

HT82J97E USB Joystick Encoder 8-Bit OTP MCU

Features

- Flexible total solution for applications that combine PS/2 and low-speed USB interface, such as mice, joysticks, and many others
- USB Specification Compliance
 - Conforms to USB specification V1.1
 - Conforms to USB HID specification V1.1
- Supports 1 Low-speed USB control endpoint and 1
 interrupt endpoint
- Each endpoint has 8×8 bytes FIFO
- Integrated USB transceiver
- 3.3V regulator output
- External 6MHz or 12MHz ceramic resonator or crystal
- 8-bit RISC microcontroller, with 2K×14 EPROM (000H~7FFH)
- 96 bytes RAM (20H~7FH)
- 6MHz/12MHz internal CPU clock
- 4-level stacks

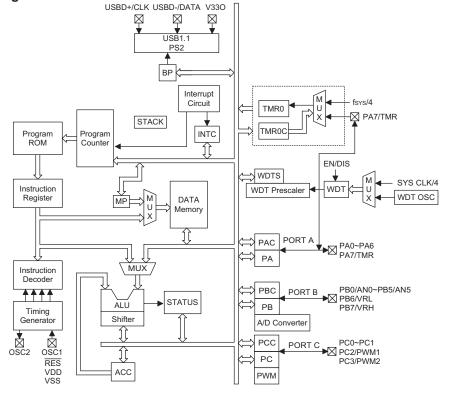
General Description

The USB MCU OTP body is suitable for USB mouse and USB joystick devices. It consists of a Holtek high

- Two 8-bit indirect addressing registers
- One 16-bit programmable timer counter with overflow interrupt (shared with PA7, vector 0CH)
- One USB interrupt input (vector 04H)
- HALT function and wake-up feature reduce power consumption
- PA0~PA7 support wake-up function
- Internal Power-On reset (POR)
- Watchdog Timer (WDT)
- 20 I/O ports (including 2-PWM output, PC2, PC3)
- 2 PWM output (PC2, PC3)
- Can produce PWM frequency range from 23Hz to 23kHz
- Built-in 8-bit Analog-to-Digital Converter, (6-channel for internal mode (PB0~PB5), 6-channel for external mode with VHL (PB7) and VRL (PB6))
- 20/28-pin SOP package

performance 8-bit MCU core for control unit, built-in USB SIE, 2K×14 EPROM and 96 bytes data RAM.

Block Diagram





Pin Assignment

		1		1
		VSS 🗆	1 28	🗆 osci
		V330 🗆	2 27	🗆 osco
		USBD+/CLK	3 26	
		USBD-/DATA 🗆	4 25	D PC3/PWM2
VSS 🗖 1	20 🗅 OSCI	RES 🗆	5 24	D PC2/PWM1
V330 🗖 2	19 🗖 OSCO	PA0 🗆	6 23	🗆 PA7
USBD+/CLK 🗖 3	18 🗖 VDD	PA1 🗆	7 22	🗆 PA6
USBD-/DATA 🗖 4	17 🗖 PA7	PC0 🗆	8 21	🗆 PA5
RES 🗆 5	16 🗖 PA6	PC1 🗆	9 20	🗆 PA4
PA0 🗖 6	15 🗖 PA5	PB0/AN0	10 19	🗆 PA3
PA1 🗖 7	14 🗖 PA4	PB1/AN1 🗆	11 12 18	🗆 PA2
РВ2 🗖 8	13 🗖 PA3	PB2/AN2	12 13 17	□ PB7/VRH
РВЗ 🗖 9	12 🗖 PA2	PB3/AN3 🗆	14 16	D PB6/VRL
РВ4 🗖 10	11 🏳 РВ7	PB4/AN4 🗆	15	D PB5/AN5
	[82J97E 0 SOP-A		HT82J97E - 28 SOP-/	=
-				-

Pin Description

Pin Name	I/O	ROM Code Option	Description
PA0~PA7	I/O	Pull-low Pull-high Wake-up CMOS/NMOS/PMOS	Bidirectional 8-bit input/output port. Each bit can be configured as a wake-up input by ROM code option. The input or output mode is con- trolled by PAC (PA control register). Pull-high resistor options: PA0~PA7 Pull-low resistor options: PA0~PA3 CMOS/NMOS/PMOS options: PA0~PA7 Wake-up options: PA0~PA7
PB0/AN0 PB1/AN1 PB2/AN2 PB3/AN3 PB5/AN5 PB6/VRL	I/O	Pull-high Analog input	Bidirectional 8-bit input/output port. Software instructions determine the CMOS output or Schmitt trigger input with pull-high resistor (determined by pull-high options). The PB can be used as analog input of the analog to digital converter (determined by options). Pull-low resistor for options: PB2, PB3
PB4/AN4 PB7/VRH	I/O	Pull-high Analog input Wake-up	Bidirectional 8-bit input/output port. Software instructions determine the CMOS output or Schmitt trigger input with pull-high resistor (determined by pull-high options). The PB can be used as analog input of the analog to digital converter (determined by options). Wake-up options: PB4, PB7
VSS	_		Negative power supply, ground
PC0~PC3	I/O	Pull-high	Bidirectional I/O lines. Software instructions determine the CMOS out- put or Schmitt trigger input with pull-high resistor (determined by pull-high options). PC2 can be used as PWM1 output PC3 can be used as PWM2 output
RES	Ι		Schmitt trigger reset input. Active low.
VDD			Positive power supply
V33O	0		3.3V regulator output
USBD+/CLK	I/O		USBD+ or PS2 CLK I/O line USB or PS2 function is controlled by software control register
USBD-/DATA	I/O		USBD- or PS2 DATA I/O line USB or PS2 function is controlled by software control register
OSCI OSCO	 0	_	OSCI, OSCO are connected to a 6MHz or 12MHz crystal/resonator (determined by software instructions) for the internal system clock.



Ta=25°C

Absolute Maximum Ratings

Supply VoltageV_SS=0.3V to V_SS+6.0V	Storage Temperature50°C to 125°C
Input VoltageV_SS-0.3V to V_DD+0.3V	Operating Temperature0°C to 70°C

Note: These are stress ratings only. Stresses exceeding the range specified under "Absolute Maximum Ratings" may cause substantial damage to the device. Functional operation of this device at other conditions beyond those listed in the specification is not implied and prolonged exposure to extreme conditions may affect device reliability.

D.C. Characteristics

Complete L	Demonster		Test Conditions		-		Unit
Symbol	Parameter	V_{DD}	Conditions	Min.	Тур.	Max.	
V _{DD}	Operating Voltage	_		4		5.5	V
I _{DD}	Operating Current (6MHz Crystal)	5V	No load, f _{SYS} =6MHz	_	7	9	mA
I _{STB}	Standby Current	5V	No load, system HALT		300	500	μA
V _{IL1}	Input Low Voltage for I/O Ports	5V		0		0.8	V
V _{IH1}	Input High Voltage for I/O Ports	5V		2		5	V
V _{IL2}	Input Low Voltage (RES)	5V		0		$0.4V_{DD}$	V
V _{IH2}	Input High Voltage (RES)	5V		0.9V _{DD}		V _{DD}	V
I _{OL}	Output Sink Current for Other Ports PA0~PA7, PB0~PB7 and PC0~PC3	5V	V _{OL} =0.4V	2	4	_	mA
I _{OH}	Output Port Source Current	5V	V _{OL} =3.4V	-2.5	-4		mA
R _{PD}	Pull-down Resistance for PA0~PA3, PB2 and PB3	5V		10	30	50	kΩ
R _{PH1}	Pull-high Resistance for CLK and DATA	_		2	4.7	6	kΩ
R _{PH2}	Pull-high Resistance for PA0~PA7, PB0~PB7 and PC0~PC3	_	_	30	50	70	kΩ
V _{LVR}	Low Voltage Reset	5V		2.4	2.7	3	V

A.C. Characteristics

Ta=25°C

Symbol	Parameter		Test Conditions	Min.	True	Max.	Unit
Symbol	Parameter		Conditions	IVIIII.	Тур.	Wax.	Unit
f _{SYS}	System Clock (Crystal OSC)	5V	—	6	_	12	MHz
f _{RCSYS}	RC Clock with 8-bit Prescaler Register			0	32	_	kHz
t _{WDT}	Watchdog Time-out Period (System Clock)	_	Without WDT prescaler	1024	_	_	t _{RCSYS}
t _{RF}	USBD+, USBD- Rising & falling Time	_		75		300	ns
t _{SST}	System Start-up Timer Period		Wake-up from HALT		1024		t _{SYS}
tosc	Crystal Setup	_			5	10	ms
f _{PWM}	PWM Cycle Frequency	_	6MHz or 12MHz	23		2300	Hz

Note: Power-on period=t_{WDT}+t_{SST}+t_{OSC}

WDT Time-out in normal mode=1/f_{RCSYS}×256×WDTS+t_{WDT}

WDT Time-out in HALT mode=1/f_{RCSYS}×256×WDTS+t_{SST}+t_{OSC}



Functional Description

Execution Flow

The system clock for the microcontroller is derived from either 6MHz or 12MHz crystal oscillator, which used a frequency that is determined by the SCLKSEL bit of the SCC Register. The default system frequency is 12MHz. The system clock is internally divided into four nonoverlapping clocks. One instruction cycle consists of four system clock cycles.

Instruction fetching and execution are pipelined in such a way that a fetch takes an instruction cycle while decoding and execution takes the next instruction cycle. However, the pipelining scheme causes each instruction to be effectively executed in a cycle. If an instruction changes the program counter, two cycles are required to complete the instruction.

Program Counter – PC

The program counter (PC) controls the sequence in which the instructions stored in the program ROM are executed and its contents specify a full range of program memory.

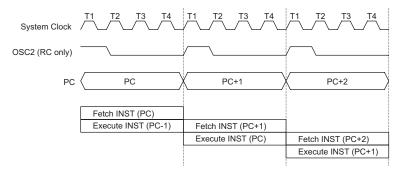
After accessing a program memory word to fetch an instruction code, the contents of the program counter are incremented by one. The program counter then points to the memory word containing the next instruction code.

When executing a jump instruction, conditional skip execution, loading to the PCL register, performing a subroutine call or return from subroutine, initial reset, internal interrupt, external interrupt or return from interrupts, the PC manipulates the program transfer by loading the address corresponding to each instruction.

The conditional skip is activated by instructions. Once the condition is met, the next instruction, fetched during the current instruction execution, is discarded and a dummy cycle replaces it to get the proper instruction. Otherwise proceed with the next instruction.

The lower byte of the program counter (PCL) is a readable and writeable register (06H). Moving data into the PCL performs a short jump. The destination will be within the current program ROM page.

When a control transfer takes place, an additional dummy cycle is required.



Mode	Program Counter										
wode	*10	*9	*8	*7	*6	*5	*4	*3	*2	*1	*0
Initial Reset	0	0	0	0	0	0	0	0	0	0	0
USB Interrupt	0	0	0	0	0	0	0	0	1	0	0
Timer/Event Counter Overflow	0	0	0	0	0	0	0	1	1	0	0
Skip						PC+2					
Loading PCL	*10	*9	*8	@7	@6	@5	@4	@3	@2	@1	@0
Jump, Call Branch	#10	#9	#8	#7	#6	#5	#4	#3	#2	#1	#0
Return from Subroutine	S10	S9	S8	S7	S6	S5	S4	S3	S2	S1	S0

Execution Flow

Program Counter

Note: *10~*0: Program counter bits

S10~S0: Stack register bits

#10~#0: Instruction code bits

@7~@0: PCL bits

Program Memory – ROM

The program memory is used to store the program instructions which are to be executed. It also contains data, table, and interrupt entries, and is organized into 2048×14 bits, addressed by the program counter and table pointer.

Certain locations in the program memory are reserved for special usage:

Location 000H

This area is reserved for program initialization. After a chip reset, the program always begins execution at location 000H.

Location 004H

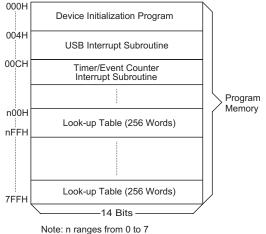
This area is reserved for the USB interrupt service program. If the USB interrupt is activated, the interrupt is enabled and the stack is not full, the program begins execution at location 004H.

Location 00CH

This location is reserved for the Timer/Event Counter interrupt service program. If a timer interrupt results from a Timer/Event Counter overflow, and the interrupt is enabled and the stack is not full, the program begins execution at location 00CH.

Table location

Any location in the program memory can be used as look-up tables. There are three method to read the



Program Memory

ROM data by two table read instructions: "TABRDC" and "TABRDL", transfer the contents of the lower-order byte to the specified data memory, and the higher-order byte to TBLH (08H).

The three methods are shown as follows:

- The instructions "TABRDC [m]" (the current page, one page=256words), where the table locations is defined by TBLP (07H) in the current page. And the ROM code option TBHP is disabled (default).
- The instructions "TABRDC [m]", where the table locations is defined by registers TBLP (07H) and TBHP (01FH). And the ROM code option TBHP is enabled.
- The instructions "TABRDL [m]", where the table locations is defined by Registers TBLP (07H) in the last page (0700H~07FFH).

Only the destination of the lower-order byte in the table is well-defined, the other bits of the table word are transferred to the lower portion of TBLH and the remaining 1-bit words are read as "0". The Table Higher-order byte register (TBLH) is read only. The table pointer (TBLP, TBHP) is a read/write register (07H, 1FH), which indicates the table location. Before accessing the table, the location must be placed in the TBLP and TBHP (If the OTP option TBHP is disabled, the value in TBHP has no effect). The TBLH is read only and cannot be restored. If the main routine and the ISR (Interrupt Service Routine) both employ the table read instruction, the contents of the TBLH in the main routine are likely to be changed by the table read instruction used in the ISR. Errors can occur. In other words, using the table read instruction in the main routine and the ISR simultaneously should be avoided. However, if the table read instruction has to be applied in both the main routine and the ISR, the interrupt should be disabled prior to the table read instruction. It will not be enabled until the TBLH has been backed up. All table related instructions require two cycles to complete the operation. These areas may function as normal program memory depending on the requirements

Once TBHP is enabled, the instruction "TABRDC [m]" reads the ROM data as defined by TBLP and TBHP value. Otherwise, the ROM code option TBHP is disabled, the instruction "TABRDC [m]" reads the ROM data as defined by TBLP and the current program counter bits.

Instruction	Table Location										
Instruction	*10	*9	*8	*7	*6	*5	*4	*3	*2	*1	*0
TABRDC [m]	P10	P9	P8	@7	@6	@5	@4	@3	@2	@1	@0
TABRDL [m]	1	1	1	@7	@6	@5	@4	@3	@2	@1	@0

Table Location

Note: *10~*0: Table location bits

@7~@0: TBLP bits

P10~P8: Current program counter bits when TBHP is disabled TBHP register bit2~bit0 when TBHP is enabled



Stack Register – STACK

This is a special part of the memory which is used to save the contents of the program counter (PC) only. The stack is organized into 8 levels and is neither part of the data nor part of the program space, and is neither readable nor writeable. The activated level is indexed by the stack pointer (SP) and is neither readable nor writeable. At a subroutine call or interrupt acknowledge signal, the contents of the program counter are pushed onto the stack. At the end of a subroutine or an interrupt routine, signaled by a return instruction (RET or RETI), the program counter is restored to its previous value from the stack. After a chip reset, the SP will point to the top of the stack.

If the stack is full and a non-masked interrupt takes place, the interrupt request flag will be recorded but the acknowledge signal will be inhibited. When the stack pointer is decremented (by RET or RETI), the interrupt will be serviced. This feature prevents stack overflow allowing the programmer to use the structure more easily. In a similar case, if the stack is full and a "CALL" is subsequently executed, stack overflow occurs and the first entry will be lost (only the most recent 8 return addresses are stored).

Data Memory - RAM for Bank 0

The data memory is designed with 96×8 bits. The data memory is divided into two functional groups: special function registers and general purpose data memory (96×8). Most are read/write, but some are read only.

The special function registers include the indirect addressing registers (R0:00H, R1:02H). Bank register (BP, 04H), PWM1 duty register (0DH), PWM2 duty register(0EH), Timer/Event Counter higher order byte register (TMRH;0FH), Timer/Event Counter lower order byte register (TMRL;10H), Timer/Event Counter control register (TMRC;11H), program counter lower-order byte register (PCL:06H), memory pointer registers (MP0;01H, MP1;03H), accumulator (ACC;05H), table pointers (TBLP;07H, TBHP;1FH), table higher-order byte register (TBLH;08H), status register (STATUS;0AH), interrupt control register (INTC;0BH), Watchdog Timer option setting register (WDTS;09H), I/O registers (PA:12H, PB:14H, PC:16H), PWM Base Period Register (18H), I/O control registers (PAC;13H, PBC;15H, PCC;17H). USB/PS2 status and control register (USC;1AH), USB endpoint interrupt status register (USR;1BH), system clock control register (SCC;1CH). A/D converter status and control register (ADSC;1DH) and A/D converter result register (ADR:1EH). The remaining space before the 20H is reserved for future ex-

Bank 0 00H Indirect Addressing Register 0 MP0 01H Indirect Addressing Register 1 02H 03H MP1 BP 04H ACC 05H PCL 06H TBLP 07H TBI H 08H WDTS 09H 0AH STATUS INTC 0BH PWM1 Duty Register 0CH PWM2 Duty Register 0DH 0EH 0FH TMRH TMRL 10H TMRC 11H 12H PA PAC 13H 14H PB 15H PBC 16H PC 17H PCC 18H PWM Base Period Register (PD) 19H USC 1AH USR 1BH SCC 1CH ADSC 1DH 1EH ADR TBHP 1FH 20H General Purpose DATA MEMORY (96 Bytes) 7FH

Bank 0 RAM Mapping

panded usage and reading these locations will get "00H". The general purpose data memory, addressed from 20H to 7FH, is used for data and control information under instruction commands.

All of the data memory areas can handle arithmetic, logic, increment, decrement and rotate operations directly. Except for some dedicated bits, each bit in the data memory can be set and reset by "SET [m].i" and "CLR [m].i". They are also indirectly accessible through memory pointer registers (MP0 or MP1).



Data Memory - RAM for Bank 1

The special function registers used in the USB interface are located in RAM Bank1. In order to access Bank1 register, only the Indirect addressing pointer MP1 can be used and the Bank register BP should be set to 1. The RAM bank 1 mapping is as shown.

	RAM Bank 1						
49H	FIFO 1						
48H	FIFO 0						
47H	FIED 0						
46H	MISC						
45H	SIES						
44H 45⊔	PIPE						
43H	STALL						
42H	AWR						
41H	Pipe_ctrl						
20H	Dia a 11						
1FH	TBHP						
1EH	ADR						
1DH	ADSC						
1CH	SCC						
1BH	USR						
1AH	USC						
19H							
18H	PWM Base Period Register (PD)						
17H	PCC						
16H	PC						
15H	PBC						
14H	PB						
13H	PAC						
12H	PA						
11H	TMRC						
10H	TMRL						
0FH	TMRH						
0EH							
0DH	PWM2 Duty Register						
0CH	PWM1 Duty Register						
0BH	INTC						
0AH	STATUS						
09H	WDTS						
08H	TBLH						
07H	TBLP						
06H	PCL						
05H	ACC						
04H	BP						
03H	MP1						
02H	Indirect Addressing Register 1						
01H	MP0						
00H	Indirect Addressing Register 0						

Address 00~1FH in RAM Bank0 and Bank1 are located in the same Registers

Indirect Addressing Register

Locations 00H and 02H are indirect addressing registers that are not physically implemented. Any read/write operation on [00H] ([02H]) will access the data memory pointed to by MP0 (MP1). Reading location 00H (02H) indirectly will return the result 00H. Writing indirectly results in no operation.

The indirect addressing pointer (MP0) always points to Bank0 RAM addresses no matter the value of Bank Register (BP).

The indirect addressing pointer (MP1) can access Bank0 or Bank1 RAM data according to the value of BP which is set to 0 or 1 respectively.

The memory pointer registers (MP0 and MP1) are 8-bit registers.

Accumulator

The accumulator is closely related to ALU operations. It is also mapped to location 05H of the data memory and can carry out immediate data operations. The data movement between two data memory locations must pass through the accumulator.

Arithmetic and Logic Unit – ALU

This circuit performs 8-bit arithmetic and logic operations. The ALU provides the following functions:

- Arithmetic operations (ADD, ADC, SUB, SBC, DAA)
- Logic operations (AND, OR, XOR, CPL)
- Rotation (RL, RR, RLC, RRC)
- Increment and Decrement (INC, DEC)
- Branch decision (SZ, SNZ, SIZ, SDZ)

The ALU not only saves the results of a data operation but also changes the status register.

Status Register – STATUS

This 8-bit register (0AH) contains the zero flag (Z), carry flag (C), auxiliary carry flag (AC), overflow flag (OV), power down flag (PDF), and watchdog time-out flag (TO). It also records the status information and controls the operation sequence.

With the exception of the TO and PDF flags, bits in the status register can be altered by instructions like most other registers. Any data written into the status register will not change the TO or PDF flag. In addition, operations related to the status register may give different results from those intended.



Labels	Bits	Function
С	0	C is set if an operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation; otherwise C is cleared. C is also affected by a rotate through carry instruction.
AC	1	AC is set if an operation results in a carry out of the low nibbles in addition or no borrow from the high nibble into the low nibble in subtraction; otherwise AC is cleared.
Z	2	Z is set if the result of an arithmetic or logic operation is zero; otherwise Z is cleared.
ov	3	OV is set if an operation results in a carry into the highest-order bit but not a carry out of the highest-order bit, or vice versa; otherwise OV is cleared.
PDF	4	PDF is cleared by a system power-up or executing the "CLR WDT" instruction. PDF is set by executing the "HALT" instruction.
то	5	TO is cleared by a system power-up or executing the "CLR WDT" or "HALT" instruction. TO is set by a WDT time-out.
	6	Unused bit, read as "0"
	7	Unused bit, read as "0"

Status Register

The TO flag can be affected only by a system power-up, a WDT time-out or executing the "CLR WDT" or "HALT" instruction. The PDF flag can be affected only by executing the "HALT" or "CLR WDT" instruction or during a system power-up.

The Z, OV, AC and C flags generally reflect the status of the latest operations.

In addition, upon entering the interrupt sequence or executing a subroutine call, the status register will not be automatically pushed onto the stack. If the contents of the status are important and if the subroutine can corrupt the status register, precautions must be taken to save it properly.

Interrupt

The device provides an external interrupt and internal timer/event counter interrupts. The Interrupt Control Register (INTC;0BH) contains the interrupt control bits to set the enable/disable and the interrupt request flags.

Once an interrupt subroutine is serviced, all the other interrupts will be blocked (by clearing the EMI bit). This scheme may prevent any further interrupt nesting. Other interrupt requests may occur during this interval but only the interrupt request flag is recorded. If a certain interrupt requires servicing within the service routine, the EMI bit and the corresponding bit of the INTC may be set to allow interrupt nesting. If the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the SP is decremented. If immediate service is desired, the stack must be prevented from becoming full.

All these kinds of interrupts have a wake-up capability. As an interrupt is serviced, a control transfer occurs by pushing the program counter onto the stack, followed by a branch to a subroutine at a specified location in the program memory. Only the program counter is pushed onto the stack. If the contents of the register or status register (STATUS) are altered by the interrupt service program which corrupts the desired control sequence, the contents should be saved in advance.

Register	Bit No.	Label	Function
	0	EMI	Controls the master (global) interrupt (1=enable; 0=disable)
	1	EUI	Controls the USB interrupt (1=enable; 0= disable)
	2	_	Unused bit, read as "0"
INTC	3	ETI	Controls the Timer/Event Counter interrupt (1=enable; 0=disable)
(0BH)	4	USBF	USB interrupt request flag (1=active; 0=inactive)
	5	_	Unused bit, read as "0"
	6 TF		Internal timer/event counter request flag (1:active; 0:inactive)
	7	_	Unused bit, read as "0"

INTC Register



The USB interrupts are triggered by the following USB events and the related interrupt request flag (USBF; bit 4 of the INTC) will be set.

- Access of the corresponding USB FIFO from PC
- suspend signal from PC
- resume signal from PC
- USB Reset signal

When the interrupt is enabled, the stack is not full and the external interrupt is active, a subroutine call to location 04H will occur. The interrupt request flag (USBF) and EMI bits will be cleared to disable other interrupts.

When the PC Host access the FIFO of the HT82J97E, the corresponding request bit of the USR is set, and a USB interrupt is triggered. So user can easily decide which FIFO is accessed. When the interrupt has been served, the corresponding bit should be cleared by firmware. When the HT82J97E receives a USB Suspend signal from the Host PC, the suspend line (bit0 of the USC) of the HT82J97E is set and a USB interrupt is also triggered.

When the HT82J97E receives a Resume signal from the Host PC, the resume line (bit3 of the USC) of the HT82J97E is set and a USB interrupt is triggered.

Whenever a USB reset signal is detected, the USB interrupt is triggered and URST_Flag bit of the USC register is set. When the interrupt has been served, the bit should be cleared by firmware.

The internal timer/even counter interrupt is initialized by setting the timer/event counter interrupt request flag (;bit 6 of the INTC), caused by a timer overflow. When the interrupt is enabled, the stack is not full and the TF is set, a subroutine call to location 0CH will occur. The related interrupt request flag (TF) will be reset and the EMI bit cleared to disable further interrupts.

During the execution of an interrupt subroutine, other interrupt acknowledge signals are held until the "RETI" instruction is executed or the EMI bit and the related interrupt control bit are set to 1 (if the stack is not full). To return from the interrupt subroutine, "RET" or "RETI" may be invoked. RETI will set the EMI bit to enable an interrupt service, but RET will not.

Interrupts, occurring in the interval between the rising edges of two consecutive T2 pulses, will be serviced on the latter of the two T2 pulses, if the corresponding interrupts are enabled. In the case of simultaneous requests the following table shows the priority that is applied. These can be masked by resetting the EMI bit.

No.	Interrupt Source	Priority	Vector
а	USB interrupt	1	04H
b	Timer/Event Counter overflow	2	0CH

The timer/event counter interrupt request flag (TF), USB interrupt request flag (USBF), enable timer/event counter interrupt bit (ETI), enable USB interrupt bit (EUI) and

enable master interrupt bit (EMI) constitute an interrupt control register (INTC) which is located at 0BH in the data memory. EMI, EUI and ETI are used to control the enabling/disabling of interrupts. These bits prevent the requested interrupt from being serviced. Once the interrupt request flags (TF, USBF) are set, they will remain in the INTC register until the interrupts are serviced or cleared by a software instruction.

It is recommended that a program does not use the "CALL subroutine" within the interrupt subroutine. Interrupts often occur in an unpredictable manner or need to be serviced immediately in some applications. If only one stack is left and enabling the interrupt is not well controlled, the original control sequence will be damaged once the "CALL" operates in the interrupt subroutine.

Oscillator Configuration

There is an oscillator circuit in the microcontroller.



System Oscillator

This oscillator is designed for system clocks. The HALT mode stops the system oscillator and ignores an external signal to conserve power.

A crystal across OSC1 and OSC2 is needed to provide the feedback and phase shift required for the oscillator. No other external components are required. In stead of a crystal, a resonator can also be connected between OSC1 and OSC2 to get a frequency reference, but two external capacitors in OSC1 and OSC2 are required.

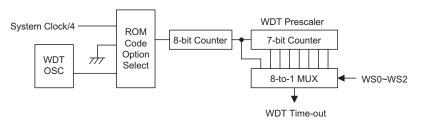
The HT82J97E can operate in 6MHz or 12MHz system clocks. In order to make sure that the USB SIE functions properly, user should correctly configure the SCLKSEL bit of the SCC Register. The default system clock is 12MHz.

The WDT oscillator is a free running on-chip RC oscillator, and no external components are required. Even if the system enters the power down mode, the system clock is stopped, but the WDT oscillator still works within a period of approximately 31µs. The WDT oscillator can be disabled by ROM code option to conserve power.

Watchdog Timer – WDT

The WDT clock source is implemented by a dedicated RC oscillator (WDT oscillator), or instruction clock (system clock divided by 4), determine by ROM code option. This timer is designed to prevent a software malfunction or sequence from jumping to an unknown location with unpredictable results. The Watchdog Timer can be disabled by ROM code option. If the Watchdog Timer is disabled by ROM code option.





Watchdog Timer

abled, all the executions related to the WDT result in no operation.

Once the internal WDT oscillator (RC oscillator with a period of 31µs/5V normally) is selected, it is first divided by 256 (8-stage) to get the nominal time-out period of 8ms/5V. This time-out period may vary with temperatures, VDD and process variations. By invoking the WDT prescaler, longer time-out periods can be realized. Writing data to WS2, WS1, WS0 (bits 2, 1, 0 of the WDTS) can give different time-out periods. If WS2, WS1, and WS0 are all equal to 1, the division ratio is up to 1:128, and the maximum time-out period is 1s/5V. If the WDT oscillator is disabled, the WDT clock may still come from the instruction clock and operates in the same manner except that in the HALT state the WDT may stop counting and lose its protecting purpose. In this situation the logic can only be restarted by external logic. The high nibble and bit 3 of the WDTS are reserved for user defined flags, which can only be set to "10000" (WDTS.7~WDTS.3).

If the device operates in a noisy environment, using the on-chip 32kHz RC oscillator (WDT OSC) is strongly recommended, since the HALT will stop the system clock.

WS2	WS1	WS0	Division Ratio
0	0	0	1:1
0	0	1	1:2
0	1	0	1:4
0	1	1	1:8
1	0	0	1:16
1	0	1	1:32
1	1	0	1:64
1	1	1	1:128

WDTS Register

The WDT overflow under normal operation will initialize a "chip reset" and set the status bit "TO". But in the HALT mode, the overflow will initialize a "warm reset" and only the PC and SP are reset to zero. To clear the contents of the WDT (including the WDT prescaler), three methods are adopted; external reset (a low level to RES), software instruction and a "HALT" instruction. The software instruction include "CLR WDT" and the other set – "CLR WDT1" and "CLR WDT2". Of these two types of instruction, only one can be active depending on the ROM code option – "CLR WDT times selec-

tion option". If the "CLR WDT" is selected (i.e. CLRWDT times is equal to one), any execution of the "CLR WDT" instruction will clear the WDT. In the case that "CLR WDT" and "CLR WDT" are chosen (i.e. CLRWDT times is equal to two), these two instructions must be executed to clear the WDT; otherwise, the WDT may reset the chip as a result of time-out.

Power Down Operation – HALT

The HALT mode is initialized by the "HALT" instruction and results in the following:

- The system oscillator will be turned off but the WDT oscillator remains running (if the WDT oscillator is selected).
- The contents of the on-chip RAM and registers remain unchanged.
- The WDT and WDT prescaler will be cleared and recounted again (if the WDT clock is from the WDT oscillator).
- All of the I/O ports remain in their original status.
- The PDF flag is set and the TO flag is cleared.

The system can leave the HALT mode by means of an external reset, an interrupt, an external falling edge signal on port A or a WDT overflow. An external reset causes a device initialization and the WDT overflow performs a "warm reset". After the TO and PDF flags are examined, the cause for chip reset can be determined. The PDF flag is cleared by a system power-up or executing the "CLR WDT" instruction and is set when executing the "HALT" instruction. The TO flag is set if the WDT time-out occurs, and causes a wake-up that only resets the PC and SP; the others remain in their original status.

The port A wake-up and interrupt methods can be considered as a continuation of normal execution. Each bit in port A can be independently selected to wake-up the device by mask option. Awakening from an I/O port stimulus, the program will resume execution of the next instruction. If it awakens from an interrupt, two sequence may occur. If the related interrupt is disabled or the interrupt is enabled but the stack is full, the program will resume execution at the next instruction. If the interrupt is enabled and the stack is not full, the regular interrupt response takes place. If an interrupt request flag is set to "1" before entering the HALT mode, the wake-up function of the related interrupt will be disabled. Once a



wake-up event occurs, it takes 1024 t_{SYS} (system clock period) to resume normal operation. In other words, a dummy period will be inserted after a wake-up. If the wake-up results from an interrupt acknowledge signal, the actual interrupt subroutine execution will be delayed by one or more cycles. If the wake-up results in the next instruction execution, this will be executed immediately after the dummy period is finished.

To minimize power consumption, all the I/O pins should be carefully managed before entering the HALT status.

Reset

There are four ways in which a reset can occur:

- RES reset during normal operation
- RES reset during HALT
- WDT time-out reset during normal operation
- USB reset

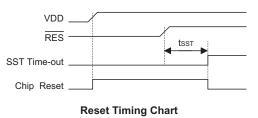
The WDT time-out during HALT is different from other chip reset conditions, since it can perform a "warm reset" that resets only the PC and SP, leaving the other circuits in their original state. Some registers remain unchanged during other reset conditions. Most registers are reset to the "initial condition" when the reset conditions are met. By examining the PDF and TO flags, the program can distinguish between different "chip resets".

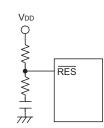
то	PDF	RESET Conditions
0	0	RES reset during power-up
u	u	RES reset during normal operation
0	1	RES wake-up HALT
1	u	WDT time-out during normal operation
1	1	WDT wake-up HALT

Note: "u" stands for "unchanged"

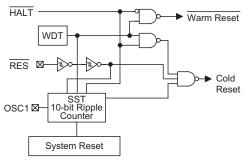
To guarantee that the system oscillator is started and stabilized, the SST (System Start-up Timer) provides an extra delay of 1024 system clock pulses when the system resets (power-up, WDT time-out or $\overline{\text{RES}}$ reset) or the system awakes from the HALT state.

When a system reset occurs, the SST delay is added during the reset period. Any wake-up from HALT will enable the SST delay.





Reset Circuit



Reset Configuration

The functional unit chip reset status are shown below.

000H
Disable
Clear
Clear. After master reset, WDT begins counting
Off
Input mode
Points to the top of the stack



Register	Reset (Power On)	WDT Time-out (Normal Operation)	RES Reset (Normal Operation)	RES Reset (HALT)	WDT Time-Out (HALT)*	USB-Reset (Normal)	USB-Reset (HALT)
TMRH	XXXX XXXX	սսսս սսսս	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน
TMRL	XXXX XXXX	սսսս սսսս	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน
TMRC	00-0 1	00-0 1	00-0 1	00-0 1	uu-u u	00-0 1	00-0 1
Program Counter	000H	000H	000H	000H	000H	000H	000H
MP0	1xxx xxxx	1սսս սսսս	1uuu uuuu	1uuu uuuu	1uuu uuuu	1uuu uuuu	1uuu uuuu
MP1	1xxx xxxx	1นนน นนนน	1นนน นนนน	1นนน นนนน	1นนน นนนน	1นนน นนนน	1uuu uuuu
ACC	XXXX XXXX	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	uuuu uuuu
TBLP	XXXX XXXX	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	uuuu uuuu
TBLH	-xxx xxxx	-uuu uuuu	-uuu uuuu	-uuu uuuu	-uuu uuuu	-uuu uuuu	-uuu uuuu
STATUS	00 xxxx	1u uuuu	uu uuuu	01 uuuu	11 uuuu	uu uuuu	01 uuuu
INTC	-000 0000	-000 0000	-000 0000	-000 0000	-uuu uuuu	-000 0000	-000 0000
WDTS	1000 0111	1000 0111	1000 0111	1000 0111	นนนน นนนน	1000 0111	1000 0111
PA	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PAC	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PB	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PBC	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PC	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PCC	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
PD	1111 1111	1111 1111	1111 1111	1111 1111	นนนน นนนน	1111 1111	1111 1111
AWR	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	0000 0000	0000 0000
PIPE	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	0000 0000	0000 0000
STALL	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	0000 0000	0000 0000
SIES	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	0000 0000	0000 0000
MISC	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	0000 0000	0000 0000
FIFO0	XXXX XXXX	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	0000 0000	0000 0000
FIFO1	XXXX XXXX	นนนน นนนน	นนนน นนนน	นนนน นนนน	นนนน นนนน	0000 0000	0000 0000
USC	11xx 0000	uuxx uuuu	11xx 0000	11xx 0000	uuxx uuuu	1100 0u00	1100 0u00
USR	0100 0000	นนนน นนนน	0100 0000	0100 0000	นนนน นนนน	u1uu 0000	u1uu 0000
SCC	0000 0000	นนนน นนนน	0000 0000	0000 0000	นนนน นนนน	uu00 u000	uu00 u000
ADSC	1000 0000	นนนน นนนน	1000 0000	1000 0000	นนนน นนนน	1000 0000	1000 0000
ADR	XXXX XXXX	นนนน นนนน	XXXX XXXX	XXXX XXXX	นนนน นนนน	XXXX XXXX	XXXX XXXX

Note: "*" stands for "warm reset"

"u" stands for "unchanged"

"x" stands for "unknown"



Timer/Event Counter

A timer/event counter (TMR) is implemented in the microcontroller.

The timer/event counter contains a 16-bit programmable count-up counter and the clock may come from an external source or from the system clock divided by 4.

Using the internal clock source, there is only 1 reference time-base for the timer/event counter. The internal clock source is coming from $f_{SYS}/4$. The external clock input allows the user to count external events, measure time intervals or pulse widths.

There are 3 registers related to the timer/event counter; TMRH (0FH), TMRL (10H), TMRC (11H). Writing TMRL will only put the written data to an internal lower-order byte buffer (8 bits) and writing TMRH will transfer the specified data and the contents of the lower-order byte buffer to TMRH and TMRL preload registers, respectively. The timer/event counter preload register is changed by each writing TMRH operations. Reading TMRH will latch the contents of TMRH and TMRL counters to the destination and the lower-order byte buffer, respectively. Reading the TMRL will read the contents of the lower-order byte buffer. The TMRC is the timer/event counter control register, which defines the operating mode, counting enable or disable and active edge.

The TM0, TM1 bits define the operating mode. The event count mode is used to count external events,

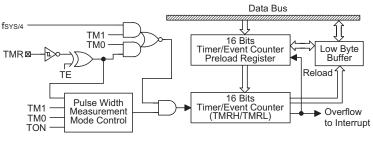
which means that the clock source comes from an external (TMR) pin. The timer mode functions as a normal timer with the clock source coming from the $f_{SYS}/4$ (Timer). The pulse width measurement mode can be used to count the high or low level duration of the external signal (TMR). The counting is based on the $f_{SYS}/4$.

In the event count or timer mode, once the timer/event counter starts counting, it will count from the current contents in the timer/event counter to FFFFH. Once overflow occurs, the counter is reloaded from the timer/event counter preload register and generates the interrupt request flag (TF; bit 6 of the INTC) at the same time.

In the pulse width measurement mode with the TON and TE bits equal to one, once the TMR has received a transient from low to high (or high to low if the TE bit is "0") it will start counting until the TMR returns to the original level and resets the TON. The measured result will remain in the timer/event counter even if the activated transient occurs again. In other words, only one cycle measurement can be done. Until setting the TON, the cycle measurement will function again as long as it receives further transient pulse. Note that, in this operating mode, the timer/event counter starts counting not according to the logic level but according to the transient edges. In the case of counter overflows, the counter is reloaded from the timer/event counter preload register and issues the interrupt request just like the other two modes. To enable the counting operation, the timer ON

Label (TMRC)	RC) Bits Function	
	0~2	Unused bit, read as "0"
TE	3 Defines the TMR active edge of the timer/event counter (0=active on low to high; 1=active on high to low)	
TON 4		Enable/disable the timer counting (0=disable; 1=enable)
5		Unused bit, read as "0"
TM0 6 TM1 7		Defines the operating mode 01=Event count mode (external clock) 10=Timer mode (internal clock) 11=Pulse width measurement mode 00=Unused

TMRC Register



Timer/Event Counter

bit (TON; bit 4 of TMRC) should be set to 1. In the pulse width measurement mode, the TON will be cleared automatically after the measurement cycle is completed. But in the other two modes the TON can only be reset by instructions. The overflow of the timer/event counter is one of the wake-up sources. No matter what the operation mode is, writing a 0 to ET can disable the corresponding interrupt services.

In the case of timer/event counter OFF condition, writing data to the timer/event counter preload register will also reload that data to the timer/event counter. But if the timer/event counter is turned on, data written to it will only be kept in the timer/event counter preload register. The timer/event counter will still operate until overflow occurs (a timer/event counter reloading will occur at the same time). When the timer/event counter (reading TMR) is read, the clock will be blocked to avoid errors. As clock blocking may result in a counting error, this must be taken into consideration by the programmer.

Input/Output Ports

There are 20 bidirectional input/output lines in the microcontroller, labeled from PA to PC, which are mapped to the data memory of [12H], [14H] and [16H] respectively. All of these I/O ports can be used for input and output operations. For input operation, these ports are non-latching, that is, the inputs must be ready at the T2 rising edge of instruction "MOV A,[m]" (m=12H, 14H or 16H). For output operation, all the data is latched and remains unchanged until the output latch is rewritten.

Each I/O line has its own control register (PAC, PBC, PCC) to control the input/output configuration. With this control register, CMOS/NMOS/PMOS output or Schmitt trigger input with or without pull-high/low resistor structures can be reconfigured dynamically under software

control. To function as an input, the corresponding latch of the control register must write a "1". The input source also depends on the control register. If the control register bit is "1", the input will read the pad state. If the control register bit is "0", the contents of the latches will move to the internal bus. The latter is possible in the "read-modify-write" instruction. For output function, CMOS/NMOS/PMOS configurations can be selected (NMOS and PMOS are available for PA only). These control registers are mapped to locations 13H, 15H and 17H.

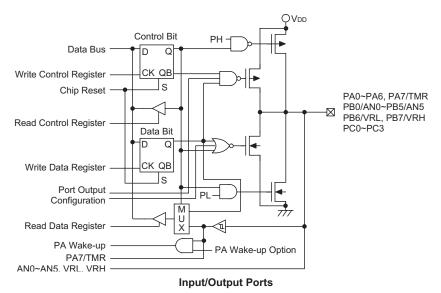
After a chip reset, these input/output lines remain at high levels or in a floating state (depending on the pull-high/low options). Each bit of these input/output latches can be set or cleared by "SET [m].i" and "CLR [m].i" (m=12H, 14H or 16H) instructions.

Some instructions first input data and then follow the output operations. For example, "SET [m].i", "CLR [m].i", "CPL [m]", "CPLA [m]" read the entire port states into the CPU, execute the defined operations (bit-operation), and then write the results back to the latches or the accumulator.

Each line of port A has the capability of waking-up the device.

There are pull-high/low (PA only) options available for I/O lines. Once the pull-high/low option of an I/O line is selected, the I/O line have pull-high/low resistor. Otherwise, the pull-high/low resistor is absent. It should be noted that a non-pull-high/low I/O line operating in input mode will cause a floating state.

It is recommended that unused or not bonded out I/O lines should be set as output pins by software instruction to avoid consuming power under input floating state.



Note: The outputs of PC2 and PC3 will be PWM outputs when PWM outputs are enabled.



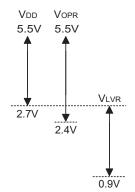
Low Voltage Reset – LVR

The microcontroller contains a low voltage reset circuit in order to monitor the supply voltage of the device. If the supply voltage of the device drops to within the range of $0.9V - V_{LVR}$ such as might occur when changing the battery, the LVR will automatically reset the device internally.

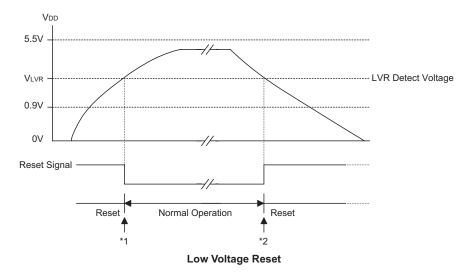
The LVR includes the following specifications:

- For a valid LVR signal, a low voltage $(0.9V \sim V_{LVR})$ must exist for more than 1ms. If the low voltage state does not exceed 1ms, the LVR will ignore it and will not perform a reset function.
- The LVR uses the "OR" function with the external RES signal to perform a chip reset.

The relationship between V_{DD} and V_{LVR} is shown below.



Note: V_{OPR} is the voltage range for proper chip operation at 6MHz or 12MHz system clock.



Note: *1: To make sure that the system oscillator has stabilized, the SST provides an extra delay of 1024 system clock pulses before entering the normal operation.

*2: A low voltage has to exist for more than 1ms, after that 1ms delay, the device enters a reset mode.



USB with MCU Interface

There are eight registers, including Pipe_ctrl, Address+Remote_WakeUp, Stall, Pipe, SIES, Misc, FIFO 0 and FIFO 1 in this buffer function.

Register Name	Pipe_ctrl	Addr.+Remote	Stall	Pipe	SIES	Misc	FIFO 0	FIFO 1
Mem. Addr.	41H	42H	43H	44H	45H	46H	48H	49H
Reserved Addr.	Bank 1, Address 40H, 4AH, 4FH							

Register Memory Mapping

Address+Remote_WakeUp register represents current address and remote wake-up function. The initial value is "00000000" from MSB to LSB.

Register Address	R/W	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
01000010B	R/W		Address value				Remote Wake-up Function 0: Not this function 1: The function exists		

Address+Remote_WakeUp Register

The PIPE_ctrl, STALL and PIPE are bitmap ones. The Pipe_ctrl Register is used for configuring IN (Bit=1) or OUT (Bit=0) Pipe. The default is defined as IN Pipe. The Pipe register represents whether the corresponding endpoint is accessed by host or not. After a USB interrupt signal is being sent out, the MCU can check which endpoint had been accessed. This register is set only after the host accessed the corresponding endpoint. The Stall register shows whether the corresponding endpoint works or not. As soon as the endpoint works improperly, the corresponding bit must be set. The bitmaps are listed as follows:

Register Name	R/W	Register Address	Bit7~Bit2 Reserved	Bit 1	Bit 0	Default Value
Pipe_ctrl	R/W	01000001B	_	Pipe 1	Pipe 0	00000011
Stall	R/W	01000011B	_	Pipe 1	Pipe 0	00000000
Pipe	R	01000100B		Pipe 1	Pipe 0	00000000

Stall and Pipe Registers

The SIES Register is used to indicate the present signal state which the USB SIE received and also determines whether the USB SIE has to change the device address automatically.

Bit No.	Function	Read/Write	Register Address
7	MNI	R/W	
6	EOT	R	
5	CRC_ERR	R/W	
4	NAK	R	040004045
3	IN	R	01000101B
2	OUT	R/W	
1	F0_ERR	R/W	
0	Adr_set	R/W	

SIES Registers Table



Function Name	Read/Write	Description			
Adr_set	R/W	This bit is used to configure the USB SIE to automatically change the device address with the value of the Address+Remote_WakeUp Register (42H). When this bit is set to 1 by F/W, the USB SIE will update the device address with the value of the Address+Remote_WakeUp Register (42H) after the PC Host has successfully read the data from the device by the IN operation. The USB SIE will clear the bit after updating the device address. Otherwise, when this bit is cleared to 0, the USB SIE will update the device address immediately after an address is written to the Address+Remote_WakeUp Register (42H).			
F0_Err	0_Err R/W This bit is used to indicate when there are some errors that occurred when the FIF accessed. This bit is set by the USB SIE and cleared by F/W.				
Out	R/W	This bit is used to indicate that there are OUT token (except for the OUT zero) that has been received. The F/W clears the bit after the OUT data has been read. Also, this bit will be cleared by the USB SIE after the next valid SETUP token is received.			
IN	R	This bit is used to indicate that the current USB receiving signal from the PC Host is IN to- ken.			
NAK	R	This bit is used to indicate that the USB SIE has transmitted the NAK signal to the Host in response to the PC Host IN or OUT token.			
CRC_err	R/W	This bit indicates that there are CRC error (bit=1). The programmer must do something to save the device and keep it alive. This bit is set by the USB SIE and cleared by F/W.			
EOT	R	End of transient flag, normal status is 1. If suspend="1" line & EOT="0" indicates that something is wrong in the USB Interface. The programmer must do something to save the device and keep it alive.			
MNI	R/W	This bit is for masking the NAK interrupt when MNI="1", the default value="0"			

SIES Function Table

The Misc register is actually a command + status to control the desired FIFO action and to show the status of the desired FIFO. Every bit's meaning and usage are listed as follows:

Bit No.	Function	Read/Write	Register Address
7	Len0	R/W	
6	Ready	R	
5	Set CMD	R/W	
4	Sel_pipe1	R/W	040004400
3	Sel_pipe0	R/W	01000110B
2	Clear	R/W	
1	Tx	R/W	
0	Request	R/W	

Misc Registers Table



Function Name	Read/Write	Description	
Request	R/W	After setting the other desired status, FIFO can be requested by setting this bit high active. After work has been done, this bit must be set low.	
Тх	R/W	Represents the direction and transition end of the MCU accesses. When being set as logic 1, the MCU wants to write data to FIFO. After work has been done, this bit must be set to logic 0 before terminating the request to represent a transition end. For reading action, this bit must be set to logic 0 to indicate that the MCU wants to read and must be set to logic 1 after work is done.	
Clear	R/W	Represents MCU clear requested FIFO, even if FIFO is not ready.	
Sel_pipe1 Sel_pipe0	H = H = H = H = H = H = H = H = H = H =		
Set (IMD) R/W		Shows that the data in FIFO is setup as command. This bit will be cleared by firmware. So, even if the MCU is busy, nothing is missed by the SETUP command from the host.	
Ready	R	Indicates that the desired FIFO is ready to work.	
Len0 R/W read action to the correspondi		Indicates that the host sent a 0-sized packet to the MCU. This bit must be cleared by a read action to the corresponding FIFO. Also, this bit will be cleared by the USB SIE after the next valid SETUP token is received.	

Misc Function Table

The HT82J97E has two 8×8 bidirectional FIFO for the two endpoints (control and Interrupt). User can easily read/write the FIFO data by accessing the corresponding FIFO pointer register (FIFO0, FIFO1). The following are two examples for reading and writing the FIFO data:

HT82J97E FIFO is read by packet. To read from FIFO, the following should be followed:

- Select one set of FIFO, set in the read mode (MISC TX bit = 0), and set the REQ bit to "1".
- Check the ready bit until the status = 1
- Read through the FIFO pointer register, and record the data number that has been read.
- Repeat steps 2 and 3 until the ready bit becomes 0 which indicates the end of the FIFO data reading.
- Set MISC TX bit = 1
- Clear the REQ bit to 0. Complete reading.

User reads the data through the FIFO pointer register, user has to record the number of bytes to be read. The

HT82J97E allows a maximum of 8 bytes of data in each packet.

The HT82J97E FIFO is written by packet. To write to FIFO, the following should be followed:

- Select a set of FIFO, set in the write mode (MISC TX bit = 1), and set the REQ bit to "1"
- Check the ready bit until the status = 1
- Write through the FIFO pointer register and take down the data number that has been written
- Repeat steps 2 and 3 until writing is complete or the ready bit becomes 0 which indicates that the FIFO no longer allows any data writing.
- Set MISC TX bit = 0
- Clear the REQ bit to 0. Complete writing.

User writes the data through the FIFO pointer register, user has to record the number of bytes that have been written. The HT82J97E allows a maximum of 8 bytes of data in each packet.



There are some timing constrains and usages illustrated here. By setting the MISC register, the MCU can perform reading, writing and clearing actions. There are some examples shown in the following table for endpoint FIFO reading, writing and clearing.

Actions	MISC Setting Flow and Status
Read FIFO0 sequence	00H \rightarrow 01H \rightarrow delay of 2µs, check 41H \rightarrow read* from FIFO0 register and check if not ready (01H) \rightarrow 03H \rightarrow 02H
Write FIFO1 sequence	0AH \rightarrow 0BH \rightarrow delay of 2µs, check 4BH \rightarrow write* to FIFO1 register and check if not ready (0BH) \rightarrow 09H \rightarrow 08H
Check whether FIFO0 can be read or not	00H \rightarrow 01H \rightarrow delay of 2µs, check 41H (if ready) or 01H (if not ready) \rightarrow 00H
Check whether FIFO1 can be written to or not	0AH $\!$
Write 0-sized packet sequence to FIFO 0	02H \rightarrow 03H \rightarrow delay of 2µs, check 43H \rightarrow 01H \rightarrow 00H

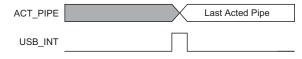
Note: *: There are 2µs gap existing between 2 reading actions or between 2 writing actions

Register Name	R/W	Register Address	Bit7~Bit0	
FIFO 0	R/W	01001000B	Data7~Data0	
FIFO 1	R/W	01001001B	Data7~Data0	

FIFO Register Address Table

USB Active Pipe Timing

The USB active pipe accessed by the host cannot be used by the MCU simultaneously. When the host finishes its work, the signal, a USB_INT will be produced to tell the MCU that the pipe can be used and the acted pipe No. will be shown in the signal, ACT_PIPE as well. The timing is illustrated in the Figure below.



USB Active Pipe Timing

Suspend Wake-Up and Remote Wake-Up

If there is no signal on the USB bus for over 3ms, the HT82J97E will go into a suspend mode. The Suspend line (bit 0 of the USC) will be set to 1 and a USB interrupt is triggered to indicate that the HT82J97E should jump to the suspend state to meet the 500μ A USB suspend current spec.

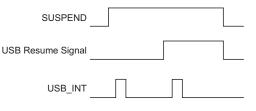
In order to meet the 500 μ A suspend current, the programmer should disable the USB clock by clearing the USBCKEN (bit3 of the SCC) to "0". The suspend current is 400 μ A.

The user can also further decrease the suspend current to 250μ A by setting the SUSP2 (bit4 of the SCC). But if the SUSP2 is set, the user has to make sure not to enable the LVR OPT option, otherwise the HT82J97E will be reset.

When the resume signal is sent out by the host, the HT82J97E will wake-up the MCU by USB interrupt and

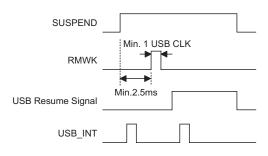
the Resume line (bit 3 of the USC) is set. In order to make the HT82J97E function properly, the programmer must set the USBCKEN (bit 3 of the SCC) to 1 and clear the SUSP2 (bit4 of the SCC). Since the Resume signal will be cleared before the Idle signal is sent out by the host and the Suspend line (bit 0 of the USC) is going to "0". So when the MCU is detecting the Suspend line (bit0 of the USC), the Resume line should be remembered and taken into consideration.

After finishing the resume signal, the suspend line will go inactive and a USB interrupt is triggered. The following is the timing diagram:





The device with remote wake-up function can wake-up the USB Host by sending a wake-up pulse through RMWK (bit 1 of USC). Once the USB Host receive the wake-up signal from the HT82J97E, it will send a Resume signal to the device. The timing is as follows:



To Configure the HT82J97E as PS2 Device

The HT82J97E can be defined as a USB interface or a PS2 interface by configuring the SPS2 (bit 4 of the USR) and SUSB (bit 5 of the USR). If SPS2=1, and SUSB=0, the HT82J97E is defined as PS2 interface, pin USBD- is now defined as PS2 Data pin and USBD+ is now defined as PS2 Clk pin. The user can easily read or write to the PS2 Data or PS2 Clk pin by accessing the corresponding bit PS2DAI (bit 4 of the USC), PS2CKI (bit 5 of the USC), PS2DAO (bit 6 of the USC) and S2CKO (bit 7 of the USC) respectively.

The user should make sure that in order to read the data properly, the corresponding output bit must be set to "1". For example, if user wants to read the PS2 Data by reading PS2DAI, the PS2DAO should be set to "1". Otherwise it always read a "0".

If SPS2=0, and SUSB=1, the HT82J97E is defined as a USB interface. Both the USBD- and USBD+ are driven by the USB SIE of the HT82J97E. User only writes or reads the USB data through the corresponding FIFO.

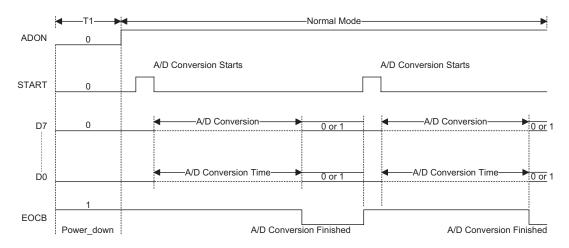
Both SPS2 and SUSB default is "0".

To Configure the ADC Block

The HT82J97E has built-in an 8-bit A/D converter with 6 channels (PB0~PB5). In order to make the A/D converter more flexible, there are two modes: External Reference voltage and Internal Reference voltage. It can be easily configured by setting the ADREF (bit 6 of the USR). For External Reference voltage, the reference voltage of the A/D converter comes from an external PB6/VRL and PB7/VRH pins. Otherwise, the reference voltage is coming from the VDD and VSS of the MCU.

PB0~PB5 is the 6-channel input of the A/D converter, it is easy to define which channel is converting by configuring ACS2~ACS0 (bit 2~0 of the ADSC). Also there are four converter clock sources to be selected by setting ADCS1 (bit 4 of the ADSC), ADCS0 (bit 3 of the ADSC).

Once the ADON (bit 6 of the ADSC) is set, it sends the start pulse through START (bit 5 of ADSC). The A/D converter will be in operation. There are EOCB (bit 7 of the ADSC) to indicate whether the A/D converter is busy or not. The EOCB is cleared when the converter is busy or not. The EOCB is cleared when the converter data by reading the register ADR. In order to meet 500μ A suspend current spec., user should disable the A/D by clearing ADON before jumping to suspend mode.



The following is an A/D converter timing diagram:

To Configure PWM Block

The HT82J97E has two PWM outputs (PWM1 and PWM2), which are shared with PC2, PC3 and can be easily enabled or disabled by the PWM1_EN or PWM2_EN bit of PORT_PC (16H) respectively.

Also there is a one 8-bit PWMBR (PWM Base Period Register, 18H) which defines both PWM output waveform cycle period.

PWM cycle period = $256 \times 1/f_{SYS} \times (PWMBR+1)$, or $256 \times 4/f_{SYS} \times (PWMBR+1)$

where $1/f_{SYS}$ or $4/f_{SYS}$ is defined by PWM_S bit of the PORT_PC (16H)

For example if PWMBR = 17, $4/f_{SYS}$ (T1) is selected and f_{SYS} = 6MHz.

So both output waveform cycle period is $256 \times 4/6 \times (17+1) =$ about $3072 \mu s (0.325 \text{kHz})$

Now user can easily define the corresponding PWM duty by configuring the PWM1DR (for PWM1) or PWM2DR (for PWM2) duty registers

PWM1 duty (high pulse) = (PWM1DR+1)/256×100%

I/O Port Special Registers Definition

PWM1 high pulse period = PWM1 duty×PWM cycle period

PWM1 Low pulse period = PWM cycle period-high pulse period

PWM2 duty (high pulse) = (PWM2DR+1)/256×100%

PWM2 high pulse period = PWM2 duty×PWM cycle period

PWM2 Low pulse period = PWM cycle period-high pulse period

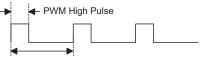
For example PWMBR=17, PWM1DR=63, 4/f_{SYS} (T1) is selected and f_{SYS} =6MHz

PWM cycle period = $256 \times 4/6 \times (17+1)$ = about $3072 \mu s$ (0.325kHz)

PWM1 duty = (63+1)/256 = 25%

PWM1 high pulse period = $25\% \times 3072\mu s$ = $768\mu s$

PWM1 low pulse period = $3072\mu s - 768\mu s = 2304\mu s$



PWM Cycle Period

• Port-A (12H) – PA								
Register	Bits	Labels	Read/Write	Option	Functions			
	0	PA0	R/W	_	I/O (R/W) has pull-low and pull-high ROM code option. Has falling edge wake-up ROM code option.			
	1	PA1	R/W	_	I/O (R/W) has pull-low and pull-high option. Has falling edge wake-up option.			
	2	PA2	R/W	_	I/O (R/W) has pull-low and pull-high option. Has falling edge and rising edge wake-up option.			
	3	PA3	R/W	_	I/O (R/W) has pull-low and pull-high option. Has falling edge and rising edge wake-up option.			
PA (12H)	4	PA4	R/W	_	I/O (R/W) has pull-high option. Has falling edge wake-up option.			
	5	PA5	R/W	_	I/O (R/W) has pull-high option. Has falling edge wake-up option.			
	6	PA6	R/W	_	I/O (R/W) has pull-high option. Has falling edge wake-up option.			
	7	PA7	R/W	_	I/O (R/W) has pull-high option. Has falling edge wake-up option, pin-shared with timer ir put pin.			

Port-A Control (13H) – PAC

This port configure the input or output mode of Port-A



• Port-B Control (14H) – PB

Register	Bits	Labels	Read/Write	Option	Functions
	0	PB0	R/W		I/O (R/W), has pull-high option, ADC input.
	1	PB1	R/W		I/O (R/W), has pull-high option, ADC input.
	2	PB2	R/W		I/O (R/W), has pull-low and pull-high option, ADC input.
	3	PB3	R/W		I/O (R/W), has pull-low and pull-high option, ADC input.
PB (14H)	4	PB4	R/W		I/O (R/W), has pull-high option, can wake-up, ADC input.
(1411)	5	PB5	R/W		I/O (R/W), has pull-high option, ADC input.
	6	PB6	R/W		I/O (R/W), has pull-high option, ADC input, VRL input for ADC external mode.
	7	PB7	R/W		I/O (R/W), has pull-high option, ADC input, VRH input for ADC external mode, has wake-up capability.

• Port-B Control (15H) – PBC

This port configures the input or output mode of Port-B for I/O mode

•	Port-C	Control ((16H)) – PC
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Register	Bits	Labels	Read/Write	Option	Functions
	0	PC0	R/W		I/O (R/W), has pull-high option
	1	PC1	R/W		I/O (R/W), has pull-high option
	2	PC2	R/W		$\ensuremath{I/O}$ (R/W), has pull-high option, can be used as PWM1 output
	3	PC3	R/W		$\ensuremath{I/O}$ (R/W), has pull-high option, can be used as PWM2 output
PC (16H)	4	PC4			Reserved bit
	5	PC5	R/W	PWM_S	PWM base period register frequency source 0= T1 (default) 1= f _{SYS}
	6	PC6	R/W	PWM1_EN	1: Internal register bit, enable PWM1 output 0: Disable (default)
	7	PC7	R/W	PWM2_EN	1: Internal register bit, enable output 0: Disable (default)

• Port-C Control (17H) - PCC

This port is used to control whether the Port-C pin is input or output pin except PC4~PC7

Register	Bits	Labels	Read/Write	Option	Functions
	0	PE0	R	SUSPEND	USB suspend mode status bit. When 1, indicates that the USB system entry is in suspend mode.
	1	PE1	W	RMOT_WK	USB remote wake-up signal. Default value is 0.
	2	PE2	R/W	URST_FLAG	USB bus reset event flag. Default value is 0.
USC	3	PE3	R	RESUME_O	When RESUME_OUT EVENT, RESUME_O is set to 1. Default value is 0.
(0X1A)	4	PE4	R	PS2_DAI	USBD-/DATA input
	5	PE5	R	PS2_CKI	USBD+/CLK input
	6	PE6	W	PS2_DAO	Output for driving USBD-/DATA pin, when working un- der 3D PS2 mouse function. Default value is 1.
	7	PE7	W	PS2_CKO	Output for driving USBD-/DATA pin, when working un- der 3D PS2 mouse function. Default value is 1.

USB/PS2 Status and Control Register USC (Address 0X1A)

Endpoint Interrupt Status Register USR (Address 0X1B)

The USR (USB endpoint interrupt status register) register is used to indicate which endpoint is accessed and to select the serial bus (PS2 or USB) and A/D converter operation modes. The endpoint request flags (EP0IF, EP1IF) are used to indicate which endpoints are accessed. If an endpoint is accessed, the related endpoint request flag will be set to "1" and a USB interrupt will occur (If a USB interrupt is enabled and the stack is not full). When the active endpoint request flag is served, the endpoint request flag has to be cleared to "0".

Register	Bits	Labels	Read/Write	Option	Functions
	0	PEC0	R/W	EP0IF	When set to "1", indicates an endpoint 0 interrupt event. Must wait for the MCU to process the interrupt event and clear this bit by firmware. This bit must be "0", then the next interrupt event will be processed. Default value is "0".
	1	PEC1	R/W	EP1IF	When set to "1", indicates an endpoint 1 interrupt event. Must wait for the MCU to process the interrupt event, then clear this bit by firmware. This bit must be "0", then the next interrupt event will be processed. Default value is "0".
	2	PEC2	R/W	—	Reserved bit, set to "0"
USR (0X1B)	3	PEC3	R/W	—	Reserved bit, set to "0"
(0,115)	4	PEC4	R/W	SELPS2	When set to "1", indicates that the chip is working under PS2 mode. Default value is "0".
	5	PEC5	R/W	SELUSB	When set to "1", indicates that the chip is working under USB mode. Default value is 0.
	6	PEC6	R/W	VRSEL	When set to "0", indicates the reference voltage of the 8-bit ADC from the external input pin. When set to "1", indicates that the reference voltage is from the internal power line. Default value is "1".
	7	PEC7	R/W	USB_flag	This flag is used to show that the MCU is in USB mode (Bit=1). This bit is R/W by FW and will be cleared to zero after power-on reset. The default is "0".



Clock Control Register SCC (Address 0X1C)

There is a system clock control register implemented to select the clock used in the MCU. This register consists of USB clock control bit (USBCKEN), second suspend mode control bit (SUSPEND2) and system clock selection (SCLKSEL).

Register	Bits	Labels	Read/Write	Option	Functions
	2~0	PF2~PF0	R/W		Reserved
	3	PF3	R/W	USBCKEN	USB clock control bit. When set to "1", indicates a USBCK ON, else USBCK OFF. Default value is "0".
	4	PF4	R/W	SUSPEND2	When set to "1", enables a 7.5k Ω resistor connected to D-pin to 5V VDD. Default value is "0".
scc	5	PF5	R/W		Reserved
(0X1C)	6	PF6	R/W	SCLKSEL	System clock 6MHz or 12MHz option, when working on external oscillator mode. Default value is "0". 0: Operating at external 12MHz mode 1: Operating at external 6MHz mode Default value is "0".
	7	PF7	R/W	PS2_flag	This flag is used to show that the MCU is in PS2 mode (Bit=1). This bit is R/W by FW and will be cleared to zero after power-on reset. The default is "0".

ADC Status and Control Register ADC (Address 0X1D)

The A/D converter implemented in the MCU is a 6-channel 8-bit A/D converter. The reference voltage (high reference voltage and low reference voltage) can be selected as coming from external pins (PB6/VRL and PB7/VRH) or internal power supplies of the MCU (VDD and VSS). The VRL and VRH are used to set the minimal and maximal boundaries of the full-scale range of the A/D converter. If an analog input, VRL or VRH is not used for A/D conversion, it can also be used as a general purpose I/O line. The ADSC (A/D converter status and control register) register is used to set the configurations and A/D clock sources of the A/D converter and controls the operation of the A/D converter.

Register	Bits	Labels	Read/Write	Option	Functions
	2~0	PFC2~PFC0	R/W	SEL_CH	These four bits selects one of the eight ADC channels for conversion. Channels 0 to 5 correspond to inputs AD0~AD5 on port pins PB0-PB5 respectively. Chan- nels 6 and 7 are the ADC reference inputs VRH and VRL, on port pins PB6 and PB7 respectively. 000: AD0 (PB0); 001: AD1 (PB1) 010: AD2 (PB2); 011: AD3 (PB3) 100: AD4 (PB4); 101: AD5 (PB5) 110: AD6 or VRL (PB6); 111: AD7 or VRH (PB7) Default value is 000'B.
ADC (0X1D)	4~3	PFC4~PFC3	R/W	SEL_CLK	Selecting ADC operating clock. 00: 6MHz (Default clock) 01: 3MHz 10: 1.5MHz 11: 0.75MHz
	5	PFC5	R/W	START	Start of ADC conversion. High active. Default value is "0"
	6	PFC6	R/W	ADON	Enable pin. ADON=1, Enable ADC block. Default value is "0".
	7	PFC7	R/W	EOCB	End of conversion. This read-only status bit is cleared when a conversion is completed, indicating that the ADC Data Register contains a valid result.



ADC High-byte Data Register ADCR (Address 0X1E)

,					,	
	Register	Bits	Labels	Read/Write	Option	Functions
	ADCR (0X1E)	7~0	PG7~PG0	R	ADCDR	The ADCDR stores the result of a valid ADC conversion bit7~bit0.

Table High Byte Pointer for Current Table Read TBHP (Address 0X1F)

Regist	er B	Bits	Labels	Read/Write	Option	Functions
TBHI (0X1F	- 9	2~0	PGC2~PG0	R	—	Store current table read bit10~bit8 data

PWM Base Period Register PWMBR (Address 0X18)

This register is used to define the base period of the PWM cycle period. The period is defined according to the following equation:

Base period = $(4/f_{SYS}) \times (PWMBR+1)$ or $(1/f_{SYS}) \times (PWMBR+1)$

Where $4/f_{SYS}$ or $1/f_{SYS}$ is defined by PWM_S bit of PORT_PC

Where PWMBR = 1~255, PWMBR=0 is not available

PWM cycle period = 256×Base period

Base period equals to 1/256 duty cycle.

Register	Bits	Labels	Read/Write	Option	Functions
PWMBR (0X18)	7~0	PD7~PD0	R		Used to define the base period of the PWM Range =2~256×Base Period Where PWMBR=1~255, PWMBR=0 is not available

PWM Duty Register PWM1DR (Address 0XCH) and PWM2DR (Address 0XDH)

This register is used to define the duty of the PWM1 output (PC2) or PWM2 output (PC3) respectively. Both PWM cycle frequency is defined according to the following equation:

Register	Bits	Read/Write	Option	Functions
PWM1DR (0XCH) PWM2DR (0XDH)	7~0	R/W	_	Used to define the PWM duty

PWM1 duty = (PWM1DR+1)/PWM cycle×100% period

Where PWM1DR= 0~255

If the PWM function is enabled by setting the corresponding bit (PWM1_EN or PWM2_EN of Port C), the PWM output (PC2 or PC3) pins always output the PWM signal whether the corresponding control register bit (PCC2 or PCC3) is defined as in input or output mode.

OTP	Options
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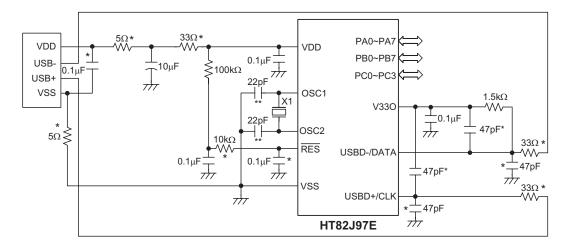
Option
WDT clock source: RC (system/4) (default: T1)
WDT clock source: enable/disable for normal mode (default: disable)
PA0~PA7, PB4, PB7 wake-up by bit (PA2, PA3 both wake-up by falling or rising edge) (default: non wake-up)
PA0~PA7 pull-high by bit (default: Pull-high)
PC0~3,PB pull-high by nibble (default: Pull-high)
2.7 V (error 0.3V) LVR enable/disable (default: enable)
PA0~PA3, PB2, PB3 Pull-low by bit (default: non pull-low 30kΩ)
"CLR WDT", 1 or 2 instructions
TBHP enable/disable (default: disable)
PA output mode (CMOS/NMOS/PMOS) by bit (default: CMOS)

The LVR voltage is define as $2.7V\pm0.3V$ and default is enable.



Application Circuits

Crystal or Ceramic Resonator for Multiple I/O Applications



Note: The resistance and capacitance for the reset circuit should be designed in such a way as to ensure that the VDD is stable and remains within a valid operating voltage range before bringing RES to high.

X1 can use 6MHz or 12MHz, X1 as close OSC1 & OSC2 as possible

Components with * are used for EMC issue.

Components with ** are used for resonator only.



Instruction Set Summary

Mnemonic	Description	Instruction Cycle	Flag Affected
Arithmetic		1	
ADD A,[m] ADDM A,[m] ADD A,x ADC A,[m] ADCM A,[m] SUB A,x SUB A,[m] SUBM A,[m] SBC A,[m] SBCM A,[m] DAA [m]	Add data memory to ACC Add ACC to data memory Add immediate data to ACC Add data memory to ACC with carry Add ACC to data memory with carry Subtract immediate data from ACC Subtract data memory from ACC Subtract data memory from ACC with result in data memory Subtract data memory from ACC with carry Subtract data memory from ACC with carry Subtract data memory from ACC with carry and result in data memory Decimal adjust ACC for addition with result in data memory	$ \begin{array}{c} 1\\ 1^{(1)}\\ 1\\ 1\\ 1^{(1)}\\ 1\\ 1^{(1)}\\ 1\\ 1^{(1)}\\ 1^{(1)}\\ 1^{(1)} \end{array} $	Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV Z,C,AC,OV C
Logic Operati	on		
AND A,[m] OR A,[m] XOR A,[m] ANDM A,[m] ORM A,[m] XORM A,[m] AND A,x OR A,x XOR A,x CPL [m] CPLA [m]	AND data memory to ACC OR data memory to ACC Exclusive-OR data memory to ACC AND ACC to data memory OR ACC to data memory Exclusive-OR ACC to data memory AND immediate data to ACC OR immediate data to ACC Exclusive-OR immediate data to ACC Complement data memory Complement data memory with result in ACC	$ \begin{array}{c c} 1 \\ 1 \\ 1 \\ 1^{(1)} \\ 1^{(1)} \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1$	Z Z Z Z Z Z Z Z Z Z Z
Increment & D	Decrement		
INCA [m] INC [m] DECA [m] DEC [m]	Increment data memory with result in ACC Increment data memory Decrement data memory with result in ACC Decrement data memory	$ \begin{array}{c c} 1 \\ 1^{(1)} \\ 1 \\ 1^{(1)} \end{array} $	Z Z Z Z
Rotate			
RRA [m] RR [m] RRCA [m] RRC [m] RLA [m] RLCA [m] RLCA [m]	Rotate data memory right with result in ACC Rotate data memory right Rotate data memory right through carry with result in ACC Rotate data memory right through carry Rotate data memory left with result in ACC Rotate data memory left Rotate data memory left Rotate data memory left through carry with result in ACC Rotate data memory left through carry	$\begin{array}{c} 1 \\ 1^{(1)} \\ 1 \\ 1^{(1)} \\ 1 \\ 1^{(1)} \\ 1 \\ 1^{(1)} \\ 1 \\ 1^{(1)} \end{array}$	None C C None None C C
Data Move			
MOV A,[m] MOV [m],A MOV A,x	Move data memory to ACC Move ACC to data memory Move immediate data to ACC	1 1 ⁽¹⁾ 1	None None None
Bit Operation			
CLR [m].i SET [m].i	Clear bit of data memory Set bit of data memory	1 ⁽¹⁾ 1 ⁽¹⁾	None None



Mnemonic	Description	Instruction Cycle	Flag Affected
Branch			
JMP addr	Jump unconditionally	2	None
SZ [m]	Skip if data memory is zero	1 ⁽²⁾	None
SZA [m]	Skip if data memory is zero with data movement to ACC	1 ⁽²⁾	None
SZ [m].i	Skip if bit i of data memory is zero	1 ⁽²⁾	None
SNZ [m].i	Skip if bit i of data memory is not zero	1 ⁽²⁾	None
SIZ [m]	Skip if increment data memory is zero	1 ⁽³⁾	None
SDZ [m]	Skip if decrement data memory is zero	1 ⁽³⁾	None
SIZA [m]	Skip if increment data memory is zero with result in ACC	1 ⁽²⁾	None
SDZA [m]	Skip if decrement data memory is zero with result in ACC	1 ⁽²⁾	None
CALL addr	Subroutine call	2	None
RET	Return from subroutine	2	None
RET A,x	Return from subroutine and load immediate data to ACC	2	None
RETI	Return from interrupt	2	None
Table Read			
TABRDC[M](5)	Read ROM code (locate by TBLPand TBHP) to data memory and TBLH	2 ⁽¹⁾	None
TABRDC[m] ⁽⁶⁾	Read ROM code (current page) to data memory and TBLH	2 ⁽¹⁾	None
TABRDL [m]	Read ROM code (last page) to data memory and TBLH	2 ⁽¹⁾	None
Miscellaneous			
NOP	No operation	1	None
CLR [m]	Clear data memory	1 ⁽¹⁾	None
SET [m]	Set data memory	1 ⁽¹⁾	None
CLR WDT	Clear Watchdog Timer	1	TO,PDF
CLR WDT1	Pre-clear Watchdog Timer	1	TO ⁽⁴⁾ ,PDF ⁽⁴⁾
CLR WDT2	Pre-clear Watchdog Timer	1	TO ⁽⁴⁾ ,PDF ⁽⁴⁾
SWAP [m]	Swap nibbles of data memory	1 ⁽¹⁾	None
SWAPA [m]	Swap nibbles of data memory with result in ACC	1	None
HALT	Enter power down mode	1	TO,PDF

Note: x: Immediate data

m: Data memory address

A: Accumulator

i: 0~7 number of bits

addr: Program memory address

 $\sqrt{}$: Flag is affected

-: Flag is not affected

⁽¹⁾: If a loading to the PCL register occurs, the execution cycle of instructions will be delayed for one more cycle (four system clocks).

⁽²⁾: If a skipping to the next instruction occurs, the execution cycle of instructions will be delayed for one more cycle (four system clocks). Otherwise the original instruction cycle is unchanged.

- $^{(3)}$: $^{(1)}$ and $^{(2)}$
- ⁽⁴⁾: The flags may be affected by the execution status. If the Watchdog Timer is cleared by executing the "CLR WDT1" or "CLR WDT2" instruction, the TO and PDF are cleared. Otherwise the TO and PDF flags remain unchanged.
- ⁽⁵⁾: "ROM code TBHP option" is enabled
- ⁽⁶⁾: "ROM code TBHP option" is disabled



Instruction Definition

ADC A,[m]	Add data	memory a	nd carry to	the accu	mulator	
Description	The conte	ents of the	specified of	data mem	ory, accum ccumulator	
Operation	$ACC \leftarrow A$	CC+[m]+0	2			
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
		_	\checkmark	\checkmark	\checkmark	\checkmark
ADCM A,[m]	Add the a	iccumulato	or and carry	/ to data r	nemory	
Description					ory, accum pecified da	
Operation	$[m] \leftarrow AC$	C+[m]+C				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
		_	\checkmark	\checkmark	\checkmark	\checkmark
ADD A,[m]	Add data	memory to	o the accur	nulator		
Description					ory and the	e accumi
Description		the accum	•			accum
Operation	$ACC \leftarrow A$	CC+[m]				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
			\checkmark	\checkmark	\checkmark	\checkmark
ADD A,x	Add imme	ediate data	a to the acc	cumulator		
ADD A,x Description		ents of the			specified o	lata are a
-	The conte	ents of the ator.			specified o	lata are a
Description	The conte accumula	ents of the ator.			specified o	lata are a
Description	The conte accumula	ents of the ator.			specified of	data are a
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Description Operation Affected flag(s)	The conte accumula ACC ← A TO	PDF	ov N	z	AC √	С
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Description Data in the accumulator and the specified data memory perforeration. The result is stored in the accumulator. Operation ACC \leftarrow ACC "AND" [m] Affected flag(s) \overline{TO} PDF OV Z AC C AND A,x Logical AND immediate data to the accumulator Description Data in the accumulator and the specified data perform a bite The result is stored in the accumulator. Operation ACC \leftarrow ACC "AND" x Affected flag(s) \overline{TO} PDF OV Z AC C $ -$ AND A,x Logical AND data memory with the accumulator. Operation ACC \leftarrow ACC "AND" x Affected flag(s) \overline{TO} PDF OV Z AC C $ -$ ANDM A,[m] Logical AND data memory with the accumulator Description Data in the specified data memory. Operation [m] \leftarrow ACC "AND" [m] Affected flag(s) \overline{TO} PDF OV Z AC C $ -$ Operation [m] \leftarrow ACC "AND" [m] Affected flag(s) \overline{TO} PDF OV Z AC C $ -$ Operation The instruction unconditionally calls a subroutine located at program counter increments once to obtain the address of the this onto the sta	AND A,[m]	Logical AND accumulator with data memory
Affected flag(s) $\hline TO PDF OV Z AC C \\ \hline - & - & - & - & - & - & - & - & - & -$		Data in the accumulator and the specified data memory perfo
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Affected flag(s) $\overrightarrow{PO} \leftarrow addr$ $\overrightarrow{TO} \overrightarrow{PDF} \overrightarrow{OV} \overrightarrow{Z} \overrightarrow{AC} \overrightarrow{C}$ $ -$	Description	program counter increments once to obtain the address of the this onto the stack. The indicated address is then loaded. F
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	Affected flag(s)	
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DescriptionThe contents of the specified data memory are cleared to 0.Operation $[m] \leftarrow 00H$ Affected flag(s)		
Operation [m] ← 00H Affected flag(s)	CLR [m]	Clear data memory
Affected flag(s)	Description	The contents of the specified data memory are cleared to 0.
	Operation	[m] ← 00H
TO PDF OV Z AC C — — — — — — —	Affected flag(s)	
		TO PDF OV Z AC C



CLR [m].i	Clear bit of	of data me	mory			
Description	The bit i c	f the spec	ified data ı	memory is	cleared to	o 0.
Operation	[m].i ← 0					
Affected flag(s)						
	то	PDF	OV	Z	AC	С
			_		_	
CLR WDT	Clear Wa	tchdog Tin	ner			
Description	The WDT cleared.	is cleared	(clears the	e WDT). TI	ne power c	lown bit (l
Operation	WDT \leftarrow 0 PDF and					
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
	0	0			_	
CLR WDT1	Preclear	Watchdog	Timer			
Description	of this inst plies this	ruction wil	WDT2, clea thout the of has been	ther precle	ar instruct	ion just se
Operation	$WDT \leftarrow C$	0H*				
	DDE and	TO / 0*				
Affected flag(s)	PDF and	TO ← 0*				
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Description Each bit of the specified data memory is logically complemented (1's complement), E which previously contained a 1 are changed to 0 and vice-versa. The complemented resiles stored in the accumulator and the contents of the data memory remain unchanged. Operation ACC \leftarrow [m] Affected flag(s) \overline{TO} \overline{PDF} \overline{OV} \overline{ACC} DAA [m] Decimal-Adjust accumulator for addition Description The accumulator value is adjusted to the BCD (Binary Coded Decimal) code. The accumulator is divided into two nibble of the accumulator is greater than 9. The BCD justment is done by adding 6 to the original value if the original value is greater than 9. The BCD is used with the ord or (ACC - ACC - Q) = 0. ACC = 1 the data memory and only the carry flag (C) may be affected. Operation If ACC.3-ACC.0 >9 or ACC=1 then [m].7-[m].4 \leftarrow ACC.7-ACC.4+AC1.0=C Affected flag(s) \overline{TO} \overline{PDF} \overline{OV} \overline{AC} \overline{C} DEC [m] Decrement data memory Decrement data memory Decrement data memory Decrement data memory and only the carry flag (C) may be affected. Decrement data in the specified data memory is decremented by 1. Operation $[f ACC.7-ACC.4+AC1.9 or C=1$ then [m].7-[m].4 $\leftarrow ACC.7-ACC.4+AC1.C=C$ $[m] \leftarrow [m] - 1$ DEC [m] Decrement data memory $[m] \leftarrow [m] - 1$ $[m] \leftarrow [m] - 1$ $[m] \leftarrow [m] - 1$	CPLA [m]	Compleme	nt data m	emorv and	d place re	sult in the	accumulat	tor	
Affected flag(s) $\hline TO PDF OV Z AC C$ DAA [m] Decimal-Adjust accumulator for addition Description The accumulator value is adjusted to the BCD (Binary Coded Decimal) code. The accumulator is divided into two nibbles. Each nibble is adjusted to the BCD code and an inter carry (AC1) will be done if the low nibble of the accumulator is greater than 9. The BCD (Justement is done by adding 6 to the original value is greater than 9. The BCD or justement is done by adding 6 to the original value is greater than 9. The BCD or carry (AC or C) is set; otherwise the original value remains unchanged. The result is store in the data memory and only the carry flag (C) may be affected. Operation If ACC.3-ACC.0 >9 or AC=1 then [m].3-[m].0 \leftarrow (ACC.3-ACC.0), AC1=0 and If ACC.7-ACC.4+AC1 >9 or C=1 then [m].7-[m].4 \leftarrow ACC.7-ACC.4+AC1,C=1 else [m].7-[m].4 \leftarrow ACC.7-ACC.4+AC1,C=1 DEC [m] Decrement data memory Decrement data memory Decremented by 1. Operation		Each bit of which previ	the spec ously cor	ified data ntained a 1	memory is are chang	s logically jed to 0 an	compleme d vice-vers	ented (1's compleme sa. The complemente	ed res
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	Operation	$ACC \leftarrow [m]$]						-
DAA [m] Decimal-Adjust accumulator for addition Description The accumulator value is adjusted to the BCD (Binary Coded Decimal) code. The accum lator is divided into two nibbles. Each nibble of the accumulator is greater than 9. The BCD (justment is done by adding 6 to the original value is mainter carry (AC1) will be done if the low nibble of the accumulator is greater than 9. The BCD (binary Coded Decimal) code. The accum is to carry (AC1) will be done if the low nibble of the accumulator is greater than 9. The BCD (binary Coded Decimal) code. The result is store in the data memory and only the carry flag (C) may be affected. Operation If ACC.3-ACC.0 >9 or AC=1 then [m].3-[m].0 \leftarrow (ACC.3-ACC.0)+6, AC1=\overline{AC} else [m].3-[m].0 \leftarrow (ACC.3-ACC.0)+6, AC1=\overline{AC} else [m].3-[m].0 \leftarrow (ACC.3-ACC.0), AC1=0 and if ACC.7-ACC.4+AC1 >9 or C=1 then [m].7-[m].4 ← ACC.7-ACC.4+AC1,C=1 else [m].7-[m].4 ← ACC.7-ACC.4+AC1,C=1 DEC [m] Decrement data memory Decrement data memory Decremented by 1. Operation [m] ← [m]-1 Affected flag(s) $\overline{TO} PDF OV Z AC C C D Decompleter data memory and place result in the accumulator Description Data in the specified data memory remain unchanged. Operation ACC ←$	Affected flag(s)								
DAA [m] Decimal-Adjust accumulator for addition Description The accumulator value is adjusted to the BCD (Binary Coded Decimal) code. The accumulator is divided into two nibbles. Each nibble is adjusted to the BCD code and an inter carry (AC1) will be done if the low nibble of the accumulator is greater than 9. The BCD justment is done by adding 6 to the original value if the original value is greater than 9 c carry (AC or C) is set; otherwise the original value (C) may be affected. Operation If ACC.3~ACC.0~9 or AC=1 then [m].3~[m].0 (- (ACC.3~ACC.0)+6, AC1=AC else [m].3~[m].0 (- (ACC.3~ACC.0), AC1=0) and If ACC.7~ACC.4+AC1 >9 or C=1 then [m].7~[m].4 (- ACC.7~ACC.4+6+AC1,C=1) else [m].7~[m].4 (- ACC.7~ACC.4+6+AC1,C=C) Affected flag(s) TO PDF OV Z AC C Description Data in the specified data memory is decremented by 1. Operation [m] (- [m]-1) Affected flag(s) TO PDF OV Z AC C DEC [m] Decrement data memory Idecremented by 1. Operation [m] (- [m]-1) DEC [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1. Operation [m] (- [m]-1) Decrement data memory and place result in the accumulator Description Data in the specified data memory is decr		ТО	PDF	OV	Z	AC	С]	
Description The accumulator value is adjusted to the BCD (Binary Coded Decimal) code. The accumulator is divided into two nibbles. Each nibble is adjusted to the BCD code and an inter carry (AC1) will be done if the low nibble of the accumulator is greater than 9. The BCD justment is done by adding 6 to the original value if the original value is greater than 9. The BCD core and an inter carry (AC or C) is set; otherwise the original value if the original value is greater than 9. The BCD core are (AC or C) is set; otherwise the original value if the original value is greater than 9. The BCD core are (AC or C) is set; otherwise the original value is mechanged. The result is store in the data memory and only the carry flag (C) may be affected. Operation If ACC.3~ACC.0 > 9 or AC=1 then [m].3~[m].0 ~ (ACC.3~ACC.0)+6, AC1=AC else [m].3~[m].0 ~ (ACC.3~ACC.0)+6, AC1=C and If ACC.7~ACC.4+AC1.9 or C=1 then [m].7~[m].4 ~ ACC.7~ACC.4+AC1,C=1 else [m].7~[m].4 ~ ACC.7~ACC.4+AC1,C=C Affected flag(s) TO PDF OV Z AC C C DEC [m] Decrement data memory Description Data in the specified data memory is decremented by 1. Operation [m] <- [m]-1			_	_	\checkmark]	
Interpretation <t< td=""><td>DAA [m]</td><td>Decimal-Ac</td><td>djust accu</td><td>umulator fo</td><td>or addition</td><td></td><td></td><td></td><td></td></t<>	DAA [m]	Decimal-Ac	djust accu	umulator fo	or addition				
then $[m].3-[m].0 \leftarrow (ACC.3-ACC.0)+6, AC1=\overline{AC}$ else $[m].3-[m].0 \leftarrow (ACC.3-ACC.0), AC1=0$ and If ACC.7-ACC.4+AC1>9 or C=1 then $[m].7-[m].4 \leftarrow ACC.7-ACC.4+6+AC1,C=1$ else $[m].7-[m].4 \leftarrow ACC.7-ACC.4+AC1,C=C$ Affected flag(s) $\boxed{TO PDF OV Z AC C}{\ \qquad - \qquad - \qquad }$ DEC $[m]$ Decrement data memory Description Data in the specified data memory is decremented by 1. Operation $[m] \leftarrow [m]-1$ Affected flag(s) $\boxed{TO PDF OV Z AC C}{\ \qquad - \qquad - \qquad }$ DEC $[m]$ Decrement data memory is decremented by 1. Operation $[m] \leftarrow [m]-1$ Affected flag(s) $\boxed{TO PDF OV Z AC C}{\ \qquad - \qquad \qquad - \qquad $	Description	lator is divid carry (AC1) justment is carry (AC o	ded into t) will be do done by a r C) is set	wo nibbles one if the lo adding 6 to ;; otherwise	s. Each nil ow nibble o o the origir e the origir	oble is adj of the accu nal value if nal value re	usted to th umulator is the origina emains un	ne BCD code and an greater than 9. The B al value is greater tha changed. The result i	interr BCD a an 9 o
TOPDFOVZACC \checkmark DEC [m]Decrement data memoryDescriptionData in the specified data memory is decremented by 1.Operation[m] \leftarrow [m]-1Affected flag(s) \overline{TO} PDFOVZACC \checkmark -DECA [m]Decrement data memory and place result in the accumulatorDescriptionData in the specified data memory is decremented by 1, leaving the result in the accumulatorDECA [m]Decrement data memory and place result in the accumulatorOperationACC \leftarrow [m]-1Affected flag(s) \overline{TO} PDFOVZACC	Operation	then [m].3~ else [m].3~ and If ACC.7~A then [m].7~	r[m].0 ← ([m].0 ← (ACC.4+AC r[m].4 ← 1	(ACC.3~A (ACC.3~A) C1 >9 or C ACC.7~A(CC.0), AC =1 CC.4+6+A	1=0 C1,C=1			
DEC [m] Decrement data memory Description Data in the specified data memory is decremented by 1. Operation [m] \leftarrow [m]-1 Affected flag(s) \overline{TO} PDF OV Z AC C $ $ $ -$ DECA [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1, leaving the result in the accumulator Description Data in the specified data memory remain unchanged. Operation ACC \leftarrow [m]-1 Affected flag(s) \overline{TO} PDF OV Z AC C \overline{TO} PDF OV Z AC C	Affected flag(s)								
DEC [m] Decrement data memory Description Data in the specified data memory is decremented by 1. Operation [m] \leftarrow [m]-1 Affected flag(s) \overline{TO} PDF OV Z AC C Decrement data memory and place result in the accumulator Decrement data memory is decremented by 1, leaving the result in the accumulator DECA [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1, leaving the result in the accumulator Operation ACC \leftarrow [m]-1 Affected flag(s) TO PDF OV Z AC C Image: The contents of the data memory remain unchanged. TO PDF OV Z AC C		ТО	PDF	OV	Z	AC	С]	
Description Data in the specified data memory is decremented by 1. Operation $[m] \leftarrow [m]-1$ Affected flag(s) \overline{TO} PDF OV Z AC C DECA [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1, leaving the result in the accumulator Description Data in the specified data memory remain unchanged. Operation ACC $\leftarrow [m]-1$ Affected flag(s) \overline{TO} PDF OV Z AC C							\checkmark		
Description Data in the specified data memory is decremented by 1. Operation $[m] \leftarrow [m]-1$ Affected flag(s) \overline{TO} PDF OV Z AC C DECA [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1, leaving the result in the accumulator Description Data in the specified data memory remain unchanged. Operation ACC $\leftarrow [m]-1$ Affected flag(s) \overline{TO} PDF OV Z AC C	DEC [m]	Decrement	data me	mory					
Affected flag(s)	Description	Data in the	specified	I data men	nory is de	cremented	l by 1.		
TOPDFOVZACC $$ DECA [m]Decrement data memory and place result in the accumulatorDescriptionData in the specified data memory is decremented by 1, leaving the result in the accumutor. The contents of the data memory remain unchanged.OperationACC \leftarrow [m]-1Affected flag(s)TOPDFOVZACC	Operation	[m] ← [m]–	1						
DECA [m]Decrement data memory and place result in the accumulatorDescriptionData in the specified data memory is decremented by 1, leaving the result in the accumutor. The contents of the data memory remain unchanged.OperationACC \leftarrow [m]-1Affected flag(s)TOPDFOVZACC	Affected flag(s)								
DECA [m] Decrement data memory and place result in the accumulator Description Data in the specified data memory is decremented by 1, leaving the result in the accumutor. The contents of the data memory remain unchanged. Operation ACC \leftarrow [m]–1 Affected flag(s) TO PDF OV Z AC C		ТО	PDF	OV	Z	AC	С]	
DescriptionData in the specified data memory is decremented by 1, leaving the result in the accumutor. The contents of the data memory remain unchanged.OperationACC \leftarrow [m]-1Affected flag(s)TOTOPDFOVZACC						_			
DescriptionData in the specified data memory is decremented by 1, leaving the result in the accumutor. The contents of the data memory remain unchanged.OperationACC \leftarrow [m]-1Affected flag(s)TOTOPDFOVZACC	DECA [m]	Decrement	data me	mory and i	place resu	ilt in the a	ccumulato	r	
Affected flag(s)		Data in the	specified	data mem	ory is deci	remented	by 1, leavir		cumu
TO PDF OV Z AC C	Operation	$ACC \gets [m]$]—1						
	Affected flag(s)								
		ТО	PDF	OV	Z	AC	С]	
		_		_]	



	- (
HALT	•	ver down i				<i>cc</i>
Description	the RAM	and registe	os program ers are reta the WDT t	ined. The	WDT and	prescaler
Operation	$PC \leftarrow PC$					
	$PDF \leftarrow 1$					
	$TO \leftarrow 0$					
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
	0	1	_			
INC [m]	Incremen	t data mer	mory			
Description	Data in th	e specifie	d data mer	mory is inc	remented	by 1
Operation	[m] ← [m]]+1				
Affected flag(s)						
	то	PDF	OV	Z	AC	С
				\checkmark		
INCA [m]			mory and p			
Description		•	d data men the data n	•		•
Operation	ACC ← [i			lemory rel		angeu.
Affected flag(s)	700 (- [i					
, moolod mag(o)	ТО	PDF	OV	Z	AC	С
				√		Ū
				v		
JMP addr	Directly ju	ımp				
Description			er are repla		he directly	-specified
	control is	passed to	this destir	nation.		
Operation	PC ←ado	lr				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
		—	_	—	—	
MOV A,[m]	Move dat	a memory	to the acc	umulator		
Description	The conte	ents of the	specified	data mem	ory are co	pied to the
Operation	ACC ← [I	n]				
Affected flag(s)						
	то	PDF	OV	Z	AC	С
		_				
		1				



MOV A,x	Move immed	liate data	to the ac	cumulator	r			
Description	The 8-bit dat	a specifie	ed by the	code is lo	aded into	the accur	nulator.	
Operation	$ACC \gets x$							
Affected flag(s)							_	
	ТО	PDF	OV	Z	AC	С		
	_	_				_]	
MOV [m],A	Move the ac	cumulato	r to data	memory				
Description	The contents memories).	s of the ac	cumulate	or are copi	ed to the	specified	data memory (one	of the dat
Operation	[m] ←ACC							
Affected flag(s)								
	ТО	PDF	OV	Z	AC	С]	
		_						
NOP	No operation					·11. 11.	d for a forward	
Description			med. Exe	ecution co	ntinues w	ith the ne	t instruction.	
Operation	$PC \leftarrow PC+1$							
Affected flag(s)							1	
	ТО	PDF	OV	Z	AC	С	-	
	_	—	_	_				
OR A,[m]	Logical OR a	accumula	tor with d	ata memo	ory			
	Data in the a	accumulat	tor and th	ie specifie	d data me		e of the data men he accumulator.	nories) pe
Description	Data in the a	accumulat e logical_	tor and th _OR oper	ie specifie	d data me		e of the data men he accumulator.	nories) pe
OR A,[m] Description Operation Affected flag(s)	Data in the a form a bitwis	accumulat e logical_	tor and th _OR oper	ie specifie	d data me			nories) pe
Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$	accumulat e logical_	tor and th _OR oper	ie specifie	d data me			nories) pe
Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$	accumulat e logical_ 2 ″OR″ [m	tor and th _OR oper ۱]	e specifie ation. The	d data me e result is	stored in		nories) pe
Description Operation Affected flag(s)	Data in the a form a bitwis $ACC \leftarrow ACC$	accumulat e logical_ C "OR" [m PDF	tor and th _OR oper n] 	e specifie ation. The Z √	AC	stored in		nories) pe
Description Operation Affected flag(s) OR A,x	Data in the a form a bitwis ACC ← ACC TO Logical OR in	accumulat e logical_ "OR" [m PDF mmediate accumula	tor and th _OR oper] OV e data to tor and ti	the specifie ration. The Z √ the accum	AC	C		
Description Operation Affected flag(s) OR A,x Description	Data in the a form a bitwis ACC ← ACC TO Logical OR in Data in the a	accumulat e logical_ C "OR" [m PDF mmediate accumula stored in	tor and th _OR oper] OV e data to tor and ti	the specifie ration. The Z √ the accum	AC	C	he accumulator.	
Description Operation Affected flag(s) OR A,x Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is	accumulat e logical_ C "OR" [m PDF mmediate accumula stored in	tor and th _OR oper] OV e data to tor and ti	the specifie ration. The Z √ the accum	AC	C	he accumulator.	
Description Operation Affected flag(s) OR A,x Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$	accumulat e logical_ C "OR" [m PDF mmediate accumula stored in	tor and th _OR oper] OV e data to tor and ti	the specifie ration. The Z √ the accum	AC	C	he accumulator.	
Description Operation Affected flag(s) OR A,x Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$	PDF mmediate accumula stored in C "OR" x	tor and th _OR oper n] OV e data to tor and th the accu	the specifie Z √ the accum ne specifie mulator.	AC	C C erform a b	he accumulator.	
Description Operation Affected flag(s) OR A,x Description Operation Affected flag(s)	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO TO	PDF PDF CrOR" [m PDF CrOR" accumula stored in PDF PDF PDF CrOR" x PDF PDF PDF	OV OV OV OV OV OV OV OV	the accum me specifie $\frac{Z}{}$ the accum mulator. $\frac{Z}{}$	AC	C C erform a b	he accumulator.	
Description Operation Affected flag(s) OR A,x Description Operation Affected flag(s) ORM A,[m]	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO TO Logical OR of	accumulative logical_ "OR" [m PDF PDF accumula stored in "OR" x PDF PDF accumula	OV OV OV OV OV OV OV OV	the accumption of the accumptic of the accumption of the accumption of the accumpti	AC AC AC AC AC AC AC AC AC	C C erform a b C	he accumulator.	e operation
Description Operation Affected flag(s) OR A,x Description Operation Affected flag(s) ORM A,[m]	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO TO Logical OR of	accumulative logical_ "OR" [m PDF PDF accumula stored in "OR" x PDF C PDF C C C C C C C C C C C C C C C C C C C	OV O	the accumption of the difference of the differe	AC AC AC AC AC AC AC AC AC AC AC AC AC A	C C erform a b C C ories) and	he accumulator.	e operation
Description Operation Affected flag(s) OR A,x Description Operation Affected flag(s) ORM A,[m] Description	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO TO Logical OR o Data in the	accumulat e logical_ ? "OR" [m PDF 	OV O	the accumption of the difference of the differe	AC AC AC AC AC AC AC AC AC AC AC AC AC A	C C erform a b C C ories) and	he accumulator.	e operation
Description Operation Affected flag(s) OR A,x Description Operation Affected flag(s) ORM A,[m] Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO TO Logical OR of Data in the bitwise logical	accumulat e logical_ ? "OR" [m PDF 	OV O	the accumption of the difference of the differe	AC AC AC AC AC AC AC AC AC AC AC AC AC A	C C erform a b C C ories) and	he accumulator.	e operation
Description Operation	Data in the a form a bitwis $ACC \leftarrow ACC$ TO Logical OR in Data in the a The result is $ACC \leftarrow ACC$ TO Logical OR of Data in the bitwise logica [m] $\leftarrow ACC$ "	accumulat e logical_ ? "OR" [m PDF 	OV O	the accumption of the difference of the differe	AC AC AC AC AC AC AC AC AC AC AC AC AC A	C C erform a b C C ories) and	he accumulator.	e operatior



RET	Return fro	m subrou	tine			
Description	The progr	am counte	er is restor	ed from th	ie stack. T	his is a 2-
Operation	$PC \leftarrow Sta$	ck				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
					_	
RET A,x	Return an	d place in	nmediate c	lata in the	accumula	tor
Description	The progr fied 8-bit i		er is restore data.	ed from the	e stack and	the accur
Operation	$PC \leftarrow Sta$ $ACC \leftarrow x$					
Affected flag(s)						
	то	PDF	OV	Z	AC	С
				_	_	
RETI	Return fro	m interrup	ot			
Description			er is restor enable ma			
Operation	PC ← Sta	ick				
Affected flog(c)	EMI ← 1					
Affected flag(s)	ТО	PDF	OV	Z	AC	С
	10	PDF	00	2	AC	
RL [m]	Rotate da	ta memor	y left			
Description	The conte	nts of the	specified d	ata memo	ry are rota	ted 1 bit le
Operation	[m].(i+1) ∢ [m].0 ← [r		n].i:bit i of t	he data m	emory (i=(0~6)
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
RLA [m]	Rotate da	ta memor	y left and p	place resu	It in the ac	cumulato
Description			data men			
·		•	accumula	•		
Operation	ACC.(i+1) ACC.0 ←		m].i:bit i of	the data r	memory (i	=0~6)
Affected flag(s)						
	то	PDF	OV	Z	AC	С
					_	
	-					



RLC [m]	Rotate da	ta memor	y left throu	gh carry			
Description			•		ory and the g is rotated	, 0	
Operation	[m].(i+1) ↔ [m].0 ← C C ← [m].	;].i:bit i of tł	ne data m	emory (i=0	~6)	
Affected flag(s)							
	то	PDF	OV	Z	AC	С	
	_	—	_	—	_	\checkmark	
RLCA [m]	Rotate lef	t through o	carry and p	lace resu	It in the ac	cumulator	
Description	carry bit a	nd the orig	ginal carry	flag is rota	e carry flag ated into bit ae data me	0 position	. The rota
Operation	ACC.(i+1) ACC.0 ← C ← [m].	С	m].i:bit i of	the data r	memory (i=	:0~6)	
Affected flag(s)						1	
	ТО	PDF	OV	Z	AC	С	
	_		—		—	\checkmark	
RR [m]	Rotate da	ta memor	y right				
Description	The conte	nts of the	specified d	ata memo	ry are rotat	ed 1 bit rig	ht with bit
Operation	[m].i ← [n	n].(i+1); [m].i:bit i of th	ne data m	emorv (i=0	~6)	
	[m].7 ← [I	n].0			oo.j (. o	0)	
Affected flag(s)	[m].7 ← [I	n].0			oo.y (0)	
Affected flag(s)	[m].7 ← [r TO	n].0 PDF	OV	Z	AC	с	
Affected flag(s)			OV				
		PDF		Z	AC		
Affected flag(s) RRA [m] Description	TO — Rotate rig Data in th	PDF — ht and pla e specified	 ce result ir d data men	Z — n the accu nory is rot	AC	C —	
RRA [m]	TO — Rotate rig Data in th the rotate	PDF — ht and pla e specified d result in t - [m].(i+1);	ce result in d data men the accumu	Z — n the accu nory is rot	AC —	C — ight with b	
RRA [m] Description	TO — Rotate rig Data in th the rotate ACC.(i) ←	PDF — ht and pla e specified d result in t - [m].(i+1);	ce result in d data men the accumu	Z — n the accu nory is rot	AC — mulator ated 1 bit r	C — ight with b	
RRA [m] Description Operation	TO — Rotate rig Data in th the rotate ACC.(i) ←	PDF — ht and pla e specified d result in t - [m].(i+1);	ce result in d data men the accumu	Z — n the accu nory is rot	AC — mulator ated 1 bit r	C — ight with b	
RRA [m] Description Operation	TO TO Rotate rig Data in th the rotate ACC.(i) \leftarrow ACC.7 \leftarrow	PDF — ht and pla e specified d result in t - [m].(i+1); [m].0	ce result ir d data men the accumu [m].i:bit i d	Z — n the accu nory is rot ulator. The of the data	AC — mulator ated 1 bit r contents c a memory (C — ight with b if the data i=0~6)	
RRA [m] Description Operation	TO TO Rotate rig Data in th the rotate ACC.(i) \leftarrow ACC.7 \leftarrow TO 	PDF — ht and pla e specified d result in t - [m].(i+1); [m].0 PDF —	ce result ir d data men the accumu [m].i:bit i d	Z — n the accu nory is rot ulator. The of the data Z —	AC — mulator ated 1 bit r contents c a memory (C — ight with b if the data i=0~6)	
RRA [m] Description Operation Affected flag(s)	TO 	PDF ht and pla e specified d result in t - [m].(i+1); [m].0 PDF 	ce result ir d data men the accumu [m].i:bit i d OV y right thro specified	Z n the accu nory is rot ulator. The of the data Z ugh carry data mem	AC — mulator ated 1 bit r contents c a memory (C ight with b if the data i i=0~6) C — ne carry fla	memory ro ng are tog
RRA [m] Description Operation Affected flag(s)	TO 	PDF — ht and pla e specified d result in t - [m].(i+1); [m].0 PDF — ta memory ents of the) replaces n].(i+1); [m];	ce result ir d data men the accumu [m].i:bit i d OV y right thro specified the carry b	Z — the accu nory is rot ulator. The of the data Z ugh carry data mem pit; the orig	AC mulator ated 1 bit r contents c a memory (AC hory and th	C ight with b if the data i i=0~6) C inter carry flat flag is rota	memory ro ng are tog
RRA [m] Description Operation Affected flag(s) RRC [m] Description	TO \neg Rotate rigData in thethe rotateACC.(i) \leftarrow ACC.7 \leftarrow TO \neg Rotate daThe conteright. Bit ([m].i \leftarrow [n[m].7 \leftarrow (C	PDF — ht and pla e specified d result in t - [m].(i+1); [m].0 PDF — ta memory ents of the) replaces n].(i+1); [m];	ce result ir d data men the accumu [m].i:bit i d OV y right thro specified the carry b	Z — the accu nory is rot ulator. The of the data Z ugh carry data mem pit; the orig	AC mulator ated 1 bit r contents c a memory (AC AC hory and th ginal carry	C ight with b if the data i i=0~6) C inter carry flat flag is rota	memory ro ng are tog
RRA [m] Description Operation Affected flag(s) RRC [m] Description Operation	TO \neg Rotate rigData in thethe rotateACC.(i) \leftarrow ACC.7 \leftarrow TO \neg Rotate daThe conteright. Bit ([m].i \leftarrow [n[m].7 \leftarrow (C	PDF — ht and pla e specified d result in t - [m].(i+1); [m].0 PDF — ta memory ents of the) replaces n].(i+1); [m];	ce result ir d data men the accumu [m].i:bit i d OV y right thro specified the carry b	Z — the accu nory is rot ulator. The of the data Z ugh carry data mem pit; the orig	AC mulator ated 1 bit r contents c a memory (AC AC hory and th ginal carry	C ight with b if the data i i=0~6) C inter carry flat flag is rota	memory ro ng are tog

HOLTEK							
	HT82J97E						
RRCA [m]	Rotate right through carry and place result in the accumulator						
Description	Data of the specified data memory and the carry flag are rotated 1 bit right. Bit 0 replace the carry bit and the original carry flag is rotated into the bit 7 position. The rotated result i stored in the accumulator. The contents of the data memory remain unchanged.						
Operation	ACC.i \leftarrow [m].(i+1); [m].i:bit i of the data memory (i=0~6) ACC.7 \leftarrow C C \leftarrow [m].0						
Affected flag(s)							
	TO PDF OV Z AC C						
SBC A,[m]	Subtract data memory and carry from the accumulator						
Description	The contents of the specified data memory and the complement of the carry flag are sub tracted from the accumulator, leaving the result in the accumulator.						
Operation	$ACC \leftarrow ACC + [\overline{m}] + C$						
Affected flag(s)							
	TO PDF OV Z AC C						
SBCM A,[m]	Subtract data memory and carry from the accumulator						
Description	The contents of the specified data memory and the complement of the carry flag are sub-						

	ACC.7 ← C ← [m].0						
Affected flag(s)							
	ТО	PDF	OV	Z	AC	С	
						\checkmark	
SBC A,[m]	Subtract d	lata memo	ory and ca	rry from th	e accumu	lator	
Description		nts of the	specified of	data memo	ory and the	e complem	nent of the carry flag are so nulator.
Operation	$ACC \leftarrow AC$	CC+[m]+C	;				
ffected flag(s)							
	ТО	PDF	OV	Z	AC	С	
			\checkmark	\checkmark	\checkmark	\checkmark	
BCM A,[m]	Subtract d	lata memo	ory and ca	rry from th	e accumu	lator	
Description	The conte	nts of the	specified of	data memo	ory and the	e complem	ent of the carry flag are s
	tracted fro	m the acc	umulator,	leaving the	e result in	the data n	nemory.
peration	$[m] \leftarrow AC$	C+[m]+C					
ffeeted flee(e)							
mected hag(s)							1
nected hag(s)	ТО	PDF	OV	Z	AC	С	
inected liag(s)	то —	PDF	OV √	Z √	AC √	C √	
	TO — Skip if dec		\checkmark	\checkmark			
6DZ [m]	Skip if dec The conte instruction	crement da nts of the s is skipper execution	√ ata memor specified d d. If the res n, is discare	√ ry is 0 ata memor sult is 0, th ded and a	√ ry are decr e following dummy cy	√ remented l g instructio cle is repla	by 1. If the result is 0, the n on, fetched during the curr aced to get the proper instr 1 cycle).
SDZ [m] Description	Skip if dec The conte instruction instruction	rement da nts of the s i is skippe executior cles). Othe	√ ata memor specified d d. If the res n, is discard rrwise proc	y is 0 ata memory sult is 0, th ded and a 0 ceed with t	√ ry are decr e following dummy cy	√ remented l g instructio cle is repla	on, fetched during the curr aced to get the proper instr
SDZ [m] Description	Skip if dec The conter instruction instruction tion (2 cyc	rement da nts of the s i is skippe executior cles). Othe	√ ata memor specified d d. If the res n, is discard rrwise proc	y is 0 ata memory sult is 0, th ded and a 0 ceed with t	√ ry are decr e following dummy cy	√ remented l g instructio cle is repla	on, fetched during the curr aced to get the proper instr
SDZ [m] Description	Skip if dec The conter instruction instruction tion (2 cyc	rement da nts of the s i is skippe executior cles). Othe	√ ata memor specified d d. If the res n, is discard rrwise proc	y is 0 ata memory sult is 0, th ded and a 0 ceed with t	√ ry are decr e following dummy cy	√ remented l g instructio cle is repla	on, fetched during the curr aced to get the proper instr
DZ [m] Description	Skip if dec The conter instruction instruction tion (2 cyc Skip if ([m	trement dants of the s nts of the s n is skippe n execution cles). Othe]–1)=0, [m	 ata memor specified d d. If the res h, is discard rwise proc rwise proc	√ ry is 0 ata memor sult is 0, th ded and a 0 ceed with t 1)	√ ry are decr e followinş dummy cy he next in	√ remented l g instructio cle is repla struction (on, fetched during the curr aced to get the proper instr
SDZ [m] Description Operation Affected flag(s)	Skip if dec The conter instruction instruction tion (2 cyc Skip if ([m	crement da nts of the s i is skippe executior executior executior les). Othe]–1)=0, [m PDF	√ ata memor specified d d. If the res n, is discare rwise proc n] ← ([m]– OV	√ y is 0 ata memoor sult is 0, th ded and a c ceed with t 1) Z 	√ ry are decr e following dummy cy he next in AC —	√ remented I g instructio cle is repla struction (C 	on, fetched during the curr aced to get the proper instr
SDZ [m] Description Operation Affected flag(s)	Skip if dec The content instruction tion (2 cyc Skip if ([m TO Decrement The content instruction unchange	rement da nts of the s i is skippe execution les). Othe]–1)=0, [m PDF 	√ ata memor specified d d. If the res rwise proc a] ← ([m] OV OV mory and specified d d. The resu sult is 0, th ded and a	√ y is 0 ata memoor sult is 0, th ded and a do ceed with t 1) Z place resurve ata memoor ult is stored e following dummy cyo	√ ry are decr e following dummy cy he next in AC 	√ remented I g instructio cle is repla struction (C C Skip if 0 remented I cumulator I n, fetched aced to ge	on, fetched during the curr aced to get the proper instr
BDZ [m] Description Operation Affected flag(s) BDZA [m] Description	Skip if dec The conten instruction tion (2 cyc Skip if ([m TO Decremen The conten instruction unchange execution,	rement da nts of the s is skippe executior executior execution [-1)=0, [m PDF 	 ata memor specified d d. If the res rwise proc rwise proc rwise proc $1] \leftarrow ([m]$ OV mory and specified d d. The resu sult is 0, the ded and a c	√ y is 0 ata memory sult is 0, th ded and a d ceed with t 1) Z place result ata memory ult is stored e following dummy cyuthe next in	√ ry are decr e following dummy cy he next in AC 	√ remented I g instructio cle is repla struction (C C Skip if 0 remented I cumulator I n, fetched aced to ge	on, fetched during the curr liced to get the proper instr 1 cycle).
SDZ [m] Description Operation Affected flag(s) SDZA [m] Description	Skip if dec The content instruction tion (2 cyc Skip if ([m TO Decrement The content instruction unchangent execution, cles). Other	rement da nts of the s is skippe executior executior execution [-1)=0, [m PDF 	 ata memor specified d d. If the res rwise proc rwise proc rwise proc $1] \leftarrow ([m]$ OV mory and specified d d. The resu sult is 0, the ded and a c	√ y is 0 ata memory sult is 0, th ded and a d ceed with t 1) Z place result ata memory ult is stored e following dummy cyuthe next in	√ ry are decr e following dummy cy he next in AC 	√ remented I g instructio cle is repla struction (C C Skip if 0 remented I cumulator I n, fetched aced to ge	on, fetched during the curr liced to get the proper instr 1 cycle).
Affected flag(s) SDZ [m] Description Operation Affected flag(s) SDZA [m] Description Operation Affected flag(s)	Skip if dec The content instruction tion (2 cyc Skip if ([m TO Decrement The content instruction unchangent execution, cles). Other	rement da nts of the s is skippe executior executior execution [-1)=0, [m PDF 	 ata memor specified d d. If the res rwise proc rwise proc rwise proc $1] \leftarrow ([m] - 1)$ OV mory and specified d d. The resu sult is 0, the ded and a conceed with	√ y is 0 ata memory sult is 0, th ded and a d ceed with t 1) Z place result ata memory ult is stored e following dummy cyuthe next in	√ ry are decr e following dummy cy he next in AC 	√ remented I g instructio cle is repla struction (C C Skip if 0 remented I cumulator I n, fetched aced to ge	on, fetched during the curr liced to get the proper instr 1 cycle).



SET [m]	Set data memory							
Description	Each bit of the speci	fied data me	emory is	set to 1.				
Operation	[m] ← FFH							
Affected flag(s)								
	TO PDF	OV	Z	AC	С			
			_	_	_			
SET [m]. i	Set bit of data memo	ory						
Description	Bit i of the specified	data memor	y is set t	o 1.				
Operation	[m].i ← 1							
Affected flag(s)								
	TO PDF	OV	Z	AC	С			
		_	_	_	_			
		I		I]			
SIZ [m]	Skip if increment dat	a memory is	s 0					
Description		•		•		by 1. If the result is 0, the fol-		
	-		-			ecution, is discarded and a es). Otherwise proceed with		
	the next instruction (()			
Operation	Skip if ([m]+1)=0, [m] ← ([m]+1)						
Affected flag(s)								
	TO PDF	OV	Z	AC	С			
		_	_	_	_			
				I				
SIZA [m]	Increment data mem	ory and place	ce result	in ACC, s	kip if 0			
Description		•		•		y 1. If the result is 0, the next ulator. The data memory re-		
	mains unchanged. If	the result is	0, the fo	llowing ins	struction, f	etched during the current in-		
	struction execution, instruction (2 cycles)					replaced to get the proper		
Operation	Skip if ([m]+1)=0, AC							
Affected flag(s)	Skip II ([III]+1)−0, AC	JC ← ([iii]+ i)					
Allected lidg(s)	TO PDF	OV	Z	AC	С			
			2	70	0			
		_	_	_	_			
SNZ [m].i	Skip if bit i of the dat	a memory is	s not 0					
Description	•					n is skipped. If bit i of the data		
		-			-	current instruction execution, instruction (2 cycles). Other-		
	wise proceed with th		•	-	ie proper	Instruction (2 cycles). Other-		
Operation	Skip if [m].i≠0		、·	. ,				
Affected flag(s)	,							
3(-)	TO PDF	OV	Z	AC	С			
			_					
				_	_			



SUB A,[m]	Subtract	data memo	ory from th	e accumu	lator	
Description	The spec		nemory is a		from the c	ontents
Operation	$ACC \leftarrow A$	CC+[m]+1				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
			\checkmark	\checkmark	\checkmark	\checkmark
SUBM A,[m]	Subtract	data memo	ory from th	e accumu	lator	
Description		ified data r he data m		subtracted	l from the c	ontents
Operation	$[m] \leftarrow AC$	C+[m]+1				
Affected flag(s)						
	то	PDF	OV	Z	AC	С
		_	\checkmark	\checkmark	\checkmark	\checkmark
SUB A,x	Subtract	mmediate	data from	the accur	nulator	
Description	The imme	ediate data	specified I	by the cod	e is subtrad	cted from
			It in the ac			
Operation	$ACC \leftarrow A$	CC+x+1				
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
		1				
			\checkmark	\checkmark	\checkmark	\checkmark
SWAP [m]	 Swap nib	bles withir	√ the data r		\checkmark	V
SWAP [m] Description	The low-o		the data r	nemory	√ the specifi	
	The low-o ries) are i	order and I	the data r nigh-order ed.	nemory		
Description	The low-o ries) are i	order and I nterchang	the data r nigh-order ed.	nemory		
Description Operation	The low-o ries) are i	order and I nterchang	the data r nigh-order ed.	nemory		
Description Operation	The low-o ries) are [m].3~[m]	order and h nterchang $.0 \leftrightarrow [m].7$	i the data r nigh-order ed. 7~[m].4	nemory nibbles of	the specifi	ed data
Description Operation Affected flag(s)	The low-ories) are [m].3~[m]	order and I nterchang .0 ↔ [m].7 PDF	i the data r nigh-order ed. '~[m].4 OV	nemory nibbles of Z	the specifi AC	ed data
Description Operation	The low-ories) are formed by the low-ories of the low-ori	a memory	the data r nigh-order ed. '~[m].4 OV and place	nemory nibbles of Z result in t	AC	C L Lator
Description Operation Affected flag(s)	The low-ories) are formed for the low-ories of the low-or	order and I nterchang .0 ↔ [m].7 PDF 	the data r nigh-order ed. '~[m].4 OV and place igh-order r	nemory nibbles of Z 	the specifi AC	C C ulator
Description Operation Affected flag(s)	The low-ories) are in [m].3~[m] TO TO Swap date The low-oring the received of the content of the low-oring the received of the content of the	order and I nterchang $.0 \leftrightarrow [m].7$ PDF a memory order and P sult to the CC.0 \leftarrow [r	a the data r nigh-order ed. '~[m].4 OV and place nigh-order r accumulat	nemory nibbles of Z 	AC	C C ulator
Description Operation Affected flag(s) SWAPA [m] Description Operation	The low-ories) are in [m].3~[m] TO TO Swap date The low-oring the received of the content of the low-oring the received of the content of the	order and H nterchang .0 ↔ [m].7 PDF a memory order and h sult to the	a the data r nigh-order ed. '~[m].4 OV and place nigh-order r accumulat	nemory nibbles of Z 	AC	C C ulator
Description Operation Affected flag(s) SWAPA [m] Description	The low-ories) are in [m].3~[m] TO	order and I nterchang $.0 \leftrightarrow [m].7$ PDF a memory order and P sult to the CC.0 \leftarrow [r CC.4 \leftarrow [r	a the data r nigh-order ed. '~[m].4 OV and place nigh-order r accumulat n].7~[m].4 n].3~[m].0	nemory nibbles of Z result in t nibbles of or. The co	AC AC he accumu the specific ontents of t	C C ulator ed data he data
Description Operation Affected flag(s) SWAPA [m] Description Operation	The low-ories) are in [m].3~[m] TO TO Swap date The low-oring the received of the content of the low-oring the received of the content of the	order and I nterchang $.0 \leftrightarrow [m].7$ PDF a memory order and P sult to the CC.0 \leftarrow [r	a the data r nigh-order ed. '~[m].4 OV and place nigh-order r accumulat	nemory nibbles of Z 	AC	C C ulator



SZ [m]	Skip if da	ta memory	is 0				
Description				data mem	ory are 0,	the followir	ng instruction, fetched during
·	the curre	nt instructi	on executi	ion, is disc	arded and	l a dummy	v cycle is replaced to get the kt instruction (1 cycle).
Operation	Skip if [m]=0					
Affected flag(s)							
	ТО	PDF	OV	Z	AC	С	
	_	_	_				
SZA [m]	Move dat	a memory	to ACC, s	kip if 0			
Description			-		rv are copi	ied to the a	occumulator. If the contents is
	0, the foll and a dur	owing inst	ruction, fei	tched durii d to get the	ng the cur	rent instrue	ction execution, is discarded 2 cycles). Otherwise proceed
Operation	Skip if [m						
Affected flag(s)] -					
	ТО	PDF	OV	Z	AC	С	
						_	
]
SZ [m].i	Skip if bit	i of the da	ta memor	y is 0			
Description		-		-		-	on, fetched during the current
						struction (aced to get the proper instruc- 1 cycle).
Operation	Skip if [m	,				,	• /
Affected flag(s)							
	то	PDF	OV	Z	AC	С	
	_	_	_	_		_	
TABRDC [m]			e (locate b	oy TBLP ar	nd TBHP)	to TBLH a	nd data memory (ROM code
	TBHP is e	,					
Description							TBLPand TBHP) is moved to FBLH directly.
Operation		OM code (le		-)			
Affected flag(s)		ROM code	e (nign byt	e)			
/ mooled mag(b)	ТО	PDF	OV	Z	AC	С	
TABRDC [m]	Move the disabled)	ROM cod	de (curren	it page) to	TBLH ar	nd data me	emory (ROM code TBHP is
Description							able pointer (TBLP) is moved o TBLH directly.
Operation		OM code (le ROM code	• •	e)			
Affected flag(s)							
	то	PDF	OV	Z	AC	С	
			_	_		_	

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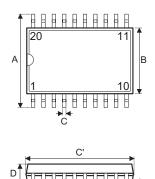
TABRDL [m]			le (last pag			
Description			M code (la: nd the high			
Operation	$[m] \leftarrow RC$	M code (I	ow byte)			
	TBLH ← I	ROM code	e (high byte	e)		
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
		_	_		_	
XOR A,[m]	Logical X	OR accun	nulator with	ı data mer	mory	
Description			ulator and t and the res			
Operation	$ACC \leftarrow A$	CC "XOR	." [m]			
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
			_			
XORM A,[m]	Logical X	OR data n	nemory wit	h the accu	umulator	
Description			ed data me The result			
Operation	$[m] \leftarrow AC$	C "XOR"	[m]			
Affected flag(s)						
	ТО	PDF	OV	Z	AC	С
			_		_	
	L				1	
XOR A,x	e		diate data t			
Description			lator and th is stored in			
Operation	ACC ← A					ie e nag ie
Affected flag(s)						
Anoted hag(s)	ТО	PDF	OV	Z	AC	С
	10	FUP		∠ √	AC	
				N		



Package Information

20-pin SOP (300mil) Outline Dimensions

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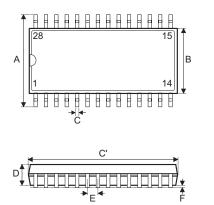
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Symbol	Dimensions in mil						
Symbol	Min.	Nom.	Max.				
A	394		419				
В	290		300				
С	14		20				
C′	490		510				
D	92		104				
E	_	50	_				
F	4						
G	32		38				
н	4		12				
α	0°		10°				



28-pin SOP (300mil) Outline Dimensions



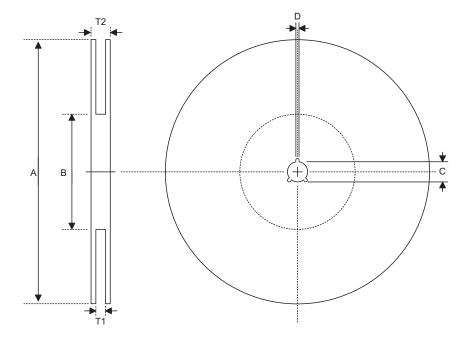


Cumula al	Dimensions in mil						
Symbol	Min.	Nom.	Max.				
А	394		419				
В	290		300				
С	14		20				
C′	697		713				
D	92		104				
E	_	50	_				
F	4		_				
G	32		38				
Н	4		12				
α	0°		10°				



Product Tape and Reel Specifications

Reel Dimensions



SOP 20W

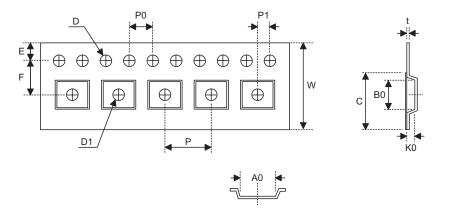
Symbol	Description	Dimensions in mm
А	Reel Outer Diameter	330±1.0
В	Reel Inner Diameter	62±1.5
С	Spindle Hole Diameter	13.0+0.5 _0.2
D	Key Slit Width	2.0±0.5
T1	Space Between Flange	24.8+0.3 _0.2
T2	Reel Thickness	30.2±0.2

SOP 28W (300mil)

Symbol	Description	Dimensions in mm
А	Reel Outer Diameter	330±1.0
В	Reel Inner Diameter	62±1.5
с	Spindle Hole Diameter	13.0+0.5 0.2
D	Key Slit Width	2.0±0.5
T1	Space Between Flange	24.8+0.3 0.2
T2	Reel Thickness	30.2±0.2



Carrier Tape Dimensions



ymbol	Description	Dimensions in mm
W	Carrier Tape Width	24.0+0.3 _0.1
Р	Cavity Pitch	12.0±0.1
E	Perforation Position	1.75±0.1
F	Cavity to Perforation (Width Direction)	11.5±0.1
D	Perforation Diameter	1.5+0.1
D1	Cavity Hole Diameter	1.5+0.25
P0	Perforation Pitch	4.0±0.1
P1	Cavity to Perforation (Length Direction)	2.0±0.1
A0	Cavity Length	10.8±0.1
B0	Cavity Width	13.3±0.1
K0	Cavity Depth	3.2±0.1
t	Carrier Tape Thickness	0.3±0.05
С	Cover Tape Width	21.3

SOP 28W (300mil)

Symbol	Description	Dimensions in mm
W	Carrier Tape Width	24.0±0.3
Р	Cavity Pitch	12.0±0.1
E	Perforation Position	1.75±0.1
F	Cavity to Perforation (Width Direction)	11.5±0.1
D	Perforation Diameter	1.5+0.1
D1	Cavity Hole Diameter	1.5+0.25
P0	Perforation Pitch	4.0±0.1
P1	Cavity to Perforation (Length Direction)	2.0±0.1
A0	Cavity Length	10.85±0.1
В0	Cavity Width	18.34±0.1
K0	Cavity Depth	2.97±0.1
t	Carrier Tape Thickness	0.35±0.01
С	Cover Tape Width	21.3



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