PAC7	PAC6	PAC5	PAC4	PAC3	PAC2	PAC1	PAC0
PAPU	PAPU7	PAPU6	PAPU5	PAPU4	PAPU3	PAPU2	PAPU1	PAPU0

Register Name	Bit							
	7	6	5	4	3	2	1	0
PAWU	PAWU7	PAWU6	PAWU5	PAWU4	PAWU3	PAWU2	PAWU1	PAWU0
PB	—	PB6	PB5	PB4	PB3	PB2	PB1	PB0
PBC	—	PBC6	PBC5	PBC4	PBC3	PBC2	PBC1	PBC0
PBPU	—	PBPU6	PBPU5	PBPU4	PBPU3	PBPU2	PBPU1	PBPU0

“—”: Unimplemented, read as “0”

### I/O Logic Function Register List

## Pull-high Resistors

Many product applications require pull-high resistors for their switch inputs usually requiring the use of an external resistor. To eliminate the need for these external resistors, all I/O pins, when configured as a digital input have the capability of being connected to an internal pull-high resistor. These pull-high resistors are selected using registers, namely PPU~PBPU, and are implemented using weak PMOS transistors.

Note that the pull-high resistor can be controlled by the relevant pull-high control register only when the pin-shared functional pin is selected as a digital input or NMOS output. Otherwise, the pull-high resistors cannot be enabled.

### • P<sub>x</sub>PU Register

Bit	7	6	5	4	3	2	1	0
Name	P <sub>x</sub> PU7	P <sub>x</sub> PU6	P <sub>x</sub> PU5	P <sub>x</sub> PU4	P <sub>x</sub> PU3	P <sub>x</sub> PU2	P <sub>x</sub> PU1	P <sub>x</sub> PU0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

**P<sub>x</sub>PU<sub>n</sub>**: I/O Port x Pin pull-high function control

0: Disable

1: Enable

The P<sub>x</sub>PU<sub>n</sub> bit is used to control the pin pull-high function. Here the “x” can be A or B. However, the actual available bits for each I/O Port may be different.

## Port A Wake-up

The HALT instruction forces the microcontroller into the SLEEP or IDLE Mode which preserves power, a feature that is important for battery and other low-power applications. Various methods exist to wake-up the microcontroller, one of which is to change the logic condition on one of the Port A pins from high to low. This function is especially suitable for applications that can be woken up via external switches. Each pin on Port A can be selected individually to have this wake-up feature using the PAWU register.

Note that the wake-up function can be controlled by the wake-up control registers only when the pin is selected as general purpose input and the MCU enters the IDLE or SLEEP mode.

### • PAWU Register

Bit	7	6	5	4	3	2	1	0
Name	PAWU7	PAWU6	PAWU5	PAWU4	PAWU3	PAWU2	PAWU1	PAWU0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **PAWU7~PAWU0**: PA7~PA0 wake-up function control

0: Disable

1: Enable

## I/O Port Control Registers

Each I/O port has its own control register known as PAC~PBC, to control the input/output configuration. With this control register, each CMOS output or input can be reconfigured dynamically under software control. Each pin of the I/O ports is directly mapped to a bit in its associated port control register. For the I/O pin to function as an input, the corresponding bit of the control register must be written as a “1”. This will then allow the logic state of the input pin to be directly read by instructions. When the corresponding bit of the control register is written as a “0”, the I/O pin will be setup as a CMOS output. If the pin is currently setup as an output, instructions can still be used to read the output register. However, it should be noted that the program will in fact only read the status of the output data latch and not the actual logic status of the output pin.

### • PxC Register

Bit	7	6	5	4	3	2	1	0
Name	PxC7	PxC6	PxC5	PxC4	PxC3	PxC2	PxC1	PxC0
R/W								
POR	1	1	1	1	1	1	1	1

**PxCn:** I/O Port x Pin type selection

0: Output

1: Input

The PxCn bit is used to control the pin type selection. Here the “x” can be A or B. However, the actual available bits for each I/O Port may be different.

## I/O Port Source Current Selection

The device supports different output source current driving capability for each I/O port. With the selection register SLEDC, specific I/O port can support four levels of the source current driving capability. These source current selection bits are available when the corresponding pin is configured as a CMOS output. Otherwise, these select bits have no effect. Users should refer to the Input/Output Characteristics section to select the desired output source current for different applications.

### • SLEDC Register

Bit	7	6	5	4	3	2	1	0
Name	SLEDC7	SLEDC6	SLEDC5	SLEDC4	SLEDC3	SLEDC2	SLEDC1	SLEDC0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 **SLEDC7~SLEDC6:** PB6~PB4 source current selection

00: Source current=Level 0 (min.)

01: Source current=Level 1

10: Source current=Level 2

11: Source current=Level 3 (max.)

Bit 5~4 **SLEDC5~SLEDC4:** PB3~PB1 source current selection

00: Source current=Level 0 (min.)

01: Source current=Level 1

10: Source current=Level 2

11: Source current=Level 3 (max.)

Bit 3~2 **SLEDC3~SLEDC2:** PA7~PA4 source current selection

00: Source current=Level 0 (min.)

01: Source current=Level 1

10: Source current=Level 2

11: Source current=Level 3 (max.)

- Bit 1~0    **SLEDC1~SLEDC0**: PA3~PA0 source current selection  
 00: Source current=Level 0 (min.)  
 01: Source current=Level 1  
 10: Source current=Level 2  
 11: Source current=Level 3 (max.)

### Pin-shared Functions

The flexibility of the microcontroller range is greatly enhanced by the use of pins that have more than one function. Limited numbers of pins can force serious design constraints on designers but by supplying pins with multi-functions, many of these difficulties can be overcome. For these pins, the desired function of the multi-function I/O pins is selected by a series of registers via the application program control.

### Pin-shared Function Selection Registers

The limited number of supplied pins in a package can impose restrictions on the amount of functions a certain device can contain. However by allowing the same pins to share several different functions and providing a means of function selection, a wide range of different functions can be incorporated into even relatively small package sizes. The device includes Port “x” Output Function Selection register “n”, labeled as PxSn, which can select the desired functions of the multi-function pin-shared pins.

The most important point to note is to make sure that the desired pin-shared function is properly selected and also deselected. For most pin-shared functions, to select the desired pin-shared function, the pin-shared function should first be correctly selected using the corresponding pin-shared control register. After that the corresponding peripheral functional setting should be configured and then the peripheral function can be enabled. However, a special point must be noted for some digital input pins, such as INTn, CTCK etc, which share the same pin-shared control configuration with their corresponding general purpose I/O functions when setting the relevant pin-shared control bit fields. To select these pin functions, in addition to the necessary pin-shared control and peripheral functional setup aforementioned, they must also be setup as an input by setting the corresponding bit in the I/O port control register. To correctly deselect the pin-shared function, the peripheral function should first be disabled and then the corresponding pin-shared function control register can be modified to select other pin-shared functions.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PAS0	PAS07	PAS06	PAS05	PAS04	PAS03	PAS02	PAS01	PAS00
PAS1	PAS17	PAS16	PAS15	PAS14	PAS13	PAS12	PAS11	PAS10
PBS0	—	—	—	—	PBS03	PBS02	—	—
PBS1	—	—	PBS15	PBS14	PBS13	PBS12	—	—

**Pin-shared Function Selection Register List**

• **PAS0 Register**

Bit	7	6	5	4	3	2	1	0
Name	PAS07	PAS06	PAS05	PAS04	PAS03	PAS02	PAS01	PAS00
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7~6     **PAS07~PAS06:** PA3 Pin-Shared function selection  
00: PA3  
01: AIN4  
10: PA3  
11: PA3
- Bit 5~4     **PAS05~PAS04:** PA2 Pin-Shared function selection  
00: PA2  
01: TX  
10: PA2  
11: PA2
- Bit 3~2     **PAS03~PAS02:** PA1 Pin-Shared function selection  
00: PA1  
01: AIN5  
10: PA1  
11: PA1
- Bit 1~0     **PAS01~PAS00:** PA0 Pin-Shared function selection  
00: PA0  
01: RX/TX  
10: PA0  
11: PA0

• **PAS1 Register**

Bit	7	6	5	4	3	2	1	0
Name	PAS17	PAS16	PAS15	PAS14	PAS13	PAS12	PAS11	PAS10
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7~6     **PAS17~PAS16:** PA7 Pin-Shared function selection  
00: PA7  
01: SCS  
10: PA7  
11: PA7
- Bit 5~4     **PAS15~PAS14:** PA6 Pin-Shared function selection  
00: PA6  
01: SDO  
10: PA6  
11: PA6
- Bit 3~2     **PAS13~PAS12:** PA5 Pin-Shared function selection  
00: PA5  
01: SDI/SDA  
10: PA5  
11: PA5
- Bit 1~0     **PAS11~PAS10:** PA4 Pin-Shared function selection  
00: PA4  
01: SCK/SCL  
10: PA4  
11: PA4

• **PBS0 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	PBS03	PBS02	—	—
R/W	—	—	—	—	R/W	R/W	—	—
POR	—	—	—	—	0	0	—	—

Bit 7~4 Unimplemented, read as “0”

Bit 3~2 **PBS03~PBS02**: PB1 Pin-Shared function selection

00: PB1

01: MODSYNC

10: PB1

11: PB1

Bit 1~0 Unimplemented, read as “0”

• **PBS1 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	PBS15	PBS14	PBS13	PBS12	—	—
R/W	—	—	R/W	R/W	R/W	R/W	—	—
POR	—	—	0	0	0	0	—	—

Bit 7~6 Unimplemented, read as “0”

Bit 5~4 **PBS15~PBS14**: PB6 Pin-Shared function selection

00: PB6

01: STP

10: PB6

11: PB6

Bit 3~2 **PBS13~PBS12**: PB5 Pin-Shared function selection

00: PB5

01: SCK/SCL

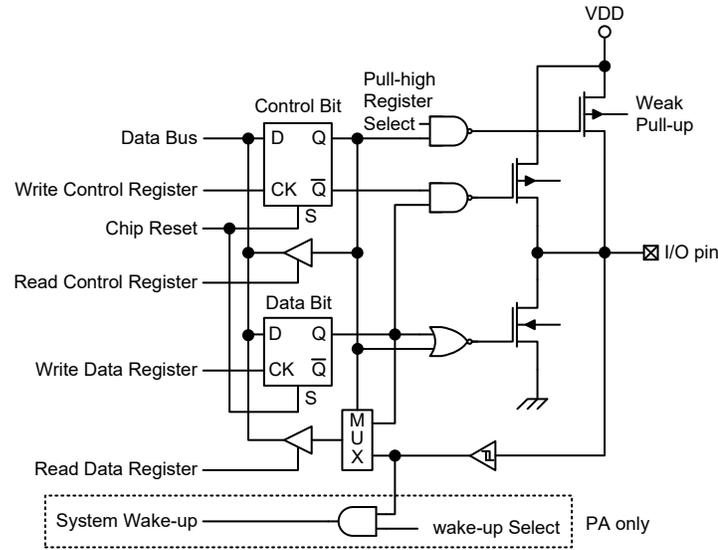
10: TX

11: PB5

Bit 1~0 Unimplemented, read as “0”

**I/O Pin Structures**

The accompanying diagram illustrates the internal structure of the I/O logic function. As the exact logical construction of the I/O pin will differ from this drawing, it is supplied as a guide only to assist with the functional understanding of the I/O logic function. The wide range of pin-shared structures does not permit all types to be shown.



**Logic Function Input/Output Structure**

### Programming Considerations

Within the user program, one of the first things to consider is port initialisation. After a reset, all of the I/O data and port control registers will be set high. This means that all I/O pins will default to an input state, the level of which depends on the other connected circuitry and whether pull-high selections have been chosen. If the port control registers are then programmed to setup some pins as outputs, these output pins will have an initial high output value unless the associated port data registers are first programmed. Selecting which pins are inputs and which are outputs can be achieved byte-wide by loading the correct values into the appropriate port control register or by programming individual bits in the port control register using the “SET [m].i” and “CLR [m].i” instructions. Note that when using these bit control instructions, a read-modify-write operation takes place. The microcontroller must first read in the data on the entire port, modify it to the required new bit values and then rewrite this data back to the output ports.

Port A has the additional capability of providing wake-up function. When the device is in the SLEEP or IDLE Mode, various methods are available to wake the device up. One of these is a high to low transition of any of the Port A pins. Single or multiple pins on Port A can be setup to have this function.

## Timer Modules – TM

One of the most fundamental functions in any microcontroller device is the ability to control and measure time. To implement time related functions the device includes a Timer Module, abbreviated to the name TM. The TM is multi-purpose timing unit and serves to provide operations such as Timer/Counter, Compare Match Output as well as being the functional unit for the generation of PWM signals. The TM has two individual interrupts. The addition of input and output pins for the TM ensures that users are provided with timing units with a wide and flexible range of features.

The general features of the Compact Type TM are described here with more detailed information provided in the individual Compact Type TM section.

### Introduction

The device contains one Compact Type TM. The main features of the CTM are summarised in the accompanying table.

TM Function	CTM
Timer/Counter	√
Input Capture	—
Compare Match Output	√
PWM Output	√
Single Pulse Output	—
PWM Alignment	Edge
PWM Adjustment Period & Duty	Duty or Period

**CTM Function Summary**

### TM Operation

The Compact Type TM offers a diverse range of functions, from simple timing operations to PWM signal generation. The key to understanding how the TM operates is to see it in terms of a free running counter whose value is then compared with the value of pre-programmed internal comparators. When the free running counter has the same value as the pre-programmed comparator, known as a compare match situation, a TM interrupt signal will be generated which can clear the counter and perhaps also change the condition of the TM output pin. The internal TM counter is driven by a user selectable clock source, which can be an internal clock or an external pin.

### TM Clock Source

The clock source which drives the main counter in the TM can originate from various sources. The selection of the required clock source is implemented using the CTCK2~CTCK0 bits in the CTM control registers. The clock source can be a ratio of the system clock  $f_{SYS}$  or the internal high clock  $f_H$ , the  $f_{SUB}$  clock source or the external CTCK pin. The CTCK pin clock source is used to allow an external signal to drive the TM as an external clock source or for event counting.

### TM Interrupts

The Compact Type TM has two internal interrupts, the internal comparator A or comparator P, which generate a TM interrupt when a compare match condition occurs. When a TM interrupt is generated, it can be used to clear the counter and also to change the state of the TM output pin.

### TM External Pins

The Compact Type TM has one TM input pin, with the label CTCK. The CTM input pin, CTCK, is essentially a clock source for the CTM and is selected using the CTCK2~CTCK0 bits in the CTMC0

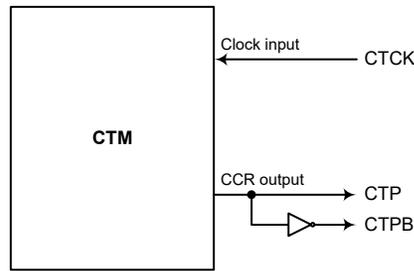
register. This external TM input pin allows an external clock source to drive the internal TM. The CTCK input pin can be chosen to have either a rising or falling active edge.

The TM has two output pins with the label CTP and CTPB. The CTPB pin outputs the inverted signal of the CTP. When the TM is in the Compare Match Output Mode, these pins can be controlled by the TM to switch to a high or low level or to toggle when a compare match situation occurs. The external CTP and CTPB output pins are also the pins where the TM generates the PWM output waveform.

As the TM input and output pins are pin-shared with other functions, the TM input and output function must first be setup using relevant pin-shared function selection register described in the Pin-shared Function section.

CTM	
Input	Output
CTCK	CTP, CTPB

**TM External Pins**

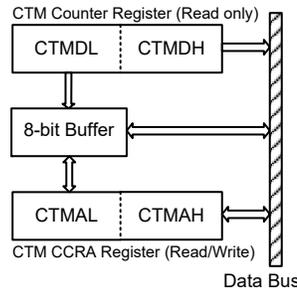


**CTM Function Pin Block Diagram**

### Programming Considerations

The TM Counter Registers and the Compare CCRA registers, all have a low and high byte structure. The high bytes can be directly accessed, but as the low bytes can only be accessed via an internal 8-bit buffer, reading or writing to these register pairs must be carried out in a specific way. The important point to note is that data transfer to and from the 8-bit buffer and its related low byte only takes place when a write or read operation to its corresponding high byte is executed.

As the CCRA registers are implemented in the way shown in the following diagram and accessing this register pair is carried out in a specific way described above, it is recommended to use the “MOV” instruction to access the CCRA low byte register, named CTMAL, in the following access procedures. Accessing the CCRA low byte register without following these access procedures will result in unpredictable values.

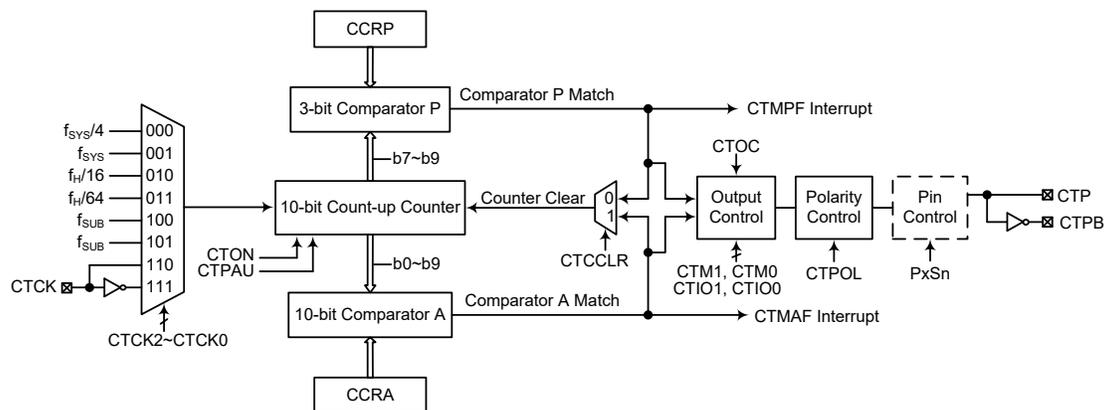


The following steps show the read and write procedures:

- Writing Data to CCRA
  - ♦ Step 1. Write data to Low Byte CTMAL
    - Note that here data is only written to the 8-bit buffer.
  - ♦ Step 2. Write data to High Byte CTMAH
    - Here data is written directly to the high byte registers and simultaneously data is latched from the 8-bit buffer to the Low Byte registers.
- Reading Data from the Counter Registers and CCRA
  - ♦ Step 1. Read data from the High Byte CTMDH or CTMAH
    - Here data is read directly from the High Byte registers and simultaneously data is latched from the Low Byte register into the 8-bit buffer.
  - ♦ Step 2. Read data from the Low Byte CTMDL or CTMAL
    - This step reads data from the 8-bit buffer.

### Compact Type TM – CTM

The Compact TM type contains three operating modes, which are Compare Match Output, Timer/Event Counter and PWM Output modes. The Compact TM can also be controlled with an external input pin and can drive two external output pins.



Note: 1. The CTM external pins are pin-shared with other functions, so before using the CTM function, the pin-shared function registers must be set properly to enable the CTM pin function. The CTCK pin, if used, must also be set as an input by setting the corresponding bits in the port control register.

2. The CTPB is the inverted signal of the CTP.

10-Bit Compact Type TM Block Diagram

### Compact TM Operation

The size of Compact Type TM is 10-bit wide and its core is a 10-bit count-up counter which is driven by a user selectable internal or external clock source. There are also two internal comparators with the names, Comparator A and Comparator P. These comparators will compare the value in the counter with CCRP and CCRA registers. The CCRP comparator is 3-bit wide whose value is compared with the highest 3 bits in the counter while the CCRA is the 10 bits and therefore compares all counter bits.

The only way of changing the value of the 10-bit counter using the application program, is to clear the counter by changing the CTON bit from low to high. The counter will also be cleared

automatically by a counter overflow or a compare match with one of its associated comparators. When these conditions occur, a CTM interrupt signal will also usually be generated. The Compact Type TM can operate in a number of different operational modes, can be driven by different clock sources including an input pin and can also control an output pin. All operating setup conditions are selected using relevant internal registers.

### Compact Type TM Register Description

Overall operation of each Compact TM is controlled using several registers. A read only register pair exists to store the internal counter 10-bit value, while a read/write register pair exists to store the internal 10-bit CCRA value. The remaining two registers are control registers which setup the different operating and control modes as well as the three CCRP bits.

Register Name	Bit							
	7	6	5	4	3	2	1	0
CTMC0	CTPAU	CTCK2	CTCK1	CTCK0	CTON	CTRP2	CTRP1	CTRP0
CTMC1	CTM1	CTM0	CTIO1	CTIO0	CTOC	CTPOL	CTDPX	CTCCLR
CTMDL	D7	D6	D5	D4	D3	D2	D1	D0
CTMDH	—	—	—	—	—	—	D9	D8
CTMAL	D7	D6	D5	D4	D3	D2	D1	D0
CTMAH	—	—	—	—	—	—	D9	D8

10-bit Compact TM Register List

#### • CTMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	CTPAU	CTCK2	CTCK1	CTCK0	CTON	CTRP2	CTRP1	CTRP0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **CTPAU**: CTM Counter Pause Control

- 0: Run
- 1: Pause

The counter can be paused by setting this bit high. Clearing the bit to zero restores normal counter operation. When in a Pause condition the CTM will remain powered up and continue to consume power. The counter will retain its residual value when this bit changes from low to high and resume counting from this value when the bit changes to a low value again.

Bit 6~4 **CTCK2~CTCK0**: Select CTM Counter clock

- 000:  $f_{SYS}/4$
- 001:  $f_{SYS}$
- 010:  $f_H/16$
- 011:  $f_H/64$
- 100:  $f_{SUB}$
- 101:  $f_{SUB}$
- 110: CTCK rising edge clock
- 111: CTCK falling edge clock

These three bits are used to select the clock source for the CTM. The external pin clock source can be chosen to be active on the rising or falling edge. The clock source  $f_{SYS}$  is the system clock, while  $f_H$  and  $f_{SUB}$  are other internal clocks, the details of which can be found in the Operating Modes and System Clocks section.

Bit 3 **CTON**: CTM Counter On/Off Control

- 0: Off
- 1: On

This bit controls the overall on/off function of the CTM. Setting the bit high enables the

counter to run, clearing the bit disables the CTM. Clearing this bit to zero will stop the counter from counting and turn off the CTM which will reduce its power consumption. When the bit changes state from low to high the internal counter value will be reset to zero, however when the bit changes from high to low, the internal counter will retain its residual value until the bit returns high again.

If the CTM is in the Compare Match Output Mode or the PWM Output Mode then the CTM output pin will be reset to its initial condition, as specified by the CTOC bit, when the CTON bit changes from low to high.

Bit 2~0 **CTRP2~CTRP0**: CTM CCRP 3-bit register, compared with the CTM Counter bit 9~bit 7

Comparator P Match Period:

- 000: 1024 CTM clocks
- 001: 128 CTM clocks
- 010: 256 CTM clocks
- 011: 384 CTM clocks
- 100: 512 CTM clocks
- 101: 640 CTM clocks
- 110: 768 CTM clocks
- 111: 896 CTM clocks

These three bits are used to setup the value on the internal CCRP 3-bit register, which are then compared with the internal counter's highest three bits. The result of this comparison can be selected to clear the internal counter if the CTCCLR bit is set to zero. Setting the CTCCLR bit to zero ensures that a compare match with the CCRP values will reset the internal counter. As the CCRP bits are only compared with the highest three counter bits, the compare values exist in 128 clock cycle multiples. Clearing all three bits to zero is in effect allowing the counter to overflow at its maximum value.

• **CTMC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	CTM1	CTM0	CTIO1	CTIO0	CTOC	CTPOL	CTDPX	CTCCLR
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 **CTM1~CTM0**: Select CTM Operating Mode

- 00: Compare Match Output Mode
- 01: Undefined
- 10: PWM Output Mode
- 11: Timer/Counter Mode

These bits setup the required operating mode for the CTM. To ensure reliable operation the CTM should be switched off before any changes are made to the CTM1 and CTM0 bits. In the Timer/Counter Mode, the CTM output pin state is undefined.

Bit 5~4 **CTIO1~CTIO0**: CTM external pin function selection

Compare Match Output Mode

- 00: No change
- 01: Output low
- 10: Output high
- 11: Toggle output

PWM Output Mode

- 00: PWM output inactive state
- 01: PWM output active state
- 10: PWM output
- 11: Undefined

Timer/counter Mode:

Unused

These two bits are used to determine how the CTM external pin changes state when a certain condition is reached. The function that these bits select depends upon in which

mode the CTM is running.

In the Compare Match Output Mode, the CTIO1~CTIO0 bits determine how the CTM output pin changes state when a compare match occurs from the Comparator A. The CTM output pin can be setup to switch high, switch low or to toggle its present state when a compare match occurs from the Comparator A. When the CTIO1~CTIO0 bits are both zero, then no change will take place on the output. The initial value of the CTM output pin should be setup using the CTOC bit. Note that the output level requested by the CTIO1~CTIO0 bits must be different from the initial value setup using the CTOC bit otherwise no change will occur on the CTM output pin when a compare match occurs. After the CTM output pin changes state it can be reset to its initial level by changing the level of the CTON bit from low to high.

In the PWM Output Mode, the CTIO1 and CTIO0 bits determine how the CTM output pin changes state when a certain compare match condition occurs. The PWM output function is modified by changing these two bits. It is necessary to change the values of the CTIO1 and CTIO0 bits only after the CTM has been switched off. Unpredictable PWM outputs will occur if the CTIO1 and CTIO0 bits are changed when the CTM is running.

- Bit 3 **CTOC**: CTM CTP Output control bit  
Compare Match Output Mode  
0: Initial low  
1: Initial high  
PWM Output Mode  
0: Active low  
1: Active high

This is the output control bit for the CTM output pin. Its operation depends upon whether CTM is being used in the Compare Match Output Mode or in the PWM Output Mode. It has no effect if the CTM is in the Timer/Counter Mode. In the Compare Match Output Mode it determines the logic level of the CTM output pin before a compare match occurs. In the PWM Output Mode it determines if the PWM signal is active high or active low.

- Bit 2 **CTPOL**: CTM CTP Output polarity Control  
0: Non-invert  
1: Invert

This bit controls the polarity of the CTM output pins. When the bit is set high the CTM output pins will be inverted and not inverted when the bit is zero. It has no effect if the CTM is in the Timer/Counter Mode.

- Bit 1 **CTDPX**: CTM PWM period/duty Control  
0: CCRP – period; CCRA – duty  
1: CCRP – duty; CCRA – period

This bit, determines which of the CCRA and CCRP registers are used for period and duty control of the PWM waveform.

- Bit 0 **CTCCLR**: CTM Counter clear condition selection  
0: CTM Comparatror P match  
1: CTM Comparatror A match

This bit is used to select the method which clears the counter. Remember that the Compact TM contains two comparators, Comparator A and Comparator P, either of which can be selected to clear the internal counter. With the CTCCLR bit set high, the counter will be cleared when a compare match occurs from the Comparator A. When the bit is low, the counter will be cleared when a compare match occurs from the Comparator P or with a counter overflow. A counter overflow clearing method can only be implemented if the CCRP bits are all cleared to zero. The CTCCLR bit is not used in the PWM Output Mode.

**• CTMDL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: CTM Counter Low Byte Register bit 7~bit 0  
 CTM 10-bit Counter bit 7~bit 0

**• CTMDH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R	R
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as “0”  
 Bit 1~0 **D9~D8**: CTM Counter High Byte Register bit 1~bit 0  
 CTM 10-bit Counter bit 9~bit 8

**• CTMAL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: CTM CCRA Low Byte Register bit 7~bit 0  
 CTM 10-bit CCRA bit 7~bit 0

**• CTMAH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as “0”  
 Bit 1~0 **D9~D8**: CTM CCRA High Byte Register bit 1~bit 0  
 CTM 10-bit CCRA bit 9~bit 8

**Compact Type TM Operating Modes**

The Compact Type TM can operate in one of three operating modes, Compare Match Output Mode, PWM Output Mode or Timer/Counter Mode. The operating mode is selected using the CTM1 and CTM0 bits in the CTMC1 register.

**Compare Match Output Mode**

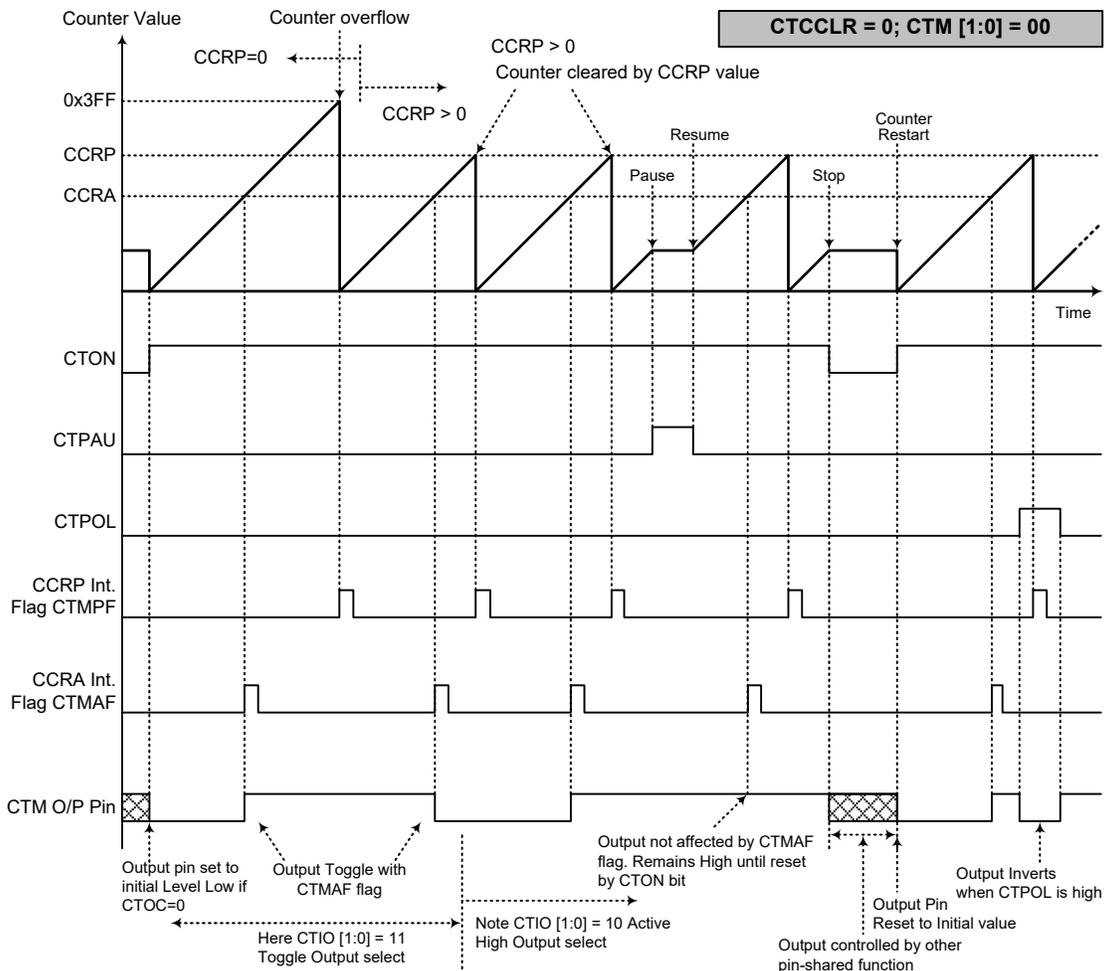
To select this mode, bits CTM1 and CTM0 in the CTMC1 register, should be set to 00 respectively. In this mode once the counter is enabled and running it can be cleared by three methods. These are a counter overflow, a compare match from Comparator A and a compare match from Comparator P. When the CTCCLR bit is low, there are two ways in which the counter can be cleared. One is when a compare match from Comparator P, the other is when the CCRP bits are all zero which allows the counter to overflow. Here both CTMAF and CTMPF interrupt request flags for Comparator A and Comparator P respectively, will both be generated.

If the CTCCLR bit in the CTMC1 register is high then the counter will be cleared when a compare

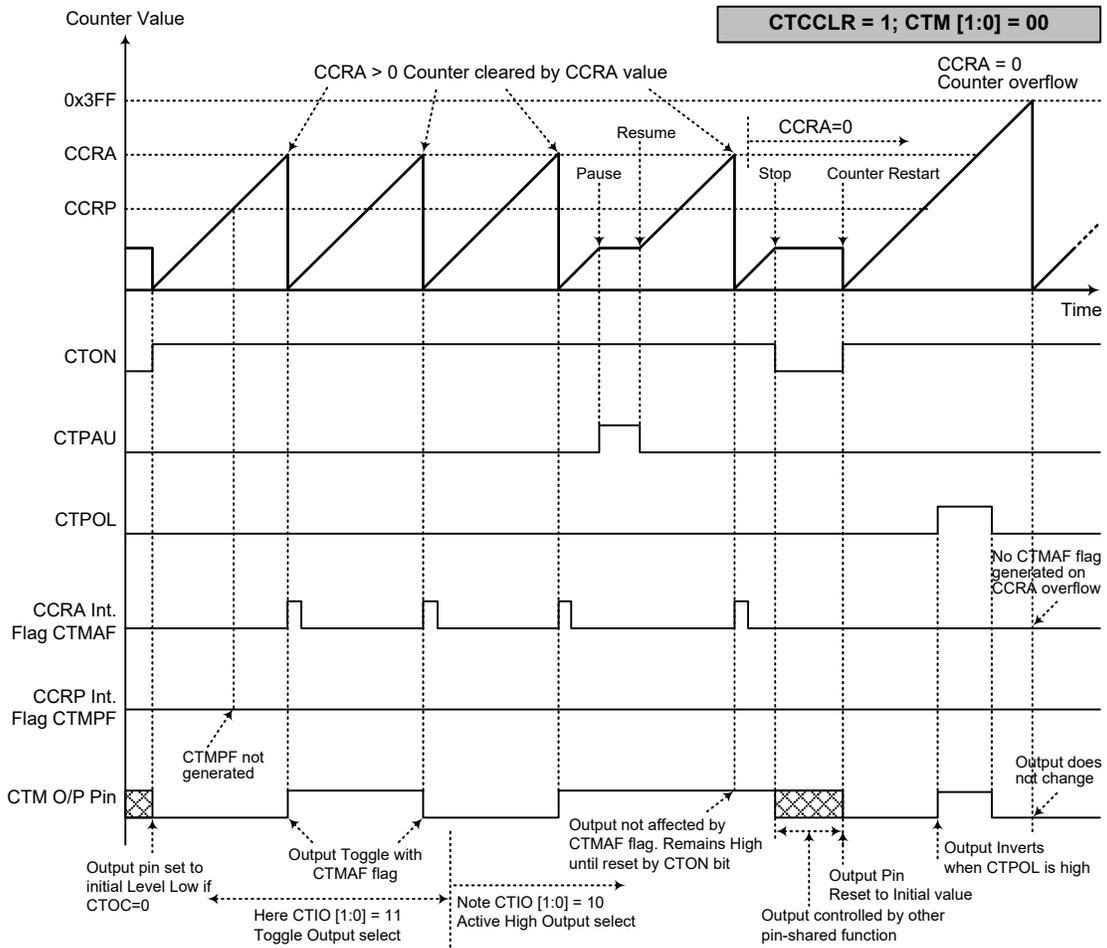
match occurs from Comparator A. However, here only the CTMAF interrupt request flag will be generated even if the value of the CCRP bits is less than that of the CCRA registers. Therefore when CTCCLR is high no CTMPF interrupt request flag will be generated.

If the CCRA bits are all zero, the counter will overflow when it reaches its maximum 10-bit, 3FF Hex, value, however here the CTMAF interrupt request flag will not be generated.

As the name of the mode suggests, after a comparison is made, the CTM output pin, will change state. The CTM output pin condition however only changes state when a CTMAF interrupt request flag is generated after a compare match occurs from Comparator A. The CTMPF interrupt request flag, generated from a compare match occurs from Comparator P, will have no effect on the CTM output pin. The way in which the CTM output pin changes state are determined by the condition of the CTIO1 and CTIO0 bits in the CTMC1 register. The CTM output pin can be selected using the CTIO1 and CTIO0 bits to go high, to go low or to toggle from its present condition when a compare match occurs from Comparator A. The initial condition of the CTM output pin, which is setup after the CTON bit changes from low to high, is setup using the CTOC bit. Note that if the CTIO1 and CTIO0 bits are zero then no pin change will take place.



- Note: 1. With CTCCLR=0, a Comparator P match will clear the counter  
 2. The CTM output pin controlled only by the CTMAF flag  
 3. The output pin reset to initial state by a CTON bit rising edge



**Compare Match Output Mode – CTCCLR=1**

- Note: 1. With CTCCLR=1, a Comparator A match will clear the counter
2. The CTM output pin controlled only by the CTMAF flag
3. The output pin reset to initial state by a CTON rising edge
4. The CTMPF flags is not generated when CTCCLR=1

**Timer/Counter Mode**

To select this mode, bits CTM1 and CTM0 in the CTMC1 register should be set to 11 respectively. The Timer/Counter Mode operates in an identical way to the Compare Match Output Mode generating the same interrupt flags. The exception is that in the Timer/Counter Mode the CTM output pin is not used. Therefore the above description and Timing Diagrams for the Compare Match Output Mode can be used to understand its function. As the CTM output pin is not used in this mode, the pin can be used as a normal I/O pin or other pin-shared function.

**PWM Output Mode**

To select this mode, bits CTM1 and CTM0 in the CTMC1 register should be set to 10 respectively. The PWM function within the CTM is useful for applications which require functions such as motor control, heating control, illumination control etc. By providing a signal of fixed frequency but of varying duty cycle on the CTM output pin, a square wave AC waveform can be generated with varying equivalent DC RMS values.

As both the period and duty cycle of the PWM waveform can be controlled, the choice of generated waveform is extremely flexible. In the PWM Output Mode, the CTCCLR bit has no effect as the PWM period. Both of the CCRA and CCRP registers are used to generate the PWM waveform, one register is used to clear the internal counter and thus control the PWM waveform frequency, while the other one is used to control the duty cycle. Which register is used to control either frequency or duty cycle is determined using the CTDPX bit in the CTMC1 register. The PWM waveform frequency and duty cycle can therefore be controlled by the values in the CCRA and CCRP registers.

An interrupt flag, one for each of the CCRA and CCRP, will be generated when a compare match occurs from either Comparator A or Comparator P. The CTOC bit In the CTMC1 register is used to select the required polarity of the PWM waveform while the two CTIO1 and CTIO0 bits are used to enable the PWM output or to force the CTM output pin to a fixed high or low level. The CTPOL bit is used to reverse the polarity of the PWM output waveform.

• **10-bit CTM, PWM Output Mode, Edge-aligned Mode, CTDPX=0**

CCRP	1~7	0
Period	CCRP×128	1024
Duty	CCRA	

If  $f_{SYS}=8\text{MHz}$ , CTM clock source is  $f_{SYS}/4$ ,  $CCRP=4$  and  $CCRA=128$ ,

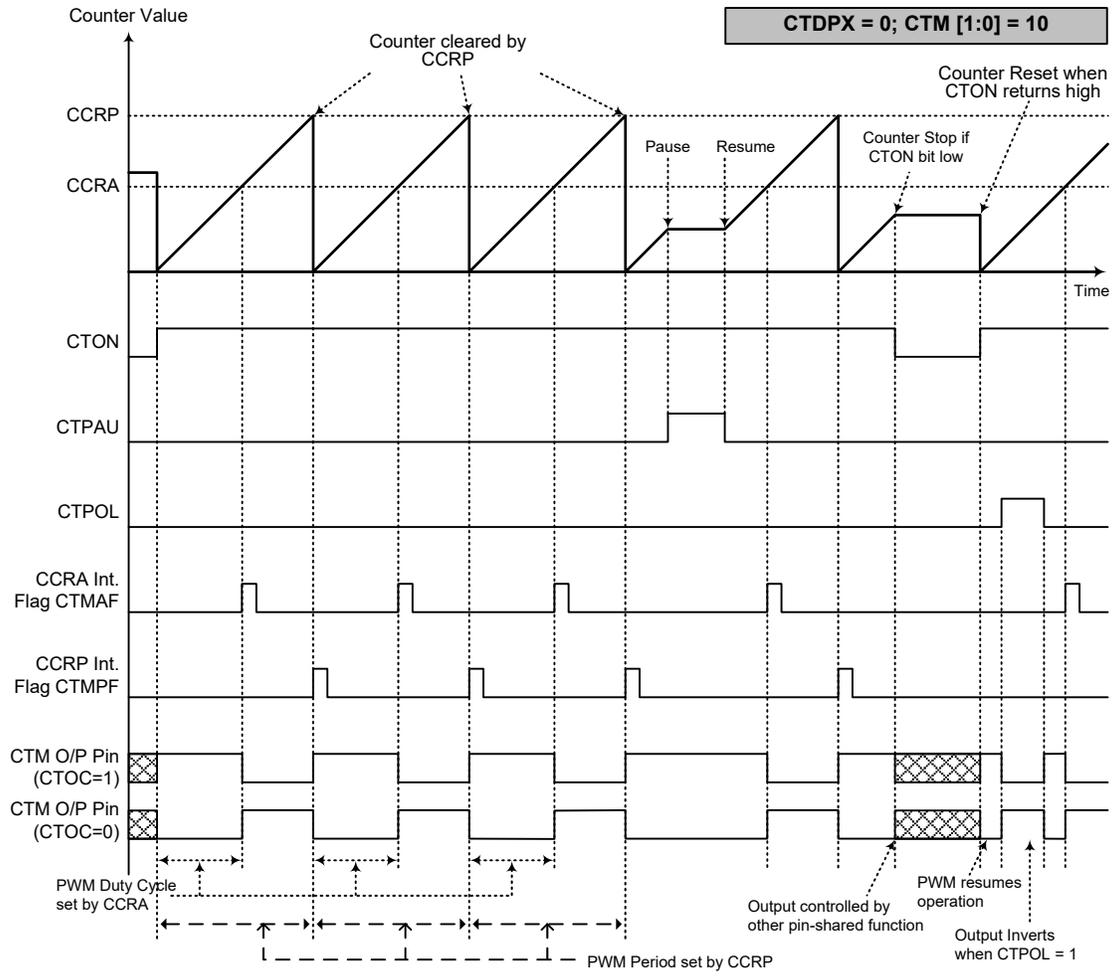
The CTM PWM output frequency= $(f_{SYS}/4)/(4\times 128)=f_{SYS}/2048=4\text{kHz}$ , duty= $128/(4\times 128)=25\%$ .

If the Duty value defined by the CCRA register is equal to or greater than the Period value, then the PWM output duty is 100%.

• **10-bit CTM, PWM Output Mode, Edge-aligned Mode, CTDPX=1**

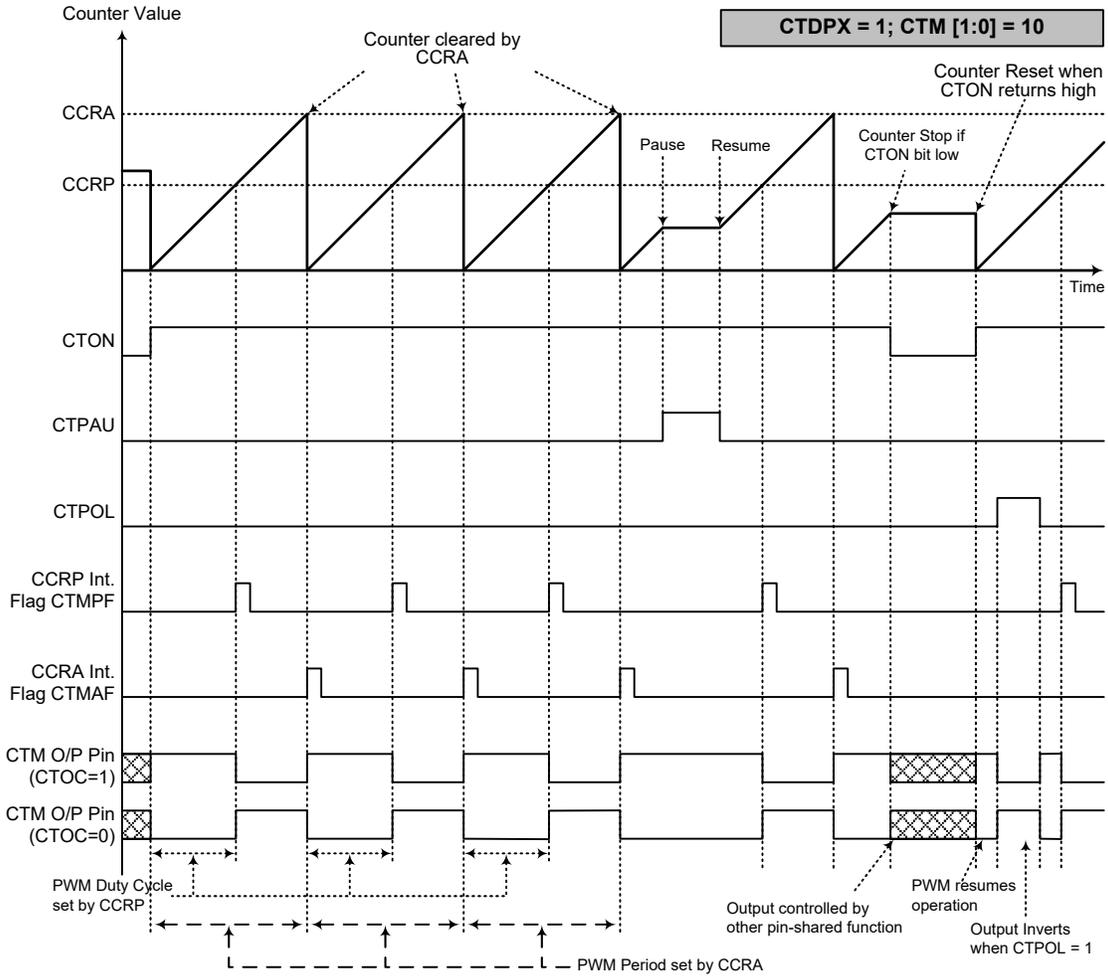
CCRP	1~7	0
Period	CCRA	
Duty	CCRP×128	1024

The PWM output period is determined by the CCRA register value together with the CTM clock while the PWM duty cycle is defined by the CCRP register value except when the CCRP value is equal to 0.



**PWM Output Mode – CTDPX=0**

- Note: 1. Here CTDPX=0 – Counter cleared by CCRP  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues running even when CTIO[1:0]=00 or 01  
 4. The CTCCLR bit has no influence on PWM operation



**PWM Output Mode – CTD PX=1**

- Note: 1. Here CTD PX=1 – Counter cleared by CCRA  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues even when CTIO[1:0]=00 or 01  
 4. The CTCCLR bit has no influence on PWM operation

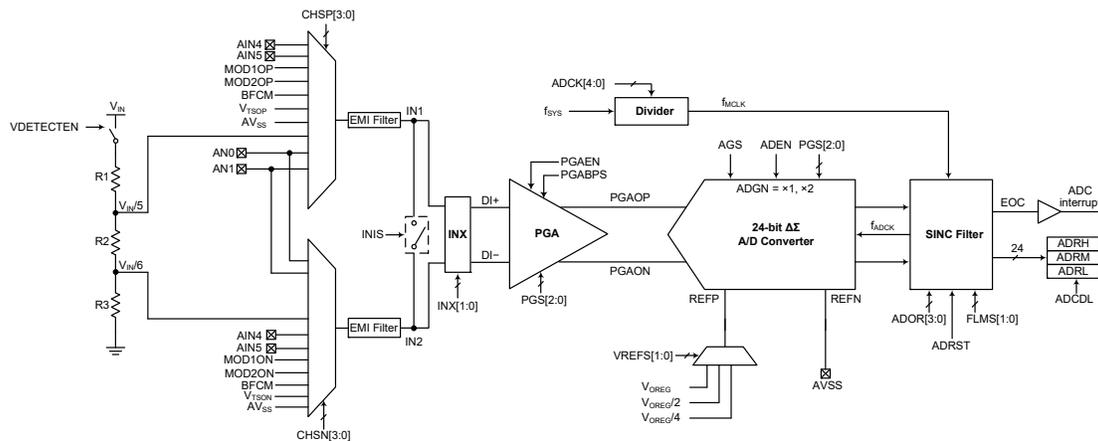
## Analog to Digital Converter

The need to interface to real world analog signals is a common requirement for many electronic systems. However, to properly process these signals by a microcontroller, they must first be converted into digital signals by A/D converters. By integrating the A/D conversion electronic circuitry into the microcontroller, the need for external components is reduced significantly with the corresponding follow-on benefits of lower costs and reduced component space requirements.

### A/D Converter Overview

The device contains a high accuracy multi-channel 24-bit Delta Sigma analog-to-digital converter which can directly interface to external analog signals, such as that from sensors or other control signals and convert these signals directly into a 24-bit digital value.

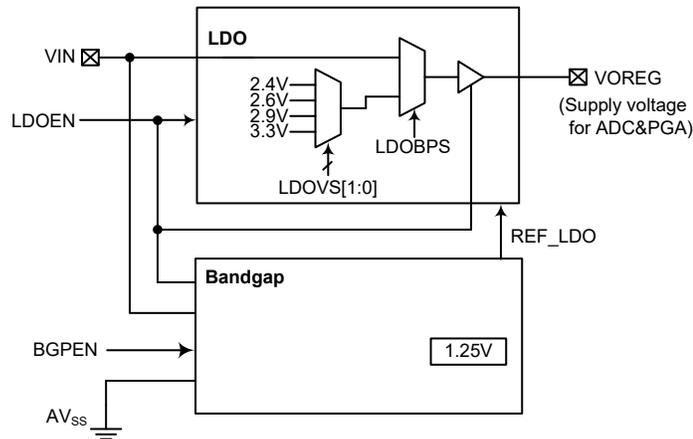
In addition, PGA gain control, A/D converter gain control and A/D converter reference gain control determine the amplification gain for A/D converter input signal. The designer can select the best gain combination for the desired amplification applied to the input signal. The following block diagram illustrates the A/D converter basic operational function. The A/D converter input channel can be arranged as 4 single-ended A/D converter input channels or 2 differential external input channels. The input signal can be amplified by PGA before entering the 24-bit Delta Sigma A/D converter. The A/D converter module will output one bit converted data to SINC filter which can transform the converted one-bit data to 24 bits and store them into the specific data registers. Additionally, the device also provides a temperature sensor to compensate the A/D converter deviation caused by the temperature. With high accuracy and performance, the device is very suitable for differential output sensor applications such as weight measurement scales and other related products.



**A/D Converter Structure**

### Internal Power Supply

The device contains an LDO for the regulated power supply. The accompanying block diagram illustrates the basic functional operation. The internal LDO can provide a fixed voltage for the PGA, A/D converter or external components. There are four LDO voltage levels, 2.4V, 2.6V, 2.9V or 3.3V, determined by the  $LDOVS1 \sim LDOVS0$  bits in the  $PWRC$  register. The LDO function can be controlled by the  $LDOEN$  and can be powered off to reduce overall power consumption.



**Internal Power Supply Block Diagram**

• **PWRC Register**

Bit	7	6	5	4	3	2	1	0
Name	LDOEN	BGPEN	—	VDETECTEN	TSEN	LDOBPS	LDOVS1	LDOVS0
R/W	R/W	R/W	—	R/W	R/W	R/W	R/W	R/W
POR	0	0	—	0	0	0	0	0

- Bit 7      **LDOEN**: LDO function control  
 0: Disable  
 1: Enable  
  
 If the LDO is disabled, there will be no power consumption and the LDO output pin will remain at a low level using a weak internal pull-low resistor.
- Bit 6      **BGPEN**: Bandgap enable control  
 0: Disable  
 1: Enable  
  
 Note: The BGPEN bit is the total band gap switch of the A/D Converter block, and all A/D Converter related functions need to be set under BGPEN=1.
- Bit 5      Unimplemented, read as “0”
- Bit 4      **VDETECTEN**: Voltage divider  $V_{IN}/5$  and  $V_{IN}/6$  enable/disable control  
 0: Disable  
 1: Enable
- Bit 3      **TSEN**: Internal temperature enable/disable control  
 0: Disable  
 1: Enable
- Bit 2      **LDOBPS**: LDO Bypass function control  
 0: Disable  
 1: Enable
- Bit 1~0    **LDOVS1~LDOVS0**: LDO output voltage selection  
 00: 2.4V  
 01: 2.6V  
 10: 2.9V  
 11: 3.3V

**A/D Converter Data Rate Definition**

The Delta Sigma A/D converter data rate can be calculated using the following equation:

Data Rate for SINC2  

$$= f_{ADCK} / (2 \times OSR)$$

$$= (f_{MCLK}/N) / (2 \times OSR)$$

$$= f_{MCLK} / (N \times 2 \times OSR)$$

Data Rate for SINC3

$$= f_{ADCK} / OSR$$

$$= (f_{MCLK}/N) / OSR$$

$$= f_{MCLK} / (N \times OSR)$$

$f_{ADCK}$ : A/D converter clock frequency, derived from  $f_{MCLK}/N$ .

$f_{MCLK}$ : A/D converter clock source, derived from  $f_{SYS}$  or  $f_{SYS}/2/(ADCK+1)$  using the ADCK bit field.

N: A constant divide factor equal to 12 or 30 determined by the FLMS bit field.

OSR: Oversampling rate determined by the ADOR bit field.

For example, if a data rate of 8Hz is desired, an  $f_{MCLK}$  clock source with a frequency of 4MHz A/D converter can be selected. Then set the FLMS field to “00” to obtain an “N” equal to 30.

For the SINC2 filter, set the ADOR field to “0010” to select an oversampling rate equal to 8192.

Therefore, the Data Rate = 4MHz / (30 × 2 × 8192) = 8Hz.

For the SINC3 filter, set the ADOR field to “0001” to select an oversampling rate equal to 16384.

Therefore, the Data Rate = 4MHz / (30 × 16384) = 8Hz.

## A/D Converter Register Description

Overall operation of the A/D converter is controlled by using a series of registers. Three read only registers exist to store the A/D converter data 24-bit value. A control register named as PWRC is used to control the required bias and supply voltages for PGA and A/D converter and is described in the “Internal Power Supply” section. The remaining registers are control registers which set up the gain selections and control functions of the A/D converter.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PWRC	LDOEN	BGPEN	—	VDETECTEN	TSEN	LDOBPS	LDOVS1	LDOVS0
PGAC0	VREFS1	VREFS0	AGS	PGAEN	PGABPS	PGS2	PGS1	PGS0
PGAC1	INIS	INX1	INX0	—	—	—	—	—
PGACS	CHSN3	CHSN2	CHSN1	CHSN0	CHSP3	CHSP2	CHSP1	CHSP0
ADRL	D7	D6	D5	D4	D3	D2	D1	D0
ADRM	D15	D14	D13	D12	D11	D10	D9	D8
ADRH	D23	D22	D21	D20	D19	D18	D17	D16
ADCR0	SINCS	D6	FLMS1	FLMS0	ADOR3	ADOR2	ADOR1	ADOR0
ADCR1	—	—	ADCDL	EOC	ADRST	ADEN	—	—
ADCS	—	—	—	ADCK4	ADCK3	ADCK2	ADCK1	ADCK0

**A/D Converter Register List**

### Programmable Gain Amplifier Registers – PGAC0, PGAC1, PGACS

There are three registers related to the programmable gain control, PGAC0, PGAC1 and PGACS. The PGAC0 register is used to select the PGA gain, A/D Converter gain and the A/D Converter reference gain. The PGAC1 register is used to define the input connection and differential input offset voltage adjustment control. In addition, the PGACS register is used to select the PGA inputs. Therefore, the input channels have to be determined by the CHSP3~CHSP0 and CHSN3~CHSN0 bits to determine which analog channel input pins, temperature detector inputs or internal power supply are actually connected to the internal differential A/D converter.

• **PGAC0 Register**

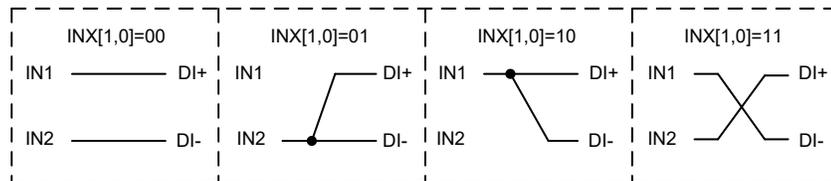
Bit	7	6	5	4	3	2	1	0
Name	VREFS1	VREFS0	AGS	PGAEN	PGABPS	PGS2	PGS1	PGS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7~6 **VREFS1~VREFS0**: A/D converter reference voltage pair selection  
 00: Internal reference voltage pair –  $V_{OREG}$  &  $AV_{SS}$   
 01: Internal reference voltage pair –  $V_{OREG} \times 1/2$  &  $AV_{SS}$   
 10: Internal reference voltage pair –  $V_{OREG} \times 1/4$  &  $AV_{SS}$   
 11: Reserved
- Bit 5 **AGS**: A/D converter PGAOP/PGAON differential input signal gain selection  
 0: ADGN = 1  
 1: ADGN = 2
- Bit 4 **PGAEN**: PGA enable control  
 0: Disable  
 1: Enable  
 This bit controls the overall enable/disable function of the PGA circuits. The related functions will only be available when the PGAEN is set high.
- Bit 3 **PGABPS**: PGA bypass control  
 0: Normal mode  
 1: Bypass PGA mode
- Bit 2~0 **PGS2~PGS0**: PGA DI+/DI- differential channel input gain selection  
 000: PGAGN = 1  
 001: PGAGN = 2  
 010: PGAGN = 4  
 011: PGAGN = 8  
 100: PGAGN = 16  
 101: PGAGN = 32  
 110: PGAGN = 64  
 111: PGAGN = 128

• **PGAC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	INIS	INX1	INX0	—	—	—	—	—
R/W	R/W	R/W	R/W	—	—	—	—	—
POR	0	0	0	—	—	—	—	—

- Bit 7 **INIS**: The selected input ends, IN1/IN2, internal connection control  
 0: Not connected  
 1: Connected
- Bit 6~5 **INX1~INX0**: The selected input ends, IN1/IN2, and the PGA differential input ends, DI+/DI-, connection control



- Bit 4~0 Unimplemented, read as “0”

• **PGACS Register**

Bit	7	6	5	4	3	2	1	0
Name	CHSN3	CHSN2	CHSN1	CHSN0	CHSP3	CHSP2	CHSP1	CHSP0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~4 **CHSN3~CHSN0**: Negative input end IN2 selection

- 0000: AN0
- 0001: AN1
- 0010~0111: Reserved
- 1000: MOD1ON
- 1001: MOD2ON
- 1010: Temperature sensor output –  $V_{TSON}$
- 1011: AIN4
- 1100: AIN5
- 1101:  $V_{IN}/6$
- 1110: BFCM
- 1111:  $AV_{SS}$

These bits are used to select the negative input, IN2. It is recommended that when the  $V_{TSON}$  signal is selected as the negative input, the  $V_{TSOP}$  signal should be selected as the positive input for proper operation.

Bit 3~0 **CHSP3~CHSP0**: Positive input end IN1 selection

- 0000: AN0
- 0001: AN1
- 0010~0111: Reserved
- 1000: MOD1OP
- 1001: MOD2OP
- 1010: Temperature sensor output –  $V_{TSOP}$
- 1011: AIN4
- 1100: AIN5
- 1101:  $V_{IN}/5$
- 1110: BFCM
- 1111:  $AV_{SS}$

These bits are used to select the positive input, IN1. It is recommended that when the  $V_{TSOP}$  signal is selected as the positive input, the  $V_{TSON}$  signal should be selected as the negative input for proper operation.

**A/D Converter Data Registers – ADRL, ADRM, ADRH**

The 24-bit Delta Sigma A/D converter requires three data registers to store the converted value. These are a high byte register, known as ADRH, a middle byte register, known as ADRM, and a low byte register, known as ADRL. After the conversion process takes place, these registers can be directly read by the microcontroller to obtain the digitised conversion value. D0~D23 are the A/D conversion result data bits.

• **ADRL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R	R	R	R	R	R	R	R
POR	x	x	x	x	x	x	x	x

“x”: unknown

Bit 7~0 **D7~D0**: A/D conversion data register bit 7~bit 0

• **ADRM Register**

Bit	7	6	5	4	3	2	1	0
Name	D15	D14	D13	D12	D11	D10	D9	D8
R/W	R	R	R	R	R	R	R	R
POR	x	x	x	x	x	x	x	x

"x": unknown

Bit 7~0     **D15~D8**: A/D conversion data register bit 15~bit 8

• **ADRH Register**

Bit	7	6	5	4	3	2	1	0
Name	D23	D22	D21	D20	D19	D18	D17	D16
R/W	R	R	R	R	R	R	R	R
POR	x	x	x	x	x	x	x	x

"x": unknown

Bit 7~0     **D23~D16**: A/D conversion data Register bit 23~bit 16

**A/D Converter Control Registers – ADCR0, ADCR1, ADCS**

To control the function and operation of the A/D converter, three control registers known as ADCR0, ADCR1 and ADCS are provided. These 8-bit registers define functions such as the selection of which reference source is used by the internal A/D converter, the A/D converter clock source, the A/D converter output data rate as well as controlling the power-up function and monitoring the A/D converter end of conversion status.

• **ADCR0 Register**

Bit	7	6	5	4	3	2	1	0
Name	SINCS	D6	FLMS1	FLMS0	ADOR3	ADOR2	ADOR1	ADOR0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7     **SINCS**: SINC selection

0: SINC3 filter

1: SINC2 filter

Bit 6     **D6**: Reserved

Bit 5~4     **FLMS1~FLMS0**: A/D converter clock  $f_{ADCK}$  selection

00:  $f_{ADCK} = f_{MCLK}/30, N = 30$

10:  $f_{ADCK} = f_{MCLK}/12, N = 12$

Others: Reserved

Bit 3~0     **ADOR3~ADOR0**: A/D conversion oversampling rate selection

0000: Oversampling rate OSR = 32768

0001: Oversampling rate OSR = 16384

0010: Oversampling rate OSR = 8192

0011: Oversampling rate OSR = 4096

0100: Oversampling rate OSR = 2048

0101: Oversampling rate OSR = 1024

0110: Oversampling rate OSR = 512

0111: Oversampling rate OSR = 256

1000: Oversampling rate OSR = 128

Others: Reserved

**• ADCR1 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	ADCDL	EOC	ADRST	ADEN	—	—
R/W	—	—	R/W	R/W	R/W	R/W	—	—
POR	—	—	0	0	0	0	—	—

Bit 7~6 Unimplemented, read as “0”

Bit 5 **ADCDL**: A/D converted data latch function enable control

0: Disable  
1: Enable

If the A/D converted data latch function is enabled, the latest converted data value will be latched and will not be updated by any subsequent converted results until this function is disabled. Although the converted data is latched into the data registers, the A/D converter circuits remain operational, but will not generate an interrupt and the EOC will not change. It is recommended that this bit should be set high before reading the converted data from the ADRL, ADRM and ADRH registers. After the converted data has been read out, the bit can then be cleared to zero to disable the A/D converter data latch function and allow further conversion values to be stored. In this way, the possibility of obtaining undesired data during A/D converter conversions can be prevented.

Bit 4 **EOC**: End of A/D conversion flag

0: A/D conversion in progress  
1: A/D conversion ended

This bit must be cleared by software.

Bit 3 **ADRST**: A/D converter software reset enable control

0: Disable  
1: Enable

This bit is used to reset the A/D converter internal digital SINC filter. This bit is cleared to zero for A/D normal operations. However, if set high, the internal digital SINC filter will be reset and the current A/D converted data will be aborted. A new A/D data conversion process will not be initiated until this bit is cleared to zero again.

Bit 2 **ADEN**: A/D converter function enable control

0: Disable  
1: Enable

This bit controls the A/D converter function enable/disable. This bit should be set high to enable the A/D converter. If the bit is cleared to zero then the A/D converter will be switched off reducing the device power consumption. As the A/D converter will consume a limited amount of power, even when not executing a conversion, this may be an important consideration in power sensitive battery powered applications.

Bit 1~0 Unimplemented, read as “0”

**• ADCS Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	ADCK4	ADCK3	ADCK2	ADCK1	ADCK0
R/W	—	—	—	R/W	R/W	R/W	R/W	R/W
POR	—	—	—	0	0	0	0	0

Bit 7~5 Unimplemented, read as “0”

Bit 4~0 **ADCK4~ADCK0**: A/D converter clock source  $f_{MCLK}$  divided ratio selection

00000~11110:  $f_{MCLK}=f_{SYS}/2 / (ADCK[4:0]+1)$   
11111:  $f_{MCLK}=f_{SYS}$

**A/D Converter Operation**

The ADRST bit in the ADCR1 register is used to start and reset the A/D converter after power on. When the microcontroller changes this bit from low to high and then low again, an analog to digital

conversion in the SINC filter will be initiated. After this setup is completed, the A/D Converter is ready for operation. This bit is used to control the overall start operation of the internal analog to digital converter.

The EOC bit in the ADCR1 register is used to indicate when the analog to digital conversion process is complete. This bit will be automatically set high by the microcontroller after a conversion cycle has ended. In addition, the corresponding A/D converter interrupt request flag will be set in the interrupt control register, and if the interrupts are enabled, an appropriate internal interrupt signal will be generated. This A/D converter internal interrupt signal will direct the program flow to the associated A/D converter internal interrupt address for processing. If the A/D converter internal interrupt is disabled, the microcontroller can poll the EOC bit in the ADCR1 register to check whether it has been set to “1” as an alternative method of detecting the end of an A/D conversion cycle. The A/D converted data will be updated continuously by the new converted data. If the A/D converted data latch function is enabled, the latest converted data will be latched and the following new converted data will be discarded until this data latch function is disabled.

The clock source for the A/D converter should be typically fixed at a value of 4MHz, which originates from the system clock  $f_{SYS}$ , and can be chosen to be either  $f_{SYS}$  or a subdivided version of  $f_{SYS}$ . The division ratio value is determined by the ADCK4~ADCK0 bits in the ADCS register to obtain a 4MHz clock source for the A/D Converter.

The differential reference voltage supply to the A/D Converter can be supplied from the  $V_{OREG}$  and  $AV_{SS}$ , or from a subdivided version of  $V_{OREG}$  and  $AV_{SS}$ . The desired selection is made using the VREFS1~VREFS0 bits in the PGAC0 register.

### **Summary of A/D Conversion Steps**

The following summarises the individual steps that should be executed in order to implement an A/D conversion process.

- Step 1  
Enable the power LDO for PGA and ADC.
- Step 2  
Select the PGA gain and the A/D converter gain as well as the A/D converter reference voltage pair by the PGAC0 register.
- Step 3  
Select the PGA settings for input connection and input offset by PGAC1 register.
- Step 4  
Select the required A/D conversion clock source by correctly programming bits ADCK4~ADCK0 in the ADCS register.
- Step 5  
Select output data rate by configuring the ADOR3~ADOR0 and FLMS1~FLMS0 bits in the ADCR0 register.
- Step 6  
Select which channel is to be connected to the internal PGA by correctly programming the CHSP3~CHSP0 and CHSN3~CHSN0 bits which are also contained in the PGACS register.
- Step 7  
Enable the A/D converter by setting the ADEN bit in the ADCR1 register high.

- Step 8  
Reset the A/D converter by setting the ADRST bit in the ADCR1 register high and then clearing this bit to zero to release the reset status.
  - Step 9  
If the A/D converter interrupt is to be used, the related interrupt control registers must be correctly configured to ensure the A/D converter interrupt function is active. The master interrupt control bit, EMI, and the A/D converter interrupt bit, ADE, must both be set high to do this.
  - Step 10  
To check when the analog to digital conversion process is complete, the EOC bit in the ADCR1 register can be polled. The conversion process is complete when this bit goes high. When this occurs the A/D converter data registers ADRL, ADRM and ADRH can be read to obtain the conversion value. As an alternative method, if the interrupts are enabled and the stack is not full, the program can wait for an A/D converter interrupt to occur.
- Note: When checking for the end of the conversion process, if the method of polling the EOC bit in the ADCR1 register is used, the interrupt enable step above can be omitted.

### Programming Considerations

During microcontroller operations where the A/D converter is not being used, the A/D converter internal circuitry can be switched off to reduce power consumption, by clearing the ADEN bit in the ADCR1 register to zero. When this happens, the internal A/D converter circuits will not consume power irrespective of what analog voltage is applied to their input lines.

### A/D Converter Transfer Function

The device contains a 24-bit Delta Sigma A/D converter and its full-scale converted digitized value is from 8388607 to -8388608 in decimal value. The converted data format is formed using a two's complement binary value. The MSB of the converted data is the signed bit. Since the full-scale analog input value is equal to the differential reference input voltage,  $\Delta VR_{\pm}$ , selected by the VREFS1~VREFS0 bits in the PGAC0 register, this gives a single bit analog input value of  $\Delta VR_{\pm}$  divided by 8388608.

$$1 \text{ LSB} = \Delta VR_{\pm} / 8388608$$

The A/D Converter input voltage value can be calculated using the following equation:

$$\Delta SI\_I = PGAGN \times ADGN \times \Delta DI_{\pm}$$

$$ADC\_Conversion\_Data = (\Delta SI\_I / \Delta VR_{\pm}) \times K$$

Where K is equal to  $2^{23}$

Note: 1. The PGAGN and ADGN values are decided by the PGS[2:0] and AGS control bits

2.  $\Delta SI\_I$ : Differential Input Signal after amplification
3. PGAGN: Programmable Gain Amplifier gain
4. ADGN: A/D Converter gain
5.  $\Delta DI_{\pm}$ : Differential input signal derived from external channels or internal signals
6.  $\Delta VR_{\pm}$ : Differential Reference voltage

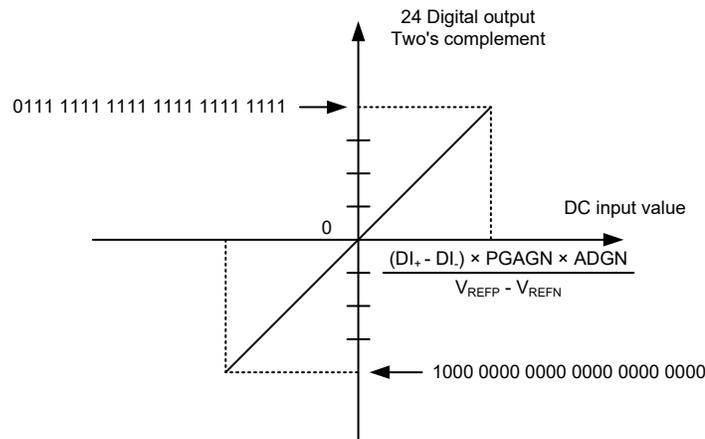
Due to the digital system design of the Delta Sigma A/D Converter, the maximum A/D converted value is 8388607 and the minimum value is -8388608. Therefore, there is a middle value of 0. The ADC\_Conversion\_Data equation illustrates this range of converted data variation.

A/D Conversion Data (2's complement, Hexadecimal)	Decimal Value
0x7FFFFFFF	8388607
0x800000	-8388608

**A/D Conversion Data Range**

The above A/D conversion data table illustrates the range of A/D conversion data.

The following diagram shows the relationship between the DC input value and the A/D converted data which is presented using Two's Complement.



**A/D Converted Data**

The A/D converted data is related to the input voltage and the PGA selections. The format of the A/D Converter output is a two's complement binary code. The length of this output code is 24 bits and the MSB is a signed bit. When the MSB is “0”, this represents a “positive” input. If the MSB is “1”, this represents a “negative” input. The maximum value is 8388607 and the minimum value is -8388608. If the input signal is greater than the maximum value, the converted data is limited to 8388607, and if the input signal is less than the minimum value, the converted data is limited to -8388608.

**A/D Converted Data to Voltage**

The converted data can be recovered using the following equations:

If MSB = 0 – Positive Converted data

$$\text{Input Voltage} = \frac{(\text{Converted\_data}) \times \text{LSB}}{\text{PGAGN} \times \text{ADGN}}$$

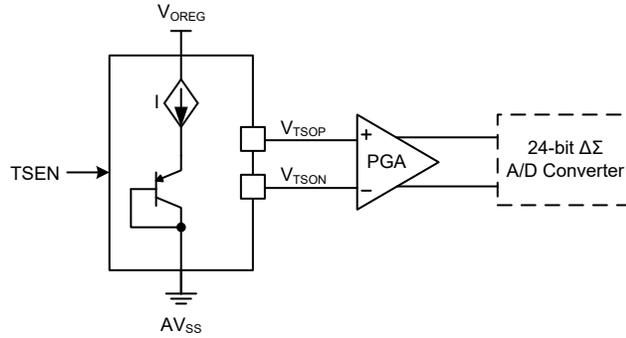
If the MSB = 1 – Negative Converted data

$$\text{Input voltage} = \frac{(\text{Two's\_complement\_of\_Converted\_data}) \times \text{LSB}}{\text{PGAGN} \times \text{ADGN}}$$

Note: Two's complement = One's complement + 1

**Temperature Sensor**

An internal temperature sensor is integrated within the device to allow compensation for temperature effects. By selecting the PGA input channels to the  $V_{TSOP}$  and  $V_{TSON}$  signals, the A/D Converter can obtain temperature information and allow compensation to be carried out on the A/D converted data. The following block diagram illustrates the functional operation for the temperature sensor.



**Temperature Sensor Structure**

## Body Fat Measurement Function

This device include a sine generator, two operational amplifiers, a filter, three comparators and three 6-bit D/A converters. It is high quality, flexibility and high integration for body fat measurement. The whole module power is from an LDO.

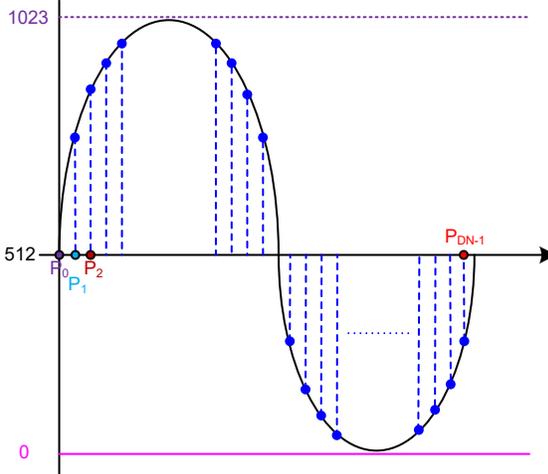
### Sine Wave Generator

The sine generator consists of a frequency divider, counter, sine pattern Data memory, 10-bit D/A converter and OP0. It offers the wide range 1kHz~500kHz sine wave generator and 64×10 bits of Data memory for sine wave pattern by software setting. The frequency divider will multiply by DN/M to generate a clock to counter. The related details refer to following formula.

- System clock / M=sine wave frequency
- System clock × (DN/M)=counter count rate
- The M must be the multiple of N and 8.
- $M=N \times DN$
- DN: sine wave cycle data number ( $DN \leq 64$ )

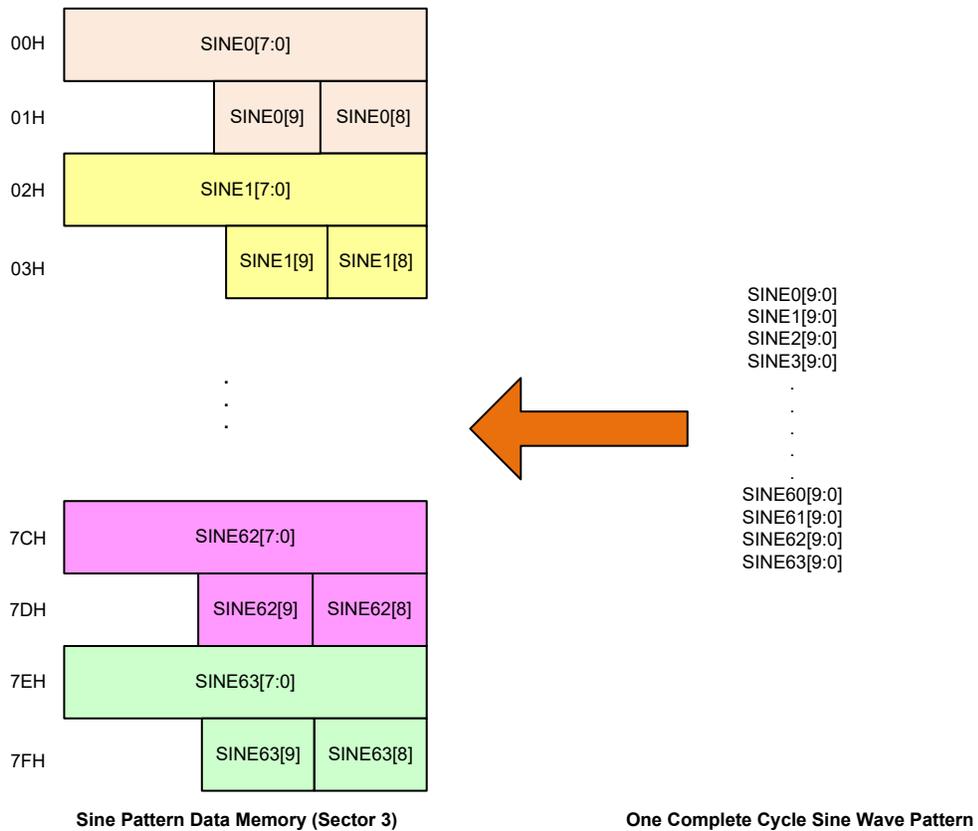
Refer to the following table and diagram for details.

System Frequency	4MHz				8MHz			
Sine Wave Frequency (kHz)	500	50	5	1	500	50	5	1
M	8	80	800	4000	16	160	1600	8000
N	1	2	20	100	1	4	25	125
DN	8	40	40	40	16	40	64	64



Users only need to generate a sine wave pattern  $P_0 \sim P_{DN-1}$  which is stored in Data memory sector 3 (00H~7FH). The sine pattern [7:0] is stored using even addresses and the sine pattern[8]/pattern[9] is stored using odd addresses. Once the sine generator is enabled, the CPU cannot write or read any data to or from this Data memory. The sine generator will then read data from this Data memory and transmit it to a 10-bit D/A converter.

The controller reads a complete cycle sine pattern from Data memory and generates a sine waveform on the SIN pin. Refer to the following diagram.



## Body Fat Measurement Registers

There are a series of registers control the overall operation of the body fat measurement function.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SGC	SGEN	BREN	—	—	—	—	EXTSYNC	SGIQRS
SGN	D7	D6	D5	D4	D3	D2	D1	D0
SGDN	—	—	D5	D4	D3	D2	D1	D0
OPAC	OPAEN	IQ_FWR	OP1BW1	OP1BW0	OP2G3	OP2G2	OP2G1	OP2G0
SWC0	SW7	SW6	SW5	SW4	SW3	SW2	SW1	SW0
SWC1	SWF	SWE	SWD	SWA	SW9A	SW8A	SW9	SW8
SWC2	SWC	SWN	SWMA	SWL	SWK	SWH	SWG	SWB
SWC3	IQMOD2S1	IQMOD2S0	IQMOD1S1	IQMOD1S0	OP2NIS1	OP2NIS0	OP2PIS1	OP2PIS0
CMPC0	CMP0EN	—	CMP0F	CMP0O	CMP1EN	—	CMP1F	CMP1O
CMPC1	—	—	—	—	CMP2EN	—	CMP2F	CMP2O
BDA0C	BDA0EN	—	D5	D4	D3	D2	D1	D0
BDA1C	BDA1EN	—	D5	D4	D3	D2	D1	D0
BDA2C	BDA2EN	—	D5	D4	D3	D2	D1	D0

**Body Fat Measurement Register List**

## Sine Wave Generator

The sine wave generator is controlled by three registers. Details are given as below.

### • SGC Register

Bit	7	6	5	4	3	2	1	0
Name	SGEN	BREN	—	—	—	—	EXTSYNC	SGIQRS
R/W	R/W	R/W	—	—	—	—	R/W	R/W
POR	0	0	—	—	—	—	0	0

Bit 7 **SGEN**: Sine generator enabled  
0: Disable  
1: Enable

When this bit is equal to “0”, the OP0 and 10-bit D/A converter will be in a power down mode.

Bit 6 **BREN**: Bias resistors control bit  
0: Disable – power down mode  
1: Enable – normal mode

Bit 5~2 Unimplemented, read as “0”

Bit 1 **EXTSYNC**: SYNC mode selection bit  
0: Internal sync mode  
1: External sync mode

Bit 0 **SGIQRS**: Sine generator & IQ circuit reset  
0: Normal  
1: Reset

The bit is normally low but if set high and then cleared low again, it will generate a reset to the IQ circuit and sine wave generator, and restart from the first data.

• **SGN Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: Sine generator data  
 The multiplier of system frequency (N)  
 The multiplier (N) is equal to  $D[7:0]+1$

• **SGDN Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	D5	D4	D3	D2	D1	D0
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	0	0	0	0

Bit 7~6 Unimplemented, read as “0”  
 Bit 5~0 **D5~D0**: Data number of sample  
 The data number of one sine wave cycle stored in Data memory Sector 3. DN is equal to  $D[5:0]+1$ .

**Operational Amplifiers and Comparators**

Several registers are associated with the operation of the amplifier part. The OPAC register is used for controlling the operational amplifier. The SWC0, SWC1, SWC2 and SWC3 registers are used for configuring the switch condition. The remaining CMPC0, CMPC1, BDA0C, BDA1C and BDA2C registers are used to determine whether the detected human body has already stood on the electrode.

• **OPAC Register**

Bit	7	6	5	4	3	2	1	0
Name	OPAEN	IQ_FWR	OP1BW1	OP1BW0	OP2G3	OP2G2	OP2G1	OP2G0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	1	0	0	0	0	0

Bit 7 **OPAEN**: Operational amplifier control bit  
 0: Disable  
 1: Enable  
 When this bit is equal to “0”, the OP1, OP2 and 6-bit D/A converter will be in a power down mode.

Bit 6 **IQ\_FWR**: IQ/FWR mode selection bit  
 0: IQ mode  
 1: FWR mode

Bit 5~4 **OP1BW1~OP1BW0**: OP1 bandwidth selection bits  
 00: 10MHz  
 01: 7.5MHz  
 10: 5MHz  
 11: 2.5MHz

Bit 3~0 **OP2G3~OP2G0**: OP2 gain control  
 0001: 1.14  
 0010: 1.31  
 0011: 1.5  
 0100: 1.73  
 0101: 2  
 0110: 2.33  
 0111: 2.75

1000: 3.285  
1001: 4  
1010: 5  
Others: 1

• **SWC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	SW7	SW6	SW5	SW4	SW3	SW2	SW1	SW0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **SW7**: Switch 7 control bit  
            0: Off  
            1: On
- Bit 6      **SW6**: Switch 6 control bit  
            0: Off  
            1: On
- Bit 5      **SW5**: Switch 5 control bit  
            0: Off  
            1: On
- Bit 4      **SW4**: Switch 4 control bit  
            0: Off  
            1: On
- Bit 3      **SW3**: Switch 3 control bit  
            0: Off  
            1: On
- Bit 2      **SW2**: Switch 2 control bit  
            0: Off  
            1: On
- Bit 1      **SW1**: Switch 1 control bit  
            0: Off  
            1: On
- Bit 0      **SW0**: Switch 0 control bit  
            0: Off  
            1: On

• **SWC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	SWF	SWE	SWD	SWA	SW9A	SW8A	SW9	SW8
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **SWF**: Switch F control bit  
            0: Off  
            1: On
- Bit 6      **SWE**: Switch E control bit  
            0: Off  
            1: On
- Bit 5      **SWD**: Switch D control bit  
            0: Off  
            1: On
- Bit 4      **SWA**: Switch A control bit  
            0: Off  
            1: On

- Bit 3      **SW9A**: Switch 9A control bit  
             0: Off  
             1: On
- Bit 2      **SW8A**: Switch 8A control bit  
             0: Off  
             1: On
- Bit 1      **SW9**: Switch 9 control bit  
             0: Off  
             1: On
- Bit 0      **SW8**: Switch 8 control bit  
             0: Off  
             1: On

• **SWC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	SWC	SWN	SWMA	SWL	SWK	SWH	SWG	SWB
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **SWC**: Switch C control bit  
             0: Off  
             1: On
- Bit 6      **SWN**: Switch N control bit  
             0: Off  
             1: On
- Bit 5      **SWMA**: Switch MA control bit  
             0: Off  
             1: On
- Bit 4      **SWL**: Switch L control bit  
             0: Off  
             1: On
- Bit 3      **SWK**: Switch K control bit  
             0: Off  
             1: On
- Bit 2      **SWH**: Switch H control bit  
             0: Off  
             1: On
- Bit 1      **SWG**: Switch G control bit  
             0: Off  
             1: On
- Bit 0      **SWB**: Switch B control bit  
             0: Off  
             1: On

• **SWC3 Register**

Bit	7	6	5	4	3	2	1	0
Name	IQMOD2S1	IQMOD2S0	IQMOD1S1	IQMOD1S0	OP2NIS1	OP2NIS0	OP2PIS1	OP2PIS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 **IQMOD2S1~IQMOD2S0**: MOD2OP & MOD2ON switch control bit

- 00: Floating
- 01: Floating
- 10: I to MOD2OP & MOD2ON
- 11: Q to MOD2OP & MOD2ON

Bit 5~4 **IQMOD1S1~IQMOD1S0**: MOD1OP & MOD1ON switch control bit

- 00: I to MOD1OP & MOD1ON
- 01: Q to MOD1OP & MOD1ON
- 10: Floating
- 11: Floating

Bit 3~2 **OP2NIS1~OP2NIS0**: OP2 negative input switch control

- 00: Floating
- 01: Pull-high to  $V_{OREG}$
- 10: Pull-low to GND
- 11: Floating

Bit 1~0 **OP2PIS1~OP2PIS0**: OP2 positive input switch control

- 00: Floating
- 01: Pull-high to  $V_{OREG}$
- 10: Pull-low to GND
- 11: Floating

• **CMPC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	CMP0EN	—	CMP0F	CMP0O	CMP1EN	—	CMP1F	CMP1O
R/W	R/W	—	R/W	R	R/W	—	R/W	R
POR	0	—	0	0	0	—	0	0

Bit 7 **CMP0EN**: Comparator 0 control bit

- 0: Disable
- 1: Enable

Bit 6 Unimplemented, read as “0”

Bit 5 **CMP0F**: Comparator 0 flag bit

- 0: Output no change
- 1: Output change (0 to 1, or 1 to 0)

Bit 4 **CMP0O**: Comparator 0 output bit

- 0: Output 0
- 1: Output 1

Bit 3 **CMP1EN**: Comparator 1 control bit

- 0: Disable
- 1: Enable

Bit 2 Unimplemented, read as “0”

Bit 1 **CMP1F**: Comparator 1 flag bit

- 0: Output no change
- 1: Output change (0 to 1, or 1 to 0)

Bit 0 **CMP1O**: Comparator 1 output bit

- 0: Output 0
- 1: Output 1

• **CMPC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	CMP2EN	—	CMP2F	CMP2O
R/W	—	—	—	—	R/W	—	R/W	R
POR	—	—	—	—	0	—	0	0

- Bit 7~4 Unimplemented, read as “0”
- Bit 3 **CMP2EN**: Comparator 2 control bit  
 0: Disable  
 1: Enable
- Bit 2 Unimplemented, read as “0”
- Bit 1 **CMP2F**: Comparator 2 flag bit  
 0: Output no change  
 1: Output change (0 to 1, or 1 to 0)
- Bit 0 **CMP2O**: Comparator 2 output bit  
 0: Output 0  
 1: Output 1

• **BDA0C Register**

Bit	7	6	5	4	3	2	1	0
Name	BDA0EN	—	D5	D4	D3	D2	D1	D0
R/W	R/W	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	—	0	0	0	0	0	0

- Bit 7 **BDA0EN**: D/A Converter 0 enable or disable control bit  
 0: Disable  
 1: Enable
- Bit 6 Unimplemented, read as “0”
- Bit 5~0 **D5~D0**: D/A Converter 0 output control code

• **BDA1C Register**

Bit	7	6	5	4	3	2	1	0
Name	BDA1EN	—	D5	D4	D3	D2	D1	D0
R/W	R/W	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	—	0	0	0	0	0	0

- Bit 7 **BDA1EN**: D/A Converter 1 enable or disable control bit  
 0: Disable  
 1: Enable
- Bit 6 Unimplemented, read as “0”
- Bit 5~0 **D5~D0**: D/A Converter 1 output control code

• **BDA2C Register**

Bit	7	6	5	4	3	2	1	0
Name	BDA2EN	—	D5	D4	D3	D2	D1	D0
R/W	R/W	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	—	0	0	0	0	0	0

- Bit 7 **BDA2EN**: D/A Converter 2 enable or disable control bit  
 0: Disable  
 1: Enable
- Bit 6 Unimplemented, read as “0”
- Bit 5~0 **D5~D0**: D/A Converter 2 output control code

## Serial Interface Module – SIM

The device contains a Serial Interface Module, which includes both the four-line SPI interface or two-line I<sup>2</sup>C interface types, to allow an easy method of communication with external peripheral hardware. Having relatively simple communication protocols, these serial interface types allow the microcontroller to interface to external SPI or I<sup>2</sup>C based hardware such as sensors, Flash or EEPROM memory, etc. The SIM interface pins are pin-shared with other I/O pins and therefore the SIM interface functional pins must first be selected using the corresponding pin-shared function selection bits. As both interface types share the same pins and registers, the choice of whether the SPI or I<sup>2</sup>C type is used is made using the SIM operating mode control bits, named SIM2~SIM0, in the SIMC0 register. These pull-high resistors of the SIM pin-shared I/O pins are selected using pull-high control registers when the SIM function is enabled and the corresponding pins are used as SIM input pins.

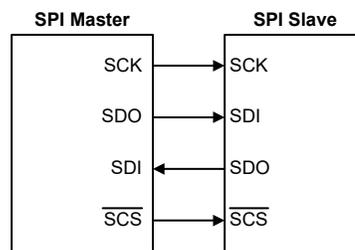
### SPI Interface

The SPI interface is often used to communicate with external peripheral devices such as sensors, Flash or EEPROM memory devices, etc. Originally developed by Motorola, the four line SPI interface is a synchronous serial data interface that has a relatively simple communication protocol simplifying the programming requirements when communicating with external hardware devices.

The communication is full duplex and operates as a slave/master type, where the device can be either master or slave. Although the SPI interface specification can control multiple slave devices from a single master, the device provided only one  $\overline{SCS}$  pin. If the master needs to control multiple slave devices from a single master, the master can use I/O pin to select the slave devices.

### SPI Interface Operation

The SPI interface is a full duplex synchronous serial data link. It is a four line interface with pin names SDI, SDO, SCK and  $\overline{SCS}$ . Pins SDI and SDO are the Serial Data Input and Serial Data Output lines, SCK is the Serial Clock line and  $\overline{SCS}$  is the Slave Select line. As the SPI interface pins are pin-shared with normal I/O pins and with the I<sup>2</sup>C function pins, the SPI interface pins must first be selected by configuring the pin-shared function selection bits and setting the correct bits in the SIMC0 and SIMC2 registers. After the desired SPI configuration has been set it can be disabled or enabled using the SIMEN bit in the SIMC0 register. Communication between devices connected to the SPI interface is carried out in a slave/master mode with all data transfer initiations being implemented by the master. The Master also controls the clock signal. As the device only contains a single  $\overline{SCS}$  pin only one slave device can be utilized. The  $\overline{SCS}$  pin is controlled by software, set CSEN bit to 1 to enable  $\overline{SCS}$  pin function, set CSEN bit to 0 the  $\overline{SCS}$  pin will be floating state.



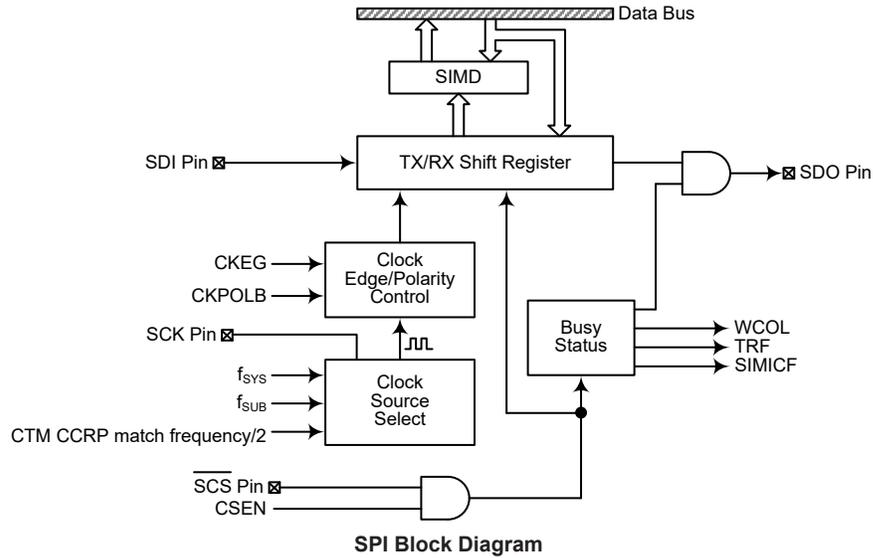
**SPI Master/Slave Connection**

The SPI function in this device offers the following features:

- Full duplex synchronous data transfer
- Both Master and Slave modes

- LSB first or MSB first data transmission modes
- Transmission complete flag
- Rising or falling active clock edge

The status of the SPI interface pins is determined by a number of factors such as whether the device is in the master or slave mode and upon the condition of certain control bits such as CSEN and SIMEN.



### SPI Registers

There are three internal registers which control the overall operation of the SPI interface. These are the SIMD data register and two registers SIMC0 and SIMC2.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SIMC0	SIM2	SIM1	SIM0	—	SIMDEB1	SIMDEB0	SIMEN	SIMICF
SIMC2	D7	D6	CKPOLB	CKEG	MLS	CSEN	WCOL	TRF
SIMD	D7	D6	D5	D4	D3	D2	D1	D0

**SPI Register List**

### SPI Data Register

The SIMD register is used to store the data being transmitted and received. The same register is used by both the SPI and I<sup>2</sup>C functions. Before the device writes data to the SPI bus, the actual data to be transmitted must be placed in the SIMD register. After the data is received from the SPI bus, the device can read it from the SIMD register. Any transmission or reception of data from the SPI bus must be made via the SIMD register.

#### • SIMD Register

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

“x”: unknown

Bit 7~0     **D7~D0**: SIM data register bit 7 ~ bit 0

### SPI Control Registers

There are also two control registers for the SPI interface, SIMC0 and SIMC2. Note that the SIMC2 register also has the name SIMA which is used by the I<sup>2</sup>C function. The SIMC1 register is not used by the SPI function, only by the I<sup>2</sup>C function. Register SIMC0 is used to control the enable/disable function and to set the data transmission clock frequency. Register SIMC2 is used for other control functions such as LSB/MSB selection, write collision flag, etc.

#### • SIMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	SIM2	SIM1	SIM0	—	SIMDEB1	SIMDEB0	SIMEN	SIMICF
R/W	R/W	R/W	R/W	—	R/W	R/W	R/W	R/W
POR	1	1	1	—	0	0	0	0

Bit 7~5 **SIM2~SIM0**: SIM Operating Mode Control  
 000: SPI master mode; SPI clock is  $f_{SYS}/4$   
 001: SPI master mode; SPI clock is  $f_{SYS}/16$   
 010: SPI master mode; SPI clock is  $f_{SYS}/64$   
 011: SPI master mode; SPI clock is  $f_{SUB}$   
 100: SPI master mode; SPI clock is CTM CCRP match frequency/2  
 101: SPI slave mode  
 110: I<sup>2</sup>C slave mode  
 111: Non SIM function

These bits setup the overall operating mode of the SIM function. As well as selecting if the I<sup>2</sup>C or SPI function, they are used to control the SPI Master/Slave selection and the SPI Master clock frequency. The SPI clock is a function of the system clock but can also be chosen to be sourced from CTM and  $f_{SUB}$ . If the SPI Slave Mode is selected then the clock will be supplied by an external Master device.

Bit 4 Unimplemented, read as “0”

Bit 3~2 **SIMDEB1~SIMDEB0**: I<sup>2</sup>C Debounce Time Selection

The SIMDEB1~SIMDEB0 bits are only used in the I<sup>2</sup>C mode and the detailed definition is described in the I<sup>2</sup>C section.

Bit 1 **SIMEN**: SIM Enable Control

0: Disable  
 1: Enable

The bit is the overall on/off control for the SIM interface. When the SIMEN bit is cleared to zero to disable the SIM interface, the SDI, SDO, SCK and  $\overline{SCS}$ , or SDA and SCL lines will lose their SPI or I<sup>2</sup>C function and the SIM operating current will be reduced to a minimum value. When the bit is high the SIM interface is enabled. If the SIM is configured to operate as an SPI interface via the SIM2~SIM0 bits, the contents of the SPI control registers will remain at the previous settings when the SIMEN bit changes from low to high and should therefore be first initialised by the application program. If the SIM is configured to operate as an I<sup>2</sup>C interface via the SIM2~SIM0 bits and the SIMEN bit changes from low to high, the contents of the I<sup>2</sup>C control bits such as HTX and TXAK will remain at the previous settings and should therefore be first initialised by the application program while the relevant I<sup>2</sup>C flags such as HCF, HAAS, HBB, SRW and RXAK will be set to their default states.

Bit 0 **SIMICF**: SIM SPI slave mode Incomplete Transfer Flag

0: SIM SPI slave mode incomplete condition not occurred  
 1: SIM SPI slave mode incomplete condition occurred

This bit is only available when the SIM is configured to operate in an SPI slave mode. If the SPI operates in the slave mode with the SIMEN and CSEN bits both being set to 1 but the  $\overline{SCS}$  line is pulled high by the external master device before the SPI data transfer is completely finished, the SIMICF bit will be set to 1 together with the TRF bit. When this condition occurs, the corresponding interrupt will occur if the interrupt function is enabled. However, the TRF bit will not be set to 1 if the SIMICF bit is set to 1 by software application program.

• **SIMC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	CKPOLB	CKEG	MLS	CSEN	WCOL	TRF
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 Undefined bits

These bits can be read or written by the application program.

Bit 5 **CKPOLB**: SPI clock line base condition selection

0: The SCK line will be high when the clock is inactive.

1: The SCK line will be low when the clock is inactive.

The CKPOLB bit determines the base condition of the clock line, if the bit is high, then the SCK line will be low when the clock is inactive. When the CKPOLB bit is low, then the SCK line will be high when the clock is inactive.

Bit 4 **CKEG**: SPI SCK clock active edge type selection

CKPOLB=0

0: SCK is high base level and data capture at SCK rising edge

1: SCK is high base level and data capture at SCK falling edge

CKPOLB=1

0: SCK is low base level and data capture at SCK falling edge

1: SCK is low base level and data capture at SCK rising edge

The CKEG and CKPOLB bits are used to setup the way that the clock signal outputs and inputs data on the SPI bus. These two bits must be configured before data transfer is executed otherwise an erroneous clock edge may be generated. The CKPOLB bit determines the base condition of the clock line, if the bit is high, then the SCK line will be low when the clock is inactive. When the CKPOLB bit is low, then the SCK line will be high when the clock is inactive. The CKEG bit determines active clock edge type which depends upon the condition of CKPOLB bit.

Bit 3 **MLS**: SPI data shift order

0: LSB first

1: MSB first

This is the data shift select bit and is used to select how the data is transferred, either MSB or LSB first. Setting the bit high will select MSB first and low for LSB first.

Bit 2 **CSEN**: SPI  $\overline{SCS}$  pin control

0: Disable

1: Enable

The CSEN bit is used as an enable/disable for the  $\overline{SCS}$  pin. If this bit is low, then the  $\overline{SCS}$  pin will be disabled and placed into a floating condition. If the bit is high, the  $\overline{SCS}$  pin will be enabled and used as a select pin.

Bit 1 **WCOL**: SPI write collision flag

0: No collision

1: Collision

The WCOL flag is used to detect whether a data collision has occurred or not. If this bit is high, it means that data has been attempted to be written to the SIMD register during a data transfer operation. This writing operation will be ignored if data is being transferred. This bit can be cleared by the application program.

Bit 0 **TRF**: SPI Transmit/Receive complete flag

0: SPI data is being transferred

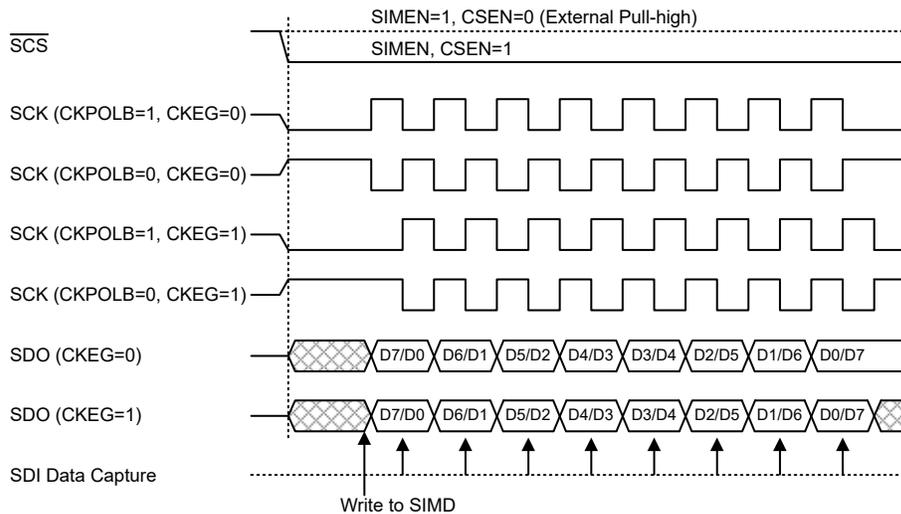
1: SPI data transfer is completed

The TRF bit is the Transmit/Receive Complete flag and is set to 1 automatically when an SPI data transfer is completed, but must be cleared to 0 by the application program. It can be used to generate an interrupt.

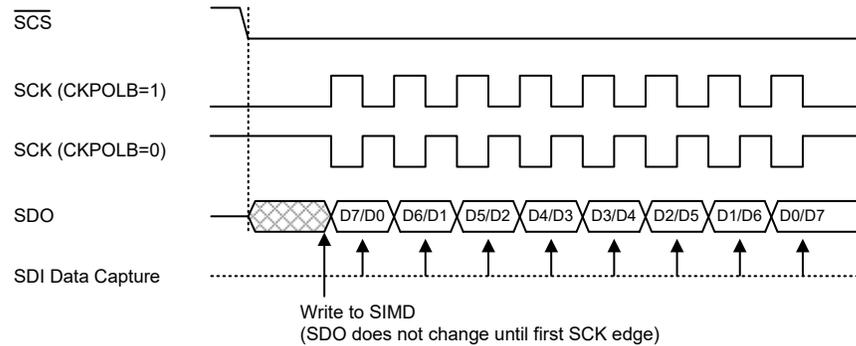
**SPI Communication**

After the SPI interface is enabled by setting the SIMEN bit high, then in the Master Mode, when data is written to the SIMD register, transmission/reception will begin simultaneously. When the data transfer is complete, the TRF flag will be set automatically, but must be cleared using the application program. In the Slave Mode, when the clock signal from the master has been received, any data in the SIMD register will be transmitted and any data on the SDI pin will be shifted into the SIMD register. The master should output a  $\overline{SCS}$  signal to enable the slave devices before a clock signal is provided. The slave data to be transferred should be well prepared at the appropriate moment relative to the SCK signal depending upon the configurations of the CKPOLB bit and CKEG bit. The accompanying timing diagram shows the relationship between the slave data and SCK signal for various configurations of the CKPOLB and CKEG bits.

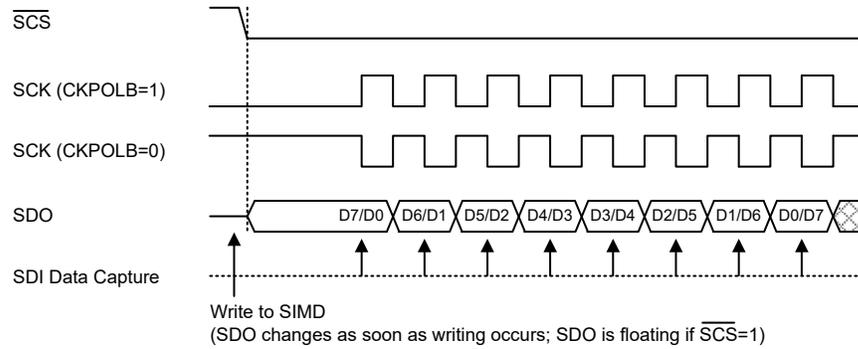
The SPI Master mode will continue to function if the SPI clock is running.



**SPI Master Mode Timing**

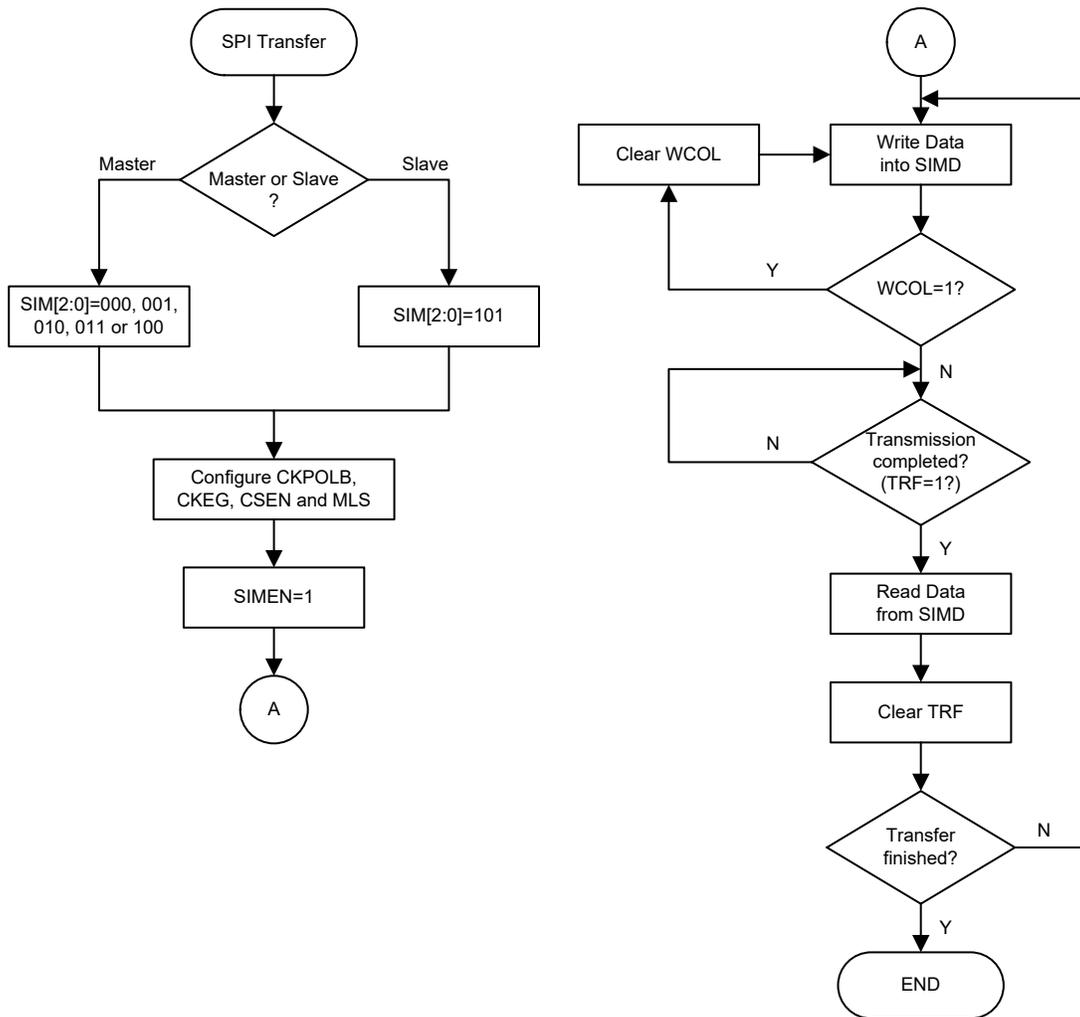


**SPI Slave Mode Timing – CKEG=0**



Note: For SPI slave mode, if SIMEN=1 and CSEN=0, SPI is always enabled and ignores the  $\overline{SCS}$  level.

**SPI Slave Mode Timing – CKEG=1**



**SPI Transfer Control Flow Chart**

**SPI Bus Enable/Disable**

To enable the SPI bus, set CSEN=1 and  $\overline{SCS}=0$ , then wait for data to be written into the SIMD (TXRX buffer) register. For the Master Mode, after data has been written to the SIMD (TXRX

buffer) register, then transmission or reception will start automatically. When all the data has been transferred, the TRF bit should be set. For the Slave Mode, when clock pulses are received on SCK, data in the TXRX buffer will be shifted out or data on SDI will be shifted in.

When the SPI bus is disabled, the SCK, SDI, SDO and  $\overline{\text{SCS}}$  can become I/O pins or other pin-shared functions using the corresponding pin-shared control bits.

#### SPI Operation Steps

All communication is carried out using the 4-line interface for either Master or Slave Mode.

The CSEN bit in the SIMC2 register controls the overall function of the SPI interface. Setting this bit high will enable the SPI interface by allowing the  $\overline{\text{SCS}}$  line to be active, which can then be used to control the SPI interface. If the CSEN bit is low, the SPI interface will be disabled and the  $\overline{\text{SCS}}$  line will be in a floating condition and can therefore not be used for control of the SPI interface. If the CSEN bit and the SIMEN bit in the SIMC0 are set high, this will place the SDI line in a floating condition and the SDO line high. If in Master Mode the SCK line will be either high or low depending upon the clock polarity selection bit CKPOLB in the SIMC2 register. If in Slave Mode the SCK line will be in a floating condition. If the SIMEN bit is low, then the bus will be disabled and the  $\overline{\text{SCS}}$ , SDI, SDO and SCK will all become I/O pins or the other functions using the corresponding pin-shared control bits. In the Master Mode the Master will always generate the clock signal. The clock and data transmission will be initiated after data has been written into the SIMD register. In the Slave Mode, the clock signal will be received from an external master device for both data transmission and reception. The following sequences show the order to be followed for data transfer in both Master and Slave Mode.

#### Master Mode

- Step 1  
Select the SPI Master mode and clock source using the SIM2~SIM0 bits in the SIMC0 control register.
- Step 2  
Setup the CSEN bit and setup the MLS bit to choose if the data is MSB or LSB first, this setting must be the same with the Slave devices.
- Step 3  
Setup the SIMEN bit in the SIMC0 control register to enable the SPI interface.
- Step 4  
For write operations: write the data to the SIMD register, which will actually place the data into the TXRX buffer. Then use the SCK and SDO lines to output the data. After this, go to step 5.  
For read operations: the data transferred in on the SDI line will be stored in the TXRX buffer until all the data has been received at which point it will be latched into the SIMD register.
- Step 5  
Check the WCOL bit if set high then a collision error has occurred so return to step 4. If equal to zero then go to the following step.
- Step 6  
Check the TRF bit or wait for a SIM SPI serial bus interrupt.
- Step 7  
Read data from the SIMD register.
- Step 8  
Clear TRF.

- Step 9  
Go to step 4.

#### **Slave Mode**

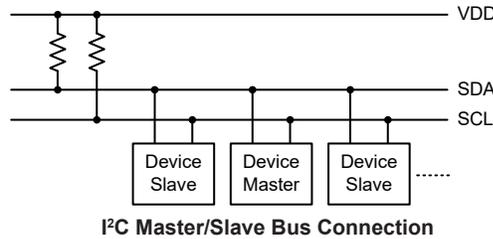
- Step 1  
Select the SPI Slave mode using the SIM2~SIM0 bits in the SIMC0 control register
- Step 2  
Setup the CSEN bit and setup the MLS bit to choose if the data is MSB or LSB first, this setting must be the same with the Master devices.
- Step 3  
Setup the SIMEN bit in the SIMC0 control register to enable the SPI interface.
- Step 4  
For write operations: write the data to the SIMD register, which will actually place the data into the TXRX buffer. Then wait for the master clock SCK and  $\overline{SCS}$  signal. After this, go to step 5.  
For read operations: the data transferred in on the SDI line will be stored in the TXRX buffer until all the data has been received at which point it will be latched into the SIMD register.
- Step 5  
Check the WCOL bit if set high then a collision error has occurred so return to step 4. If equal to zero then go to the following step.
- Step 6  
Check the TRF bit or wait for a SIM SPI serial bus interrupt.
- Step 7  
Read data from the SIMD register.
- Step 8  
Clear TRF.
- Step 9  
Go to step 4.

#### **Error Detection**

The WCOL bit in the SIMC2 register is provided to indicate errors during data transfer. The bit is set by the SPI serial Interface but must be cleared by the application program. This bit indicates that a data collision has occurred which happens if a write to the SIMD register takes place during a data transfer operation and will prevent the write operation from continuing.

#### **I<sup>2</sup>C Interface**

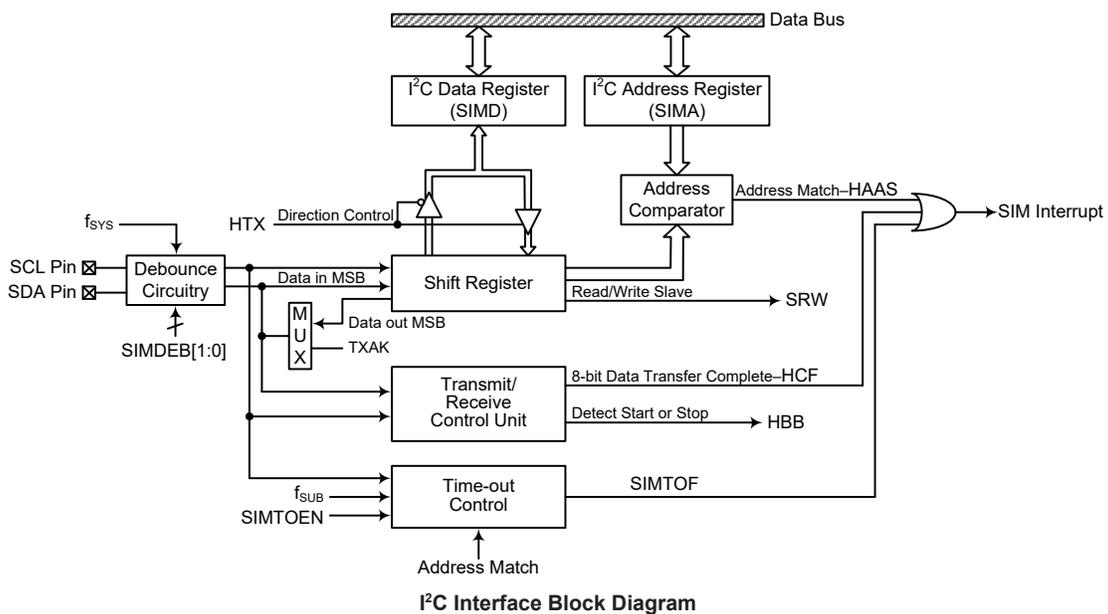
The I<sup>2</sup>C interface is used to communicate with external peripheral devices such as sensors, EEPROM memory etc. Originally developed by Philips, it is a two line low speed serial interface for synchronous serial data transfer. The advantage of only two lines for communication, relatively simple communication protocol and the ability to accommodate multiple devices on the same bus has made it an extremely popular interface type for many applications.

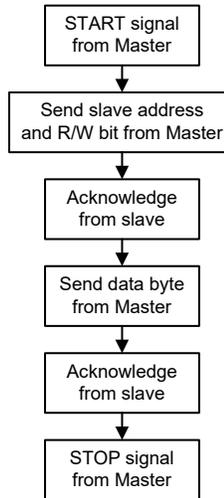


**I2C Interface Operation**

The I2C serial interface is a two line interface, a serial data line, SDA, and serial clock line, SCL. As many devices may be connected together on the same bus, their outputs are both open drain types. For this reason it is necessary that external pull-high resistors are connected to these outputs. Note that no chip select line exists, as each device on the I2C bus is identified by a unique address which will be transmitted and received on the I2C bus.

When two devices communicate with each other on the bidirectional I2C bus, one is known as the master device and one as the slave device. Both master and slave can transmit and receive data; however, it is the master device that has overall control of the bus. For the device, which only operates in slave mode, there are two methods of transferring data on the I2C bus, the slave transmit mode and the slave receive mode. The pull-high control function pin-shared with SCL/SDA pin is still applicable even if I2C device is activated and the related internal pull-high function could be controlled by its corresponding pull-high control register.





**I<sup>2</sup>C Interface Operation**

The SIMDEB1 and SIMDEB0 bits determine the debounce time of the I<sup>2</sup>C interface. This uses the system clock to in effect add a debounce time to the external clock to reduce the possibility of glitches on the clock line causing erroneous operation. The debounce time, if selected, can be chosen to be either 2 or 4 system clocks. To achieve the required I<sup>2</sup>C data transfer speed, there exists a relationship between the system clock,  $f_{SYS}$ , and the I<sup>2</sup>C debounce time. For either the I<sup>2</sup>C Standard or Fast mode operation, users must take care of the selected system clock frequency and the configured debounce time to match the criterion shown in the following table.

I <sup>2</sup> C Debounce Time Selection	I <sup>2</sup> C Standard Mode (100kHz)	I <sup>2</sup> C Fast Mode (400kHz)
No Debounce	$f_{SYS} > 2\text{MHz}$	$f_{SYS} > 4\text{MHz}$
2 system clock debounce	$f_{SYS} > 4\text{MHz}$	$f_{SYS} > 8\text{MHz}$
4 system clock debounce	$f_{SYS} > 4\text{MHz}$	$f_{SYS} > 8\text{MHz}$

**I<sup>2</sup>C Minimum  $f_{SYS}$  Frequency Requirement**

**I<sup>2</sup>C Registers**

There are three control registers associated with the I<sup>2</sup>C bus, SIMC0, SIMC1 and SIMTOC, one address register SIMA and one data register, SIMD.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SIMC0	SIM2	SIM1	SIM0	—	SIMDEB1	SIMDEB0	SIMEN	SIMICF
SIMC1	HCF	HAAS	HBB	HTX	TXAK	SRW	IAMWU	RXAK
SIMA	SIMA6	SIMA5	SIMA4	SIMA3	SIMA2	SIMA1	SIMA0	D0
SIMD	D7	D6	D5	D4	D3	D2	D1	D0
SIMTOC	SIMTOEN	SIMTOF	SIMTOS5	SIMTOS4	SIMTOS3	SIMTOS2	SIMTOS1	SIMTOS0

**I<sup>2</sup>C Register List**

**I<sup>2</sup>C Data Register**

The SIMD register is used to store the data being transmitted and received. The same register is used by both the SPI and I<sup>2</sup>C functions. Before the device writes data to the I<sup>2</sup>C bus, the actual data to be transmitted must be placed in the SIMD register. After the data is received from the I<sup>2</sup>C bus, the device can read it from the SIMD register. Any transmission or reception of data from the I<sup>2</sup>C bus must be made via the SIMD register.

**• SIMD Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

"x": unknown

 Bit 7~0 **D7~D0**: SIM data register bit 7 ~ bit 0

**I<sup>2</sup>C Address Register**

The SIMA register is also used by the SPI interface but has the name SIMC2. The SIMA register is the location where the 7-bit slave address of the slave device is stored. Bits 7~1 of the SIMA register define the device slave address. Bit 0 is not implemented.

When a master device, which is connected to the I<sup>2</sup>C bus, sends out an address, which matches the slave address in the SIMA register, the slave device will be selected. Note that the SIMA register locates at the same register address as SIMC2 which is used by the SPI interface.

**• SIMA Register**

Bit	7	6	5	4	3	2	1	0
Name	SIMA6	SIMA5	SIMA4	SIMA3	SIMA2	SIMA1	SIMA0	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

 Bit 7~1 **SIMA6~SIMA0**: I<sup>2</sup>C slave address

 SIMA6~SIMA0 is the I<sup>2</sup>C slave address bit 6~bit 0.

 Bit 0 **D0**: Reserved bit, can be read or written

**I<sup>2</sup>C Control Registers**

There are three control registers for the I<sup>2</sup>C interface, SIMC0, SIMC1 and SIMTOC. The register SIMC0 is used to control the enable/disable function and to select the I<sup>2</sup>C slave mode and debounce time. The SIMC1 register contains the relevant flags which are used to indicate the I<sup>2</sup>C communication status. Another register, SIMTOC, is used to control the I<sup>2</sup>C time-out function and is described in the corresponding section.

**• SIMC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	SIM2	SIM1	SIM0	—	SIMDEB1	SIMDEB0	SIMEN	SIMICF
R/W	R/W	R/W	R/W	—	R/W	R/W	R/W	R/W
POR	1	1	1	—	0	0	0	0

 Bit 7~5 **SIM2~SIM0**: SIM Operating Mode Control

 000: SPI master mode; SPI clock is  $f_{SYS}/4$ 

 001: SPI master mode; SPI clock is  $f_{SYS}/16$ 

 010: SPI master mode; SPI clock is  $f_{SYS}/64$ 

 011: SPI master mode; SPI clock is  $f_{SUB}$ 

100: SPI master mode; SPI clock is CTM CCRP match frequency/2

101: SPI slave mode

 110: I<sup>2</sup>C slave mode

111: Non SIM function

These bits setup the overall operating mode of the SIM function. As well as selecting if the I<sup>2</sup>C or SPI function, they are used to control the SPI Master/Slave selection and the SPI Master clock frequency. The SPI clock is a function of the system clock but can also be chosen to be sourced from CTM and  $f_{SUB}$ . If the SPI Slave Mode is selected then the clock will be supplied by an external Master device.

- Bit 4 Unimplemented, read as “0”
- Bit 3~2 **SIMDEB1~SIMDEB0**: I<sup>2</sup>C Debounce Time Selection  
 00: No debounce  
 01: 2 system clock debounce  
 1x: 4 system clock debounce  
 These bits are used to select the I<sup>2</sup>C debounce time when the SIM is configured as the I<sup>2</sup>C interface function by setting the SIM2~SIM0 bits to “110”.
- Bit 1 **SIMEN**: SIM Enable Control  
 0: Disable  
 1: Enable  
 The bit is the overall on/off control for the SIM interface. When the SIMEN bit is cleared to zero to disable the SIM interface, the SDI, SDO, SCK and  $\overline{SCS}$ , or SDA and SCL lines will lose their SPI or I<sup>2</sup>C function and the SIM operating current will be reduced to a minimum value. When the bit is high the SIM interface is enabled. The SIM configuration option must have first enabled the SIM interface for this bit to be effective. If the SIM is configured to operate as an SPI interface via the SIM2~SIM0 bits, the contents of the SPI control registers will remain at the previous settings when the SIMEN bit changes from low to high and should therefore be first initialised by the application program. If the SIM is configured to operate as an I<sup>2</sup>C interface via the SIM2~SIM0 bits and the SIMEN bit changes from low to high, the contents of the I<sup>2</sup>C control bits such as HTX and TXAK will remain at the previous settings and should therefore be first initialised by the application program while the relevant I<sup>2</sup>C flags such as HCF, HAAS, HBB, SRW and RXAK will be set to their default states.
- Bit 0 **SIMICF**: SIM SPI Incomplete Flag  
 The SIMICF bit is only used in the SPI mode and the detailed definition is described in the SPI section.

• **SIMC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	HCF	HAAS	HBB	HTX	TXAK	SRW	IAMWU	RXAK
R/W	R	R	R	R/W	R/W	R	R/W	R
POR	1	0	0	0	0	0	0	1

- Bit 7 **HCF**: I<sup>2</sup>C Bus data transfer completion flag  
 0: Data is being transferred  
 1: Completion of an 8-bit data transfer  
 The HCF flag is the data transfer completion flag. This flag will be zero when data is being transferred. Upon completion of an 8-bit data transfer the flag will go high and an interrupt will be generated.
- Bit 6 **HAAS**: I<sup>2</sup>C Bus data transfer completion flag  
 0: Not address match  
 1: Address match  
 The HAAS flag is the address match flag. This flag is used to determine if the slave device address is the same as the master transmit address. If the addresses match then this bit will be high, if there is no match then the flag will be low.
- Bit 5 **HBB**: I<sup>2</sup>C Bus busy flag  
 0: I<sup>2</sup>C Bus is not busy  
 1: I<sup>2</sup>C Bus is busy  
 The HBB flag is the I<sup>2</sup>C busy flag. This flag will be “1” when the I<sup>2</sup>C bus is busy which will occur when a START signal is detected. The flag will be cleared to “0” when the bus is free which will occur when a STOP signal is detected.
- Bit 4 **HTX**: I<sup>2</sup>C slave device transmitter/receiver selection  
 0: Slave device is the receiver  
 1: Slave device is the transmitter

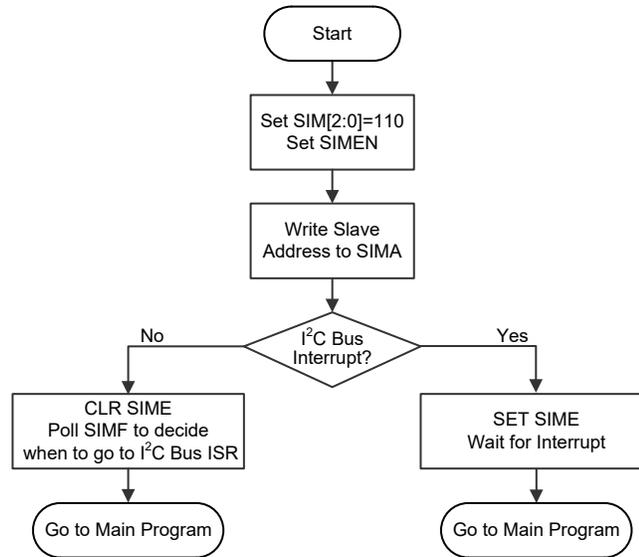
- Bit 3      **TXAK**: I<sup>2</sup>C bus transmit acknowledge flag  
             0: Slave sends acknowledge flag  
             1: Slave does not send acknowledge flag  
 The TXAK flag is the transmit acknowledge flag. After the slave device has received 8 bits of data, this flag will be transmitted to the bus on the 9<sup>th</sup> clock from the slave device. The slave device must always set the TXAK bit to “0” before further data is received.
- Bit 2      **SRW**: I<sup>2</sup>C slave read/write flag  
             0: Slave device should be in receive mode  
             1: Slave device should be in transmit mode  
 The SRW flag is the I<sup>2</sup>C Slave Read/Write flag. This flag determines whether the master device wishes to transmit or receive data from the I<sup>2</sup>C bus. When the transmitted address and slave address is match, that is when the HAAS flag is set high, the slave device will check the SRW flag to determine whether it should be in transmit mode or receive mode. If the SRW flag is high, the master is requesting to read data from the bus, so the slave device should be in transmit mode. When the SRW flag is zero, the master will write data to the bus, therefore the slave device should be in receive mode to read this data.
- Bit 1      **IAMWU**: I<sup>2</sup>C Address Match Wake-Up control  
             0: Disable  
             1: Enable – must be cleared by the application program after wake-up  
 This bit should be set to 1 to enable the I<sup>2</sup>C address match wake-up from the SLEEP or IDLE Mode. If the IAMWU bit has been set before entering either the SLEEP or IDLE mode to enable the I<sup>2</sup>C address match wake-up, then this bit must be cleared by the application program after wake-up to ensure correction device operation.
- Bit 0      **RXAK**: I<sup>2</sup>C bus receive acknowledge flag  
             0: Slave receives acknowledge flag  
             1: Slave does not receive acknowledge flag  
 The RXAK flag is the receiver acknowledge flag. When the RXAK flag is “0”, it means that a acknowledge signal has been received at the 9<sup>th</sup> clock, after 8 bits of data have been transmitted. When the slave device is in the transmit mode, the slave device checks the RXAK flag to determine if the master receiver wishes to receive the next byte. The slave transmitter will therefore continue sending out data until the RXAK flag is “1”. When this occurs, the slave transmitter will release the SDA line to allow the master to send a STOP signal to release the I<sup>2</sup>C Bus.

### I<sup>2</sup>C Bus Communication

Communication on the I<sup>2</sup>C bus requires four separate steps, a START signal, a slave device address transmission, a data transmission and finally a STOP signal. When a START signal is placed on the I<sup>2</sup>C bus, all devices on the bus will receive this signal and be notified of the imminent arrival of data on the bus. The first seven bits of the data will be the slave address with the first bit being the MSB. If the address of the slave device matches that of the transmitted address, the HAAS bit in the SIMC1 register will be set and an SIM interrupt will be generated. After entering the interrupt service routine, the slave device must first check the condition of the HAAS and SIMTOF bits to determine whether the interrupt source originates from either an address match or the completion of an 8-bit data transfer or the I<sup>2</sup>C bus time-out occurrence. During a data transfer, note that after the 7-bit slave address has been transmitted, the following bit, which is the 8<sup>th</sup> bit, is the read/write bit whose value will be placed in the SRW bit. This bit will be checked by the slave device to determine whether to go into transmit or receive mode. Before any transfer of data to or from the I<sup>2</sup>C bus, the microcontroller must initialise the bus; the following are steps to achieve this:

- Step 1  
     Set the SIM2~SIM0 bits to “110” and SIMEN bit to “1” in the SIMC0 register to enable the I<sup>2</sup>C bus.

- Step 2  
Write the slave address of the device to the I<sup>2</sup>C bus address register SIMA.
- Step 3  
Set the SIME interrupt enable bit of the interrupt control register to enable the SIM interrupt.



**I<sup>2</sup>C Bus Initialisation Flow Chart**

### I<sup>2</sup>C Bus Start Signal

The START signal can only be generated by the master device connected to the I<sup>2</sup>C bus and not by the slave device. This START signal will be detected by all devices connected to the I<sup>2</sup>C bus. When detected, this indicates that the I<sup>2</sup>C bus is busy and therefore the HBB bit will be set. A START condition occurs when a high to low transition on the SDA line takes place when the SCL line remains high.

### I<sup>2</sup>C Slave Address

The transmission of a START signal by the master will be detected by all devices on the I<sup>2</sup>C bus. To determine which slave device the master wishes to communicate with, the address of the slave device will be sent out immediately following the START signal. All slave devices, after receiving this 7-bit address data, will compare it with their own 7-bit slave address. If the address sent out by the master matches the internal address of the microcontroller slave device, then an internal SIM I<sup>2</sup>C bus interrupt signal will be generated. The next bit following the address, which is the 8<sup>th</sup> bit, defines the read/write status and will be saved to the SRW bit of the SIMC1 register. The slave device will then transmit an acknowledge bit, which is a low level, as the 9<sup>th</sup> bit. The slave device will also set the status flag HAAS when the addresses match.

As an SIM I<sup>2</sup>C bus interrupt can come from three sources, when the program enters the interrupt subroutine, the HAAS and SIMTOF bits should be examined to see whether the interrupt source has come from either a matching slave address, the completion of a data byte transfer or the I<sup>2</sup>C bus time-out occurrence. When a slave address is matched, the device must be placed in either the transmit mode and then write data to the SIMD register, or in the receive mode where it must implement a dummy read from the SIMD register to release the SCL line.

### **I<sup>2</sup>C Bus Read/Write Signal**

The SRW bit in the SIMC1 register defines whether the master device wishes to read data from the I<sup>2</sup>C bus or write data to the I<sup>2</sup>C bus. The slave device should examine this bit to determine if it is to be a transmitter or a receiver. If the SRW flag is “1” then this indicates that the master device wishes to read data from the I<sup>2</sup>C bus, therefore the slave device must be setup to send data to the I<sup>2</sup>C bus as a transmitter. If the SRW flag is “0” then this indicates that the master wishes to send data to the I<sup>2</sup>C bus, therefore the slave device must be setup to read data from the I<sup>2</sup>C bus as a receiver.

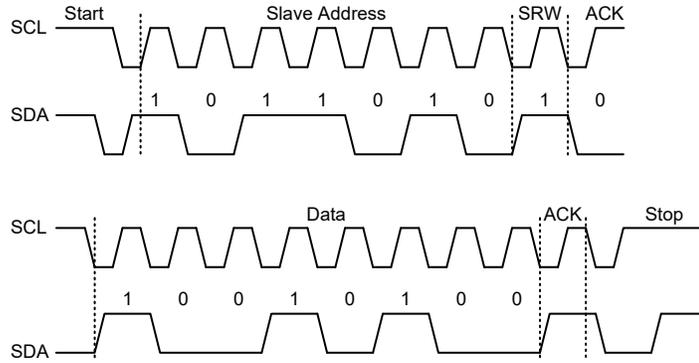
### **I<sup>2</sup>C Bus Slave Address Acknowledge Signal**

After the master has transmitted a calling address, any slave device on the I<sup>2</sup>C bus, whose own internal address matches the calling address, must generate an acknowledge signal. The acknowledge signal will inform the master that a slave device has accepted its calling address. If no acknowledge signal is received by the master then a STOP signal must be transmitted by the master to end the communication. When the HAAS flag is high, the addresses have matched and the slave device must check the SRW flag to determine if it is to be a transmitter or a receiver. If the SRW flag is high, the slave device should be setup to be a transmitter so the HTX bit in the SIMC1 register should be set to “1”. If the SRW flag is low, then the microcontroller slave device should be setup as a receiver and the HTX bit in the SIMC1 register should be cleared to “0”.

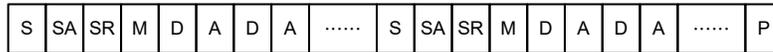
### **I<sup>2</sup>C Bus Data and Acknowledge Signal**

The transmitted data is 8-bit wide and is transmitted after the slave device has acknowledged receipt of its slave address. The order of serial bit transmission is the MSB first and the LSB last. After receipt of 8 bits of data, the receiver must transmit an acknowledge signal, level “0”, before it can receive the next data byte. If the slave transmitter does not receive an acknowledge bit signal from the master receiver, then the slave transmitter will release the SDA line to allow the master to send a STOP signal to release the I<sup>2</sup>C Bus. The corresponding data will be stored in the SIMD register. If setup as a transmitter, the slave device must first write the data to be transmitted into the SIMD register. If setup as a receiver, the slave device must read the transmitted data from the SIMD register.

When the slave receiver receives the data byte, it must generate an acknowledge bit, known as TXAK, on the 9<sup>th</sup> clock. The slave device, which is setup as a transmitter will check the RXAK bit in the SIMC1 register to determine if it is to send another data byte, if not then it will release the SDA line and await the receipt of a STOP signal from the master.

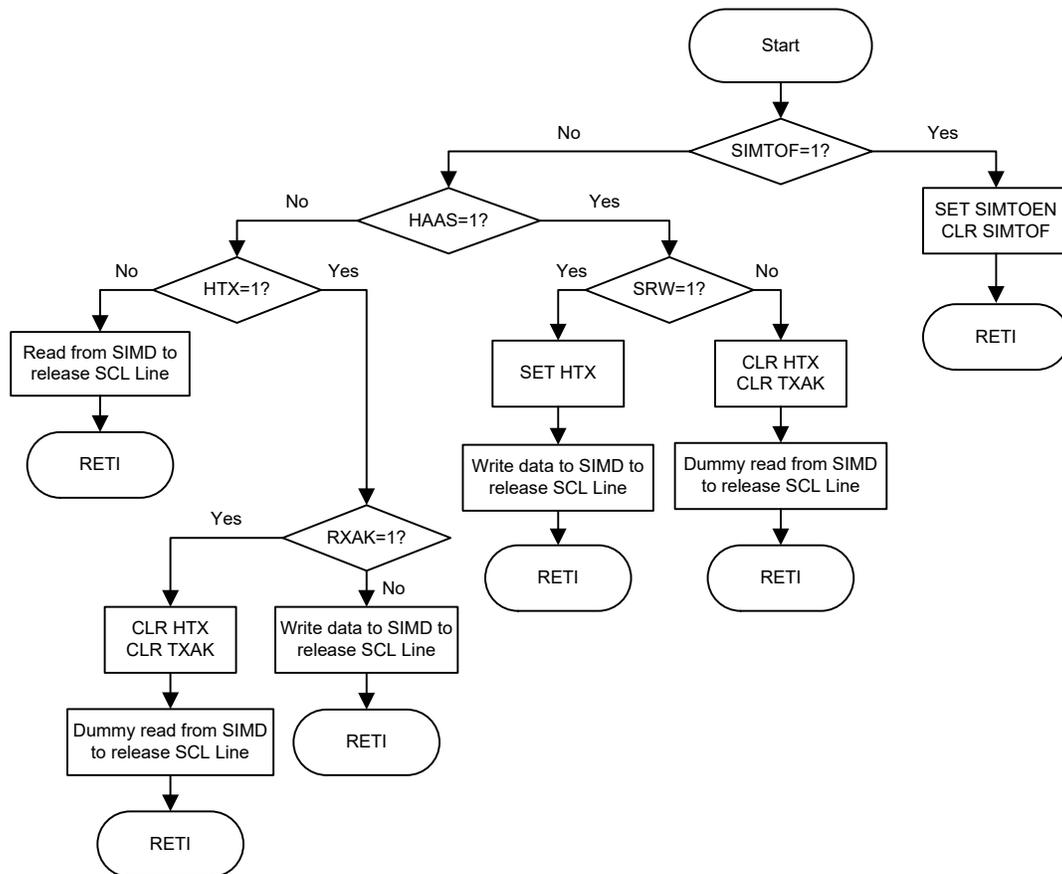


S=Start (1 bit)  
 SA=Slave Address (7 bits)  
 SR=SRW bit (1 bit)  
 M=Slave device send acknowledge bit (1 bit)  
 D=Data (8 bits)  
 A=ACK (RXAK bit for transmitter, TXAK bit for receiver, 1 bit)  
 P=Stop (1 bit)



Note: When a slave address is matched, the device must be placed in either the transmit mode and then write data to the SIMD register, or in the receive mode where it must implement a dummy read from the SIMD register to release the SCL line.

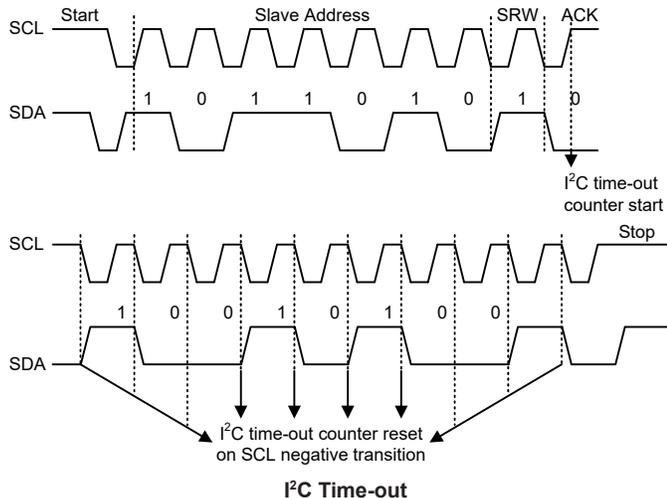
**I<sup>2</sup>C Communication Timing Diagram**



**I<sup>2</sup>C Bus ISR Flow Chart**

**I<sup>2</sup>C Time-out Control**

In order to reduce the I<sup>2</sup>C lockup problem due to reception of erroneous clock sources, a time-out function is provided. If the clock source connected to the I<sup>2</sup>C bus is not received for a while, then the I<sup>2</sup>C circuitry and registers will be reset after a certain time-out period. The time-out counter starts to count on an I<sup>2</sup>C bus “START” & “address match” condition, and is cleared by an SCL falling edge. Before the next SCL falling edge arrives, if the time elapsed is greater than the time-out period specified by the SIMTOC register, then a time-out condition will occur. The time-out function will stop when an I<sup>2</sup>C “STOP” condition occurs.



When an I<sup>2</sup>C time-out counter overflow occurs, the counter will stop and the SIMTOEN bit will be cleared to zero and the SIMTOF bit will be set high to indicate that a time-out condition has occurred. The time-out condition will also generate an interrupt which uses the SIM interrupt vector. When an I<sup>2</sup>C time-out occurs, the I<sup>2</sup>C internal circuitry will be reset and the registers will be reset into the following condition:

Registers	After I <sup>2</sup> C Time-out
SIMD, SIMA, SIMC0	No change
SIMC1	Reset to POR condition

**I<sup>2</sup>C Registers after Time-out**

The SIMTOF flag can be cleared by the application program. There are 64 time-out period selections which can be selected using the SIMTOS5~SIMTOS0 bits in the SIMTOC register. The time-out duration is calculated by the formula:  $((1\sim64) \times (32/f_{SUB}))$ . This gives a time-out period which ranges from about 1ms to 64ms.

**• SIMTOC Register**

Bit	7	6	5	4	3	2	1	0
Name	SIMTOEN	SIMTOF	SIMTOS5	SIMTOS4	SIMTOS3	SIMTOS2	SIMTOS1	SIMTOS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **SIMTOEN**: SIM I<sup>2</sup>C Time-out control

0: Disable

1: Enable

Bit 6 **SIMTOF**: SIM I<sup>2</sup>C Time-out flag

0: No time-out occurred

1: Time-out occurred

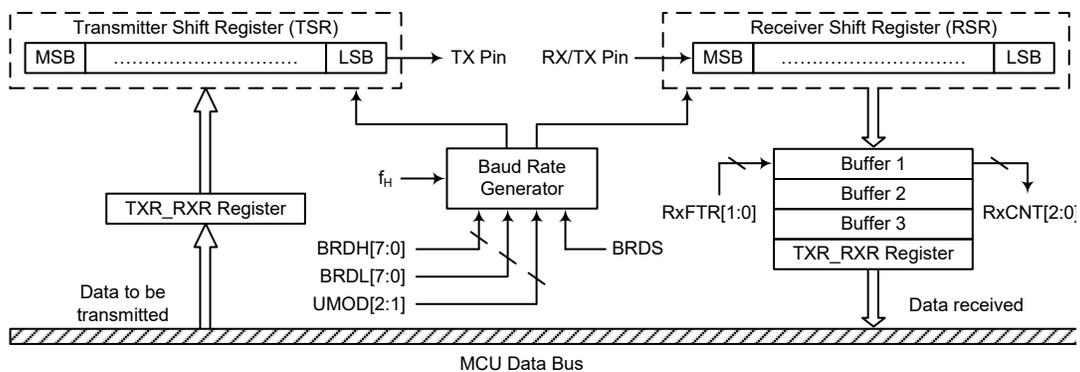
Bit 5~0 **SIMTOS5~SIMTOS0**: SIM I<sup>2</sup>C Time-out period selection  
 I<sup>2</sup>C Time-out clock source is  $f_{SUB}/32$ .  
 I<sup>2</sup>C Time-out period is equal to  $(SIMTOS[5:0]+1) \times (32/f_{SUB})$ .

## UART Interface

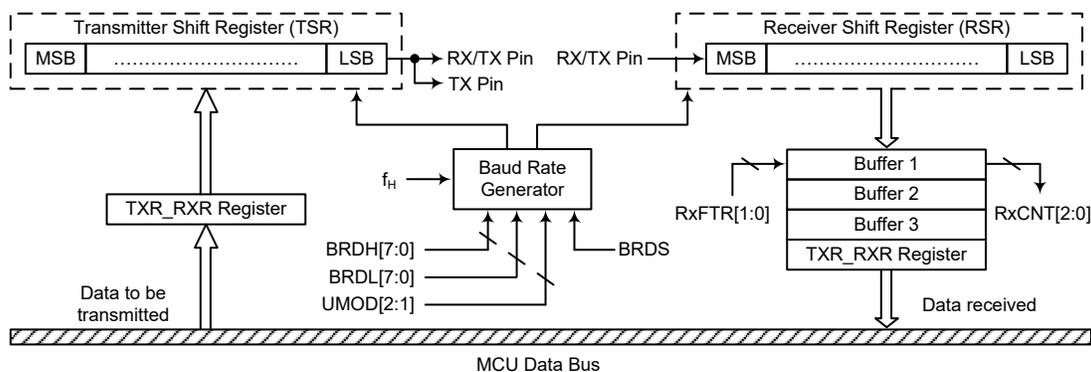
The device contains an integrated full-duplex or half-duplex asynchronous serial communications UART interface that enables communication with external devices that contain a serial interface. The UART function has many features and can transmit and receive data serially by transferring a frame of data with eight or nine data bits per transmission as well as being able to detect errors when the data is overwritten or incorrectly framed. The UART function possesses its own internal interrupt which can be used to indicate when a reception occurs or when a transmission terminates.

The integrated UART function contains the following features:

- Full-duplex or half-duplex (single wire mode), asynchronous communication
- 8 or 9 bits character length
- Even, odd, mark, space or no parity options
- One or two stop bits configurable for receiver
- Two stop bits for transmitter
- Baud rate generator with 16-bit prescaler
- Parity, framing, noise and overrun error detection
- Support for interrupt on address detect (last character bit=1)
- Separately enabled transmitter and receiver
- 4-byte Deep FIFO Receive Data Buffer
- 1-byte Deep FIFO Transmit Data Buffer
- RX/TX pin wake-up function
- Transmit and receive interrupts
- Interrupts can be triggered by the following conditions:
  - ♦ Transmitter Empty
  - ♦ Transmitter Idle
  - ♦ Receiver reaching FIFO trigger level
  - ♦ Receiver Overrun
  - ♦ Address Mode Detect



**UART Data Transfer Block Diagram – SWM=0**



UART Data Transfer Block Diagram – SWM=1

### UART External Pins

To communicate with an external serial interface, the internal UART has two external pins known as TX and RX/TX, which are pin-shared with I/O or other pin functions. The TX and RX/TX pin function should first be selected by the corresponding pin-shared function selection register before the UART function is used. Along with the UARTEN bit, the TXEN and RXEN bits, if set, will configure these pins to transmitter output and receiver input conditions. At this time the internal pull-high resistor related to the transmitter output pin will be disabled, while the internal pull-high resistor related to the receiver input pin is controlled by the corresponding I/O pull-high function control bit. When the TX or RX/TX pin function is disabled by clearing the UARTEN, TXEN or RXEN bit, the TX or RX/TX pin will be set to a floating state. At this time whether the internal pull-high resistor is connected to the TX or RX/TX pin or not is determined by the corresponding I/O pull-high function control bit.

### UART Single Wire Mode

The UART function also supports a Single Wire Mode communication which is selected using the SWM bit in the UCR3 register. When the SWM bit is set high, the UART function will be in the single wire mode. In the single wire mode, a single RX/TX pin can be used to transmit and receive data depending upon the corresponding control bits. When the RXEN bit is set high, the RX/TX pin is used as a receiver pin. When the RXEN bit is cleared to zero and the TXEN bit is set high, the RX/TX pin will act as a transmitter pin.

It is recommended not to set both the RXEN and TXEN bits high in the single wire mode. If both the RXEN and TXEN bits are set high, the RXEN bit will have the priority and the UART will act as a receiver.

It is important to note that the functional description in this UART chapter, which is described from the full-duplex communication standpoint, also applies to the half-duplex (single wire mode) communication except the pin usage. In the single wire mode, the TX pin mentioned in this chapter should be replaced by the RX/TX pin to understand the whole UART single wire mode function.

In the single wire mode, the data can also be transmitted on the TX pin in a transmission operation with proper software configurations. Therefore, the data will be output on the RX/TX and TX pins.

### UART Data Transfer Scheme

The UART Data Transfer Block Diagram shows the overall data transfer structure arrangement for the UART. The actual data to be transmitted from the MCU is first transferred to the TXR\_RXR register by the application program. The data will then be transferred to the Transmit Shift Register from where it will be shifted out, LSB first, onto the TX pin at a rate controlled by the Baud Rate

Generator. Only the TXR\_RXR register is mapped onto the MCU Data Memory, the Transmit Shift Register is not mapped and is therefore inaccessible to the application program.

Data to be received by the UART is accepted on the external RX/TX pin, from where it is shifted in, LSB first, to the Receiver Shift Register at a rate controlled by the Baud Rate Generator. When the shift register is full, the data will then be transferred from the shift register to the internal TXR\_RXR register, where it is buffered and can be manipulated by the application program. Only the TXR\_RXR register is mapped onto the MCU Data Memory, the Receiver Shift Register is not mapped and is therefore inaccessible to the application program.

It should be noted that the actual register for data transmission and reception only exists as a single shared register, TXR\_RXR, in the Data Memory.

### UART Status and Control Registers

There are nine control registers associated with the UART function. The SWM bit in the UCR3 register is used to enable/disable the UART Single Wire Mode. The USR, UCR1, UCR2, UFCR and RxCNT registers control the overall function of the UART, while the BRDH and BRDL registers control the Baud rate. The actual data to be transmitted and received on the serial interface is managed through the TXR\_RXR data register.

Register Name	Bit							
	7	6	5	4	3	2	1	0
USR	PERR	NF	FERR	OERR	RIDLE	RXIF	TIDLE	TXIF
UCR1	UARTEN	BNO	PREN	PRT1	PRT0	TXBRK	RX8	TX8
UCR2	TXEN	RXEN	STOPS	ADDEN	WAKE	RIE	TIIE	TEIE
UCR3	—	—	—	—	—	—	—	SWM
TXR_RXR	TXRX7	TXRX6	TXRX5	TXRX4	TXRX3	TXRX2	TXRX1	TXRX0
BRDH	D7	D6	D5	D4	D3	D2	D1	D0
BRDL	D7	D6	D5	D4	D3	D2	D1	D0
UFCR	—	—	UMOD2	UMOD1	UMOD0	BRDS	RxFTR1	RxFTR0
RxCNT	—	—	—	—	—	D2	D1	D0

**UART Register List**

#### • USR Register

The USR register is the status register for the UART, which can be read by the program to determine the present status of the UART. All flags within the USR register are read only. Further explanation on each of the flags is given below:

Bit	7	6	5	4	3	2	1	0
Name	PERR	NF	FERR	OERR	RIDLE	RXIF	TIDLE	TXIF
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	1	0	1	1

Bit 7 **PERR**: Parity error flag  
 0: No parity error is detected  
 1: Parity error is detected

The PERR flag is the parity error flag. When this read only flag is “0”, it indicates a parity error has not been detected. When the flag is “1”, it indicates that the parity of the received word is incorrect. This error flag is applicable only if the parity is enabled and the parity type (odd, even, mark or space) is selected. The flag can also be cleared by a software sequence which involves a read to the status register USR followed by an access to the TXR\_RXR data register.

- Bit 6      **NF:** Noise flag  
             0: No noise is detected  
             1: Noise is detected
- The NF flag is the noise flag. When this read only flag is “0”, it indicates no noise condition. When the flag is “1”, it indicates that the UART has detected noise on the receiver input. The NF flag is set during the same cycle as the RXIF flag but will not be set in the case of an overrun. The NF flag can be cleared by a software sequence which will involve a read to the status register USR followed by an access to the TXR\_RXR data register.
- Bit 5      **FERR:** Framing error flag  
             0: No framing error is detected  
             1: Framing error is detected
- The FERR flag is the framing error flag. When this read only flag is “0”, it indicates that there is no framing error. When the flag is “1”, it indicates that a framing error has been detected for the current character. The flag can also be cleared by a software sequence which will involve a read to the status register USR followed by an access to the TXR\_RXR data register.
- Bit 4      **OERR:** Overrun error flag  
             0: No overrun error is detected  
             1: Overrun error is detected
- The OERR flag is the overrun error flag which indicates when the receiver buffer has overflowed. When this read only flag is “0”, it indicates that there is no overrun error. When the flag is “1”, it indicates that an overrun error occurs which will inhibit further transfers to the TXR\_RXR receive data register. The flag is cleared by a software sequence, which is a read to the status register USR followed by an access to the TXR\_RXR data register.
- Bit 3      **RIDLE:** Receiver status  
             0: Data reception is in progress (Data being received)  
             1: No data reception is in progress (Receiver is idle)
- The RIDLE flag is the receiver status flag. When this read only flag is “0”, it indicates that the receiver is between the initial detection of the start bit and the completion of the stop bit. When the flag is “1”, it indicates that the receiver is idle. Between the completion of the stop bit and the detection of the next start bit, the RIDLE bit is “1” indicating that the UART receiver is idle and the RX/TX pin stays in logic high condition.
- Bit 2      **RXIF:** Receive TXR\_RXR data register status  
             0: TXR\_RXR data register is empty  
             1: TXR\_RXR data register has available data and reach Receiver FIFO trigger level
- The RXIF flag is the receive data register status flag. When this read only flag is “0”, it indicates that the TXR\_RXR read data register is empty. When the flag is “1”, it indicates that the TXR\_RXR read data register contains new data and reaches the Receiver FIFO trigger level. When the contents of the shift register are transferred to the TXR\_RXR register and reach Receiver FIFO trigger level, an interrupt is generated if RIE=1 in the UCR2 register. If one or more errors are detected in the received word, the appropriate receive-related flags NF, FERR, and/or PERR are set within the same clock cycle. The RXIF flag is cleared when the USR register is read with RXIF set, followed by a read from the TXR\_RXR register, and if the TXR\_RXR register has no data available.
- Bit 1      **TIDLE:** Transmission idle  
             0: Data transmission is in progress (Data being transmitted)  
             1: No data transmission is in progress (Transmitter is idle)
- The TIDLE flag is known as the transmission complete flag. When this read only flag is “0”, it indicates that a transmission is in progress. This flag will be set high when the TXIF flag is “1” and when there is no transmit data or break character being transmitted. When TIDLE is equal to “1”, the TX pin becomes idle with the pin state in logic high condition. The TIDLE flag is cleared by reading the USR register with

TIDLE set and then writing to the TXR\_RXR register. The flag is not generated when a data character or a break is queued and ready to be sent.

Bit 0 **TXIF**: Transmit TXR\_RXR data register status

0: Character is not transferred to the transmit shift register

1: Character has transferred to the transmit shift register (TXR\_RXR data register is empty)

The TXIF flag is the transmit data register empty flag. When this read only flag is “0”, it indicates that the character is not transferred to the transmitter shift register. When the flag is “1”, it indicates that the transmitter shift register has received a character from the TXR\_RXR data register. The TXIF flag is cleared by reading the UART status register (USR) with TXIF set and then writing to the TXR\_RXR data register. Note that when the TXEN bit is set, the TXIF flag will also be set since the transmit data register is not yet full.

• **UCR1 Register**

The UCR1 register together with the UCR2 and UCR3 register are the three UART control registers that are used to set the various options for the UART function, such as overall on/off control, parity control, data transfer bit length, single wire mode communication etc. Further explanation on each of the bits is given below:

Bit	7	6	5	4	3	2	1	0
Name	UARTEN	BNO	PREN	PRT1	PRT0	TXBRK	RX8	TX8
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R	W
POR	0	0	0	0	0	0	x	0

“x”: unknown

Bit 7 **UARTEN**: UART function enable control

0: Disable UART. TX and RX/TX pins are in a floating state

1: Enable UART. TX and RX/TX pins function as UART pins

The UARTEN bit is the UART enable bit. When this bit is equal to “0”, the UART will be disabled and the RX/TX pin as well as the TX pin will be set in a floating state. When the bit is equal to “1”, the UART will be enabled and the TX and RX/TX pins will function as defined by the SWM mode selection bit together with the TXEN and RXEN enable control bits.

When the UART is disabled, it will empty the buffer so any character remaining in the buffer will be discarded. In addition, the value of the baud rate counter will be reset. If the UART is disabled, all error and status flags will be reset. Also the TXEN, RXEN, TXBRK, RXIF, OERR, FERR, PERR and NF bits as well as the RxCNT register will be cleared, while the TIDLE, TXIF and RIDLE bits will be set. Other control bits in UCR1, UCR2, UCR3, UFCR, BRDH and BRDL registers will remain unaffected. If the UART is active and the UARTEN bit is cleared, all pending transmissions and receptions will be terminated and the module will be reset as defined above. When the UART is re-enabled, it will restart in the same configuration.

Bit 6 **BNO**: Number of data transfer bits selection

0: 8-bit data transfer

1: 9-bit data transfer

This bit is used to select the data length format, which can have a choice of either 8-bit or 9-bit format. When this bit is equal to “1”, a 9-bit data length format will be selected. If the bit is equal to “0”, then an 8-bit data length format will be selected. If 9-bit data length format is selected, then bits RX8 and TX8 will be used to store the 9th bit of the received and transmitted data respectively.

Note that the 9th bit of data if BNO=1, or the 8th bit of data if BNO=0, which is used as the parity bit, does not transfer to RX8 or TXRX7 respectively when the parity function is enabled.

- Bit 5     **PREN**: Parity function enable control  
           0: Parity function is disabled  
           1: Parity function is enabled  
 This is the parity enable bit. When this bit is equal to “1”, the parity function will be enabled. If the bit is equal to “0”, then the parity function will be disabled. Replace the most significant bit position with a parity bit.
- Bit 4~3   **PRT1~PRT0**: Parity type selection bits  
           00: Even parity for parity generator  
           01: Odd parity for parity generator  
           10: Mark parity for parity generator  
           11: Space parity for parity generator  
 These bits are the parity type selection bits. When these bits are equal to 00b, even parity type will be selected. If these bits are equal to 01b, then odd parity type will be selected. If these bits are equal to 10b, then a 1 (Mark) in the parity bit location will be selected. If these bits are equal to 11b, then a 0 (Space) in the parity bit location will be selected.
- Bit 2     **TXBRK**: Transmit break character  
           0: No break character is transmitted  
           1: Break characters transmit  
 The TXBRK bit is the Transmit Break Character bit. When this bit is “0”, there are no break characters and the TX pin operates normally. When the bit is “1”, there are transmit break characters and the transmitter will send logic zeros. When this bit is equal to “1”, after the buffered data has been transmitted, the transmitter output is held low for a minimum of a 13-bit length and until the TXBRK bit is reset.
- Bit 1     **RX8**: Receive data bit 8 for 9-bit data transfer format (read only)  
 This bit is only used if 9-bit data transfers are used, in which case this bit location will store the 9th bit of the received data known as RX8. The BNO bit is used to determine whether data transfers are in 8-bit or 9-bit format.
- Bit 0     **TX8**: Transmit data bit 8 for 9-bit data transfer format (write only)  
 This bit is only used if 9-bit data transfers are used, in which case this bit location will store the 9th bit of the transmitted data known as TX8. The BNO bit is used to determine whether data transfers are in 8-bit or 9-bit format.

• **UCR2 Register**

The UCR2 register is the second of the two UART control registers and serves several purposes. One of its main functions is to control the basic enable/disable operation of the UART Transmitter and Receiver as well as enabling the various UART interrupt sources. The register also serves to control the receiver STOP bit number selection, receiver wake-up enable and the address detect enable. Further explanation on each of the bits is given below:

Bit	7	6	5	4	3	2	1	0
Name	TXEN	RXEN	STOPS	ADDEN	WAKE	RIE	TIIE	TEIE
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7     **TXEN**: UART Transmitter enabled control  
           0: UART transmitter is disabled  
           1: UART transmitter is enabled  
 The bit named TXEN is the Transmitter Enable Bit. When this bit is equal to “0”, the transmitter will be disabled with any pending data transmissions being aborted. In addition the buffers will be reset. In this situation the TX pin will be set in a floating state.  
 If the TXEN bit is equal to “1” and the UARTEN bit is also equal to “1”, the transmitter will be enabled and the TX pin will be controlled by the UART. Clearing the TXEN bit during a transmission will cause the data transmission to be aborted and

will reset the transmitter. If this situation occurs, the TX pin will be set in a floating state.

- Bit 6 **RXEN**: UART Receiver enabled control  
0: UART receiver is disabled  
1: UART receiver is enabled  
The bit named RXEN is the Receiver Enable Bit. When this bit is equal to “0”, the receiver will be disabled with any pending data receptions being aborted. In addition the receive buffers will be reset. In this situation the RX/TX pin will be set in a floating state. If the RXEN bit is equal to “1” and the UARTEN bit is also equal to “1”, the receiver will be enabled and the RX/TX pin will be controlled by the UART. Clearing the RXEN bit during a reception will cause the data reception to be aborted and will reset the receiver. If this situation occurs, the RX/TX pin will be set in a floating state.
- Bit 5 **STOPS**: Number of Stop bits selection for receiver  
0: One stop bit format is used  
1: Two stop bits format is used  
This bit determines if one or two stop bits are to be used for receiver. When this bit is equal to “1”, two stop bits are used. If this bit is equal to “0”, then only one stop bit is used. Two stop bits are used for transmitter.
- Bit 4 **ADDEN**: Address detect function enable control  
0: Address detect function is disabled  
1: Address detect function is enabled  
The bit named ADDEN is the address detect function enable control bit. When this bit is equal to “1”, the address detect function is enabled. When it occurs, if the 8th bit, which corresponds to TXRX7 if BNO=0 or the 9th bit, which corresponds to RX8 if BNO=1, has a value of “1”, then the received word will be identified as an address, rather than data. If the corresponding interrupt is enabled, an interrupt request will be generated each time the received word has the address bit set, which is the 8th or 9th bit depending on the value of BNO. If the address bit known as the 8th or 9th bit of the received word is “0” with the address detect function being enabled, an interrupt will not be generated and the received data will be discarded.
- Bit 3 **WAKE**: RX/TX pin wake-up UART function enable control  
0: RX/TX pin wake-up UART function is disabled  
1: RX/TX pin wake-up UART function is enabled  
This bit is used to control the wake-up UART function when a falling edge on the RX/TX pin occurs. Note that this bit is only available when the UART clock ( $f_{H}$ ) is switched off. There will be no RX/TX pin wake-up UART function if the UART clock ( $f_{H}$ ) exists. If the WAKE bit is set to 1 as the UART clock ( $f_{H}$ ) is switched off, a UART wake-up request will be initiated when a falling edge on the RX/TX pin occurs. When this request happens and the corresponding interrupt is enabled, an RX/TX pin wake-up UART interrupt will be generated to inform the MCU to wake up the UART function by switching on the UART clock ( $f_{H}$ ) via the application program. Otherwise, the UART function cannot resume even if there is a falling edge on the RX/TX pin when the WAKE bit is cleared to 0.
- Bit 2 **RIE**: Receiver interrupt enable control  
0: Receiver related interrupt is disabled  
1: Receiver related interrupt is enabled  
This bit enables or disables the receiver interrupt. If this bit is equal to “1” and when the receiver overrun flag OERR or receive data available flag RXIF is set, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt request flag will not be influenced by the condition of the OERR or RXIF flags.
- Bit 1 **TIE**: Transmitter Idle interrupt enable control  
0: Transmitter idle interrupt is disabled  
1: Transmitter idle interrupt is enabled  
This bit enables or disables the transmitter idle interrupt. If this bit is equal to “1” and when the transmitter idle flag TIDLE is set, due to a transmitter idle condition, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt

request flag will not be influenced by the condition of the TIDLE flag.

Bit 0 **TEIE**: Transmitter Empty interrupt enable control  
 0: Transmitter empty interrupt is disabled  
 1: Transmitter empty interrupt is enabled

This bit enables or disables the transmitter empty interrupt. If this bit is equal to “1” and when the transmitter empty flag TXIF is set, due to a transmitter empty condition, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt request flag will not be influenced by the condition of the TXIF flag.

• **UCR3 Register**

The UCR3 register is used to enable the UART Single Wire Mode communication. As the name suggests in the single wire mode the UART communication can be implemented in one single line, RX/TX, together with the control of the RXEN and TXEN bits in the UCR2 register.

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	—	SWM
R/W	—	—	—	—	—	—	—	R/W
POR	—	—	—	—	—	—	—	0

Bit 7~1 Unimplemented, read as “0”

Bit 0 **SWM**: Single Wire Mode enable control  
 0: Disable, the RX/TX pin is used as UART receiver function only  
 1: Enable, the RX/TX pin can be used as UART receiver or transmitter function controlled by the RXEN and TXEN bits

Note that when the Single Wire Mode is enabled, if both the RXEN and TXEN bits are high, the RX/TX pin will just be used as UART receiver input.

• **TXR\_RXR Register**

The TXR\_RXR register is the data register which is used to store the data to be transmitted on the TX pin or being received from the RX/TX pin.

Bit	7	6	5	4	3	2	1	0
Name	TXRX7	TXRX6	TXRX5	TXRX4	TXRX3	TXRX2	TXRX1	TXRX0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

“x”: unknown

Bit 7~0 **TXRX7~TXRX0**: UART Transmit/Receive Data bit 7 ~ bit 0

• **BRDH Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: Baud rate divider high byte

The baud rate divider BRD (BRDH/BRDL) defines the UART clock divider ratio.  
 $Baud\ Rate = f_H / (BRD + UMOD/8)$   
 BRD=16~65535 or 8~65535 depending on BRDS

Note: 1. The BRD value should not be set to less than 16 when BRDS=0 or less than 8 when BRDS=1, otherwise errors may occur.  
 2. The BRDL must be written first and then BRDH, otherwise errors may occur.  
 3. The BRDH register should not be modified during data transmission process.

• **BRDL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: Baud rate divider low byte  
 The baud rate divider BRD (BRDH/BRDL) defines the UART clock divider ratio.  
 $\text{Baud Rate} = f_{\text{H}} / (\text{BRD} + \text{UMOD}/8)$   
 $\text{BRD} = 16 \sim 65535$  or  $8 \sim 65535$  depending on BRDS  
 Note: 1. The BRD value should not be set to less than 16 when BRDS=0 or less than 8 when BRDS=1, otherwise errors may occur.  
 2. The BRDL must be written first and then BRDH, otherwise errors may occur.  
 3. The BRDL register should not be modified during data transmission process.

• **UFCR Register**

The UFCR register is the FIFO control register which is used for UART modulation control, BRD range selection and trigger level selection for RXIF and interrupt.

Bit	7	6	5	4	3	2	1	0
Name	—	—	UMOD2	UMOD1	UMOD0	BRDS	RxFTR1	RxFTR0
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	0	0	0	0

Bit 7~6 Unimplemented, read as “0”  
 Bit 5~3 **UMOD2~UMOD0**: UART Modulation Control bits  
 The modulation control bits are used to correct the baud rate of the received or transmitted UART signal. These bits determine if the extra UART clock cycle should be added in a UART bit time. The UMOD2~UMOD0 will be added to internal accumulator for every UART bit time. Until a carry to bit 3, the corresponding UART bit time increases a UART clock cycle.  
 Bit 2 **BRDS**: BRD range selection  
 0: BRD range is from 16 to 65535  
 1: BRD range is from 8 to 65535  
 The BRDS is used to control the sampling point in a UART bit time. If the BRDS bit is cleared to zero, the sampling point will be  $\text{BRD}/2$ ,  $\text{BRD}/2 + 1 \times f_{\text{H}}$ , and  $\text{BRD}/2 + 2 \times f_{\text{H}}$  in a UART bit time. If the BRDS bit is set high, the sampling point will be  $\text{BRD}/2 - 1 \times f_{\text{H}}$ ,  $\text{BRD}/2$ , and  $\text{BRD}/2 + 2 \times f_{\text{H}}$  in a UART bit time.  
 Note that the BRDS bit should not be modified during data transmission process.  
 Bit 1~0 **RxFTR1~RxFTR0**: Receiver FIFO trigger level (bytes)  
 00: 4 bytes in Receiver FIFO  
 01: 1 or more bytes in Receiver FIFO  
 10: 2 or more bytes in Receiver FIFO  
 11: 3 or more bytes in Receiver FIFO  
 For the receiver these bits define the number of received data bytes in the Receiver FIFO that will trigger the RXIF bit being set high, an interrupt will also be generated if the RIE bit is enabled. To prevent OERR from being set high, the receiver FIFO trigger level can be set to 2 bytes, avoiding an overrun state that cannot be processed by the program in time when more than 4 data bytes are received. After the reset the Receiver FIFO is empty.

• **RxCNT Register**

The RxCNT register is the counter used to indicate the number of received data bytes in the Receiver FIFO which have not been read by the MCU. This register is read only.

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	D2	D1	D0
R/W	—	—	—	—	—	R	R	R
POR	—	—	—	—	—	0	0	0

Bit 7~3 Unimplemented, read as “0”

Bit 2~0 **D2~D0**: Receiver FIFO counter

The RxCNT register is the counter used to indicate the number of received data bytes in the Receiver FIFO which is not read by the MCU. When Receiver FIFO receives one byte data, the RxCNT will increase by one; when the MCU reads one byte data from the Receiver FIFO, the RxCNT will decrease by one. If there are 4 bytes of data in the Receiver FIFO, the 5th data will be saved in the shift register. If there is 6th data, the 6th data will be saved in the shift register. But the RxCNT remains the value of 4. The RxCNT will be cleared when reset occurs or UARTEN=1. This register is read only.

**Baud Rate Generator**

To setup the speed of the serial data communication, the UART function contains its own dedicated baud rate generator. The baud rate is controlled by its own internal free running 16-bit timer, the period of which is determined by two factors. The first of these is the value placed in the BRDH/BRDL register and the second is the UART modulation control bits UMOD2~UMOD0. To prevent accumulated error of the receiver baud rate frequency, it is recommended to use two stop bits for resynchronization after each byte is received. If a baud rate BR is required with UART clock  $f_H$ .

$$f_H/BR = \text{Integer Part} + \text{Fractional Part}$$

The integer part is loaded into BRD (BRDH/BRDL). The fractional part is multiplied by 8 and rounded, then loaded into the UMOD bit field below:

$$BRD = \text{TRUNC}(f_H/BR)$$

$$UMOD = \text{ROUND}[\text{MOD}(f_H/BR) \times 8]$$

Therefore, the actual baud rate is calculated as follows:

$$\text{Baud rate} = f_H/[BRD+(UMOD/8)]$$

**Calculating the Baud Rate and Error Values**

For a clock frequency of 4MHz, determine the BRDH/BRDL register value, the actual baud rate and the error value for a desired baud rate of 230400.

From the above formula, the  $BRD = \text{TRUNC}(f_H/BR) = \text{TRUNC}(17.36111) = 17$

The  $UMOD = \text{ROUND}[\text{MOD}(f_H/BR) \times 8] = \text{ROUND}(0.36111 \times 8) = \text{ROUND}(2.88888) = 3$

The actual Baud Rate =  $f_H/[BRD+(UMOD/8)] = 230215.83$

Therefore the error is equal to  $(230215.83-230400)/230400 = -0.08\%$

**Modulation Control Example**

To get the best-fitting bit sequence for UART modulation control bits UMOD2~UMOD0, the following algorithm can be used: Firstly, the fractional part of the theoretical division factor is multiplied by 8. Then the product will be rounded and UMOD2~UMOD0 bits will be filled with the rounded value. The UMOD2~UMOD0 will be added to internal accumulator for every UART bit time. Until a carry to bit 3, the corresponding UART bit time increases a UART clock cycle.

The following is an example using the fraction 0.36111 previously calculated:  $UMOD[2:0] = ROUND(0.36111 \times 8) = 011b$ .

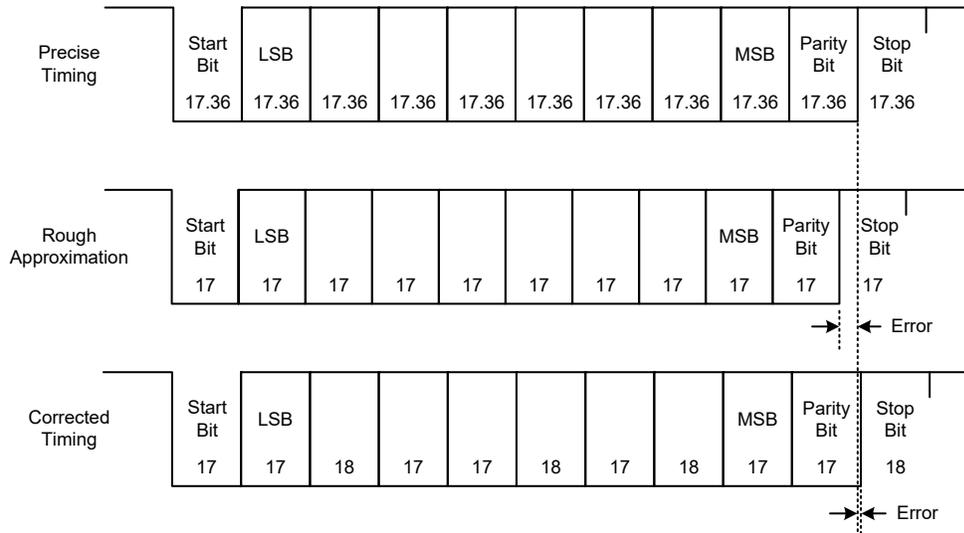
Fraction Addition	Carry to Bit 3	UART Bit Time Sequence	Extra UART Clock Cycle
0000b + 0011b = 0011b	No	Start bit	No
0011b + 0011b = 0110b	No	D0	No
0110b + 0011b = 1001b	Yes	D1	Yes
1001b + 0011b = 1100b	No	D2	No
1100b + 0011b = 1111b	No	D3	No
1111b + 0011b = 0010b	Yes	D4	Yes
0010b + 0011b = 0101b	No	D5	No
0101b + 0011b = 1000b	Yes	D6	Yes
1000b + 0011b = 1011b	No	D7	No
1011b + 0011b = 1110b	No	Parity bit	No
1110b + 0011b = 0001b	Yes	Stop bit	Yes

### Baud Rate Correction Example

The following figure presents an example using a baud rate of 230400 generated with UART clock  $f_H$ . The data format for the following figure is: eight data bits, parity enabled, no address bit, two stop bits.

The following figure shows three different frames:

- The upper frame is the correct one, with a bit-length of 17.36  $f_H$  cycles ( $4000000/230400=17.36$ ).
- The middle frame uses a rough estimate, with 17  $f_H$  cycles for the bit length.
- The lower frame shows a corrected frame using the best fit for the UART modulation control bits UMOD2~UMOD0.



### UART Setup and Control

For data transfer, the UART function utilizes a non-return-to-zero, more commonly known as NRZ, format. This is composed of one start bit, eight or nine data bits, and one or two stop bits. Parity is supported by the UART hardware, and can be setup to be even, odd, mark, space or no parity. For the most common data format, 8 data bits along with no parity and one stop bit, denoted as 8, N, 1, is used as the default setting, which is the setting at power-on. The number of data bits along with

the parity are setup by programming the BNO, PRT1~PRT0 and PREN bits. The transmitter always uses two stop bits while the receiver uses one or two stop bits which is determined by the STOPS bit. The baud rate used to transmit and receive data is setup using the internal 16-bit baud rate generator, while the data is transmitted and received LSB first. Although the UART transmitter and receiver are functionally independent, they both use the same data format and baud rate. In all cases stop bits will be used for data transmission.

### Enabling/Disabling the UART Interface

The basic on/off function of the internal UART function is controlled using the UARTEN bit in the UCR1 register. If the UARTEN, TXEN and RXEN bits are set, then these two UART pins will act as normal TX output pin and RX/TX input pin respectively. If no data is being transmitted on the TX pin, then it will default to a logic high value.

Clearing the UARTEN bit will disable the TX and RX/TX pins and allow these two pins to be used as normal I/O or other pin-shared functional pins by configuring the corresponding pin-shared control bits. When the UART function is disabled the buffer will be reset to an empty condition, at the same time discarding any remaining residual data. Disabling the UART will also reset the error and status flags with bits TXEN, RXEN, TXBRK, RXIF, OERR, FERR, PERR and NF as well as register RxCNT being cleared while bits TIDLE, TXIF and RIDLE will be set. The remaining control bits in the UCR1, UCR2, UCR3, UFCR, BRDH and BRDL registers will remain unaffected. If the UARTEN bit in the UCR1 register is cleared while the UART is active, then all pending transmissions and receptions will be immediately suspended and the UART will be reset to a condition as defined above. If the UART is then subsequently re-enabled, it will restart again in the same configuration.

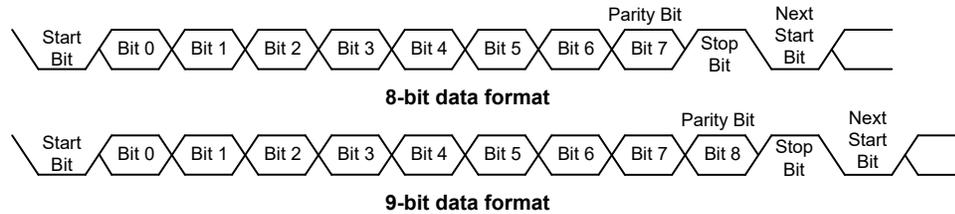
### Data, Parity and Stop Bit Selection

The format of the data to be transferred is composed of various factors such as data bit length, parity on/off, parity type, address bits and the number of stop bits. These factors are determined by the setup of various bits within the UCR1 and UCR2 registers. The BNO bit controls the number of data bits which can be set to either 8 or 9, the PRT1~PRT0 bits control the choice of odd, even, mark or space parity, the PREN bit controls the parity on/off function and the STOPS bit decides whether one or two stop bits are to be used for the receiver, while the transmitter always uses two stop bits. The following table shows various formats for data transmission. The address bit, which is the MSB of the data byte, identifies the frame as an address character or data if the address detect function is enabled. The number of stop bits, which can be either one or two, is independent of the data length and only configurable for the receiver. The transmitter uses two stop bits.

Start Bit	Data Bits	Address Bit	Parity Bit	Stop Bit
<b>Example of 8-bit Data Formats</b>				
1	8	0	0	1 or 2
1	7	0	1	1 or 2
1	7	1	0	1 or 2
<b>Example of 9-bit Data Formats</b>				
1	9	0	0	1 or 2
1	8	0	1	1 or 2
1	8	1	0	1 or 2

**Transmitter Receiver Data Format**

The following diagram shows the transmit and receive waveforms for both 8-bit and 9-bit data formats.



## UART Transmitter

Data word lengths of either 8 or 9 bits can be selected by programming the BNO bit in the UCR1 register. When BNO bit is set, the word length will be set to 9 bits. In this case the 9th bit, which is the MSB, needs to be stored in the TX8 bit in the UCR1 register. At the transmitter core lies the Transmitter Shift Register, more commonly known as the TSR, whose data is obtained from the transmit data register, which is known as the TXR\_RXR register. The data to be transmitted is loaded into this TXR\_RXR register by the application program. The TSR register is not written to with new data until the stop bit from the previous transmission has been sent out. As soon as this stop bit has been transmitted, the TSR can then be loaded with new data from the TXR\_RXR register, if it is available. It should be noted that the TSR register, unlike many other registers, is not directly mapped into the Data Memory area and as such is not available to the application program for direct read/write operations. An actual transmission of data will normally be enabled when the TXEN bit is set, but the data will not be transmitted until the TXR\_RXR register has been loaded with data and the baud rate generator has defined a shift clock source. However, the transmission can also be initiated by first loading data into the TXR\_RXR register, after which the TXEN bit can be set. When a transmission of data begins, the TSR is normally empty, in which case a transfer to the TXR\_RXR register will result in an immediate transfer to the TSR. If during a transmission the TXEN bit is cleared, the transmission will immediately cease and the transmitter will be reset. The TX output pin can then be configured as the I/O or other pin-shared functions by configuring the corresponding pin-shared control bits.

### Transmitting Data

When the UART is transmitting data, the data is shifted on the TX pin from the shift register, with the least significant bit first. In the transmit mode, the TXR\_RXR register forms a buffer between the internal bus and the transmitter shift register. It should be noted that if 9-bit data format has been selected, then the MSB will be taken from the TX8 bit in the UCR1 register. The steps to initiate a data transfer can be summarized as follows:

- Make the correct selection of the BNO, PRT1~PRT0 and PREN bits to define the required word length and parity type. Two stop bits are used for the transmitter.
- Setup the BRDH and BRDL registers and the UMOD2~UMOD0 bits to select the desired baud rate.
- Set the TXEN bit to ensure that the TX pin is used as a UART transmitter pin.
- Access the USR register and write the data that is to be transmitted into the TXR\_RXR register. Note that this step will clear the TXIF bit.

This sequence of events can now be repeated to send additional data.

It should be noted that when TXIF=0, data will be inhibited from being written to the TXR\_RXR register. Clearing the TXIF flag is always achieved using the following software sequence:

1. A USR register access
2. A TXR\_RXR register write execution

The read-only TXIF flag is set by the UART hardware and if set indicates that the TXR\_RXR register is empty and that other data can now be written into the TXR\_RXR register without overwriting the previous data. If the TEIE bit is set then the TXIF flag will generate an interrupt.

During a data transmission, a write instruction to the TXR\_RXR register will place the data into the TXR\_RXR register, which will be copied to the shift register at the end of the present transmission. When there is no data transmission in progress, a write instruction to the TXR\_RXR register will place the data directly into the shift register, resulting in the commencement of data transmission, and the TXIF bit being immediately set. When a frame transmission is complete, which happens after stop bits are sent or after the break frame, the TIDLE bit will be set. To clear the TIDLE bit the following software sequence is used:

1. A USR register access
2. A TXR\_RXR register write execution

Note that both the TXIF and TIDLE bits are cleared by the same software sequence.

### Transmitting Break

If the TXBRK bit is set and the state keeps for a time greater than  $(BRD+1) \times t_{TH}$  while TIDLE=1, then break characters will be sent on the next transmission. Break character transmission consists of a start bit, followed by  $13 \times N$  '0' bits and stop bits, where  $N=1, 2, \text{etc.}$  If a break character is to be transmitted then the TXBRK bit must be first set by the application program, and then cleared to generate the stop bits. Transmitting a break character will not generate a transmit interrupt. Note that a break condition length is at least 13 bits long. If the TXBRK bit is continually kept at a logic high level then the transmitter circuitry will transmit continuous break characters. After the application program has cleared the TXBRK bit, the transmitter will finish transmitting the last break character and subsequently send out two stop bits. The automatic logic highs at the end of the last break character will ensure that the start bit of the next frame is recognized.

### UART Receiver

The UART is capable of receiving word lengths of either 8 or 9 bits. If the BNO bit is set, the word length will be set to 9 bits with the MSB being stored in the RX8 bit of the UCR1 register. At the receiver core lies the Receive Serial Shift Register, commonly known as the RSR. The data which is received on the RX/TX external input pin is sent to the data recovery block. The data recovery block operating speed is 16 times that of the baud rate, while the main receive serial shifter operates at the baud rate. After the RX/TX pin is sampled for the stop bit, the received data in RSR is transferred to the receive data register, if the register is empty. The data which is received on the external RX/TX input pin is sampled three times by a majority detect circuit to determine the logic level that has been placed onto the RX/TX pin. It should be noted that the RSR register, unlike many other registers, is not directly mapped into the Data Memory area and as such is not available to the application program for direct read/write operations.

### Receiving Data

When the UART receiver is receiving data, the data is serially shifted in on the external RX/TX input pin, LSB first. In the read mode, the TXR\_RXR register forms a buffer between the internal bus and the receiver shift register. The TXR\_RXR register is a four byte deep FIFO data buffer, where four bytes can be held in the FIFO while a fifth byte can continue to be received. Note that the application program must ensure that the data is read from TXR\_RXR before the fifth byte has been completely shifted in, otherwise this fifth byte will be discarded and an overrun error OERR will be subsequently indicated. For continuous multi-byte data transmission, it is strongly recommended that the receiver uses two stop bits to avoid a receiving error caused by the accumulated error of the

receiver baud rate frequency.

The steps to initiate a data transfer can be summarized as follows:

- Make the correct selection of BNO, PRT1~PRT0, PREN and STOPS bits to define the word length, parity type and number of stop bits.
- Setup the BRDH and BRDL registers and the UMOD2~UMOD0 bits to select the desired baud rate.
- Set the RXEN bit to ensure that the RX/TX pin is used as a UART receiver pin.

At this point the receiver will be enabled which will begin to look for a start bit.

When a character is received the following sequence of events will occur:

- The RXIF bit in the USR register will be set when the TXR\_RXR register has data available, the number of the available data bytes can be checked by polling the RxCNT register content.
- When the contents of the shift register have been transferred to the TXR\_RXR register and reach Receiver FIFO trigger level if the RIE bit is set, then an interrupt will be generated.
- If during reception, a frame error, noise error, parity error, or an overrun error has been detected, then the error flags can be set.

The RXIF bit can be cleared using the following software sequence:

1. A USR register access
2. A TXR\_RXR register read execution

### **Receiving Break**

Any break character received by the UART will be managed as a framing error. The receiver will count and expect a certain number of bit times as specified by the values programmed into the BNO plus one or two stop bits. If the break is much longer than 13 bit times, the reception will be considered as complete after the number of bit times specified by BNO plus one or two stop bits. The RXIF bit is set, FERR is set, zeros are loaded into the receive data register, interrupts are generated if appropriate and the RIDLE bit is set. A break is regarded as a character that contains only zeros with the FERR flag set. If a long break signal has been detected, the receiver will regard it as a data frame including a start bit, data bits and the invalid stop bit and the FERR flag will be set. The receiver must wait for a valid stop bit before looking for the next start bit. The receiver will not make the assumption that the break condition on the line is the next start bit. The break character will be loaded into the buffer and no further data will be received until one or two stop bits are received. It should be noted that the RIDLE read only flag will go high when the stop bits have not yet been received. The reception of a break character on the UART registers will result in the following:

- The framing error flag, FERR, will be set.
- The receive data register, TXR\_RXR, will be cleared.
- The OERR, NF, PERR, RIDLE or RXIF flags will possibly be set.

### **Idle Status**

When the receiver is reading data, which means it will be in between the detection of a start bit and the reading of a stop bit, the receiver status flag in the USR register, otherwise known as the RIDLE flag, will have a zero value. In between the reception of a stop bit and the detection of the next start bit, the RIDLE flag will have a high value, which indicates the receiver is in an idle condition.

### Receiver Interrupt

The read only receive interrupt flag RXIF in the USR register is set by an edge generated by the receiver. An interrupt is generated if RIE=1, when a word is transferred from the Receive Shift Register, RSR, to the Receive Data Register, TXR\_RXR. An overrun error can also generate an interrupt if RIE=1.

When a subroutine will be called with an execution time longer than the time for UART to receive five data bytes, if the UART received data could not be read in time during the subroutine execution, clear the RXEN bit to zero in advance to suspend data reception. If the UART interrupt could not be served in time to process the overrun error during the subroutine execution, ensure that both EMI and RXEN bits are disabled during this period, and then enable EMI and RXEN again after the subroutine execution has been completed to continue the UART data reception.

### Managing Receiver Errors

Several types of reception errors can occur within the UART module, the following section describes the various types and how they are managed by the UART.

#### Overrun Error – OERR

The TXR\_RXR register is composed of a four byte deep FIFO data buffer, where four bytes can be held in the FIFO register, while a fifth byte can continue to be received. Before this fifth byte has been entirely shifted in, the data should be read from the TXR\_RXR register. If this is not done, the overrun error flag OERR will be consequently indicated.

In the event of an overrun error occurring, the following will happen:

- The OERR flag in the USR register will be set.
- The TXR\_RXR contents will not be lost.
- The shift register will be overwritten.
- An interrupt will be generated if the RIE bit is set.

When the OERR flag is set to “1”, it is necessary to read five data bytes from the four-byte deep receiver FIFO and the shift register immediately to avoid unexpected errors, such as the UART is unable to receive data. If such an error occurs, clear the RXEN bit to “0” then set it to “1” again to continue data reception.

The OERR flag can be cleared by an access to the USR register followed by a read to the TXR\_RXR register.

#### Noise Error – NF

Over-sampling is used for data recovery to identify valid incoming data and noise. If noise is detected within a frame the following will occur:

- The read only noise flag, NF, in the USR register will be set on the rising edge of the RXIF bit.
- Data will be transferred from the Shift register to the TXR\_RXR register.
- No interrupt will be generated. However this bit rises at the same time as the RXIF bit which itself generates an interrupt.

Note that the NF flag is reset by a USR register read operation followed by a TXR\_RXR register read operation.

#### Framing Error – FERR

The read only framing error flag, FERR, in the USR register, is set if a zero is detected instead of stop bits. If two stop bits are selected, both stop bits must be high; otherwise the FERR flag will

be set. The FERR flag and the received data will be recorded in the USR and TXR\_RXR registers respectively, and the flag is cleared in any reset.

### **Parity Error – PERR**

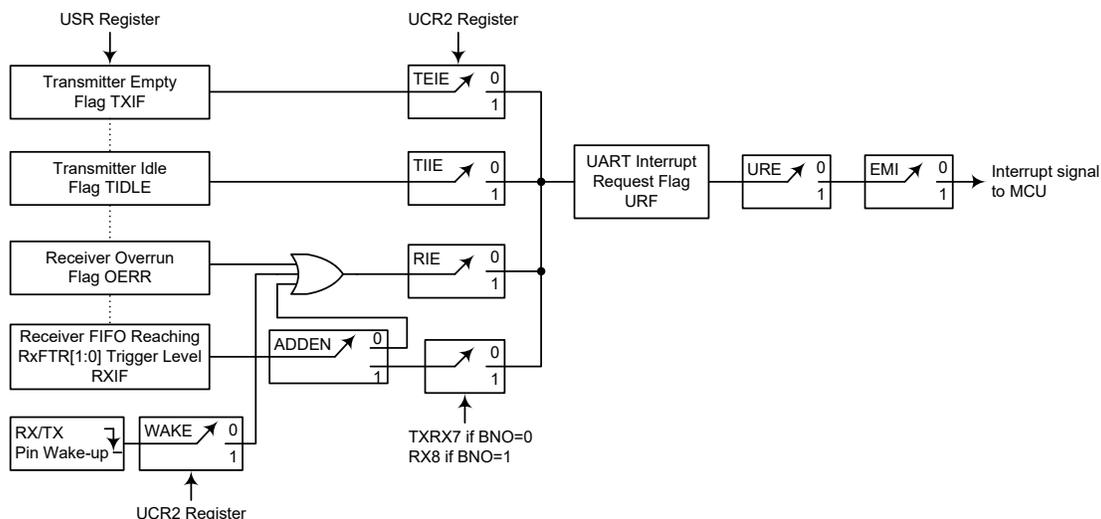
The read only parity error flag, PERR, in the USR register, is set if the parity of the received word is incorrect. This error flag is only applicable if the parity is enabled, PREN=1, and if the parity type, odd, even, mark or space, is selected. The read only PERR flag and the received data will be recorded in the USR and TXR\_RXR registers respectively. It is cleared on any reset, it should be noted that the flags, FERR and PERR, in the USR register should first be read by the application program before reading the data word.

### **UART Interrupt Structure**

Several individual UART conditions can generate a UART interrupt. When these conditions exist, a low pulse will be generated to get the attention of the microcontroller. These conditions are a transmitter data register empty, transmitter idle, receiver reaching FIFO trigger level, receiver overrun, address detect and an RX/TX pin wake-up. When any of these conditions are created, if the global interrupt enable bit and its corresponding interrupt control bit are enabled and the stack is not full, the program will jump to its corresponding interrupt vector where it can be serviced before returning to the main program. Four of these conditions have the corresponding USR register flags which will generate a UART interrupt if its associated interrupt enable control bit in the UCR2 register is set. The two transmitter interrupt conditions have their own corresponding enable control bits, while the two receiver interrupt conditions have a shared enable control bit. These enable bits can be used to mask out individual UART interrupt sources.

The address detect condition, which is also a UART interrupt source, does not have an associated flag, but will generate a UART interrupt when an address detect condition occurs if its function is enabled by setting the ADDEN bit in the UCR2 register. An RX/TX pin wake-up, which is also a UART interrupt source, does not have an associated flag, but will generate a UART interrupt if the UART clock ( $f_{H}$ ) source is switched off and the WAKE and RIE bits in the UCR2 register are set when a falling edge on the RX/TX pin occurs.

Note that the USR register flags are read only and cannot be cleared or set by the application program, neither will they be cleared when the program jumps to the corresponding interrupt servicing routine, as is the case for some of the other interrupts. The flags will be cleared automatically when certain actions are taken by the UART, the details of which are given in the UART register section. The overall UART interrupt can be disabled or enabled by the related interrupt enable control bits in the interrupt control registers of the microcontroller to decide whether the interrupt requested by the UART module is masked out or allowed.



UART Interrupt Structure

### Address Detect Mode

Setting the Address Detect Mode bit, ADDEN, in the UCR2 register, enables this special mode. If this bit is enabled then an additional qualifier will be placed on the generation of a Receiver Data Available interrupt, which is requested by the RXIF flag. If the ADDEN bit is enabled, then when data is available, an interrupt will only be generated, if the highest received bit has a high value. Note that the URE and EMI interrupt enable bits must also be enabled for correct interrupt generation. This highest address bit is the 9th bit if BNO=1 or the 8th bit if BNO=0. If this bit is high, then the received word will be defined as an address rather than data. A Data Available interrupt will be generated every time the last bit of the received word is set. If the ADDEN bit is not enabled, then a Receiver Data Available interrupt will be generated each time the RXIF flag is set, irrespective of the data last bit status. The address detect mode and parity enable are mutually exclusive functions. Therefore if the address detect mode is enabled, then to ensure correct operation, the parity function should be disabled by resetting the parity enable bit PREN to zero.

ADDEN	9th Bit if BNO=1, 8th Bit if BNO=0	UART Interrupt Generated
0	0	√
	1	√
1	0	×
	1	√

ADDEN Bit Function

### UART Power Down and Wake-up

When the UART clock ( $f_{H}$ ) is off, the UART will cease to function, all clock sources to the module are shutdown. If the UART clock ( $f_{H}$ ) is off while a transmission is still in progress, then the transmission will be paused until the UART clock source derived from the microcontroller is activated. In a similar way, if the MCU enters the IDLE or SLEEP mode while receiving data, then the reception of data will likewise be paused. When the MCU enters the IDLE or SLEEP mode, note that the USR, UCR1, UCR2, UCR3, UFCR, RxCNT, TXR\_RXR as well as the BRDH and BRDL registers will not be affected. It is recommended to make sure first that the UART data transmission or reception has been finished before the microcontroller enters the IDLE or SLEEP mode.

The UART function contains a receiver RX/TX pin wake-up function, which is enabled or disabled

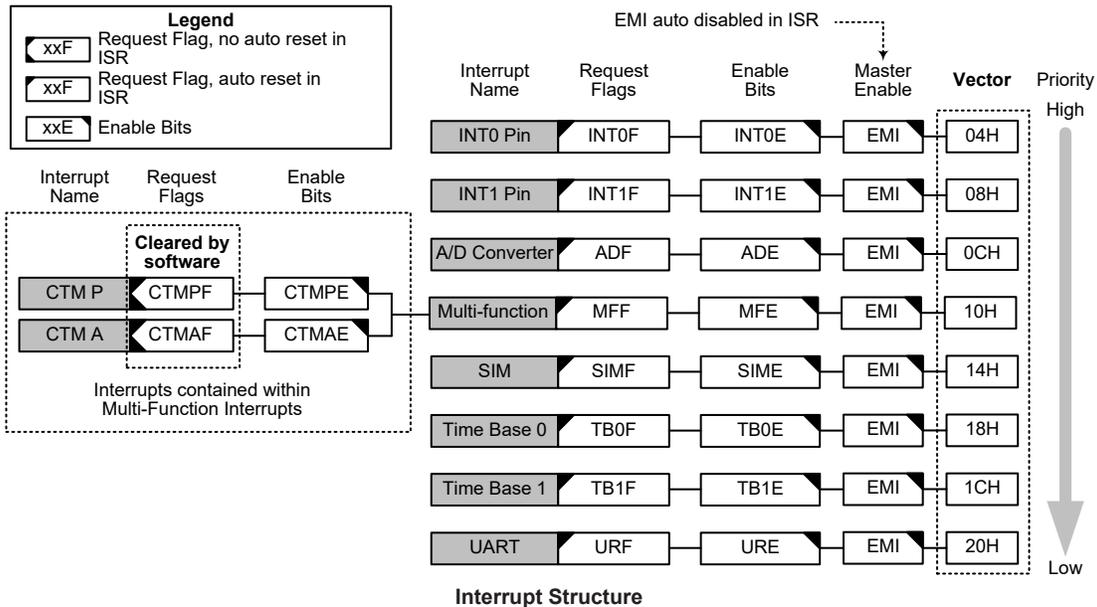
by the WAKE bit in the UCR2 register. If this bit, along with the UART enable bit, UARTEN, the receiver enable bit, RXEN and the receiver interrupt bit, RIE, are all set when the UART clock ( $f_{H}$ ) is off, then a falling edge on the RX/TX pin will trigger an RX/TX pin wake-up UART interrupt. Note that as it takes certain system clock cycles after a wake-up, before normal microcontroller operation resumes, any data received during this time on the RX/TX pin will be ignored.

For a UART wake-up interrupt to occur, in addition to the bits for the wake-up being set, the global interrupt enable bit, EMI, and the UART interrupt enable bit, URE, must be set. If the EMI and URE bits are not set then only a wake up event will occur and no interrupt will be generated. Note also that as it takes certain system clock cycles after a wake-up before normal microcontroller resumes, the UART interrupt will not be generated until after this time has elapsed.

## Interrupts

Interrupts are an important part of any microcontroller system. When an external event or an internal function such as a Timer Module or an A/D converter requires microcontroller attention, their corresponding interrupt will enforce a temporary suspension of the main program allowing the microcontroller to direct attention to their respective needs. The device contains several external interrupt and internal interrupt functions. The external interrupts are generated by the action of the external INT0~INT1 pins, while the internal interrupts are generated by various internal functions such as the CTM, Time Bases, SIM, UART and the A/D converter, etc.

The various interrupt enable bits, together with their associated request flags, are shown in the accompanying diagrams with their order of priority. Some interrupt sources have their own individual vector while others share the same multi-function interrupt vector.



## Interrupt Registers

Overall interrupt control, which basically means the setting of request flags when certain microcontroller conditions occur and the setting of interrupt enable bits by the application program, is controlled by a series of registers, located in the Special Purpose Data Memory, as shown in the accompanying table. The number of registers falls into three categories. The first is the INTC0~INTC2 registers which setup the primary interrupts, the second is the MFI register which

setup the Multi-function interrupt. Finally there is an INTEG register to setup the external interrupt trigger edge type.

Each register contains a number of enable bits to enable or disable individual registers as well as interrupt flags to indicate the presence of an interrupt request. The naming convention of these follows a specific pattern. First is listed an abbreviated interrupt type, then the (optional) number of that interrupt followed by either an “E” for enable/disable bit or “F” for request flag.

Function	Enable Bit	Request Flag	Notes
Global	EMI	—	—
INTn Pin	INTnE	INTnF	n=0~1
A/D Converter	ADE	ADF	—
Multi-function	MFE	MFF	—
CTM	CTMPE	CTMPF	—
	CTMAE	CTMAF	—
Time Base	TBnE	TBnF	n=0~1
SIM	SIME	SIMF	—
UART	URE	URF	—

**Interrupt Register Bit Naming Conventions**

Register Name	Bit							
	7	6	5	4	3	2	1	0
INTEG	—	—	—	—	INT1S1	INT1S0	INT0S1	INT0S0
INTC0	—	ADF	INT1F	INT0F	ADE	INT1E	INT0E	EMI
INTC1	TB1F	TB0F	SIMF	MFF	TB1E	TB0E	SIME	MFE
INTC2	—	—	—	URF	—	—	—	URE
MFI	—	—	CTMAF	CTMPF	—	—	CTMAE	CTMPE

**Interrupt Register List**

• **INTEG Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	INT1S1	INT1S0	INT0S1	INT0S0
R/W	—	—	—	—	R/W	R/W	R/W	R/W
POR	—	—	—	—	0	0	0	0

Bit 7~4 Unimplemented, read as “0”

Bit 3~2 **INT1S1~INT1S0**: Interrupt edge control for INT1 pin  
 00: Disable  
 01: Rising edge  
 10: Falling edge  
 11: Rising and falling edges

Bit 1~0 **INT0S1~INT0S0**: Interrupt edge control for INT0 pin  
 00: Disable  
 01: Rising edge  
 10: Falling edge  
 11: Rising and falling edges

• **INTC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	ADF	INT1F	INT0F	ADE	INT1E	INT0E	EMI
R/W	—	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	0	0	0	0	0	0	0

- Bit 7      Unimplemented, read as “0”
- Bit 6      **ADF**: A/D Converter interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 5      **INT1F**: INT1 interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 4      **INT0F**: INT0 interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 3      **ADE**: A/D Converter interrupt control  
             0: Disable  
             1: Enable
- Bit 2      **INT1E**: INT1 interrupt control  
             0: Disable  
             1: Enable
- Bit 1      **INT0E**: INT0 interrupt control  
             0: Disable  
             1: Enable
- Bit 0      **EMI**: Global interrupt control  
             0: Disable  
             1: Enable

• **INTC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	TB1F	TB0F	SIMF	MFF	TB1E	TB0E	SIME	MFE
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **TB1F**: Time Base 1 interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 6      **TB0F**: Time Base 0 interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 5      **SIMF**: SIM interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 4      **MFF**: Multi-function interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 3      **TB1E**: Time Base 1 interrupt control  
             0: Disable  
             1: Enable
- Bit 2      **TB0E**: Time Base 0 interrupt control  
             0: Disable  
             1: Enable

- Bit 1      **SIME**: SIM interrupt control  
0: Disable  
1: Enable
- Bit 0      **MFE**: Multi-function interrupt control  
0: Disable  
1: Enable

• **INTC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	URF	—	—	—	URE
R/W	—	—	—	R/W	—	—	—	R/W
POR	—	—	—	0	—	—	—	0

- Bit 7~5      Unimplemented, read as “0”
- Bit 4      **URF**: UART interrupt request flag  
0: No request  
1: Interrupt request
- Bit 3~1      Unimplemented, read as “0”
- Bit 0      **URE**: UART interrupt control  
0: Disable  
1: Enable

• **MFI Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	CTMAF	CTMPF	—	—	CTMAE	CTMPE
R/W	—	—	R/W	R/W	—	—	R/W	R/W
POR	—	—	0	0	—	—	0	0

- Bit 7~6      Unimplemented, read as “0”
- Bit 5      **CTMAF**: CTM Comparator A match interrupt request flag  
0: No request  
1: Interrupt request  
  
Note that this bit must be cleared to zero by the application program when the interrupt is serviced.
- Bit 4      **CTMPF**: CTM Comparator P match interrupt request flag  
0: No request  
1: Interrupt request  
  
Note that this bit must be cleared to zero by the application program when the interrupt is serviced.
- Bit 3~2      Unimplemented, read as “0”
- Bit 1      **CTMAE**: CTM Comparator A match interrupt control  
0: Disable  
1: Enable
- Bit 0      **CTMPE**: CTM Comparator P match interrupt control  
0: Disable  
1: Enable

**Interrupt Operation**

When the conditions for an interrupt event occur, such as a TM Comparator P or Comparator A match or A/D conversion completion etc., the relevant interrupt request flag will be set. Whether the request flag actually generates a program jump to the relevant interrupt vector is determined by the condition of the interrupt enable bit. If the enable bit is set high then the program will jump to its relevant vector; if the enable bit is zero then although the interrupt request flag is set an actual interrupt will not be generated and the program will not jump to the relevant interrupt vector. The

global interrupt enable bit, if cleared to zero, will disable all interrupts.

When an interrupt is generated, the Program Counter, which stores the address of the next instruction to be executed, will be transferred onto the stack. The Program Counter will then be loaded with a new address which will be the value of the corresponding interrupt vector. The microcontroller will then fetch its next instruction from this interrupt vector. The instruction at this vector will usually be a “JMP” which will jump to another section of program which is known as the interrupt service routine. Here is located the code to control the appropriate interrupt. The interrupt service routine must be terminated with a “RETI”, which retrieves the original Program Counter address from the stack and allows the microcontroller to continue with normal execution at the point where the interrupt occurred.

Once an interrupt subroutine is serviced, all the other interrupts will be blocked, as the global interrupt enable bit, EMI bit will be cleared automatically. This will prevent any further interrupt nesting from occurring. However, if other interrupt requests occur during this interval, although the interrupt will not be immediately serviced, the request flag will still be recorded.

If an interrupt requires immediate servicing while the program is already in another interrupt service routine, the EMI bit should be set after entering the routine, to allow interrupt nesting. If the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the Stack Pointer is decremented. If immediate service is desired, the stack must be prevented from becoming full. In case of simultaneous requests, the interrupt structure diagram shows the priority that is applied. All of the interrupt request flags when set will wake up the device if it is in SLEEP or IDLE Mode, however to prevent a wake-up from occurring the corresponding flag should be set before the device is in SLEEP or IDLE Mode.

## **External Interrupts**

The external interrupts are controlled by signal transitions on the pins INT0 and INT1. An external interrupt request will take place when the external interrupt request flags, INT0F~INT1F, are set, which will occur when a transition, whose type is chosen by the edge select bits, appears on the external interrupt pins. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and respective external interrupt enable bits, INT0E~INT1E, must first be set. Additionally the correct interrupt edge type must be selected using the INTEG register to enable the external interrupt function and to choose the trigger edge type. As the external interrupt pins are pin-shared with I/O pins, they can only be configured as external interrupt pins if their external interrupt enable bit in the corresponding interrupt register has been set and the external interrupt pin is selected by the corresponding pin-shared function selection bits. The pin must also be setup as an input by setting the corresponding bit in the port control register. When the interrupt is enabled, the stack is not full and the correct transition type appears on the external interrupt pin, a subroutine call to the external interrupt vector, will take place. When the interrupt is serviced, the external interrupt request flags, INT0F~INT1F, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts. Note that any pull-high resistor selections on the external interrupt pins will remain valid even if the pin is used as an external interrupt input.

The INTEG register is used to select the type of active edge that will trigger the external interrupt. A choice of either rising or falling or both edge types can be chosen to trigger an external interrupt. Note that the INTEG register can also be used to disable the external interrupt function.

## **A/D Converter Interrupt**

An A/D Converter Interrupt request will take place when the A/D Converter Interrupt request flag, ADF, is set, which occurs when the A/D conversion process finishes. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and A/D Interrupt

enable bit, ADE, must first be set. When the interrupt is enabled, the stack is not full and the A/D conversion process has ended, a subroutine call to the A/D Interrupt vector, will take place. When the A/D Converter Interrupt is serviced, the A/D Interrupt flag, ADF, will be automatically cleared. The EMI bit will also be automatically cleared to disable other interrupts.

### **Multi-function Interrupt**

Within the device there is one Multi-function interrupt. Unlike the other independent interrupts, this interrupt has no independent source, but rather is formed from other existing interrupt source, namely the CTM interrupt.

A Multi-function interrupt request will take place when the Multi-function interrupt request flag MFF is set. The Multi-function interrupt flag will be set when any of their included functions generate an interrupt request flag. When the Multi-function interrupt is enabled and the stack is not full, and one of the interrupts contained within each of Multi-function interrupt occurs, a subroutine call to the relevant Multi-function interrupt vector will take place. When the interrupt is serviced, the related Multi-Function request flag will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

However, it must be noted that, although the Multi-function Interrupt request flags will be automatically reset when the interrupt is serviced, the request flags from the original source of the Multi-function interrupt, namely the TM interrupt request flags will not be automatically reset and must be manually reset by the application program.

### **TM Interrupts**

The Timer Module Interrupts are contained within the Multi-function Interrupt. The Compact Type TM has two interrupts, one comes from the comparator A match situation and the other comes from the comparator P match situation. All of the Compact Type TM interrupts are contained within the Multi-function Interrupts. For the TM there are two interrupt request flags and two enable control bits. A TM interrupt request will take place when any of the TM request flags are set, a situation which occurs when a TM comparator P or A match situation happens.

To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, respective TM Interrupt enable bit, and relevant Multi-function Interrupt enable bit, MFE, must first be set. When the interrupt is enabled, the stack is not full and a TM comparator match situation occurs, a subroutine call to the relevant Multi-function Interrupt vector locations, will take place. When the TM interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts, however only the related MFF flag will be automatically cleared. As the TM interrupt request flags will not be automatically cleared, they have to be cleared by the application program.

### **Serial Interface Module Interrupt**

The Serial Interface Module Interrupt is also known as the SIM interrupt. A SIM Interrupt request will take place when the SIM Interrupt request flag, SIMF, is set, which occurs when a byte of data has been received or transmitted by the SIM interface, an I<sup>2</sup>C slave address match or I<sup>2</sup>C bus time-out occurs. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and SIM Interrupt enable bit, SIME, must first be set. When the interrupt is enabled, the stack is not full and any of the above described situations occurs, a subroutine call to the corresponding SIM Interrupt vector, will take place. When the interrupt is serviced, the SIMF flag will be automatically cleared and the EMI bit will be automatically cleared to disable other interrupts.

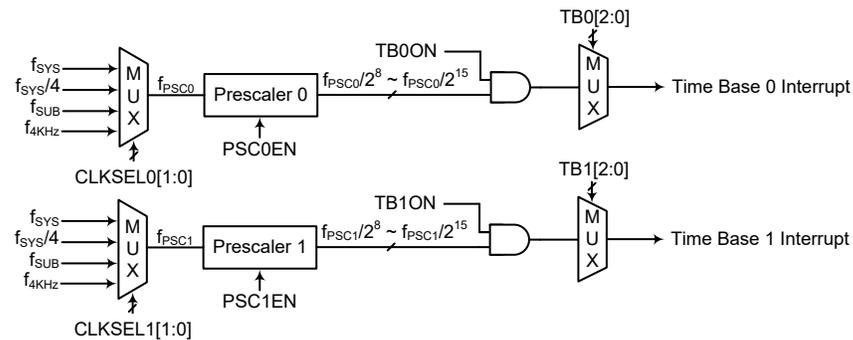
## UART Interrupt

The UART Interrupt is controlled by several UART transfer conditions. When one of these conditions occurs, an interrupt pulse will be generated to get the attention of the microcontroller. These conditions are a transmitter data register empty, transmitter idle, receiver reaching FIFO trigger level, receiver overrun, address detect and an RX/TX pin wake-up. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and UART Interrupt enable bit, URE, must first be set. When the interrupt is enabled, the stack is not full and any of the conditions described above occurs, a subroutine call to the corresponding UART Interrupt vector, will take place. When the interrupt is serviced, the UART Interrupt flag, URF, will be automatically cleared. The EMI bit will also be automatically cleared to disable other interrupts. However, the USR register flags will only be cleared when certain actions are taken by the UART, the details of which are given in the UART section.

## Time Base Interrupts

The function of the Time Base Interrupts is to provide regular time signal in the form of an internal interrupt. They are controlled by the overflow signals from their respective timer functions. When these happens their respective interrupt request flags, TB0F or TB1F will be set. To allow the program to branch to their respective interrupt vector addresses, the global interrupt enable bit, EMI and Time Base enable bits, TB0E or TB1E, must first be set. When the interrupt is enabled, the stack is not full and the Time Base overflows, a subroutine call to their respective vector locations will take place. When the interrupt is serviced, the respective interrupt request flag, TB0F or TB1F, will be automatically reset and the EMI bit will be cleared to disable other interrupts.

The purpose of the Time Base Interrupt is to provide an interrupt signal at fixed time periods. Their clock sources,  $f_{PSC0}$  or  $f_{PSC1}$ , originate from the internal clock source  $f_{SYS}$ ,  $f_{SYS}/4$  or  $f_{SUB}$  and then passes through a divider, the division ratio of which is selected by programming the appropriate bits in the TB0C and TB1C registers to obtain longer interrupt periods whose value ranges. The clock source that generate  $f_{PSC0}$  or  $f_{PSC1}$ , which in turn controls the Time Base interrupt period, is selected using the CLKSEL0[1:0] and CLKSEL1[1:0] bits in the PSC0R or PSC1R register respectively.



**Time Base Interrupts**

**• PSC0R Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	CLKSEL01	CLKSEL00
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as “0”

Bit 1~0 **CLKSEL01~CLKSEL00**: Prescaler 0 clock source selection  
 00:  $f_{SYS}$   
 01:  $f_{SYS}/4$   
 1x:  $f_{SUB}$

**• PSC1R Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	CLKSEL11	CLKSEL10
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as “0”

Bit 1~0 **CLKSEL11~CLKSEL10**: Prescaler 1 clock source selection  
 00:  $f_{SYS}$   
 01:  $f_{SYS}/4$   
 1x:  $f_{SUB}$

**• TB0C Register**

Bit	7	6	5	4	3	2	1	0
Name	TB0ON	—	—	—	—	TB02	TB01	TB00
R/W	R/W	—	—	—	—	R/W	R/W	R/W
POR	0	—	—	—	—	0	0	0

Bit 7 **TB0ON**: Time Base 0 Control  
 0: Disable  
 1: Enable

Bit 6~3 Unimplemented, read as “0”

Bit 2~0 **TB02~TB00**: Select Time Base 0 Time-out Period  
 000:  $2^8/f_{PSC0}$   
 001:  $2^9/f_{PSC0}$   
 010:  $2^{10}/f_{PSC0}$   
 011:  $2^{11}/f_{PSC0}$   
 100:  $2^{12}/f_{PSC0}$   
 101:  $2^{13}/f_{PSC0}$   
 110:  $2^{14}/f_{PSC0}$   
 111:  $2^{15}/f_{PSC0}$

**• TB1C Register**

Bit	7	6	5	4	3	2	1	0
Name	TB1ON	—	—	—	—	TB12	TB11	TB10
R/W	R/W	—	—	—	—	R/W	R/W	R/W
POR	0	—	—	—	—	0	0	0

Bit 7 **TB1ON**: Time Base 1 Control  
 0: Disable  
 1: Enable

Bit 6~3 Unimplemented, read as “0”

Bit 2~0	<b>TB12~TB10:</b> Select Time Base 1 Time-out Period
	000: $2^8/f_{PSC1}$
	001: $2^9/f_{PSC1}$
	010: $2^{10}/f_{PSC1}$
	011: $2^{11}/f_{PSC1}$
	100: $2^{12}/f_{PSC1}$
	101: $2^{13}/f_{PSC1}$
	110: $2^{14}/f_{PSC1}$
	111: $2^{15}/f_{PSC1}$

### Interrupt Wake-up Function

Each of the interrupt functions has the capability of waking up the microcontroller when in the SLEEP or IDLE Mode. A wake-up is generated when an interrupt request flag changes from low to high and is independent of whether the interrupt is enabled or not. Therefore, even though the device is in the SLEEP or IDLE Mode and its system oscillator stopped, situations such as external edge transitions on the external interrupt pins or a low power supply voltage may cause their respective interrupt flag to be set high and consequently generate an interrupt. Care must therefore be taken if spurious wake-up situations are to be avoided. If an interrupt wake-up function is to be disabled then the corresponding interrupt request flag should be set high before the device enters the SLEEP or IDLE Mode. The interrupt enable bits have no effect on the interrupt wake-up function.

### Programming Considerations

By disabling the relevant interrupt enable bits, a requested interrupt can be prevented from being serviced, however, once an interrupt request flag is set, it will remain in this condition in the interrupt register until the corresponding interrupt is serviced or until the request flag is cleared by the application program.

Where a certain interrupt is contained within a Multi-function interrupt, then when the interrupt service routine is executed, as only the Multi-function interrupt request flags, MFF, will be automatically cleared, the individual request flag for the function needs to be cleared by the application program.

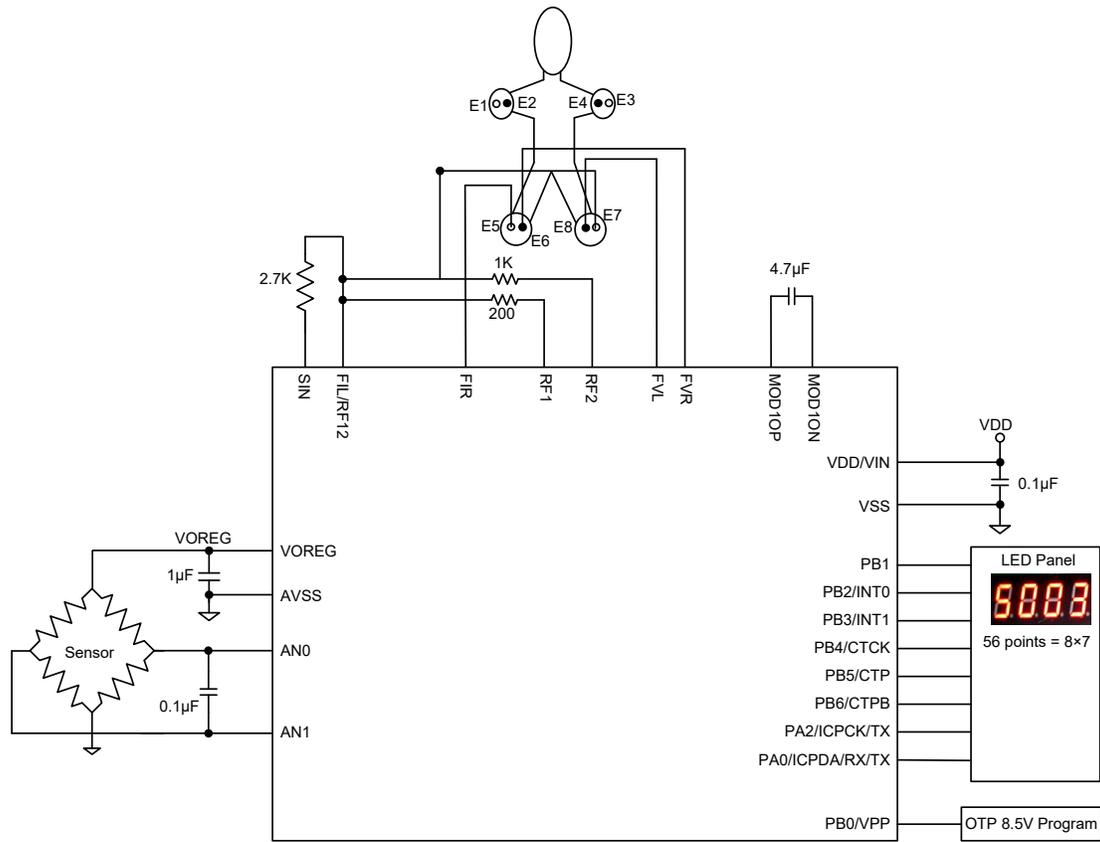
It is recommended that programs do not use the “CALL” instruction within the interrupt service subroutine. Interrupts often occur in an unpredictable manner or need to be serviced immediately. If only one stack is left and the interrupt is not well controlled, the original control sequence will be damaged once a CALL subroutine is executed in the interrupt subroutine.

Every interrupt has the capability of waking up the microcontroller when it is in the SLEEP or IDLE Mode, the wake up being generated when the interrupt request flag changes from low to high. If it is required to prevent a certain interrupt from waking up the microcontroller then its respective request flag should be first set high before enter SLEEP or IDLE Mode.

As only the Program Counter is pushed onto the stack, then when the interrupt is serviced, if the contents of the accumulator, status register or other registers are altered by the interrupt service program, their contents should be saved to the memory at the beginning of the interrupt service routine.

To return from an interrupt subroutine, either a RET or RETI instruction may be executed. The RETI instruction in addition to executing a return to the main program also automatically sets the EMI bit high to allow further interrupts. The RET instruction however only executes a return to the main program leaving the EMI bit in its present zero state and therefore disabling the execution of further interrupts.

**Application Circuits**



## Instruction Set

### Introduction

Central to the successful operation of any microcontroller is its instruction set, which is a set of program instruction codes that directs the microcontroller to perform certain operations. In the case of Holtek microcontroller, a comprehensive and flexible set of over 60 instructions is provided to enable programmers to implement their application with the minimum of programming overheads.

For easier understanding of the various instruction codes, they have been subdivided into several functional groupings.

### Instruction Timing

Most instructions are implemented within one instruction cycle. The exceptions to this are branch, call, or table read instructions where two instruction cycles are required. One instruction cycle is equal to 4 system clock cycles, therefore in the case of an 8MHz system oscillator, most instructions would be implemented within 0.5 $\mu$ s and branch or call instructions would be implemented within 1 $\mu$ s. Although instructions which require one more cycle to implement are generally limited to the JMP, CALL, RET, RETI and table read instructions, it is important to realize that any other instructions which involve manipulation of the Program Counter Low register or PCL will also take one more cycle to implement. As instructions which change the contents of the PCL will imply a direct jump to that new address, one more cycle will be required. Examples of such instructions would be “CLR PCL” or “MOV PCL, A”. For the case of skip instructions, it must be noted that if the result of the comparison involves a skip operation then this will also take one more cycle, if no skip is involved then only one cycle is required.

### Moving and Transferring Data

The transfer of data within the microcontroller program is one of the most frequently used operations. Making use of several kinds of MOV instructions, data can be transferred from registers to the Accumulator and vice-versa as well as being able to move specific immediate data directly into the Accumulator. One of the most important data transfer applications is to receive data from the input ports and transfer data to the output ports.

### Arithmetic Operations

The ability to perform certain arithmetic operations and data manipulation is a necessary feature of most microcontroller applications. Within the Holtek microcontroller instruction set are a range of add and subtract instruction mnemonics to enable the necessary arithmetic to be carried out. Care must be taken to ensure correct handling of carry and borrow data when results exceed 255 for addition and less than 0 for subtraction. The increment and decrement instructions such as INC, INCA, DEC and DECA provide a simple means of increasing or decreasing by a value of one of the values in the destination specified.

## Logical and Rotate Operation

The standard logical operations such as AND, OR, XOR and CPL all have their own instruction within the Holtek microcontroller instruction set. As with the case of most instructions involving data manipulation, data must pass through the Accumulator which may involve additional programming steps. In all logical data operations, the zero flag may be set if the result of the operation is zero. Another form of logical data manipulation comes from the rotate instructions such as RR, RL, RRC and RLC which provide a simple means of rotating one bit right or left. Different rotate instructions exist depending on program requirements. Rotate instructions are useful for serial port programming applications where data can be rotated from an internal register into the Carry bit from where it can be examined and the necessary serial bit set high or low. Another application which rotate data operations are used is to implement multiplication and division calculations.

## Branches and Control Transfer

Program branching takes the form of either jumps to specified locations using the JMP instruction or to a subroutine using the CALL instruction. They differ in the sense that in the case of a subroutine call, the program must return to the instruction immediately when the subroutine has been carried out. This is done by placing a return instruction “RET” in the subroutine which will cause the program to jump back to the address right after the CALL instruction. In the case of a JMP instruction, the program simply jumps to the desired location. There is no requirement to jump back to the original jumping off point as in the case of the CALL instruction. One special and extremely useful set of branch instructions are the conditional branches. Here a decision is first made regarding the condition of a certain data memory or individual bits. Depending upon the conditions, the program will continue with the next instruction or skip over it and jump to the following instruction. These instructions are the key to decision making and branching within the program perhaps determined by the condition of certain input switches or by the condition of internal data bits.

## Bit Operations

The ability to provide single bit operations on Data Memory is an extremely flexible feature of all Holtek microcontrollers. This feature is especially useful for output port bit programming where individual bits or port pins can be directly set high or low using either the “SET [m].i” or “CLR [m].i” instructions respectively. The feature removes the need for programmers to first read the 8-bit output port, manipulate the input data to ensure that other bits are not changed and then output the port with the correct new data. This read-modify-write process is taken care of automatically when these bit operation instructions are used.

## Table Read Operations

Data storage is normally implemented by using registers. However, when working with large amounts of fixed data, the volume involved often makes it inconvenient to store the fixed data in the Data Memory. To overcome this problem, Holtek microcontrollers allow an area of Program Memory to be setup as a table where data can be directly stored. A set of easy to use instructions provides the means by which this fixed data can be referenced and retrieved from the Program Memory.

## Other Operations

In addition to the above functional instructions, a range of other instructions also exist such as the “HALT” instruction for Power-down operations and instructions to control the operation of the Watchdog Timer for reliable program operations under extreme electric or electromagnetic environments. For their relevant operations, refer to the functional related sections.

## Instruction Set Summary

The instructions related to the data memory access in the following table can be used when the desired data memory is located in Data Memory sector 0.

### Table Conventions

- x: Bits immediate data
- m: Data Memory address
- A: Accumulator
- i: 0~7 number of bits
- addr: Program memory address

Mnemonic	Description	Cycles	Flag Affected
<b>Arithmetic</b>			
ADD A,[m]	Add Data Memory to ACC	1	Z, C, AC, OV, SC
ADDM A,[m]	Add ACC to Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC
ADD A,x	Add immediate data to ACC	1	Z, C, AC, OV, SC
ADC A,[m]	Add Data Memory to ACC with Carry	1	Z, C, AC, OV, SC
ADCM A,[m]	Add ACC to Data memory with Carry	1 <sup>Note</sup>	Z, C, AC, OV, SC
SUB A,x	Subtract immediate data from the ACC	1	Z, C, AC, OV, SC, CZ
SUB A,[m]	Subtract Data Memory from ACC	1	Z, C, AC, OV, SC, CZ
SUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
SBC A,x	Subtract immediate data from ACC with Carry	1	Z, C, AC, OV, SC, CZ
SBC A,[m]	Subtract Data Memory from ACC with Carry	1	Z, C, AC, OV, SC, CZ
SBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
DAA [m]	Decimal adjust ACC for Addition with result in Data Memory	1 <sup>Note</sup>	C
<b>Logic Operation</b>			
AND A,[m]	Logical AND Data Memory to ACC	1	Z
OR A,[m]	Logical OR Data Memory to ACC	1	Z
XOR A,[m]	Logical XOR Data Memory to ACC	1	Z
ANDM A,[m]	Logical AND ACC to Data Memory	1 <sup>Note</sup>	Z
ORM A,[m]	Logical OR ACC to Data Memory	1 <sup>Note</sup>	Z
XORM A,[m]	Logical XOR ACC to Data Memory	1 <sup>Note</sup>	Z
AND A,x	Logical AND immediate Data to ACC	1	Z
OR A,x	Logical OR immediate Data to ACC	1	Z
XOR A,x	Logical XOR immediate Data to ACC	1	Z
CPL [m]	Complement Data Memory	1 <sup>Note</sup>	Z
CPLA [m]	Complement Data Memory with result in ACC	1	Z
<b>Increment &amp; Decrement</b>			
INCA [m]	Increment Data Memory with result in ACC	1	Z
INC [m]	Increment Data Memory	1 <sup>Note</sup>	Z
DECA [m]	Decrement Data Memory with result in ACC	1	Z
DEC [m]	Decrement Data Memory	1 <sup>Note</sup>	Z
<b>Rotate</b>			
RRA [m]	Rotate Data Memory right with result in ACC	1	None
RR [m]	Rotate Data Memory right	1 <sup>Note</sup>	None
RRCA [m]	Rotate Data Memory right through Carry with result in ACC	1	C
RRC [m]	Rotate Data Memory right through Carry	1 <sup>Note</sup>	C
RLA [m]	Rotate Data Memory left with result in ACC	1	None
RL [m]	Rotate Data Memory left	1 <sup>Note</sup>	None
RLCA [m]	Rotate Data Memory left through Carry with result in ACC	1	C
RLC [m]	Rotate Data Memory left through Carry	1 <sup>Note</sup>	C

Mnemonic	Description	Cycles	Flag Affected
<b>Data Move</b>			
MOV A,[m]	Move Data Memory to ACC	1	None
MOV [m],A	Move ACC to Data Memory	1 <sup>Note</sup>	None
MOV A,x	Move immediate data to ACC	1	None
<b>Bit Operation</b>			
CLR [m].i	Clear bit of Data Memory	1 <sup>Note</sup>	None
SET [m].i	Set bit of Data Memory	1 <sup>Note</sup>	None
<b>Branch Operation</b>			
JMP addr	Jump unconditionally	2	None
SZ [m]	Skip if Data Memory is zero	1 <sup>Note</sup>	None
SZA [m]	Skip if Data Memory is zero with data movement to ACC	1 <sup>Note</sup>	None
SZ [m].i	Skip if bit i of Data Memory is zero	1 <sup>Note</sup>	None
SNZ [m]	Skip if Data Memory is not zero	1 <sup>Note</sup>	None
SNZ [m].i	Skip if bit i of Data Memory is not zero	1 <sup>Note</sup>	None
SIZ [m]	Skip if increment Data Memory is zero	1 <sup>Note</sup>	None
SDZ [m]	Skip if decrement Data Memory is zero	1 <sup>Note</sup>	None
SIZA [m]	Skip if increment Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
SDZA [m]	Skip if decrement Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
CALL addr	Subroutine call	2	None
RET	Return from subroutine	2	None
RET A,x	Return from subroutine and load immediate data to ACC	2	None
RETI	Return from interrupt	2	None
<b>Table Read Operation</b>			
TABRD [m]	Read table (specific page) to TBLH and Data Memory	2 <sup>Note</sup>	None
TABRDL [m]	Read table (last page) to TBLH and Data Memory	2 <sup>Note</sup>	None
ITABRD [m]	Increment table pointer TBLP first and Read table (specific page) to TBLH and Data Memory	2 <sup>Note</sup>	None
ITABRDL [m]	Increment table pointer TBLP first and Read table (last page) to TBLH and Data Memory	2 <sup>Note</sup>	None
<b>Miscellaneous</b>			
NOP	No operation	1	None
CLR [m]	Clear Data Memory	1 <sup>Note</sup>	None
SET [m]	Set Data Memory	1 <sup>Note</sup>	None
CLR WDT	Clear Watchdog Timer	1	TO, PDF
SWAP [m]	Swap nibbles of Data Memory	1 <sup>Note</sup>	None
SWAPA [m]	Swap nibbles of Data Memory with result in ACC	1	None
HALT	Enter power down mode	1	TO, PDF

Note: 1. For skip instructions, if the result of the comparison involves a skip then two cycles are required, if no skip takes place only one cycle is required.

2. Any instruction which changes the contents of the PCL will also require 2 cycles for execution.

### Extended Instruction Set

The extended instructions are used to support the full range address access for the data memory. When the accessed data memory is located in any data memory sector except sector 0, the extended instruction can be used to directly access the data memory instead of using the indirect addressing access. This can not only reduce the use of Flash memory space but also improve the CPU execution efficiency.

Mnemonic	Description	Cycles	Flag Affected
<b>Arithmetic</b>			
LADD A,[m]	Add Data Memory to ACC	2	Z, C, AC, OV, SC
LADDM A,[m]	Add ACC to Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC
LADC A,[m]	Add Data Memory to ACC with Carry	2	Z, C, AC, OV, SC
LADCM A,[m]	Add ACC to Data memory with Carry	2 <sup>Note</sup>	Z, C, AC, OV, SC
LSUB A,[m]	Subtract Data Memory from ACC	2	Z, C, AC, OV, SC, CZ
LSUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
LSBC A,[m]	Subtract Data Memory from ACC with Carry	2	Z, C, AC, OV, SC, CZ
LSBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
LDAA [m]	Decimal adjust ACC for Addition with result in Data Memory	2 <sup>Note</sup>	C
<b>Logic Operation</b>			
LAND A,[m]	Logical AND Data Memory to ACC	2	Z
LOR A,[m]	Logical OR Data Memory to ACC	2	Z
LXOR A,[m]	Logical XOR Data Memory to ACC	2	Z
LANDM A,[m]	Logical AND ACC to Data Memory	2 <sup>Note</sup>	Z
LORM A,[m]	Logical OR ACC to Data Memory	2 <sup>Note</sup>	Z
LXORM A,[m]	Logical XOR ACC to Data Memory	2 <sup>Note</sup>	Z
LCPL [m]	Complement Data Memory	2 <sup>Note</sup>	Z
LCPLA [m]	Complement Data Memory with result in ACC	2	Z
<b>Increment &amp; Decrement</b>			
LINCA [m]	Increment Data Memory with result in ACC	2	Z
LINC [m]	Increment Data Memory	2 <sup>Note</sup>	Z
LDECA [m]	Decrement Data Memory with result in ACC	2	Z
LDEC [m]	Decrement Data Memory	2 <sup>Note</sup>	Z
<b>Rotate</b>			
LRRA [m]	Rotate Data Memory right with result in ACC	2	None
LRR [m]	Rotate Data Memory right	2 <sup>Note</sup>	None
LRRCA [m]	Rotate Data Memory right through Carry with result in ACC	2	C
LRRC [m]	Rotate Data Memory right through Carry	2 <sup>Note</sup>	C
LRLA [m]	Rotate Data Memory left with result in ACC	2	None
LRL [m]	Rotate Data Memory left	2 <sup>Note</sup>	None
LRLCA [m]	Rotate Data Memory left through Carry with result in ACC	2	C
LRLC [m]	Rotate Data Memory left through Carry	2 <sup>Note</sup>	C
<b>Data Move</b>			
LMOV A,[m]	Move Data Memory to ACC	2	None
LMOV [m],A	Move ACC to Data Memory	2 <sup>Note</sup>	None
<b>Bit Operation</b>			
LCLR [m].i	Clear bit of Data Memory	2 <sup>Note</sup>	None
LSET [m].i	Set bit of Data Memory	2 <sup>Note</sup>	None

Mnemonic	Description	Cycles	Flag Affected
<b>Branch</b>			
LSZ [m]	Skip if Data Memory is zero	2 <sup>Note</sup>	None
LSZA [m]	Skip if Data Memory is zero with data movement to ACC	2 <sup>Note</sup>	None
LSNZ [m]	Skip if Data Memory is not zero	2 <sup>Note</sup>	None
LSZ [m].i	Skip if bit i of Data Memory is zero	2 <sup>Note</sup>	None
LSNZ [m].i	Skip if bit i of Data Memory is not zero	2 <sup>Note</sup>	None
LSIZ [m]	Skip if increment Data Memory is zero	2 <sup>Note</sup>	None
LSDZ [m]	Skip if decrement Data Memory is zero	2 <sup>Note</sup>	None
LSIZA [m]	Skip if increment Data Memory is zero with result in ACC	2 <sup>Note</sup>	None
LSDZA [m]	Skip if decrement Data Memory is zero with result in ACC	2 <sup>Note</sup>	None
<b>Table Read</b>			
LTABRD [m]	Read table (specific page) to TBLH and Data Memory	3 <sup>Note</sup>	None
LTABRDL [m]	Read table (last page) to TBLH and Data Memory	3 <sup>Note</sup>	None
LITABRD [m]	Increment table pointer TBLP first and Read table (specific page) to TBLH and Data Memory	3 <sup>Note</sup>	None
LITABRDL [m]	Increment table pointer TBLP first and Read table (last page) to TBLH and Data Memory	3 <sup>Note</sup>	None
<b>Miscellaneous</b>			
LCLR [m]	Clear Data Memory	2 <sup>Note</sup>	None
LSET [m]	Set Data Memory	2 <sup>Note</sup>	None
LSWAP [m]	Swap nibbles of Data Memory	2 <sup>Note</sup>	None
LSWAPA [m]	Swap nibbles of Data Memory with result in ACC	2	None

Note: 1. For these extended skip instructions, if the result of the comparison involves a skip then three cycles are required, if no skip takes place two cycles is required.

2. Any extended instruction which changes the contents of the PCL register will also require three cycles for execution.

## Instruction Definition

<b>ADC A,[m]</b>	Add Data Memory to ACC with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADCM A,[m]</b>	Add ACC to Data Memory with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADD A,[m]</b>	Add Data Memory to ACC
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADD A,x</b>	Add immediate data to ACC
Description	The contents of the Accumulator and the specified immediate data are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + x$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADDM A,[m]</b>	Add ACC to Data Memory
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>AND A,[m]</b>	Logical AND Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
<b>AND A,x</b>	Logical AND immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bit wise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } x$
Affected flag(s)	Z

<b>ANDM A,[m]</b>	Logical AND ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow \text{ACC} \text{ "AND" } [m]$
Affected flag(s)	Z
<b>CALL addr</b>	Subroutine call
Description	Unconditionally calls a subroutine at the specified address. The Program Counter then increments by 1 to obtain the address of the next instruction which is then pushed onto the stack. The specified address is then loaded and the program continues execution from this new address. As this instruction requires an additional operation, it is a two cycle instruction.
Operation	Stack $\leftarrow$ Program Counter + 1 Program Counter $\leftarrow$ addr
Affected flag(s)	None
<b>CLR [m]</b>	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	$[m] \leftarrow 00\text{H}$
Affected flag(s)	None
<b>CLR [m].i</b>	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	$[m].i \leftarrow 0$
Affected flag(s)	None
<b>CLR WDT</b>	Clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared.
Operation	WDT cleared TO $\leftarrow$ 0 PDF $\leftarrow$ 0
Affected flag(s)	TO, PDF
<b>CPL [m]</b>	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	$[m] \leftarrow [m]$
Affected flag(s)	Z
<b>CPLA [m]</b>	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC $\leftarrow [m]$
Affected flag(s)	Z

<b>DAA [m]</b>	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	[m] ← ACC + 00H or [m] ← ACC + 06H or [m] ← ACC + 60H or [m] ← ACC + 66H
Affected flag(s)	C
<b>DEC [m]</b>	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	[m] ← [m] - 1
Affected flag(s)	Z
<b>DECA [m]</b>	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	ACC ← [m] - 1
Affected flag(s)	Z
<b>HALT</b>	Enter power down mode
Description	This instruction stops the program execution and turns off the system clock. The contents of the Data Memory and registers are retained. The WDT and prescaler are cleared. The power down flag PDF is set and the WDT time-out flag TO is cleared.
Operation	TO ← 0 PDF ← 1
Affected flag(s)	TO, PDF
<b>INC [m]</b>	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	[m] ← [m] + 1
Affected flag(s)	Z
<b>INCA [m]</b>	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	ACC ← [m] + 1
Affected flag(s)	Z

<b>JMP addr</b>	Jump unconditionally
Description	The contents of the Program Counter are replaced with the specified address. Program execution then continues from this new address. As this requires the insertion of a dummy instruction while the new address is loaded, it is a two cycle instruction.
Operation	Program Counter $\leftarrow$ addr
Affected flag(s)	None
<b>MOV A,[m]</b>	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	ACC $\leftarrow$ [m]
Affected flag(s)	None
<b>MOV A,x</b>	Move immediate data to ACC
Description	The immediate data specified is loaded into the Accumulator.
Operation	ACC $\leftarrow$ x
Affected flag(s)	None
<b>MOV [m],A</b>	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	[m] $\leftarrow$ ACC
Affected flag(s)	None
<b>NOP</b>	No operation
Description	No operation is performed. Execution continues with the next instruction.
Operation	No operation
Affected flag(s)	None
<b>OR A,[m]</b>	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC $\leftarrow$ ACC “OR” [m]
Affected flag(s)	Z
<b>OR A,x</b>	Logical OR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC $\leftarrow$ ACC “OR” x
Affected flag(s)	Z
<b>ORM A,[m]</b>	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	[m] $\leftarrow$ ACC “OR” [m]
Affected flag(s)	Z

<b>RET</b>	Return from subroutine
Description	The Program Counter is restored from the stack. Program execution continues at the restored address.
Operation	Program Counter $\leftarrow$ Stack
Affected flag(s)	None
<b>RET A,x</b>	Return from subroutine and load immediate data to ACC
Description	The Program Counter is restored from the stack and the Accumulator loaded with the specified immediate data. Program execution continues at the restored address.
Operation	Program Counter $\leftarrow$ Stack ACC $\leftarrow$ x
Affected flag(s)	None
<b>RETI</b>	Return from interrupt
Description	The Program Counter is restored from the stack and the interrupts are re-enabled by setting the EMI bit. EMI is the master interrupt global enable bit. If an interrupt was pending when the RETI instruction is executed, the pending Interrupt routine will be processed before returning to the main program.
Operation	Program Counter $\leftarrow$ Stack EMI $\leftarrow$ 1
Affected flag(s)	None
<b>RL [m]</b>	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	[m].(i+1) $\leftarrow$ [m].i; (i=0~6) [m].0 $\leftarrow$ [m].7
Affected flag(s)	None
<b>RLA [m]</b>	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.(i+1) $\leftarrow$ [m].i; (i=0~6) ACC.0 $\leftarrow$ [m].7
Affected flag(s)	None
<b>RLC [m]</b>	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	[m].(i+1) $\leftarrow$ [m].i; (i=0~6) [m].0 $\leftarrow$ C C $\leftarrow$ [m].7
Affected flag(s)	C

<b>RLCA [m]</b>	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim 6)$ $ACC.0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>RR [m]</b>	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim 6)$ $[m].7 \leftarrow [m].0$
Affected flag(s)	None
<b>RRA [m]</b>	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory is rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim 6)$ $ACC.7 \leftarrow [m].0$
Affected flag(s)	None
<b>RRC [m]</b>	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim 6)$ $[m].7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>RRCA [m]</b>	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim 6)$ $ACC.7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>SBC A,[m]</b>	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>SBC A, x</b>	Subtract immediate data from ACC with Carry
Description	The immediate data and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SBCM A,[m]</b>	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SDZ [m]</b>	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] - 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>SDZA [m]</b>	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] - 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>SET [m]</b>	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
<b>SET [m].i</b>	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None

<b>SIZ [m]</b>	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>SIZA [m]</b>	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>SNZ [m].i</b>	Skip if bit i of Data Memory is not 0
Description	If bit i of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None
<b>SNZ [m]</b>	Skip if Data Memory is not 0
Description	The contents of the specified Data Memory are read out and then written back to the specified Data Memory again. If the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m] \neq 0$
Affected flag(s)	None
<b>SUB A,[m]</b>	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SUBM A,[m]</b>	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>SUB A,x</b>	Subtract immediate data from ACC
Description	The immediate data specified by the code is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - x$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SWAP [m]</b>	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	$[m].3\sim[m].0 \leftrightarrow [m].7\sim[m].4$
Affected flag(s)	None
<b>SWAPA [m]</b>	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC.3\sim ACC.0 \leftarrow [m].7\sim[m].4$ $ACC.7\sim ACC.4 \leftarrow [m].3\sim[m].0$
Affected flag(s)	None
<b>SZ [m]</b>	Skip if Data Memory is 0
Description	The contents of the specified Data Memory are read out and then written back to the specified Data Memory again. If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if $[m]=0$
Affected flag(s)	None
<b>SZA [m]</b>	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m]$ Skip if $[m]=0$
Affected flag(s)	None
<b>SZ [m].i</b>	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if $[m].i=0$
Affected flag(s)	None

<b>TABRD [m]</b>	Read table (specific page) to TBLH and Data Memory
Description	The low byte of the program code (specific page) addressed by the table pointer (TBLP and TBHP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>TABRDL [m]</b>	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>ITABRD [m]</b>	Increment table pointer low byte first and read table (specific page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the program code (specific page) addressed by the table pointer (TBHP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>ITABRDL [m]</b>	Increment table pointer low byte first and read table (last page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>XOR A,[m]</b>	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC “XOR” [m]
Affected flag(s)	Z
<b>XORM A,[m]</b>	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC “XOR” [m]
Affected flag(s)	Z
<b>XOR A,x</b>	Logical XOR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC “XOR” x
Affected flag(s)	Z

### Extended Instruction Definition

The extended instructions are used to directly access the data stored in any data memory sections.

**LADC A,[m]** Add Data Memory to ACC with Carry

Description The contents of the specified Data Memory, Accumulator and the carry flag are added.  
The result is stored in the Accumulator.

Operation  $ACC \leftarrow ACC + [m] + C$

Affected flag(s) OV, Z, AC, C, SC

**LADCM A,[m]** Add ACC to Data Memory with Carry

Description The contents of the specified Data Memory, Accumulator and the carry flag are added.  
The result is stored in the specified Data Memory.

Operation  $[m] \leftarrow ACC + [m] + C$

Affected flag(s) OV, Z, AC, C, SC

**LADD A,[m]** Add Data Memory to ACC

Description The contents of the specified Data Memory and the Accumulator are added.  
The result is stored in the Accumulator.

Operation  $ACC \leftarrow ACC + [m]$

Affected flag(s) OV, Z, AC, C, SC

**LADDM A,[m]** Add ACC to Data Memory

Description The contents of the specified Data Memory and the Accumulator are added.  
The result is stored in the specified Data Memory.

Operation  $[m] \leftarrow ACC + [m]$

Affected flag(s) OV, Z, AC, C, SC

**LAND A,[m]** Logical AND Data Memory to ACC

Description Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.

Operation  $ACC \leftarrow ACC \text{ "AND" } [m]$

Affected flag(s) Z

**LANDM A,[m]** Logical AND ACC to Data Memory

Description Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.

Operation  $[m] \leftarrow ACC \text{ "AND" } [m]$

Affected flag(s) Z

<b>LCLR [m]</b>	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	$[m] \leftarrow 00H$
Affected flag(s)	None
<b>LCLR [m].i</b>	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	$[m].i \leftarrow 0$
Affected flag(s)	None
<b>LCPL [m]</b>	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	$[m] \leftarrow [m]$
Affected flag(s)	Z
<b>LCPLA [m]</b>	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m]$
Affected flag(s)	Z
<b>LDAA [m]</b>	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	$[m] \leftarrow ACC + 00H$ or $[m] \leftarrow ACC + 06H$ or $[m] \leftarrow ACC + 60H$ or $[m] \leftarrow ACC + 66H$
Affected flag(s)	C
<b>LDEC [m]</b>	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	$[m] \leftarrow [m] - 1$
Affected flag(s)	Z

<b>LDECA [m]</b>	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Z
<b>LINC [m]</b>	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	$[m] \leftarrow [m] + 1$
Affected flag(s)	Z
<b>LINCA [m]</b>	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] + 1$
Affected flag(s)	Z
<b>LMOV A,[m]</b>	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	$ACC \leftarrow [m]$
Affected flag(s)	None
<b>LMOV [m],A</b>	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	$[m] \leftarrow ACC$
Affected flag(s)	None
<b>LOR A,[m]</b>	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "OR" } [m]$
Affected flag(s)	Z
<b>LORM A,[m]</b>	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "OR" } [m]$
Affected flag(s)	Z

<b>LRL [m]</b>	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	$[m].(i+1) \leftarrow [m].i; (i=0\sim6)$ $[m].0 \leftarrow [m].7$
Affected flag(s)	None
<b>LRLA [m]</b>	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim6)$ $ACC.0 \leftarrow [m].7$
Affected flag(s)	None
<b>LRLC [m]</b>	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	$[m].(i+1) \leftarrow [m].i; (i=0\sim6)$ $[m].0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>LRLCA [m]</b>	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim6)$ $ACC.0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>LRR [m]</b>	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim6)$ $[m].7 \leftarrow [m].0$
Affected flag(s)	None

<b>LRRRA [m]</b>	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory is rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim6)$ $ACC.7 \leftarrow [m].0$
Affected flag(s)	None
<b>LRRRC [m]</b>	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim6)$ $[m].7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>LRRCA [m]</b>	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim6)$ $ACC.7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>LSBC A,[m]</b>	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>LSBCM A,[m]</b>	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>LSDZ [m]</b>	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] - 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>LSDZA [m]</b>	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] - 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>LSET [m]</b>	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
<b>LSET [m].i</b>	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None
<b>LSIZ [m]</b>	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if $[m]=0$
Affected flag(s)	None

<b>LSIZA [m]</b>	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if ACC=0
Affected flag(s)	None
<b>LSNZ [m].i</b>	Skip if bit i of Data Memory is not 0
Description	If bit i of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None
<b>LSNZ [m]</b>	Skip if Data Memory is not 0
Description	The contents of the specified Data Memory are read out and then written to the specified Data Memory again. If the content of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m] \neq 0$
Affected flag(s)	None
<b>LSUB A,[m]</b>	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>LSUBM A,[m]</b>	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>LSWAP [m]</b>	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	$[m].3\sim[m].0 \leftrightarrow [m].7\sim[m].4$
Affected flag(s)	None
<b>LSWAPA [m]</b>	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC.3\sim ACC.0 \leftarrow [m].7\sim[m].4$ $ACC.7\sim ACC.4 \leftarrow [m].3\sim[m].0$
Affected flag(s)	None
<b>LSZ [m]</b>	Skip if Data Memory is 0
Description	The contents of the specified Data Memory are read out and then written to the specified Data Memory again. If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if $[m]=0$
Affected flag(s)	None
<b>LSZA [m]</b>	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m]$ Skip if $[m]=0$
Affected flag(s)	None
<b>LSZ [m].i</b>	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a three cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if $[m].i=0$
Affected flag(s)	None
<b>LTABRD [m]</b>	Read table (specific page) to TBLH and Data Memory
Description	The low byte of the program code (specific page) addressed by the table pointer (TBHP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	$[m] \leftarrow$ program code (low byte) $TBLH \leftarrow$ program code (high byte)
Affected flag(s)	None

<b>LTABRDL [m]</b>	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LITABRD [m]</b>	Increment table pointer low byte first and read table (specific page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the program code (specific page) addressed by the table pointer (TBHP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LITABRDL [m]</b>	Increment table pointer low byte first and read table (last page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LXOR A,[m]</b>	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC “XOR” [m]
Affected flag(s)	Z
<b>LXORM A,[m]</b>	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC “XOR” [m]
Affected flag(s)	Z

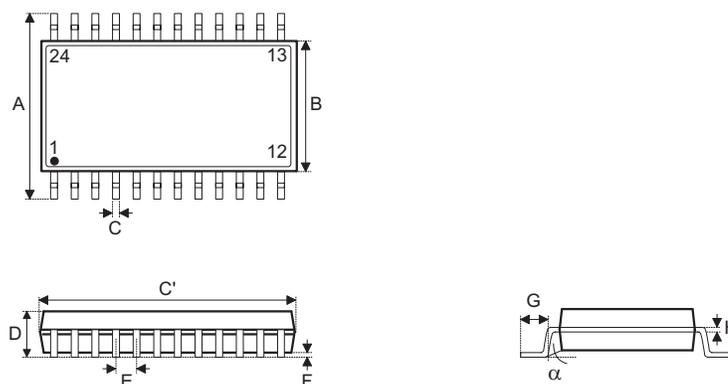
## Package Information

Note that the package information provided here is for consultation purposes only. As this information may be updated at regular intervals users are reminded to consult the [Holtek website](#) for the latest version of the [Package/Carton Information](#).

Additional supplementary information with regard to packaging is listed below. Click on the relevant section to be transferred to the relevant website page.

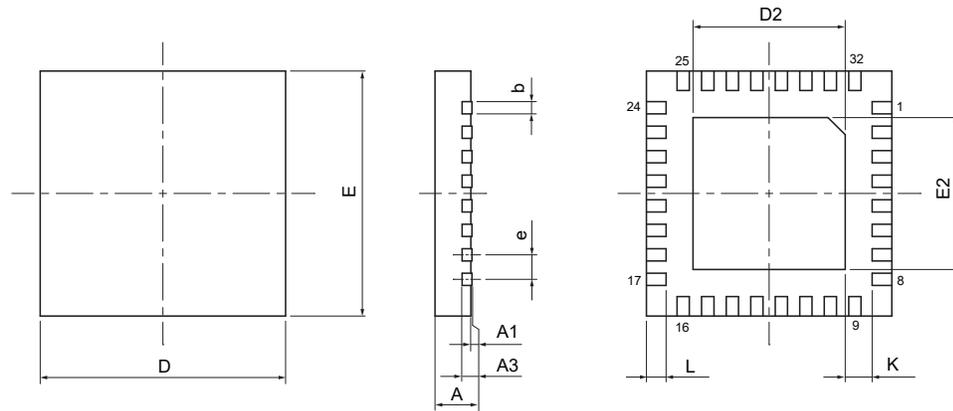
- Package Information (include Outline Dimensions, Product Tape and Reel Specifications)
- The Operation Instruction of Packing Materials
- Carton information

**24-pin SSOP (150mil) Outline Dimensions**



Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.236 BSC		
B	0.154 BSC		
C	0.008	—	0.012
C'	0.341 BSC		
D	—	—	0.069
E	0.025 BSC		
F	0.004	—	0.010
G	0.016	—	0.050
H	0.004	—	0.010
α	0°	—	8°

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	6.00 BSC		
B	3.90 BSC		
C	0.20	—	0.30
C'	8.66 BSC		
D	—	—	1.75
E	0.635 BSC		
F	0.10	—	0.25
G	0.41	—	1.27
H	0.10	—	0.25
α	0°	—	8°

**SAW Type 32-pin QFN (4mm×4mm×0.75mm) Outline Dimensions**


Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.028	0.030	0.031
A1	0.000	0.001	0.002
A3	0.008 REF		
b	0.006	0.008	0.010
D	0.157 BSC		
E	0.157 BSC		
e	0.016 BSC		
D2	0.100	—	0.108
E2	0.100	—	0.108
L	0.010	—	0.018
K	0.008	—	—

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	0.70	0.75	0.80
A1	0.00	0.02	0.05
A3	0.203 REF		
b	0.15	0.20	0.25
D	4.00 BSC		
E	4.00 BSC		
e	0.40 BSC		
D2	2.55	—	2.75
E2	2.55	—	2.75
L	0.25	—	0.45
K	0.20	—	—

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