



**A/D Flash MCU with EEPROM**

**HT66F488/HT66F489**

Revision: V1.30 Date: August 17, 2021

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**Note that the HT66F488 device, although mentioned in this datasheet, has already been phased out and is presently no longer available.**

## Features

### CPU Features

- Operating Voltage
  - ♦  $f_{SYS} = 8\text{MHz}$ : 2.2V~5.5V
  - ♦  $f_{SYS} = 10\text{MHz}$ : 2.7V~5.5V
  - ♦  $f_{SYS} = 12\text{MHz}$ : 3.3V~5.5V
  - ♦  $f_{SYS} = 16\text{MHz}$ : 4.5V~5.5V
- Up to 0.25 $\mu\text{s}$  instruction cycle with 16MHz system clock at  $V_{DD}=5\text{V}$
- Power down and wake-up functions to reduce power consumption
- Four Oscillators:
  - External High Frequency Crystal – HXT
  - External 32.768kHz Crystal-- LXT
  - Internal High Frequency RC -- HIRC
  - Internal 32kHz -- LIRC
- Fully intergrated internal 8MHz oscillator requires no external components
- Multi-mode operation: NORMAL, SLOW, IDLE and SLEEP
- All instructions executed in one to three instruction cycles
- Table read instructions
- 115 powerful instructions
- 8-level subroutine nesting
- Bit manipulation instruction

### Peripheral Features

- Flash Program Memory: 8K $\times$ 16 (HT66F489)/4K $\times$ 16 (HT66F488)
- RAM Data Memory: 384 $\times$ 8
- True EEPROM Memory: 64 $\times$ 8
- Watchdog Timer function
- 30 bidirectional I/O lines
- LED driver
- 1/3 Bias Software LCD
- Serial Interface Module – SIM for SPI or I<sup>2</sup>C
- UART Function
- Six pin-shared external interrupts
- Multiple Timer Module for time measure, input capture, compare match output, PWM output function or single pulse output function
- Dual Time-Base functions for generation of fixed time interrupt signals
- 8-channel 12-bit resolution A/D converter
- Low voltage reset function
- Low voltage detector function
- Package: 28-pin SOP/SSOP

## General Description

The devices are 8-bit high performance RISC architecture microcontroller, designed especially for applications that interface directly to analog signals, such as those from sensor. Offering users the convenience of Flash Memory multi-programming features, this device also includes a wide range of functions and features. Other memory includes an area of RAM Data Memory as well as an area of true EEPROM memory for storage of non-volatile data such as serial numbers, calibration data etc.

Analog features include a multi-channel 12-bit A/D converter. Multiple and extremely flexible Timer Modules provide timing, pulse generation and PWM generation functions. Communication with the outside world catered for by including fully integrated SPI, I<sup>2</sup>C and UART interface functions, three popular interfaces which provide designers with a means of easy communication with external peripheral hardware. Protective features such as an internal Watchdog Timer, Low Voltage Reset and Low Voltage Detector coupled with excellent noise immunity and ESD protection ensure that reliable operation is maintained in hostile electrical environments.

A full choice of HXT, LXT, HIRC and LIRC oscillator functions are provided including a fully integrated system oscillator which requires no external components for its implementation. The ability to operate and switch dynamically between a range of operating modes using different clock sources gives users the ability to optimise microcontroller operation and minimize power consumption.

The inclusion of flexible I/O programming features, Time-Base functions along with many other features ensure that the device will find excellent use in applications such as sensor signal processing, chargers, motor driving, industrial control, consumer product, subsystem controller, etc.

## Selection Table

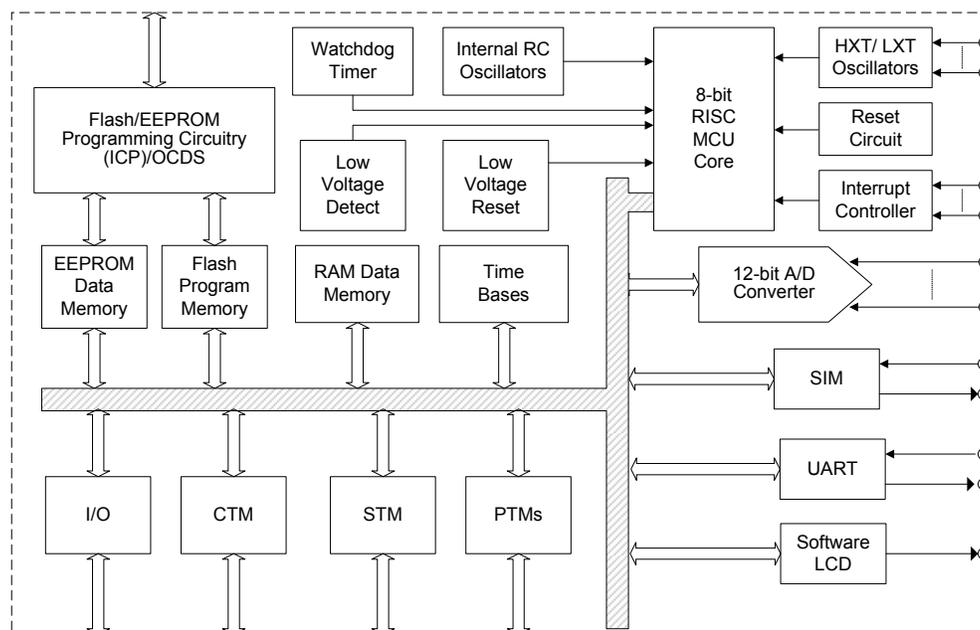
Most features are common to all devices, the main feature distinguishing them are Memory capacity and whether EEPROM or not. The following table summarises the main features of each device.

Part No.	Program Memory	Data Memory	Data EEPROM	I/O	R-type LCD (1/3 Bias)	High Current LED Output
HT66F488	4K×16	384×8	64×8	26	6(SCOM/SSEG) +20SSEG	26
HT66F489	8K×16					

Part No.	Timer Module	RTC	A/D	Interface	Stack	Package
HT66F488	10-bit CTM×1 16-bit STM×1	√	12-bit×8	SPI/I <sup>2</sup> C, UART	8	28SOP/SSOP
HT66F489	10-bit PTM×2					

## Block Diagram



## Pin Assignment

PC4/AN4/SSEG22	1	28	PC5/AN5/INT3/SSEG21
PC3/AN3/SSEG23	2	27	PC6/AN6/VREF/INT2/SSEG20
PC0/AN2/SSEG26	3	26	PB1/PTCK0/INT1/AN7/SSEG17
PB7/STP/RX/AN1/SSEG27	4	25	PB2/CTP/INT0/SSEG16
PB6/STCK/TX/AN0/SSEG28	5	24	PB3/OSC2/XT2/SSEG15
PB5/SSEG29	6	23	PB4/OSC1/XT1/ICPCK/SSEG14
VSS/AVSS	7	22	VDD/AVDD
PD0/ICPDA/SCOM0/SSEG0	8	21	PA7/CTCK/SSEG13
PD1/SCOM1/SSEG1	9	20	PA6/PTCK1/SCS/SSEG12
PD2/SCOM2/SSEG2	10	19	PA5/PTP1/SCK/SCL/SSEG11
PD3/SCOM3/SSEG3	11	18	PA4/SDI/SDA/SSEG10
PD4/SCOM4/SSEG4	12	17	PA3/SDO/SSEG9
PD5/SCOM5/SSEG5	13	16	PA2/OCDSCK/SSEG8
PA0/OCSDSA/SSEG6	14	15	PA1/SSEG7

**HT66F488/489**  
**HT66V488/489**  
**28 SOP-A/SSOP-A**

- Note: 1. If the pin-shared pin functions have multiple outputs simultaneously, the function used is determined by the corresponding software control bit or configuration option.  
 2. AVDD&VDD means the VDD and AVDD are the double bonding. VSS&AVSS means the VSS and AVSS are the double bonding.  
 3. HT66V488/489 is the EV chip, and the OCSDSA, OCDSCK pins are for EV chip only.

## Pin Descriptions

Pin Name	Function	OPT	I/T	O/T	Description
PA0/OCDSDA/ SSEG6	PA0	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	OCDSDA	—	ST	CMOS	OCDS Data/Address pin, for EV chip only
	SSEG6	SLCDC2	—	CMOS	Software LCD SEG output
PA1/SSEG7	PA1	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	SSEG7	SLCDC2	—	CMOS	Software LCD SEG output
PA2/OCDSCK/ SSEG8	PA2	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	OCDSCK	—	ST	—	OCDS Clock pin, for EV chip only.
	SSEG8	SLCDC2	—	CMOS	Software LCD SEG output
PA3/SDO/SSEG9	PA3	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	SDO	—	—	CMOS	SPI Data output
	SSEG9	SLCDC2	—	CMOS	Software LCD SEG output
PA4/SDI/SDA/ SSEG10	PA4	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	SDI	—	ST	—	SPI Data input
	SDA	—	ST	NMOS	I <sup>2</sup> C Data
	SSEG10	SLCDC2	—	CMOS	Software LCD SEG output
PA5/PTP1/SCK/ SCL/SSEG11	PA5	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	PTP1	TMPC0	ST	CMOS	PTM1 input/output
	SCK	—	ST	CMOS	SPI Serial Clock
	SCL	—	ST	NMOS	I <sup>2</sup> C Clock
	SSEG11	SLCDC2	—	CMOS	Software LCD SEG output
PA6/PTCK1/ $\overline{\text{SCS}}$ / SSEG12	PA6	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	PTCK1	—	ST	—	PTM1 input
	$\overline{\text{SCS}}$	—	ST	CMOS	SPI slaver selection
	SSEG12	SLCDC2	—	CMOS	Software LCD SEG output
PA7/CTCK/ SSEG13	PA7	PAWU PAPU	ST	CMOS	General purpose I/O. Register enable pull-up and wake-up
	CTCK	—	ST	—	CTM input
	SSEG13	SLCDC2	—	CMOS	Software LCD SEG output
PB0/PTP0/INT4/ SSEG18	PB0	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	PTP0	TMPC0	ST	CMOS	PTM0 output
	INT4	—	ST	—	External interrupt input
	SSEG18	SLCDC3	—	CMOS	Software LCD SEG output
PB1/PTCK0/ INT1/AN7/ SSEG17	PB1	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	PTCK0	—	ST	—	PTM0 input
	INT1	—	ST	—	External interrupt input
	AN7	ACERL	AN	—	ADC input
	SSEG17	SLCDC3	—	CMOS	Software LCD SEG output
PB2/CTP/INT0/ SSEG16	PB2	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	CTP	TMPC0	ST	CMOS	CTM output
	INT0	—	ST	—	External interrupt input
	SSEG16	SLCDC3	—	CMOS	Software LCD SEG output

Pin Name	Function	OPT	I/T	O/T	Description
PB3/OSC2/XT2/ SSEG15	PB3	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	OSC2	CO	—	HXT	HXT pin
	XT2	CO	—	LXT	LXT pin
	SSEG15	SLCDC3	—	CMOS	Software LCD SEG output
PB4/OSC1/ XT1/ICPCK/ SSEG14	PB4	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	OSC1	CO	HXT	—	HXT pin
	XT1	CO	LXT	—	LXT pin
	ICPCK	—	ST	—	ICP Clock pin
PB5/SSEG29	PB5	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	SSEG29	SLCDC4	—	CMOS	Software LCD SEG output
PB6/STCK/TX/ AN0/SSEG28	PB6	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	STCK	—	ST	—	STM input
	TX	—	—	CMOS	UART data transmission output
	AN0	ACERL	AN	—	ADC input
PB7/STP/RX/ AN1/SSEG27	PB7	PBPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	STP	TMPC0	ST	CMOS	STM input/output
	RX	—	ST	—	UART data received input
	AN1	ACERL	AN	—	ADC input
PC0/AN2/ SSEG26	PC0	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	AN2	ACERL	AN	—	ADC input
	SSEG26	SLCDC4	—	CMOS	Software LCD SEG output
PC1/SSEG25	PC1	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	SSEG25	SLCDC4	—	CMOS	Software LCD SEG output
PC2/SSEG24	PC2	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	SSEG24	SLCDC4	—	CMOS	Software LCD SEG output
PC3/AN3/ SSEG23	PC3	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	AN3	ACERL	AN	—	ADC input
	SSEG23	SLCDC4	—	CMOS	Software LCD SEG output
PC4/AN4/ SSEG22	PC4	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	AN4	ACERL	AN	—	ADC input
	SSEG22	SLCDC4	—	CMOS	Software LCD SEG output
PC5/AN5/INT3/ SSEG21	PC5	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	AN5	ACERL	AN	—	ADC input
	INT3	—	ST	—	External interrupt input
	SSEG21	SLCDC3	—	CMOS	Software LCD SEG output
PC6/AN6/VREF/ INT2/SSEG20	PC6	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	AN6	ACERL	AN	—	ADC input
	VREF	ADCR1	AN	—	ADC Reference voltage input
	INT2	—	ST	—	External interrupt input
	SSEG20	SLCDC3	—	CMOS	Software LCD SEG output
PC7/INT5/ SSEG19	PC7	PCPU	ST	CMOS	General purpose I/O. Register enable pull-up.
	INT5	—	ST	—	External interrupt input
	SSEG19	SLCDC3	—	CMOS	Software LCD SEG output

Pin Name	Function	OPT	I/T	O/T	Description
PD0/ICPDA/ SCOM0/SSEG0	PD0	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	ICPDA	—	ST	CMOS	ICP Data/Address pin
	SCOM0	SLCDC0 SLCDC1	—	CMOS	Software LCD COM output
	SSEG0	SLCDC0 SLCDC1	—	CMOS	Software LCD SEG output
PD1/SCOM1/ SSEG1	PD1	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	SCOM1	SLCDC0 SLCDC1	—	CMOS	Software LCD COM output
	SSEG1	SLCDC0 SLCDC1	—	CMOS	Software LCD SEG output
PD2/SCOM2/ SSEG2	PD2	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	SCOM2	SLCDC0 SLCDC1	—	CMOS	Software LCD COM output
	SSEG2	SLCDC0 SLCDC1	—	CMOS	Software LCD SEG output
PD3/SCOM3/ SSEG3	PD3	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	SCOM3	SLCDC0 SLCDC1	—	CMOS	Software LCD COM output
	SSEG3	SLCDC0 SLCDC1	—	CMOS	Software LCD SEG output
PD4/SCOM4/ SSEG4	PD4	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	SCOM4	SLCDC1	—	CMOS	Software LCD COM output
	SSEG4	SLCDC1	—	CMOS	Software LCD SEG output
PD5/SCOM5/ SSEG5	PD5	PDPUP	ST	CMOS	General purpose I/O. Register enable pull-up.
	SCOM5	SLCDC1	—	CMOS	Software LCD COM output
	SSEG5	SLCDC1	—	CMOS	Software LCD SEG output
VDD/AVDD	VDD	—	PWR	—	Positive Power Supply
	AVDD	—	PWR	—	Analog Positive Power Supply
VSS/AVSS	VSS	—	PWR	—	Negative Power Supply. Ground
	AVSS	—	PWR	—	Analog Negative Power Supply

Note: I/T: Input type;

O/T: Output type

OP: Optional by configuration option (CO) or register option

PWR: Power;

CO: Configuration option; ST: Schmitt Trigger input

CMOS: CMOS output;

AN: Analog input pin

HXT: High frequency crystal oscillator

LXT: Low frequency crystal oscillator

\*: VDD is the device power supply while AVDD is the ADC power supply. The AVDD pin is bonded together internally with VDD.

\*\* : VSS is the device ground pin while AVSS is the ADC ground pin. The AVSS pin is bonded together internally with VSS.

Note: 1. This Pin Description Summary table describes all resources which may not be connected to external pins.

2. For 28-SKDIP/SOP package PB0, PC1, PC2 and PC7 are not bonded to external pins and should be properly configured to avoid a floating condition which will result in additional current leakage.

## Absolute Maximum Ratings

Supply Voltage .....	$V_{SS} = -0.3V$ to $V_{SS} + 6.0V$
Input Voltage .....	$V_{SS} = -0.3V$ to $V_{DD} + 0.3V$
Storage Temperature .....	$-50^{\circ}C$ to $125^{\circ}C$
Operating Temperature .....	$-40^{\circ}C$ to $85^{\circ}C$
$I_{OH}$ Total .....	$-150mA$
$I_{OL}$ Total .....	$100mA$
Total Power Dissipation .....	$500mW$

Note: These are stress ratings only. Stresses exceeding the range specified under “Absolute Maximum Ratings” may cause substantial damage to these devices. Functional operation of these devices at other conditions beyond those listed in the specification is not implied and prolonged exposure to extreme conditions may affect devices reliability.

## D.C. Characteristics

$T_a = 25^{\circ}C$

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		$V_{DD}$	Conditions				
$V_{DD}$	Operating Voltage (HXT)	—	$f_{SYS} = 8MHz$	2.2	—	5.5	V
			$f_{SYS} = 10MHz$	2.7	—	5.5	V
			$f_{SYS} = 12MHz$	3.3	—	5.5	V
			$f_{SYS} = 16MHz$	4.5	—	5.5	V
	Operating Voltage (HIRC)	—	$f_{SYS} = 8MHz$	2.2	—	5.5	V
$I_{DD}$	Operating Current, Normal Mode, $f_{SYS} = f_H$ (HXT), $f_S = f_{LIRC}$	3V	No load, $f_{SYS} = f_H = 8MHz$	—	1.0	1.5	mA
		5V	ADC off, WDT enable	—	2.5	4.0	mA
		3V	No load, $f_{SYS} = f_H = 10MHz$	—	1.2	2.0	mA
		5V	ADC off, WDT enable	—	2.8	4.5	mA
		5V	No load, $f_{SYS} = f_H = 12MHz$ ADC off, WDT enable	—	3.5	5.5	mA
		5V	No load, $f_{SYS} = f_H = 16MHz$ ADC off, WDT enable	—	4.5	7.0	mA
	Operating Current, Normal Mode, $f_{SYS} = f_H$ (HIRC), $f_S = f_{LXT}$ or $f_{LIRC}$	3V	No load, $f_{SYS} = f_H = 8MHz$	—	1.2	2.0	mA
		5V	ADC off, WDT enable	—	2.8	4.5	mA
	Operating Current, Slow Mode 1 $f_H = 12MHz$ (HXT), $f_S = f_{LIRC}$	5V	No load, $f_{SYS} = f_H/2$ , ADC off, WDT enable	—	2.1	3.3	mA
		5V	No load, $f_{SYS} = f_H/4$ , ADC off, WDT enable	—	1.6	2.5	mA
		5V	No load, $f_{SYS} = f_H/8$ , ADC off, WDT enable	—	1.2	2.5	mA
		5V	No load, $f_{SYS} = f_H/16$ ADC off, WDT enable	—	1.1	2.0	mA
		5V	No load, $f_{SYS} = f_H/32$ , ADC off, WDT enable	—	1.0	2.0	mA
		5V	No load, $f_{SYS} = f_H/64$ , ADC off, WDT enable	—	1.0	2.0	mA

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
I <sub>DD</sub>	Operating Current, Slow Mode 0 f <sub>sys</sub> =f <sub>s</sub> =f <sub>LXT</sub>	3V	No load, ADC off, WDT enable, LXTLP=0	—	15	25	μA
		5V	WDT enable, LXTLP=0	—	30	50	μA
		3V	No load, ADC off, WDT enable, LXTLP=1	—	10	20	μA
		5V	WDT enable, LXTLP=1	—	25	45	μA
	Operating Current, Slow Mode 0 f <sub>sys</sub> =f <sub>s</sub> =f <sub>LIRC</sub>	3V	No load, ADC off, WDT enable	—	10	20	μA
		5V	WDT enable	—	20	40	μA
I <sub>STB</sub>	Standby Current, IDLE Mode (f <sub>sys</sub> =off, f <sub>s</sub> =f <sub>SUB</sub> = f <sub>LIRC</sub> )	3V	No load, system HALT, IDLEN=1, WDT enable	—	1.3	3.0	μA
		5V	IDLEN=1, WDT enable	—	2.2	5.0	μA
	Standby Current, SLEEP Mode (f <sub>sys</sub> =off, f <sub>s</sub> =f <sub>SUB</sub> = f <sub>LIRC</sub> )	3V	No load, system HALT, IDLEN=0, WDT enable	—	1.3	3.0	μA
		5V	IDLEN=0, WDT enable	—	2.2	5.0	μA
	Standby Current, SLEEP Mode f <sub>sys</sub> =12MHz (HXT), f <sub>sys</sub> =off, f <sub>s</sub> =f <sub>SUB</sub> = f <sub>LIRC</sub>	—	No load, system HALT, ADC off, WDT enable	—	65	95	μA
V <sub>IL</sub>	Input Low Voltage for I/O Ports	5V	—	0	—	1.5	V
		—	—	0	—	0.2V <sub>DD</sub>	V
V <sub>IH</sub>	Input High Voltage for I/O Ports	5V	—	3.5	—	5	V
		—	—	0.8V <sub>DD</sub>	—	V <sub>DD</sub>	V
I <sub>OL</sub>	I/O Port Sink Current Large Sink I/O (PD)	5V	V <sub>OL</sub> =1.5V	120	200	—	mA
		3V	V <sub>OH</sub> =0.1V <sub>DD</sub>	15	30	—	mA
		5V	V <sub>OH</sub> =0.1V <sub>DD</sub>	30	60	—	mA
	I/O Port Sink Current (PA, PB, PC)	3V	V <sub>OL</sub> =0.1V <sub>DD</sub>	15	30	—	mA
		5V	V <sub>OL</sub> =0.1V <sub>DD</sub>	30	60	—	mA
I <sub>OH</sub>	I/O Port Source Current (PA, PB, PC, PD)	3V	V <sub>OH</sub> =0.9V <sub>DD</sub> , select full source current option	-3.0	-6.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 10/22 full source current option	-1.5	-3.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 7/22 full source current option	-1.0	-2.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 4/22 full source current option	-0.8	-1.5	—	mA
		5V	V <sub>OH</sub> =0.9V <sub>DD</sub> , select full source current option	-9.0	-18.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 10/22 full source current option	-4.0	-8.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 7/22 full source current option	-3.0	-6.0	—	mA
			V <sub>OH</sub> =0.9V <sub>DD</sub> , select 4/22 full source current option	-2.0	-3.6	—	mA
R <sub>PH</sub>	Pull-high Resistance	3V	—	20	60	100	kΩ
		5V	—	10	30	50	kΩ
I <sub>BIAS</sub>	Bias current for LCD	5V	ISEL[1:0]=00	4.2	8.3	13	μA
			ISEL[1:0]=01	8.3	16.7	25	μA
			ISEL[1:0]=10	25	50	75	μA
			ISEL[1:0]=11	50	100	150	μA
V <sub>SEG_H</sub>	1/3 bias LCD SEG output (SEG_H)	2.2~5.5V	No load	0.645	0.67	0.695	V <sub>DD</sub>
V <sub>SEG_L</sub>	1/3 bias LCD SEG output (SEG_L)	2.2~5.5V	No load	0.305	0.33	0.355	V <sub>DD</sub>

## A.C. Characteristics

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
f <sub>sys</sub>	System clock (HXT)	2.2V~5.5V	—	0.4	—	8	MHz
		2.7V~5.5V		0.4	—	10	MHz
		3.3V~5.5V		0.4	—	12	MHz
		4.5V~5.5V		0.4	—	16	MHz
	System clock (HIRC)	5V	Ta=0°C~70°C	-2%	8	2%	MHz
	System clock (LXT)	—	—	—	32768	—	Hz
System clock (LIRC)	5V	Ta=25°C	-10%	32	+10%	kHz	
	2.2V~5.5V	Ta=-40°C~85°C	-50%	32	+60%	kHz	
t <sub>TIMER</sub>	xTCKn Input Pin Minimum Pulse width	—	—	0.06	0.15	0.3	µs
t <sub>SST</sub>	System Start-up Timer Period (wake up from HALT)	—	f <sub>sys</sub> =HXT	128	—	—	t <sub>sys</sub>
			f <sub>sys</sub> =LXT	128	—	—	t <sub>sys</sub>
			f <sub>sys</sub> =HIRC	16	—	—	t <sub>sys</sub>
			f <sub>sys</sub> =LIRC	2	—	—	t <sub>sys</sub>
System Start-up Timer Period (wake up from HALT, f <sub>sys</sub> on at HALT)	—	—	2	—	—	t <sub>sys</sub>	
t <sub>RSTD</sub>	System Reset Delay Time (POR, LVR, LVR S/W, WDT S/W reset)	—	—	25	50	100	ms
	System Reset Delay Time (WDT reset)	—	—	8.3	16.7	33.3	ms
t <sub>INT</sub>	Interrupt Minimum Pulse Width	—	—	1	3.3	5	µs
t <sub>LVR</sub>	Low Voltage width to reset	—	—	120	240	480	µs
t <sub>SRESET</sub>	Software Reset width to reset	—	—	45	90	120	µs
t <sub>EERD</sub>	EEPROM Read Time	—	—	1	2	4	t <sub>sys</sub>
t <sub>EEWR</sub>	EEPROM Write Time	—	—	1	2	4	ms

Note: t<sub>sys</sub>=1/f<sub>sys</sub>

**LVR & LVD Characteristics**

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>LVR</sub>	Low Voltage Reset Voltage	—	LVR Enable, 2.10V option	-5%	2.10	+5%	V
			LVR Enable, 2.55V option		2.55		V
			LVR Enable, 3.15V option		3.15		V
			LVR Enable, 3.80V option		3.80		V
V <sub>LVD</sub>	Low Voltage Detector Voltage	—	LVDEN = 1, V <sub>LVD</sub> = 2.0V	-5%	2.0	+5%	V
			LVDEN = 1, V <sub>LVD</sub> = 2.2V		2.2		V
			LVDEN = 1, V <sub>LVD</sub> = 2.4V		2.4		V
			LVDEN = 1, V <sub>LVD</sub> = 2.7V		2.7		V
			LVDEN = 1, V <sub>LVD</sub> = 3.0V		3.0		V
			LVDEN = 1, V <sub>LVD</sub> = 3.3V		3.3		V
			LVDEN = 1, V <sub>LVD</sub> = 3.6V		3.6		V
			LVDEN = 1, V <sub>LVD</sub> = 4.0V		4.0		V
I <sub>LVR</sub>	Additional Power Consumption if LVR is Used	3V	LVR disable→LVR enable	—	30	45	μA
		5V		—	60	90	μA
I <sub>LVD</sub>	Additional Power Consumption if LVD is Used	3V	LVD disable→LVD enable (LVR disable)	—	40	60	μA
		5V		—	75	115	μA
		3V	LVD disable→LVD enable (LVR enable)	—	30	45	μA
		5V		—	60	90	μA
t <sub>LVR</sub>	Low Voltage Width to Reset	—	—	120	240	480	μs
t <sub>LVD</sub>	Low Voltage Width to Interrupt	—	—	65	125	250	μs
t <sub>LVDs</sub>	LVDO stable time	—	LVD disable→LVD enable (LVR enable)	—	—	15	μs
		—	LVD disable→LVD enable (LVR disable)	—	—	200	μs

## A/D Converter Electrical Characteristics

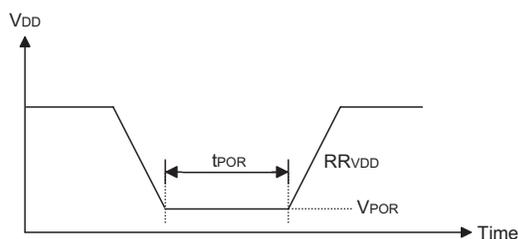
Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
AV <sub>DD</sub>	A/D Converter Operating Voltage	—	—	2.2	—	5.5	V
V <sub>AD</sub>	ADC Input Voltage	—	—	0	—	AV <sub>DD</sub> /V <sub>REF</sub>	V
V <sub>REF</sub>	ADC Reference Voltage	—	—	2	—	AV <sub>DD</sub> +0.1	V
t <sub>ADC</sub>	A/D Conversion Time	—	12-bit ADC	—	16	—	t <sub>AD</sub>
t <sub>AD</sub>	A/D Clock Period	2.2V~2.7V	—	8	—	10	μs
		2.7V~5.5V	—	0.5	—	10	μs
t <sub>ON2ST</sub>	A/D Converter On-to-Start Time	—	—	2	—	—	μs
DNL1	A/D Differential Non-linearity	2.2V~2.7V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =8μs Ta=25°C	—	±15	—	LSB
		2.7V~5.5V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =0.5μs Ta=25°C	-3	—	+3	LSB
DNL2	A/D Differential Non-linearity	3V/5V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =0.5μs Ta=-40°C ~ 85°C	-4	—	+4	LSB
INL1	A/D Integral Non-linearity	2.2V~2.7V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =8μs Ta=25°C	—	±16	—	LSB
		2.7V~5.5V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =0.5μs Ta=25°C	-4	—	+4	LSB
INL2	A/D Integral Non-linearity	3V/5V	V <sub>REF</sub> =V <sub>DD</sub> , t <sub>AD</sub> =0.5μs Ta=-40°C ~ 85°C	-6	—	+6	LSB
I <sub>ADC</sub>	Additional Power consumption if A/D Converter is used	3V	No load, t <sub>AD</sub> =0.5μs	—	1.0	2.0	mA
		5V		—	1.5	3.0	mA
t <sub>BGS</sub>	V <sub>BG</sub> turn on stable time	—	—	10	—	—	ms
V <sub>BG</sub>	Bandgap reference with buffer voltage	—	—	-5%	1.25	+5%	V
I <sub>BG</sub>	Additional power consumption if reference with buffer is used	—	—	—	200	300	μA

## Power on Reset Electrical Characteristics

Ta = 25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>POR</sub>	V <sub>DD</sub> Start Voltage to Ensure Power-on Reset	—	—	—	—	100	mV
RR <sub>VDD</sub>	V <sub>DD</sub> Rising Rate to Ensure Power-on Reset	—	—	0.035	—	—	V/ms
t <sub>POR</sub>	Minimum time for V <sub>DD</sub> Stays at V <sub>POR</sub> to ensure Power-on Reset	—	—	10	—	—	ms

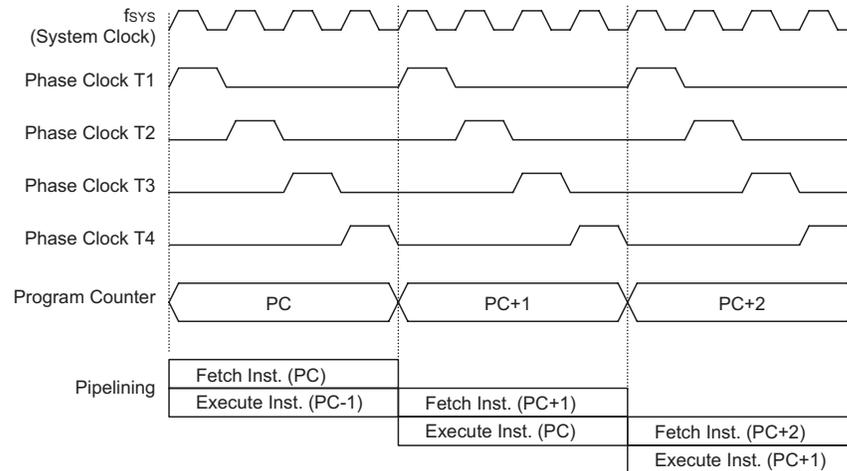


## System Architecture

A key factor in the high-performance features of the Holtek range of microcontrollers is attributed to their internal system architecture. The device takes advantage of the usual features found within RISC microcontrollers providing increased speed of operation and Periodic performance. The pipelining scheme is implemented in such a way that instruction fetching and instruction execution are overlapped, hence instructions are effectively executed in one cycle, with the exception of branch or call instructions. An 8-bit wide ALU is used in practically all instruction set operations, which carries out arithmetic operations, logic operations, rotation, increment, decrement, branch decisions, etc. The internal data path is simplified by moving data through the Accumulator and the ALU. Certain internal registers are implemented in the Data Memory and can be directly or indirectly addressed. The simple addressing methods of these registers along with additional architectural features ensure that a minimum of external components is required to provide a functional I/O and A/D control system with maximum reliability and flexibility. This makes the device suitable for low-cost, high-volume production for controller applications.

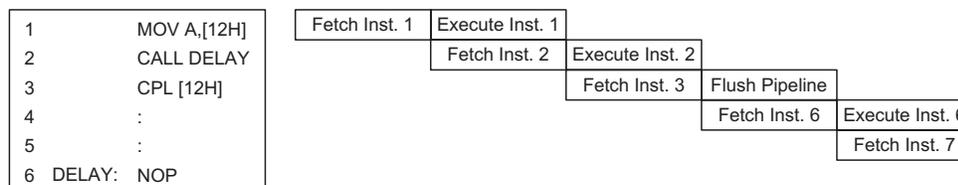
### Clocking and Pipelining

The main system clock, derived from either a HXT, LXT, HIRC or LIRC oscillator is subdivided into four internally generated non-overlapping clocks, T1~T4. The Program Counter is incremented at the beginning of the T1 clock during which time a new instruction is fetched. The remaining T2~T4 clocks carry out the decoding and execution functions. In this way, one T1~T4 clock cycle forms one instruction cycle. Although the fetching and execution of instructions takes place in consecutive instruction cycles, the pipelining structure of the microcontroller ensures that instructions are effectively executed in one instruction cycle. The exception to this are instructions where the contents of the Program Counter are changed, such as subroutine calls or jumps, in which case the instruction will take one more instruction cycle to execute.



**System Clock and Pipelining**

For instructions involving branches, such as jump or call instructions, two machine cycles are required to complete instruction execution. An extra cycle is required as the program takes one cycle to first obtain the actual jump or call address and then another cycle to actually execute the branch. The requirement for this extra cycle should be taken into account by programmers in timing sensitive applications.



**Instruction Fetching**

### Program Counter

During program execution, the Program Counter is used to keep track of the address of the next instruction to be executed. It is automatically incremented by one each time an instruction is executed except for instructions, such as “JMP” or “CALL” that demand a jump to a non-consecutive Program Memory address. Only the lower 8 bits, known as the Program Counter Low Register, are directly addressable by the application program.

When executing instructions requiring jumps to non-consecutive addresses such as a jump instruction, a subroutine call, interrupt or reset, etc., the microcontroller manages program control by loading the required address into the Program Counter. For conditional skip instructions, once the condition has been met, the next instruction, which has already been fetched during the present instruction execution, is discarded and a dummy cycle takes its place while the correct instruction is obtained.

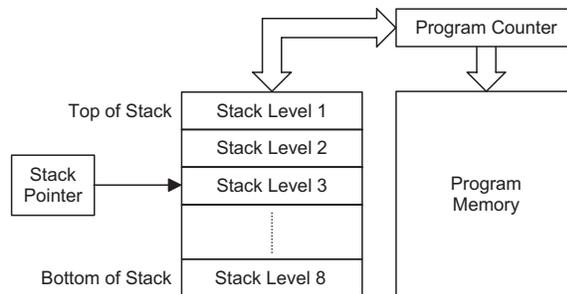
Device	Program Counter	
	Program Counter High byte	PCL Register
HT66F488	PC11~PC8	PCL7~PCL0
HT66F489	PC12~PC8	PCL7~PCL0

The lower byte of the Program Counter, known as the Program Counter Low register or PCL, is available for program control and is a readable and writable register. By transferring data directly into this register, a short program jump can be executed directly, however, as only this low byte is available for manipulation, the jumps are limited to the present page of memory, that is 256 locations. When such program jumps are executed it should also be noted that a dummy cycle will be inserted. Manipulating the PCL register may cause program branching, so an extra cycle is needed to pre-fetch.

## Stack

This is a special part of the memory which is used to save the contents of the Program Counter only. The stack is neither part of the data nor part of the program space, and is neither readable nor writeable. The activated level is indexed by the Stack Pointer, and is neither readable nor writeable. At a subroutine call or interrupt acknowledge signal, the contents of the Program Counter are pushed onto the stack. At the end of a subroutine or an interrupt routine, signaled by a return instruction, RET or RETI, the Program Counter is restored to its previous value from the stack. After a device reset, the Stack Pointer will point to the top of the stack.

If the stack is full and an enabled interrupt takes place, the interrupt request flag will be recorded but the acknowledge signal will be inhibited. When the Stack Pointer is decremented, by RET or RETI, the interrupt will be serviced. This feature prevents stack overflow allowing the programmer to use the structure more easily. However, when the stack is full, a CALL subroutine instruction can still be executed which will result in a stack overflow. Precautions should be taken to avoid such cases which might cause unpredictable program branching. If the stack is overflow, the first Program Counter save in the stack will be lost.



## Arithmetic and Logic Unit – ALU

The arithmetic-logic unit or ALU is a critical area of the microcontroller that carries out arithmetic and logic operations of the instruction set. Connected to the main microcontroller data bus, the ALU receives related instruction codes and performs the required arithmetic or logical operations after which the result will be placed in the specified register. As these ALU calculation or operations may result in carry, borrow or other status changes, the status register will be correspondingly updated to reflect these changes. The ALU supports the following functions:

- Arithmetic operations: ADD, ADDM, ADC, ADCM, SUB, SUBM, SBC, SBCM, DAA, LADD, LADDM, LADC, LADCM, LSUB, LSUBM, LSBC, LSBCM, LDAA
- Logic operations: AND, OR, XOR, ANDM, ORM, XORM, CPL, CPLA, LAND, LANDM, LOR, LORM, LXOR, LXORM, LCPL, LCPLA
- Rotation: RRA, RR, RRCA, RRC, RLA, RL, RLCA, RLC, LRR, LRRCA, LRRCA, LRLA, LRLCA, LRLC
- Increment and Decrement: INCA, INC, DECA, DEC, LINCA, LINC, LDECA, LDEC
- Branch decision: JMP, SZ, SZA, SNZ, SIZ, SDZ, SIZA, SDZA, CALL, RET, RETI, LSNZ, LSZ, LSZA, LSIZ, LSDZ, LSIZA, LSDZA

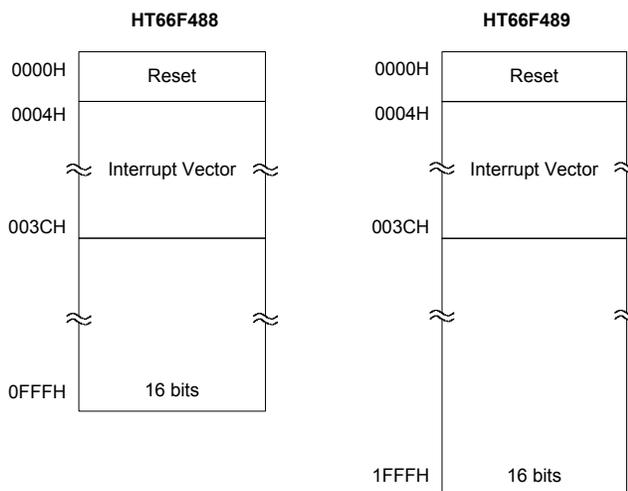
## Flash Program Memory

The Program Memory is the location where the user code or program is stored. For this device the Program Memory is Flash type, which means it can be programmed and re-programmed a large number of times, allowing the user the convenience of code modification on the same device. By using the appropriate programming tools, this Flash device offers users the flexibility to conveniently debug and develop their applications while also offering a means of field programming and updating.

### Structure

The Program Memory has a capacity of 4K×16 bits to 8K×16 bits. The Program Memory is addressed by the Program Counter and also contains data, table information and interrupt entries. Table data, which can be setup in any location within the Program Memory, is addressed by a separate table pointer register.

Device	Capacity
HT66F488	4K×16
HT66F489	8K×16



**Program Memory Structure**

### Special Vectors

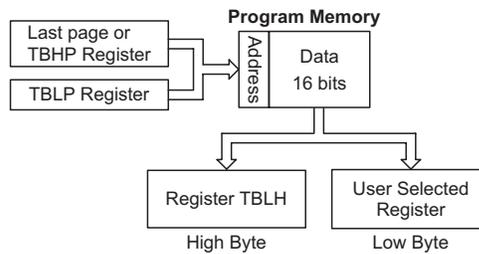
Within the Program Memory, certain locations are reserved for the reset and interrupts. The location 000H is reserved for use by the device reset for program initialisation. After a device reset is initiated, the program will jump to this location and begin execution.

## Look-up Table

Any location within the Program Memory can be defined as a look-up table where programmers can store fixed data. To use the look-up table, the table pointer must first be setup by placing the address of the look up data to be retrieved in the table pointer register, TBLP and TBHP. These registers define the total address of the look-up table.

After setting up the table pointer, the table data can be retrieved from the Program Memory using the “TABRD [m]” or “TABRDL [m]” instructions respectively when the memory [m] is located in sector 0. If the memory [m] is located in other sectors, the table data can be retrieved from the Program Memory using the “LTABRD [m]” or “LTABRDL [m]” instructions respectively. When the instruction is executed, the lower order table byte from the Program Memory will be transferred to the user defined Data Memory register [m] as specified in the instruction. The higher order table data byte from the Program Memory will be transferred to the TBLH special register. Any unused bits in this transferred higher order byte will be read as 0.

The accompanying diagram illustrates the addressing data flow of the look-up table.



## Table Program Example

The following example shows how the table pointer and table data is defined and retrieved from the microcontroller. This example uses raw table data located in the Program Memory which is stored there using the ORG statement. The value at this ORG statement is “F00H” which refers to the start address of the last page within the 4K words Program Memory of the HT66F488 device. The table pointer is setup here to have an initial value of “06H”. This will ensure that the first data read from the data table will be at the Program Memory address “F06H” or 6 locations after the start of the present page. Note that the value for the table pointer is referenced to the first address of the last page if the “TABRDL [m]” or “LTABRDL [m]” instruction is being used. The high byte of the table data which in this case is equal to zero will be transferred to the TBLH register automatically when the “TABRDL [m]” or “LTABRDL [m]” instruction is executed.

Because the TBLH register a read/write register and can be restored, care should be taken to ensure its protection if both the main routine and Interrupt Service Routine use table read instructions. If using the table read instructions, the Interrupt Service Routines may change the value of the TBLH and subsequently cause errors if used again by the main routine. As a rule it is recommended that simultaneous use of the table read instructions should be avoided. However, in situations where simultaneous use cannot be avoided, the interrupts should be disabled prior to the execution of any main routine table-read instructions. Note that all table related instructions require two instruction cycles to complete their operation.

**Table Read Program Example**

```

tempreg1 db ?      ; temporary register #1
tempreg2 db ?      ; temporary register #2
:
:
mov a,06h          ; initialise low table pointer - note that this address is referenced
mov tblp,a         ; to the last page
:
:
tabrdl tempreg1    ; transfers value in table referenced by table pointer data at program
                  ; memory address "F06H" transferred to tempreg1 and TBLH
dec tblp           ; reduce value of table pointer by one
tabrdl tempreg2    ; transfers value in table referenced by table pointer
                  ; data at program memory address "F05H" transferred to
                  ; tempreg2 and TBLH
                  ; in this example the data "1AH" is transferred to tempreg1 and data "0FH"
                  ; to register tempreg2
:
:
org F00h           ; sets initial address of program memory
dc 00Ah, 00Bh, 00Ch, 00Dh, 00Eh, 00Fh, 01Ah, 01Bh
:
:

```

**In Circuit Programming**

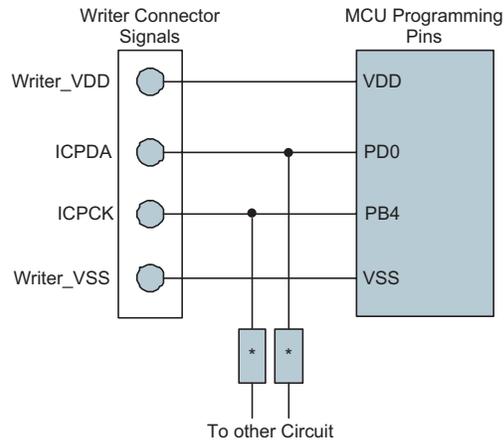
The provision of Flash type Program Memory provides the user with a means of convenient and easy upgrades and modifications to their programs on the same device. As an additional convenience, Holtek has provided a means of programming the microcontroller in-circuit using a 4-pin interface. This provides manufacturers with the possibility of manufacturing their circuit boards complete with a programmed or un-programmed microcontroller, and then programming or upgrading the program at a later stage. This enables product manufacturers to easily keep their manufactured products supplied with the latest program releases without removal and re-insertion of the device.

The Holtek Flash MCU to Writer Programming Pin correspondence table is as follows:

Holtek Write Pins	MCU Programming Pins	Function
ICPDA	PD0	Programming Serial Data
ICPCK	PB4	Programming Serial Clock
VDD	VDD	Power Supply
VSS	VSS	Ground

The Program Memory and EEPROM data memory can both be programmed serially in-circuit using this 4-wire interface. Data is downloaded and uploaded serially on a single pin with an additional line for the clock. Two additional lines are required for the power supply and ground. The technical details regarding the in-circuit programming of the device are beyond the scope of this document and will be supplied in supplementary literature.

During the programming process, the user must take care to ensure that no other outputs are connected to the PD0 and PB4 pins.



Note: \* may be resistor or capacitor. The resistance of \* must be greater than 1k or the capacitance of \* must be less than 1nF.

### On-Chip Debug Support – OCDS

There is an EV chip, HT66V488/489, which is used to emulate the HT66F488/HT66F489 device. This EV chip device also provides an “On-Chip Debug” function to debug the device during the development process. The EV chip and the actual MCU devices are almost functionally compatible except for the “On-Chip Debug” function. Users can use the EV chip device to emulate the real chip device behavior by connecting the OCSDSA and OCDSCK pins to the Holtek HT-IDE development tools. The OCSDSA pin is the OCDS Data/Address input/output pin while the OCDSCK pin is the OCDS clock input pin. When users use the EV chip for debugging, other functions which are shared with the OCSDSA and OCDSCK pins in the actual MCU device will have no effect in the EV chip. For a more detailed OCDS description, refer to the corresponding document named “Holtek e-Link for 8-bit MCU OCDS User’s Guide”.

Holtek e-Link Pins	EV Chip Pins	Pin Description
OCSDSA	OCSDSA	OCDS Data/Address input/output
OCDSCK	OCDSCK	OCDS Clock input
VDD	VDD	Power Supply
GND	VSS	Ground

## RAM Data Memory

The Data Memory is an 8-bit wide RAM internal memory and is the location where temporary information is stored.

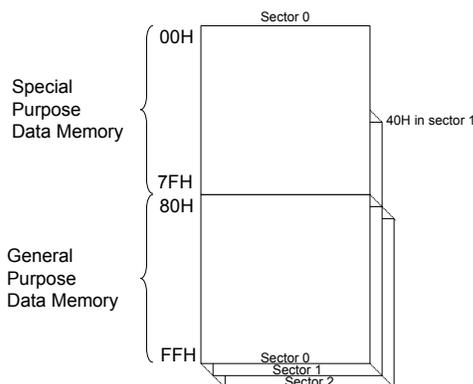
Divided into two types, the first of Data Memory is an area of RAM where special function registers are located. These registers have fixed locations and are necessary for correct operation of the device. Many of these registers can be read from and written to directly under program control, however, some remain protected from user manipulation. The second area of Data Memory is reserved for general purpose use. All locations within this area are read and write accessible under program control.

### Structure

The Data Memory is divided into several sectors, all of which are implemented in 8-bit wide Memory. Each of the Data Memory sectors is categorized into two types, the Special Purpose Data Memory and the General Purpose Data Memory.

The start address of the Special Purpose Data Memory for all devices is the address 00H while the start address of the General Purpose Data Memory is the address 80H. The Special Purpose Data Memory registers are accessible in all sectors, with the exception of the EEC register at address 40H, which is only accessible in Sector 1.

Device	Capacity	Sectors
HT66F488 HT66F489	General Purpose: 384×8	0: 80H~FFH 1: 80H~FFH 2: 80H~FFH



**Data Memory Structure**

### General Purpose Data Memory

There are 384 bytes of general purpose memory which are arranged in 80H~FFH of Sector 0~Sector 2. All microcontroller programs require an area of read/write memory where temporary data can be stored and retrieved for use later. It is this area of RAM memory that is known as General Purpose Data Memory. The general purpose data memory is fully accessible by the user program for both read and writing operations. By using the "SET [m].i" and "CLR [m].i" instructions individual bits can be set or reset under program control giving the user a large range of flexibility for bit manipulation in the Data Memory.

### Special Purpose Data Memory

This area of Data Memory is where registers, necessary for the correct operation of the microcontroller, are stored. They are overlapped in any sector. Most of the registers are both readable and writable but some are protected and are readable only, the details of which are located under the relevant Special Function Register section. Note that for locations that are unused before 80H, any read instruction to these addresses will return the value "00H".

Sector 0~2		Sector 0, 2		Sector 1
00H	IAR0	40H	Unused	EEC
01H	MP0	41H		EEA
02H	IAR1	42H		EED
03H	MP1L	43H		PTMOC0
04H	MP1H	44H		PTMOC1
05H	ACC	45H		PTMODL
06H	PCL	46H		PTMODH
07H	TBLP	47H		PTMOAL
08H	TBLH	48H		PTMOAH
09H	TBHP	49H		PTMORPL
0AH	STATUS	4AH		PTMORPH
0BH	SMOD	4BH		PTM1C0
0CH	IAR2	4CH		PTM1C1
0DH	MP2L	4DH		PTM1DL
0EH	MP2H	4EH		PTM1DH
0FH	Unused	4FH		PTM1AL
10H	INTC0	50H		PTM1AH
11H	INTC1	51H		PTM1RPL
12H	INTC2	52H		PTM1RPH
13H	INTC3	53H		INTEG0
14H	MF10	54H		INTEG1
15H	MF11	55H		WDTC
16H	MF12	56H		TBC
17H	Unused	57H		CTRL
18H	PAWU	58H		LVRC
19H	PAPU	59H		LVDC
1AH	PA	5AH	Unused	
1BH	PAC	5BH		USR
1CH	PBPU	5CH		UCR1
1DH	PB	5DH		UCR2
1EH	PBC	5EH		BRG
1FH	PCPU	5FH		TXR/RXR
20H	PC	60H		SLCDC0
21H	PCC	61H		SLCDC1
22H	PDPU	62H		SLCDC2
23H	PD	63H		SLCDC3
24H	PDC	64H		SLCDC4
25H	IOHR0	65H	Unused	
26H	IOHR1	66H	Unused	
27H	Unused	67H	Unused	
28H	ADRL	68H		IOHR2
29H	ADRH	69H		IOHR3
2AH	ADCR0	6AH	Unused	
2BH	ADCR1	6BH	Unused	
2CH	ACERL	6CH	Unused	
2DH	TMPC0	6DH	Unused	
2EH	CTMC0	6EH	Unused	
2FH	CTMC1	6FH	Unused	
30H	CTMDL	70H	Unused	
31H	CTMDH	71H	Unused	
32H	CTMAL	72H	Unused	
33H	CTMAH	73H	Unused	
34H	STMC0	74H	Unused	
35H	STMC1	75H	Unused	
36H	STMDL	76H	Unused	
37H	STMDH	77H	Unused	
38H	STMAL	78H	Unused	
39H	STMAH	79H	Unused	
3AH	STMRP	7AH	Unused	
3BH	SIMTOC	7BH	Unused	
3CH	SIMC0	7CH	Unused	
3DH	SIMC1	7DH	Unused	
3EH	SIMD	7EH	Unused	
3FH	SIMC2/SIMA	7FH	Unused	

□ : Unused, read as 00H

### Special Purpose Data Memory

## Special Function Register Description

Most of the Special Function Register details will be described in the relevant functional sections, however several registers require a separate description in this section.

### Indirect Addressing Register – IAR0, IAR1, IAR2

The Indirect Addressing Registers, IAR0, IAR1 and IAR2, although having their locations in normal RAM register space, do not actually physically exist as normal registers. The method of indirect addressing for RAM data manipulation uses these Indirect Addressing Registers and Memory Pointers, in contrast to direct memory addressing, where the actual memory address is specified. Actions on the IAR0, IAR1 and IAR2 registers will result in no actual read or write operation to these registers but rather to the memory location specified by their corresponding Memory Pointers, MP0, MP1L/MP1H or MP2L/MP2H. Acting as a pair, IAR0 and MP0 can together access data only from Sector 0 while the IAR1 register together with MP1L/MP1H register pair and IAR2 register together with MP2L/MP2H register pair can access data from any Data Memory sector. As the Indirect Addressing Registers are not physically implemented, reading the Indirect Addressing Registers indirectly will return a result of “00H” and writing to the registers indirectly will result in no operation.

### Memory Pointers – MP0, MP1L, MP1H, MP2L, MP2H

Five Memory Pointers, known as MP0, MP1L, MP1H, MP2L and MP2H, are provided. These Memory Pointers are physically implemented in the Data Memory and can be manipulated in the same way as normal registers providing a convenient way with which to address and track data. When any operation to the relevant Indirect Addressing Registers is carried out, the actual address that the microcontroller is directed to is the address specified by the related Memory Pointer. MP0, together with Indirect Addressing Register, IAR0, are used to access data from Sector 0, while MP1L/MP1H together with IAR1 and MP2L/MP2H together with IAR2 are used to access data from all data sectors according to the corresponding MP1H or MP2H register. Direct Addressing can be used in all data sectors using the corresponding instruction which can address all available data memory space.

The following example shows how to clear a sector of four Data Memory locations already defined as locations `adres1` to `adres4`.

**Indirect Addressing Program Example 1**

```
data .section 'data'
adres1 db ?
adres2 db ?
adres3 db ?
adres4 db ?
block db ?
code .section at 0 code
org 00h
start:
mov a, 04h ; setup size of block
mov block, a
mov a, offset adres1 ; Accumulator loaded with first RAM address
mov mp0, a ; setup memory pointer with first RAM address
loop:
clr IAR0 ; clear the data at address defined by MP0
inc mp0 ; increment memory pointer
sdz block ; check if last memory location has been cleared
jmp loop
continue:
```

**Indirect Addressing Program Example 2**

```
data .section at 01F0H 'data'
adres1 db ?
adres2 db ?
adres3 db ?
adres4 db ?
block db ?
code .section at 0 'code'
org00h
start:
mov a,04h ; setup size of block
mov block,a
mov a, 01h ; setup the memory sector
mov mplh, a
mov a,offset adres1 ; Accumulator loaded with first RAM address
mov mp1l,a ; setup memory pointer with first RAM address
loop:
clr IAR1 ; clear the data at address defined by MP1L
inc mp1l ; increment memory pointer MP1L
sdz block ; check if last memory location has been cleared
jmp loop
continue:
```

The important point to note here is that in the example shown above, no reference is made to specific Data Memory addresses.

### **Accumulator – ACC**

The Accumulator is central to the operation of any microcontroller and is closely related with operations carried out by the ALU. The Accumulator is the place where all intermediate results from the ALU are stored. Without the Accumulator it would be necessary to write the result of each calculation or logical operation such as addition, subtraction, shift, etc., to the Data Memory resulting in higher programming and timing overheads. Data transfer operations usually involve the temporary storage function of the Accumulator; for example, when transferring data between one user-defined register and another, it is necessary to do this by passing the data through the Accumulator as no direct transfer between two registers is permitted.

### **Program Counter Low Register – PCL**

To provide additional program control functions, the low byte of the Program Counter is made accessible to programmers by locating it within the Special Purpose area of the Data Memory. By manipulating this register, direct jumps to other program locations are easily implemented. Loading a value directly into this PCL register will cause a jump to the specified Program Memory location, however, as the register is only 8-bit wide, only jumps within the current Program Memory page are permitted. When such operations are used, note that a dummy cycle will be inserted.

### **Look-up Table Registers – TBLP, TBHP, TBLH**

These three special function registers are used to control operation of the look-up table which is stored in the Program Memory. TBLP and TBHP are the table pointers and indicate the location where the table data is located. Their value must be setup before any table read commands are executed. Their value can be changed, for example using the “INC” or “DEC” instructions, allowing for easy table data pointing and reading. TBLH is the location where the high order byte of the table data is stored after a table read data instruction has been executed. Note that the lower order table data byte is transferred to a user defined location.

### **Status Register – STATUS**

This 8-bit register contains the SC flag, CZ flag, zero flag (Z), carry flag (C), auxiliary carry flag (AC), overflow flag (OV), power down flag (PDF), and watchdog time-out flag (TO). These arithmetic/logical operation and system management flags are used to record the status and operation of the microcontroller.

With the exception of the TO and PDF flags, bits in the status register can be altered by instructions like most other registers. Any data written into the status register will not change the TO or PDF flag. In addition, operations related to the status register may give different results due to the different instruction operations. The TO flag can be affected only by a system power-up, a WDT time-out or by executing the “CLR WDT” or “HALT” instruction. The PDF flag is affected only by executing the “HALT” or “CLR WDT” instruction or during a system power-up.

The Z, OV, AC, C, SC and CZ flags generally reflect the status of the latest operations.

- SC is the result of the “XOR” operation which is performed by the OV flag and the MSB of the current instruction operation result.
- CZ is the operational result of different flags for different instructions. Refer to register definitions for more details.
- C is set if an operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation; otherwise C is cleared. C is also affected by a rotate through carry instruction.

- AC is set if an operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction; otherwise AC is cleared.
- Z is set if the result of an arithmetic or logical operation is zero; otherwise Z is cleared.
- OV is set if an operation results in a carry into the highest-order bit but not a carry out of the highest-order bit, or vice versa; otherwise OV is cleared.
- PDF is cleared by a system power-up or executing the “CLR WDT” instruction. PDF is set by executing the “HALT” instruction.
- TO is cleared by a system power-up or executing the “CLR WDT” or “HALT” instruction. TO is set by a WDT time-out.

In addition, on entering an interrupt sequence or executing a subroutine call, the status register will not be pushed onto the stack automatically. If the contents of the status registers are important and if the subroutine can corrupt the status register, precautions must be taken to correctly save it.

### STATUS Register

Bit	7	6	5	4	3	2	1	0
Name	SC	CZ	TO	PDF	OV	Z	AC	C
R/W	R/W	R/W	R	R	R/W	R/W	R/W	R/W
POR	x	x	0	0	x	x	x	x

“x” unknown

- Bit 7      **SC**: The result of the “XOR” operation which is performed by the OV flag and the MSB of the instruction operation result.
- Bit 6      **CZ**: The the operational result of different flags for different instructions.  
 For SUB/SUBM/LSUB/LSUBM instructions, the CZ flag is equal to the Z flag.  
 For SBC/SBCM/LSBC/LSBCM instructions, the CZ flag is the “AND” operation result which is performed by the previous operation CZ flag and current operation zero flag.  
 For other instructions, the CZ flag will not be affected.
- Bit 5      **TO**: Watchdog Time-Out flag  
 0: After power up or executing the “CLR WDT” or “HALT” instruction  
 1: A watchdog time-out occurred.
- Bit 4      **PDF**: Power down flag  
 0: After power up or executing the “CLR WDT” instruction  
 1: By executing the “HALT” instruction
- Bit 3      **OV**: Overflow flag  
 0: No overflow  
 1: An operation results in a carry into the highest-order bit but not a carry out of the highest-order bit or vice versa.
- Bit 2      **Z**: Zero flag  
 0: The result of an arithmetic or logical operation is not zero  
 1: The result of an arithmetic or logical operation is zero
- Bit 1      **AC**: Auxiliary flag  
 0: No auxiliary carry  
 1: An operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction
- Bit 0      **C**: Carry flag  
 0: No carry-out  
 1: An operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation  
 C is also affected by a rotate through carry instruction.

## EEPROM Data Memory

One of the special features in the device is its internal EEPROM Data Memory. EEPROM, which stands for Electrically Erasable Programmable Read Only Memory, is by its nature a non-volatile form of memory, with data retention even when its power supply is removed. By incorporating this kind of data memory, a whole new host of application possibilities are made available to the designer. The availability of EEPROM storage allows information such as product identification numbers, calibration values, specific user data, system setup data or other product information to be stored directly within the product microcontroller. The process of reading and writing data to the EEPROM memory has been reduced to a very trivial affair.

### EEPROM Data Memory Structure

The EEPROM Data Memory capacity is 64×8 bits for the device. Unlike the Program Memory and RAM Data Memory, the EEPROM Data Memory is not directly mapped into memory space and is therefore not directly addressable in the same way as the other types of memory. Read and Write operations to the EEPROM are carried out in single byte operations using an address and data register in Sector 0 and a single control register in Sector 1.

### EEPROM Registers

Three registers control the overall operation of the internal EEPROM Data Memory. These are the address register, EEA, the data register, EED and a single control register, EEC. As both the EEA and EED registers are located in all sectors, they can be directly accessed in the same way as any other Special Function Register. The EEC register however, being located in Sector 1, cannot be directly addressed directly and can only be read from or written to indirectly using the MP1L/MP1H or MP2L/MP2H Memory Pointer and Indirect Addressing Register, IAR1 or IAR2. Because the EEC control register is located at address 40H in Sector 1, the MP1L or MP2L Memory Pointer low byte must first be set to the value 40H and the MP1H or MP2H Memory Pointer high byte set to the value 01H before any operations on the EEC register are executed.

Register Name	Bit							
	7	6	5	4	3	2	1	0
EEA	—	—	D5	D4	D3	D2	D1	D0
EED	D7	D6	D5	D4	D3	D2	D1	D0
EEC	—	—	—	—	WREN	WR	RDEN	RD

EEPROM Control Registers List

#### EEA Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	D5	D4	D3	D2	D1	D0
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	0	0	0	0

- Bit 7~6 Unimplemented, read as “0”
- Bit 5~0 **D5~D0**: Data EEPROM address  
Data EEPROM address bit 5 ~ bit 0

**EED Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0     **D7~D0**: Data EEPROM data  
 Data EEPROM data bit 7 ~ bit 0

**EED Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	WREN	WR	RDEN	RD
R/W	—	—	—	—	R/W	R/W	R/W	R/W
POR	—	—	—	—	0	0	0	0

Bit 7~4     Unimplemented, read as “0”

Bit 3     **WREN**: Data EEPROM Write Enable  
 0: Disable  
 1: Enable

This is the Data EEPROM Write Enable Bit which must be set high before Data EEPROM write operations are carried out. Clearing this bit to zero will inhibit Data EEPROM write operations.

Bit 2     **WR**: EEPROM Write Control  
 0: Write cycle has finished  
 1: Activate a write cycle

This is the Data EEPROM Write Control Bit and when set high by the application program will activate a write cycle. This bit will be automatically reset to zero by the hardware after the write cycle has finished. Setting this bit high will have no effect if the WREN has not first been set high.

Bit 1     **RDEN**: Data EEPROM Read Enable  
 0: Disable  
 1: Enable

This is the Data EEPROM Read Enable Bit which must be set high before Data EEPROM read operations are carried out. Clearing this bit to zero will inhibit Data EEPROM read operations.

Bit 0     **RD**: EEPROM Read Control  
 0: Read cycle has finished  
 1: Activate a read cycle

This is the Data EEPROM Read Control Bit and when set high by the application program will activate a read cycle. This bit will be automatically reset to zero by the hardware after the read cycle has finished. Setting this bit high will have no effect if the RDEN has not first been set high.

Note: The WREN, WR, RDEN and RD can not be set to “1” at the same time in one instruction. The WR and RD can not be set to “1” at the same time.

## Reading Data from the EEPROM

To read data from the EEPROM, the read enable bit, RDEN, in the EEC register must first be set high to enable the read function. The EEPROM address of the data to be read must then be placed in the EEA register. If the RD bit in the EEC register is now set high, a read cycle will be initiated. Setting the RD bit high will not initiate a read operation if the RDEN bit has not been set. When the read cycle terminates, the RD bit will be automatically cleared to zero, after which the data can be read from the EED register. The data will remain in the EED register until another read or write operation is executed. The application program can poll the RD bit to determine when the data is valid for reading.

## Writing Data to the EEPROM

The EEPROM address of the data to be written must first be placed in the EEA register and the data placed in the EED register. To write data to the EEPROM, the write enable bit, WREN, in the EEC register must first be set high to enable the write function. After this, the WR bit in the EEC register must be immediately set high to initiate a write cycle. These two instructions must be executed consecutively. The global interrupt bit EMI should also first be cleared before implementing any write operations, and then set again after the write cycle has started. Setting the WR bit high will not initiate a write cycle if the WREN bit has not been set. As the EEPROM write cycle is controlled using an internal timer whose operation is asynchronous to microcontroller system clock, a certain time will elapse before the data will have been written into the EEPROM. Detecting when the write cycle has finished can be implemented either by polling the WR bit in the EEC register or by using the EEPROM interrupt. When the write cycle terminates, the WR bit will be automatically cleared to zero by the microcontroller, informing the user that the data has been written to the EEPROM. The application program can therefore poll the WR bit to determine when the write cycle has ended.

## Write Protection

Protection against inadvertent write operation is provided in several ways. After the device is powered-on the Write Enable bit in the control register will be cleared preventing any write operations. Also at power-on MP1L/MP1H and MP2L/MP2H will be reset to zero, which means that Data Memory Sector 0 will be selected. As the EEPROM control register is located in Sector 1, this adds a further measure of protection against spurious write operations. During normal program operation, ensuring that the Write Enable bit in the control register is cleared will safeguard against incorrect write operations.

## EEPROM Interrupt

The EEPROM write interrupt is generated when an EEPROM write cycle has ended. The EEPROM interrupt must first be enabled by setting the DEE bit in the relevant interrupt register. When an EEPROM write cycle ends, the DEF request flag will be set. If the global and EEPROM interrupts are enabled and the stack is not full, a jump to the associated Interrupt vector will take place. When the interrupt is serviced the EEPROM interrupt flag will be automatically reset. More details can be obtained in the Interrupt section.

## Programming Considerations

Care must be taken that data is not inadvertently written to the EEPROM. Protection can be Periodic by ensuring that the Write Enable bit is normally cleared to zero when not writing. Also the Bank Pointer could be normally cleared to zero as this would inhibit access to Bank 1 where the EEPROM control register exist. Although certainly not necessary, consideration might be given in the application program to the checking of the validity of new write data by a simple read back process. When writing data the WR bit must be set high immediately after the WREN bit has been set high, to ensure the write cycle executes correctly. The global interrupt bit EMI should also be cleared before a write cycle is executed and then re-enabled after the write cycle starts. Note that the device should not enter the IDLE or SLEEP mode until the EEPROM read or write operation is totally complete. Otherwise, the EEPROM read or write operation will fail.

## Programming Examples

### Reading data from the EEPROM - polling method

```
MOV A, EEPROM_ADRES      ; user defined address
MOV EEA, A
MOV A, 040H              ; setup memory pointer MP1L
MOV MP1L, A              ; MP1 points to EEC register
MOV A, 01H               ; setup memory pointer MP1H
MOV MP1H, A
SET IAR1.1               ; set RDEN bit, enable read operations
SET IAR1.0               ; start Read Cycle - set RD bit
BACK:
SZ IAR1.0                ; check for read cycle end
JMP BACK
CLR IAR1                  ; disable EEPROM read/write
CLR BP
MOV A, EED               ; move read data to register
MOV READ_DATA, A
```

### Writing Data to the EEPROM - polling method

```
MOV A, EEPROM_ADRES      ; user defined address
MOV EEA, A
MOV A, EEPROM_DATA       ; user defined data
MOV EED, A
MOV A, 040H              ; setup memory pointer MP1L
MOV MP1L, A              ; MP1 points to EEC register
MOV A, 01H               ; setup memory pointer MP1H
MOV MP1H, A
CLR EMI                  ; disable global interrupt
SET IAR1.3               ; set WREN bit, enable write operations
SET IAR1.2               ; start Write Cycle - set WR bit
SET EMI                  ; enable global interrupt
BACK:
SZ IAR1.2                ; check for write cycle end
JMP BACK
CLR IAR1                  ; disable EEPROM write
CLR MP1H
```

## Oscillator

Various oscillator options offer the user a wide range of functions according to their various application requirements. The flexible features of the oscillator functions ensure that the best optimisation can be achieved in terms of speed and power saving. Oscillator selections and operation are selected through a combination of configuration options and registers.

### Oscillator Overview

In addition to being the source of the main system clock the oscillators also provide clock sources for the Watchdog Timer and Time Base Interrupts. The external oscillators requiring some external components as well as fully integrated internal oscillators requiring no external components are provided to form a wide range of both fast and slow system oscillators. All oscillator options are selected through the configuration options. The higher frequency oscillator provides higher performance but carry with it the disadvantage of higher power requirements, while the opposite is of course true for the lower frequency oscillators. With the capability of dynamically switching between fast and slow system clock, the device has the flexibility to optimize the performance/power ratio, a feature especially important in power sensitive portable applications.

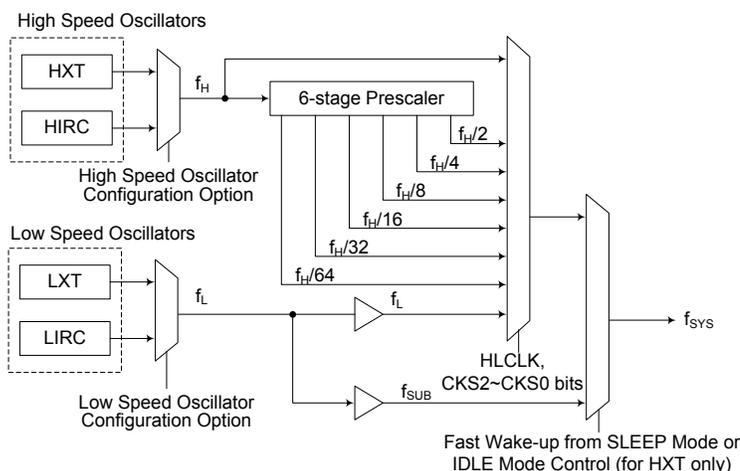
Type	Name	Freq.	Pins
External Crystal	HXT	400kHz~20MHz	OSC1/OSC2
Internal High Speed RC	HIRC	8MHz	—
External low speed Crystal	LXT	32.768kHz	XT1/XT2
Internal Low Speed RC	LIRC	32kHz	—

**Oscillator Types**

### System Clock Configurations

There are four methods of generating the system clock, two high speed oscillators and two low speed oscillators. The high speed oscillators are the external crystal oscillator and the internal 8MHz RC oscillator. The low speed oscillators are the external 32.768kHz crystal oscillator and internal 32kHz RC oscillator. Selecting whether the low or high speed oscillator is used as the system oscillator is implemented using the HLCLK bit and CKS2 ~ CKS0 bits in the SMOD register and as the system clock can be dynamically selected.

The actual source clock used for each of the high speed and the low speed oscillators is chosen via configuration options. The frequency of the slow speed or high speed system clock is also determined using the HLCLK bit and CKS2 ~ CKS0 bits in the SMOD register. Note that two oscillator selections must be made namely one high speed and one low speed system oscillators. It is not possible to choose a no-oscillator selection for either the high or low speed oscillator.



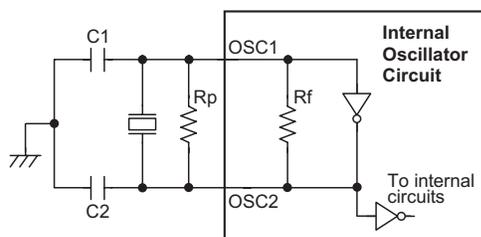
Note: The OSC1/XT1 and OSC2/XT2 share the same pins, so the HXT oscillator and the LXT oscillator can not be selected at the same time.

**System Clock Configurations**

**External Crystal/Ceramic Oscillator – HXT**

The External Crystal/Ceramic System Oscillator is one of the high frequency oscillator choices, which is selected via configuration option. For most crystal oscillator configurations, the simple connection of a crystal across OSC1 and OSC2 will create the necessary phase shift and feedback for oscillation, without requiring external capacitors. However, for some crystal types and frequencies, to ensure oscillation, it may be necessary to add two small value capacitors, C1 and C2. Using a ceramic resonator will usually require two small value capacitors, C1 and C2, to be connected as shown for oscillation to occur. The values of C1 and C2 should be selected in consultation with the crystal or resonator manufacturer’s specification.

For oscillator stability and to minimise the effects of noise and crosstalk, it is important to ensure that the crystal and any associated resistors and capacitors along with interconnecting lines are all located as close to the MCU as possible.



**Note:** 1. Rp is normally not required. C1 and C2 are required.  
 2. Although not shown OSC1/OSC2 pins have a parasitic capacitance of around 7pF.

**Crystal/Resonator Oscillator – HXT**

Crystal Oscillator C1 and C2 Values		
Crystal Frequency	C1	C2
12MHz	0pF	0pF
8MHz	0pF	0pF
4MHz	0pF	0pF
1MHz	100pF	100pF

Note: C1 and C2 values are for guidance only.

**Crystal Recommended Capacitor Values**

**Internal RC Oscillator – HIRC**

The internal RC oscillator is a fully integrated system oscillator requiring no external components. The internal RC oscillator has a fixed frequency of 8MHz. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised. As a result, at a power supply of 5V and at temperature of 25°C degrees, the fixed oscillation frequency of the HIRC will have a tolerance within 2%. PB3 and PB4 can be used as normal I/O pins.

**External 32.768kHz Crystal Oscillator – LXT**

The External 32.768kHz Crystal System Oscillator is one of the low frequency oscillator choices, which is selected via configuration option. This clock source has a fixed frequency of 32.768kHz and requires a 32.768kHz crystal to be connected between pins XT1 and XT2. The external resistor and capacitor components connected to the 32.768kHz crystal are necessary to provide oscillation. For applications where precise frequencies are essential, these components may be required to provide frequency compensation due to different crystal manufacturing tolerances. During power-up there is a time delay associated with the LXT oscillator waiting for it to start-up.

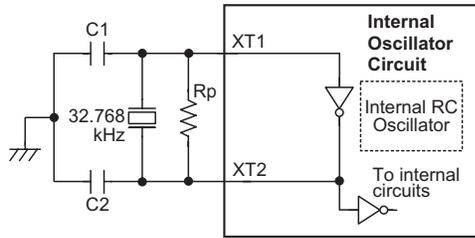
When the microcontroller enters the SLEEP or IDLE Mode, the system clock is switched off to stop microcontroller activity and to conserve power. However, in many microcontroller applications it may be necessary to keep the internal timers operational even when the microcontroller is in the SLEEP or IDLE Mode. To do this, another clock, independent of the system clock, must be provided.

However, for some crystals, to ensure oscillation and accurate frequency generation, it is necessary to add two small value external capacitors, C1 and C2. The exact values of C1 and C2 should be selected in consultation with the crystal or resonator manufacturer specification. The external parallel feedback resistor, Rp, is required.

Some configuration options determine if the XT1/XT2 pins are used for the LXT/HXT oscillator or as I/O pins.

- If the LXT oscillator is not used for any clock source, the XT1/XT2 pins can be used as normal I/O pins or HXT oscillator.
- If the LXT oscillator is used for any clock source, the 32.768kHz crystal should be connected to the XT1/XT2 pins.

For oscillator stability and to minimise the effects of noise and crosstalk, it is important to ensure that the crystal and any associated resistors and capacitors along with interconnecting lines are all located as close to the MCU as possible.



**Note:** 1. Rp, C1 and C2 are required.  
 2. Although not shown pins have parasitic capacitance of around 7pF.

**External LXT Oscillator**

LXT Oscillator C1 and C2 Values		
Crystal Frequency	C1	C2
32.768kHz	10pF	10pF
Note: 1. C1 and C2 values are for guidance only. 2. Rp=5MΩ~10MΩ is recommended.		

**32.768kHz Crystal Recommended Capacitor Values**

**LXT Oscillator Low Power Function**

The LXT oscillator can function in one of two modes, the Quick Start Mode and the Low Power Mode. The mode selection is executed using the LXTLP bit in the TBC register.

LXTLP	LXT Mode
0	Quick Start Mode
1	Low Power Mode

After power on, the LXTLP bit will be automatically cleared to zero ensuring that the LXT oscillator is in the Quick Start operating mode. In the Quick Start Mode the LXT oscillator will power up and stabilise quickly. However, after the LXT oscillator has fully powered up it can be placed into the Low-power mode by setting the LXTLP bit high. The oscillator will continue to run but with reduced current consumption, as the higher current consumption is only required during the LXT oscillator start-up. In power sensitive applications, such as battery applications, where power consumption must be kept to a minimum, it is therefore recommended that the application program sets the LXTLP bit high about 2 seconds after power-on.

It should be noted that, no matter what condition the LXTLP bit is set to, the LXT oscillator will always function normally, the only difference is that it will take more time to start up if in the Low-power mode.

**Internal 32kHz Oscillator – LIRC**

The Internal 32kHz System Oscillator is the low frequency oscillator. It is a fully integrated RC oscillator with a typical frequency of 32kHz at 5V, requiring no external components for its implementation. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised. As a result, at a power supply of 5V and at a temperature of 25°C degrees, the fixed oscillation frequency of 32kHz will have a tolerance within 10%.

**Supplementary Oscillator**

The low speed oscillator, in addition to providing a system clock source is also used to provide a clock source to two other device functions. These are the Watchdog Timer and the Time Base Interrupts.

## Operating Modes and System Clocks

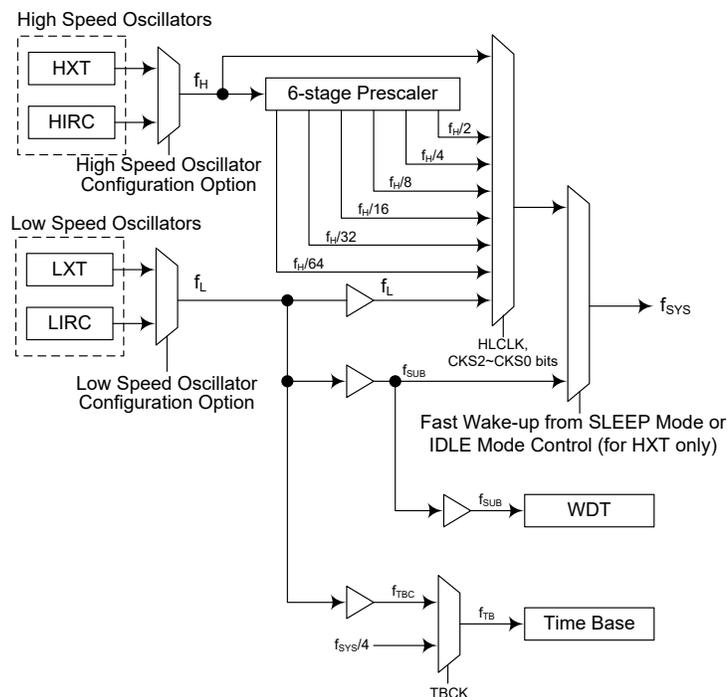
Present day applications require that their microcontrollers have high performance but often still demand that they consume as little power as possible, conflicting requirements that are especially true in battery powered portable applications. The fast clocks required for high performance will by their nature increase current consumption and of course vice-versa, lower speed clocks reduce current consumption. As Holtek has provided this device with both high and low speed clock sources and the means to switch between them dynamically, the user can optimise the operation of their microcontroller to achieve the best performance/power ratio.

### System Clocks

The device has many different clock sources for both the CPU and peripheral function operation. By providing the user with a wide range of clock options using configuration options and register programming, a clock system can be configured to obtain maximum application performance.

The main system clock, can come from either a high frequency,  $f_H$ , or a low frequency,  $f_L$ , and is selected using the HLCLK bit and CKS2~CKS0 bits in the SMOD register. The high speed system clock can be sourced from either an HXT or HIRC oscillator, selected via a configuration option. The low speed system clock source can be sourced from internal clock  $f_L$ . If  $f_L$  is selected then it can be sourced by either the LXT or LIRC oscillator, selected via a configuration option. The other choice, which is a divided version of the high speed system oscillator has a range of  $f_H/2 \sim f_H/64$ .

There are two additional internal clocks for the peripheral circuits, the substitute clock,  $f_{SUB}$ , and the Time Base clock,  $f_{TBC}$ . Each of these internal clocks is sourced by the LIRC or LXT oscillator. The  $f_{SUB}$  clock is used to provide a substitute clock for the microcontroller just after a wake-up has occurred to enable faster wake-up times. The  $f_{SUB}$  is used for the Watchdog Timer while the  $f_{TBC}$  provides the clock for the Time base and TMs.



**System Clock Configurations**

Note: When the system clock source  $f_{SYS}$  is switched to  $f_L$  from  $f_H$ , the high speed oscillation will stop to conserve the power. Thus there is no  $f_H \sim f_H/64$  for peripheral circuit to use.

## System Operation Modes

There are 6 different modes of operation for the microcontroller, each one with its own special characteristics and which can be chosen according to the specific performance and power requirements of the application. There are three modes allowing normal operation of the microcontroller, the NORMAL, SLOW0 and SLOW1 Mode. The remaining three modes, the SLEEP, IDLE0 and IDLE1 Modes are used when the microcontroller CPU is switched off to conserve power.

Operating Mode	Description			
	CPU	f <sub>sys</sub>	f <sub>SUB</sub>	f <sub>TBC</sub>
NORMAL mode	On	f <sub>H</sub> ~f <sub>H</sub> /64	On	On
SLOW0 mode	On	f <sub>L</sub>	On	On
SLOW1 mode	On	f <sub>H</sub> /2~f <sub>H</sub> /64	On	On
IDLE0 mode	Off	Off	On	On
IDLE1 mode	Off	On	On	On
SLEEP mode	Off	Off	On	Off

### NORMAL Mode

As the name suggests this is one of the main operating modes where the microcontroller has all of its functions operational and where the system clock is provided by one of the high speed oscillators. This mode operates allowing the microcontroller to operate normally with a clock source will come from one of the high speed oscillators, HXT or HIRC.

### SLOW0 Mode

This is also a mode where the microcontroller operates normally although now with a slower speed clock source. The clock source used will be from the low speed oscillator LIRC or LXT. Running the microcontroller in this mode allows it to run with much lower operating currents. In the SLOW Mode 0, the f<sub>H</sub> is off.

### SLOW1 Mode

This is also a mode where the microcontroller operates normally although now with a slower speed clock source. The high speed oscillator will however first be divided by a ratio ranging from 2 to 64, the actual ratio being selected by the CKS2~CKS0 bits in the SMOD register. Although a high speed oscillator is used, running the microcontroller at a divided clock ratio reduces the operating current.

### SLEEP Mode

The SLEEP Mode is entered when an HALT instruction is executed and the IDLEN bit in the SMOD register is low. In the SLEEP mode the CPU will be stopped. However the f<sub>SUB</sub> clock will continue to operate as the Watchdog Timer clock.

### IDLE0 Mode

The IDLE0 Mode is entered when a HALT instruction is executed and when the IDLEN bit in the SMOD register is high and the FSYSON bit in the CTRL register is low. In the IDLE0 Mode the system oscillator will be inhibited from driving the CPU but some peripheral functions will remain operational such as the Watchdog Timer and TMs. In the IDLE0 Mode, the system oscillator will be stopped.

### IDLE1 Mode

The IDLE1 Mode is entered when a HALT instruction is executed and when the IDLEN bit in the SMOD register is high and the FSYSOEN bit in the CTRL register is high. In the IDLE1 Mode the system oscillator will be inhibited from driving the CPU but may continue to provide a clock source to keep some peripheral functions operational such as the Watchdog Timer and TMs. In the IDLE1 Mode, the system oscillator will continue to run, and this system oscillator may be high speed or low speed system oscillator. In the IDLE1 Mode the Watchdog Timer clock,  $f_{SUB}$ , will be on.

### Control Register

A single register, SMOD, is used for overall control of the internal clocks within the device.

### SMOD Register

Bit	7	6	5	4	3	2	1	0
Name	CKS2	CKS1	CKS0	FSTEN	LTO	HTO	IDLEN	HLCLK
R/W	R/W	R/W	R/W	R/W	R	R	R/W	R/W
POR	0	0	0	0	0	0	1	1

Bit 7 ~ 5 **CKS2 ~ CKS0**: The system clock selection bits when HLCLK=0

- 000:  $f_L$  (LIRC or LXT)
- 001:  $f_L$  (LIRC or LXT)
- 010:  $f_H/64$
- 011:  $f_H/32$
- 100:  $f_H/16$
- 101:  $f_H/8$
- 110:  $f_H/4$
- 111:  $f_H/2$

These three bits are used to select which clock is used as the system clock source. In addition to the system clock source, which can be the LIRC or LXT, a divided version of the high speed system oscillator can also be chosen as the system clock source.

Bit 4 **FSTEN**: Fast Wake-up Control (only for HXT)

- 0: Disable
- 1: Enable

This is the Fast Wake-up Control bit which determines if the  $f_{SUB}$  clock source is initially used after the device wakes up. When the bit is high, the  $f_{SUB}$  clock source can be used as a temporary system clock to provide a faster wake up time as the  $f_{SUB}$  clock is available.

Bit 3 **LTO**: Low speed system oscillator ready flag

- 0: Not ready
- 1: Ready

This is the low speed system oscillator ready flag which indicates when the low speed system oscillator is stable after power on reset.

Bit 2 **HTO**: High speed system oscillator ready flag

- 0: Not ready
- 1: Ready

This is the high speed system oscillator ready flag which indicates when the high speed system oscillator is stable. This flag is cleared to "0" by hardware when the device is powered on and then changes to a high level after the high speed system oscillator is stable. Therefore this flag will always be read as "1" by the application program after device power-on. The flag will be low when in the SLEEP, IDLE0 or SLOW0 Mode, but after power on reset or a wake-up has occurred, the flag will change to a high level after 128 clock cycles if the HXT oscillator is used and after 15~16 clock cycles if the HIRC oscillator is used.

- Bit 1      **IDLEN**: IDLE Mode Control  
             0: Disable  
             1: Enable  
 This is the IDLE Mode Control bit and determines what happens when the HALT instruction is executed. If this bit is high, when a HALT instruction is executed the device will enter the IDLE Mode. In the IDLE1 Mode the CPU will stop running but the system clock will continue to keep the peripheral functions operational, if FSYSON bit is high. If FSYSON bit is low, the CPU and the system clock will all stop in IDLE0 mode. If the bit is low the device will enter the SLEEP Mode when a HALT instruction is executed.
- Bit 0      **HLCLK**: System Clock Selection  
             0:  $f_H/2 \sim f_H/64$  or  $f_L$   
             1:  $f_H$   
 This bit is used to select if the  $f_H$  clock or the  $f_H/2 \sim f_H/64$  or  $f_L$  clock is used as the system clock. When the bit is high the  $f_H$  clock will be selected and if low the  $f_H/2 \sim f_H/64$  or  $f_L$  clock will be selected. When system clock switches from the  $f_H$  clock to the  $f_L$  clock and the  $f_H$  clock will be automatically switched off to conserve power.

### CTRL Register

Bit	7	6	5	4	3	2	1	0
Name	FSYSON	—	—	—	—	LVRF	LRF	WRF
R/W	R/W	—	—	—	—	R/W	R/W	R/W
POR	0	—	—	—	—	0	0	0

- Bit 7      **FSYSON**:  $f_{SYS}$  Control IDLE Mode  
             0: Disable  
             1: Enable
- Bit 6~3    Unimplemented, read as “0”
- Bit 2      **LVRF**: LVR function reset flag  
 Describe elsewhere
- Bit 1      **LRF**: LVRC register software reset flag  
 Describe elsewhere
- Bit 0      **WRF**: WDTC register software reset flag  
 Describe elsewhere

### Fast Wake-up

To minimise power consumption the device can enter the SLEEP or IDLE0 Mode, where the system clock source to the device will be stopped. However when the device is woken up again, it can take a considerable time for the original system oscillator to restart, stabilise and allow normal operation to resume. To ensure the device is up and running as fast as possible a Fast Wake-up function is provided, which allows  $f_{SUB}$ , to act as a temporary clock to first drive the system until the original system oscillator has stabilised. The temporary clock is sourced by the LIRC or LXT oscillator. As the clock source for the Fast Wake-up function is  $f_{SUB}$ , the Fast Wake-up function is only available in the SLEEP and IDLE0 modes. The Fast Wake-up enable/disable function is controlled using the FSTEN bit in the SMOD register.

If the HXT oscillator is selected as the NORMAL Mode system clock, and if the Fast Wake-up function is enabled, then it will take one to two  $t_{SUB}$  clock cycles for the system to wake-up. The system will then initially run under the  $f_{SUB}$  clock source until 128 HXT clock cycles have elapsed, at which point the HTO flag will switch high and the system will switch over to operating from the HXT oscillator.

If the HIRC oscillators or LIRC oscillator is used as the system oscillator then it will take 15~16 clock cycles or 1~2 cycles to wake up the system from the SLEEP or IDLE0 Mode. The Fast Wake-up bit, FSTEN will have no effect in these cases.

System Oscillator	FSTEN Bit	Wake-up Time (SLEEP Mode)	Wake-up Time (IDLE0 Mode)	Wake-up Time (IDLE1 Mode)
HXT	0	128 HXT cycles		1~2 HXT cycles
	1	1~2 $f_{SUB}$ cycles (System runs with $f_{SUB}$ first for 128 HXT cycles and then switches over to run with the HXT clock)		1~2 HXT cycles
HIRC	x	15~16 HIRC cycles		1~2 HIRC cycles
LIRC	x	1~2 LIRC cycles		1~2 LIRC cycles
LXT	x	128 LXT cycles		1~2 LXT cycles

**Wake-up Times**

### Operating Mode Switching

The device can switch between operating modes dynamically allowing the user to select the best performance/power ratio for the present task in hand. In this way microcontroller operations that do not require high performance can be executed using slower clocks thus requiring less operating current and prolonging battery life in portable applications.

In simple terms, Mode Switching between the NORMAL Mode and SLOW Mode is executed using the HLCLK bit and CKS2~CKS0 bits in the SMOD register while Mode Switching from the NORMAL/SLOW Modes to the SLEEP/IDLE Modes is executed via the HALT instruction. When a HALT instruction is executed, whether the device enters the IDLE Mode or the SLEEP Mode is determined by the condition of the IDLEN bit in the SMOD register and FSYSON in the CTRL register.

When the HLCLK bit switches to a low level, which implies that clock source is switched from the high speed clock source,  $f_{H}$ , to the clock source,  $f_{H}/2 \sim f_{H}/64$  or  $f_{L}$ . If the clock is from the  $f_{L}$ , the high speed clock source will stop running to conserve power. When this happens it must be noted that the  $f_{H}/16$  and  $f_{H}/64$  internal clock sources will also stop running, which may affect the operation of other internal functions such as the TMs. The accompanying flowchart shows what happens when the device moves between the various operating modes.

#### NORMAL Mode to SLOW0 Mode Switching

When running in the NORMAL Mode, which uses the high speed system oscillator, and therefore consumes more power, the system clock can switch to run in the SLOW0 Mode by setting the HLCLK bit to “0” and setting the CKS2~CKS0 bits to “000” or “001” in the SMOD register. This will then use the low speed system oscillator which will consume less power. Users may decide to do this for certain operations which do not require high performance and can subsequently reduce power consumption. The SLOW Mode is sourced from the LIRC or LXT oscillator and therefore requires these oscillators to be stable before full mode switching occurs. This is monitored using the LTO bit in the SMOD register.

### **NORMAL Mode to SLOW1 Mode Switching**

When running in the NORMAL Mode, which uses the high speed system oscillator, and therefore consumes more power, the system clock can switch to run in the SLOW1 Mode by setting the HLCLK bit to “0” and setting the CKS2~CKS0 bits to “010”, “011”, “100”, “101”, “110” or “111” in the SMOD register. This will then use a divided clock of the high speed system oscillator which can reduce the operating current. The SLOW1 Mode is still sourced from the HIRC or HXT oscillator and therefore requires less time for full mode switching.

### **SLOW0 Mode to NORMAL Mode Switching**

In SLOW0 Mode the system uses LIRC or LXT low speed system oscillator. To switch back to the NORMAL Mode, where the high speed system oscillator is used, the HLCLK bit should be set to “1”. As a certain amount of time will be required for the high frequency clock to stabilise, the status of the HTO bit is checked. The amount of time required for high speed system oscillator stabilization depends upon which high speed system oscillator type is used.

### **SLOW1 Mode to NORMAL Mode Switching**

In SLOW1 Mode the system still uses high speed system oscillator. To switch back to the NORMAL Mode, where also the high speed system oscillator is used, the HLCLK bit should be set to “1”. As the two modes both use the high speed system oscillator, therefore requires less time for full mode switching.

### **Entering the SLEEP Mode**

There is only one way for the device to enter the SLEEP Mode and that is to execute the “HALT” instruction in the application program with the IDLEN bit in SMOD register equal to “0” and the WDT on. When this instruction is executed under the conditions described above, the following will occur:

- The system clock and Time Base clock will be stopped and the application program will stop at the “HALT” instruction, but the WDT will remain with the clock source coming from the  $f_{SUB}$  clock.
- The Data Memory contents and registers will maintain their present condition.
- The WDT will be cleared and resume counting.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.

### **Entering the IDLE0 Mode**

There is only one way for the device to enter the IDLE0 Mode and that is to execute the “HALT” instruction in the application program with the IDLEN bit in SMOD register equal to “1” and the FSYSON bit in CTRL register equal to “0”. When this instruction is executed under the conditions described above, the following will occur:

- The system clock will be stopped and the application program will stop at the “HALT” instruction, but the Time Base clock  $f_{TBC}$  and the  $f_{SUB}$  clock will be on.
- The Data Memory contents and registers will maintain their present condition.
- The WDT will be cleared and resume counting.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.

### **Entering the IDLE1 Mode**

There is only one way for the device to enter the IDLE1 Mode and that is to execute the “HALT” instruction in the application program with the IDLEN bit in SMOD register equal to “1” and the FSYSON bit in CTRL register equal to “1”. When this instruction is executed under the conditions described above, the following will occur:

- The system clock and Time Base clock and  $f_{SUB}$  will be on and the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition.
- The WDT will be cleared and resume counting.
- The I/O ports will maintain their present conditions.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.

### **Standby Current Considerations**

As the main reason for entering the SLEEP or IDLE Mode is to keep the current consumption of the device to as low a value as possible, perhaps only in the order of several micro-amps except in the IDLE1 Mode, there are other considerations which must also be taken into account by the circuit designer if the power consumption is to be minimised. Special attention must be made to the I/O pins on the device. All high-impedance input pins must be connected to either a fixed high or low level as any floating input pins could create internal oscillations and result in increased current consumption. This also applies to devices which have different package types, as there may be unbonded pins. These must either be setup as outputs or if setup as inputs must have pull-high resistors connected.

Care must also be taken with the loads, which are connected to I/O pins, which are setup as outputs. These should be placed in a condition in which minimum current is drawn or connected only to external circuits that do not draw current, such as other CMOS inputs. In the IDLE1 Mode the system oscillator is on, if the system oscillator is from the high speed system oscillator, the additional standby current will also be perhaps in the order of several hundred micro-amps.

### **Wake-up**

After the system enters the SLEEP or IDLE Mode, it can be woken up from one of various sources listed as follows:

- An external falling edge on Port A
- A system interrupt
- A WDT overflow

If the device is woken up by a WDT overflow, a Watchdog Timer reset will be initiated. Although both of these wake-up methods will initiate a reset operation, the actual source of the wake-up can be determined by examining the TO and PDF flags. The PDF flag is cleared by a system power-up or executing the clear Watchdog Timer instructions and is set when executing the “HALT” instruction. The TO flag is set if a WDT time-out occurs, and causes a wake-up that only resets the Program Counter and Stack Pointer, the other flags remain in their original status.

Each pin on Port A can be setup using the PAWU register to permit a negative transition on the pin to wake-up the system. When a Port A pin wake-up occurs, the program will resume execution at the instruction following the “HALT” instruction. If the system is woken up by an interrupt, then two possible situations may occur. The first is where the related interrupt is disabled or the interrupt is enabled but the stack is full, in which case the program will resume execution at the instruction following the “HALT” instruction. In this situation, the interrupt which woke-up the device will not be immediately serviced, but will rather be serviced later when the related interrupt is finally enabled or when a stack level becomes free. The other situation is where the related interrupt is enabled and the stack is not full, in which case the regular interrupt response takes place. If an interrupt request flag is set high before entering the SLEEP or IDLE Mode, the wake-up function of the related interrupt will be disabled.

### **Programming Considerations**

The high speed and low speed oscillators both use the same SST counter. For example, if the system is woken up from the SLEEP Mode and the HIRC oscillators need to start-up from an off state.

- If the device is woken up from the SLEEP Mode to NORMAL Mode, and the system clock source is from HXT oscillator and FSTEN is “1”, the system clock can be switched to the LIRC or LXT oscillator after wake up.
- There are peripheral functions, such as TMs, for which the  $f_{SYS}$  is used. If the system clock source is switched from  $f_H$  to  $f_L$ , the clock source to the peripheral functions mentioned above will change accordingly.

### **Watchdog Timer**

The Watchdog Timer is provided to prevent program malfunctions or sequences from jumping to unknown locations, due to certain uncontrollable external events such as electrical noise.

#### **Watchdog Timer Clock Source**

The Watchdog Timer clock source is provided by the  $f_{SUB}$  clock. The  $f_{SUB}$  clock can be sourced from either the LIRC or LXT oscillator selected by a configuration option. The LIRC internal oscillator has an approximate frequency of 32kHz and this specified internal clock period can vary with  $V_{DD}$ , temperature and process variations. The LXT oscillator is supplied by an external 32.768 kHz crystal. The Watchdog Timer source clock is then subdivided by a ratio of  $2^8$  to  $2^{18}$  to give longer timeouts, the actual value being chosen using the WS2~WS0 bits in the WDTC register.

#### **Watchdog Timer Control Register**

A single register, WDTC, controls the required timeout period as well as the enable or reset operation.

### WDTC Register

Bit	7	6	5	4	3	2	1	0
Name	WE4	WE3	WE2	WE1	WE0	WS2	WS1	WS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	1	0	1	0	0	1	1

Bit 7~3 **WE4 ~ WE0**: WDT function software control  
 10101/01010: WDT enable  
 Other values: Reset MCU  
 When these bits are changed to any other values by the environmental noise to reset the microcontroller, the reset operation will be activated after 2~3  $t_{SUB}$  clock cycles and the WRF bit in the CTRL register will be set to 1 to indicate the reset source.

Bit 2~0 **WS2 ~ WS0**: WDT Time-out period selection  
 000:  $2^8/f_{SUB}$   
 001:  $2^{10}/f_{SUB}$   
 010:  $2^{12}/f_{SUB}$   
 011:  $2^{14}/f_{SUB}$   
 100:  $2^{15}/f_{SUB}$   
 101:  $2^{16}/f_{SUB}$   
 110:  $2^{17}/f_{SUB}$   
 111:  $2^{18}/f_{SUB}$   
 These three bits determine the division ratio of the Watchdog Timer source clock, which in turn determines the timeout period.

### CTRL Register

Bit	7	6	5	4	3	2	1	0
Name	FSYSON	—	—	—	—	LVRF	LRF	WRF
R/W	R/W	—	—	—	—	R/W	R/W	R/W
POR	0	—	—	—	—	0	0	0

Bit 7 **FSYSON**:  $f_{SYS}$  Control IDLE Mode  
 Describe elsewhere

Bit 6~3 Unimplemented, read as “0”

Bit 2 **LVRF**: LVR function reset flag  
 Describe elsewhere

Bit 1 **LRF**: LVRC register software reset flag  
 Describe elsewhere

Bit 0 **WRF**: WDTC register software reset flag  
 0: Not occur  
 1: Occurred  
 This bit is set to 1 by the WDTC Control register software reset and cleared by the application program. Note that this bit can only be cleared to 0 by the application program.

## Watchdog Timer Operation

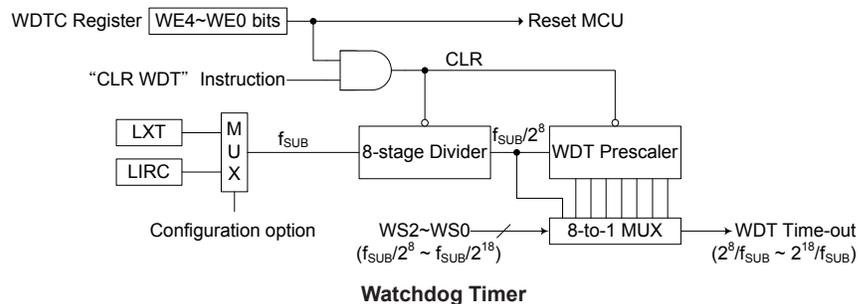
The Watchdog Timer operates by providing a device reset when its timer overflows. This means that in the application program and during normal operation the user has to strategically clear the Watchdog Timer before it overflows to prevent the Watchdog Timer from executing a reset. This is done using the clear watchdog instructions. If the program malfunctions for whatever reason, jumps to an unknown location, or enters an endless loop, the clear WDT instructions will not be executed in the correct manner, in which case the Watchdog Timer will overflow and reset the device. With regard to the Watchdog Timer enable/reset function, there are five bits, WE4~WE0, in the WDTC register to additional enable and reset control of the Watchdog Timer.

WE4 ~ WE0 Bits	WDT Function
10101B/01010B	Enable
Any other value	Reset MCU

### Watchdog Timer Enable/Reset Control

Under normal program operation, a Watchdog Timer time-out will initialise a device reset and set the status bit TO. However, if the system is in the SLEEP or IDLE Mode, when a Watchdog Timer time-out occurs, the TO bit in the status register will be set and only the Program Counter and Stack Pointer will be reset. Three methods can be adopted to clear the contents of the Watchdog Timer. The first is a WDT reset, which means a value other than 01010B and 10101B is written into the WE4~WE0 bit locations, the second is using the Watchdog Timer software clear instructions and the third is via a HALT instruction. There is only one method of using software instruction to clear the Watchdog Timer. That is to use the single “CLR WDT” instruction to clear the WDT.

The maximum time-out period is when the  $2^{18}$  division ratio is selected. As an example, with a 32 kHz LIRC oscillator as its source clock, this will give a maximum watchdog period of around 8 seconds for the  $2^{18}$  division ratio, and a minimum timeout of 7.8ms for the  $2^8$  division ratio.



## Reset and Initialisation

A reset function is a fundamental part of any microcontroller ensuring that the device can be set to some predetermined condition irrespective of outside parameters. The most important reset condition is after power is first applied to the microcontroller. In this case, internal circuitry will ensure that the microcontroller, after a short delay, will be in a well defined state and ready to execute the first program instruction. After this power-on reset, certain important internal registers will be set to defined states before the program commences. One of these registers is the Program Counter, which will be reset to zero forcing the microcontroller to begin program execution from the lowest Program Memory address.

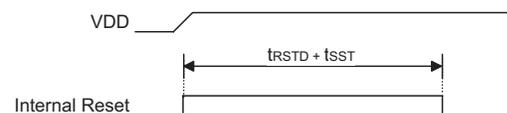
Another type of reset is when the Watchdog Timer overflows and resets the microcontroller. All types of reset operations result in different register conditions being setup. Another reset exists in the form of a Low Voltage Reset, LVR, where a full reset is implemented in situations where the power supply voltage falls below a certain threshold.

### Reset Functions

There are several ways in which a microcontroller reset can occur, through events occurring both internally and externally:

#### Power-on Reset

The most fundamental and unavoidable reset is the one that occurs after power is first applied to the microcontroller. As well as ensuring that the Program Memory begins execution from the first memory address, a power-on reset also ensures that certain other registers are preset to known conditions. All the I/O port and port control registers will power up in a high condition ensuring that all pins will be first set to inputs.

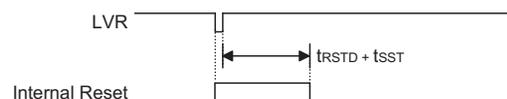


Note:  $t_{RSTD}$  is power-on delay, typical time=50ms

**Power-On Reset Timing Chart**

#### Low Voltage Reset – LVR

The microcontroller contains a low voltage reset circuit in order to monitor the supply voltage of the device, which is always enabled selected by the configuration option. If the supply voltage of the device drops to within a range of  $0.9V \sim V_{LVR}$  such as might occur when changing the battery, the LVR will automatically reset the device internally and the LVRF bit in the CTRL register will also be set to 1. For a valid LVR signal, a low voltage, i.e., a voltage in the range between  $0.9V \sim V_{LVR}$  must exist for greater than the value  $t_{LVR}$  specified in the A.C. characteristics. If the low voltage state does not exceed this value, the LVR will ignore the low supply voltage and will not perform a reset function. The actual  $V_{LVR}$  can be selected by the LVS7~LVS0 bits in the LVRC register. If the LVS7~LVS0 bits are changed to some certain values by the environmental noise, the LVR will reset the device after 2~3 LIRC/LXT clock cycles. When this happens, the LRF bit in the CTRL register will be set to 1. After power on the register will have the value of 01010101B. Note that the LVR function will be automatically disabled when the device enters the power down mode.



Note:  $t_{RSTD}$  is power-on delay, typical time=50ms

**Low Voltage Reset Timing Chart**

• **LVRC Register**

Bit	7	6	5	4	3	2	1	0
Name	LVS7	LVS6	LVS5	LVS4	LVS3	LVS2	LVS1	LVS0
R/W	R/W	R/W	R/W	R/W	R	R	R/W	R/W
POR	0	1	0	1	0	1	0	1

Bit 7 ~ 0 **LVS7 ~ LVS0**: LVR Voltage Select control

01010101: 2.1V  
 00110011: 2.55V  
 10011001: 3.15V  
 10101010: 3.8V

Other values: MCU reset – register is reset to POR value

When an actual low voltage condition occurs, as specified by one of the four defined LVR voltage values above, an MCU reset will be generated. The reset operation will be activated after 2~3 LIRC/LXT clock cycles. In this situation this register contents will remain the same after such a reset occurs.

• **CTRL Register**

Bit	7	6	5	4	3	2	1	0
Name	FSYSON	—	—	—	—	LVRF	LRF	WRF
R/W	R/W	—	—	—	—	R/W	R/W	R/W
POR	0	—	—	—	—	0	0	0

Bit 7 **FSYSON**:  $f_{SYS}$  Control IDLE Mode

Describe elsewhere

Bit 6~3 Unimplemented, read as “0”

Bit 2 **LVRF**: LVR function reset flag

0: Not occur  
 1: Occurred

This bit is set to 1 when a specific Low Voltage Reset situation condition occurs. This bit can only be cleared to zero by the application program.

Bit 1 **LRF**: LVRC register software reset flag

0: Not occur  
 1: Occurred

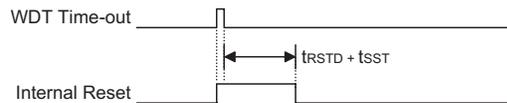
This bit is set to 1 if the LVRC register contains any non defined LVR voltage register values. This in effect acts like a software reset function. This bit can only be cleared to zero by the application program.

Bit 0 **WRF**: WDTC register software reset flag

Describe elsewhere

**Watchdog Time-out Reset during Normal Operation**

The Watchdog time-out Reset during normal operation is the same as an LVR reset except that the Watchdog time-out flag TO will be set to “1”.

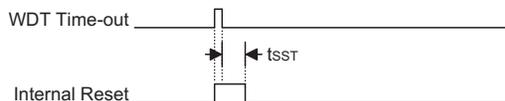


Note:  $t_{RSTD}$  is power-on delay, typical time=16.7ms

**WDT Time-out Reset during Normal Operation Timing Chart**

### Watchdog Time-out Reset during SLEEP or IDLE Mode

The Watchdog time-out Reset during SLEEP or IDLE Mode is a little different from other kinds of reset. Most of the conditions remain unchanged except that the Program Counter and the Stack Pointer will be cleared to “0” and the TO flag will be set to “1”. Refer to the A.C. Characteristics for  $t_{SST}$  details.



Note: The  $t_{SST}$  is 15~16 clock cycles if the system clock source is provided by the HIRC.  
 The  $t_{SST}$  is 128 clock cycles for HXT/LXT.  
 The  $t_{SST}$  is 1~2 clock cycles for the LIRC.

### WDT Time-out Reset during SLEEP or IDLE Timing Chart

### Reset Initial Conditions

The different types of reset described affect the reset flags in different ways. These flags, known as PDF and TO are located in the status register and are controlled by various microcontroller operations, such as the SLEEP or IDLE Mode function or Watchdog Timer. The reset flags are shown in the table:

TO	PDF	RESET Conditions
0	0	Power-on reset
u	u	LVR reset during NORMAL or SLOW Mode operation
1	u	WDT time-out reset during NORMAL or SLOW Mode operation
1	1	WDT time-out reset during IDLE or SLEEP Mode operation

Note: “u” stands for unchanged

The following table indicates the way in which the various components of the microcontroller are affected after a power-on reset occurs.

Item	Condition After RESET
Program Counter	Reset to zero
Interrupts	All interrupts will be disabled
WDT	Clear after reset, WDT begins counting
Timer Modules	Timer Modules will be turned off
Input/Output Ports	I/O ports will be setup as inputs
Stack Pointer	Stack Pointer will point to the top of the stack

The different kinds of resets all affect the internal registers of the microcontroller in different ways. To ensure reliable continuation of normal program execution after a reset occurs, it is important to know what condition the microcontroller is in after a particular reset occurs. The following table describes how each type of reset affects each of the microcontroller internal registers. Note that where more than one package type exists the table will reflect the situation for the larger package type.

Register	Reset (Power On)	WDT Time-out (Normal Operation)	LVR Reset	WDT Time-out (HALT)
IAR0	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
MP0	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
IAR1	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
MP1L	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
MP1H	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
IAR2	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
MP2L	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
MP2H	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
ACC	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
PCL	0000 0000	0000 0000	0000 0000	0000 0000
TBLP	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
TBLH	xxxx xxxx	uuuu uuuu	uuuu uuuu	uuuu uuuu
TBHP	---x xxxx	---u uuuu	---u uuuu	---u uuuu
STATUS	xx00 xxxx	uu1u uuuu	uuuu uuuu	uu11 uuuu
SMOD	0000 0011	0000 0011	0000 0011	uuuu uuuu
INTEG0	0000 0000	0000 0000	0000 0000	uuuu uuuu
INTEG1	---- 0000	---- 0000	---- 0000	---- uuuu
WDTC	0101 0011	0101 0011	0101 0011	uuuu uuuu
TBC	0011 0111	0011 0111	0011 0111	uuuu uuuu
INTC0	-000 0000	-000 0000	-000 0000	-uuu uuuu
INTC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
INTC2	0000 0000	0000 0000	0000 0000	uuuu uuuu
INTC3	0000 0000	0000 0000	0000 0000	uuuu uuuu
MFIO	--00 --00	--00 --00	--00 --00	--uu --uu
MF11	0-00 --00	0-00 --00	0-00 --00	u-uu --uu
MF12	0000 0000	0000 0000	0000 0000	uuuu uuuu
PAWU	0000 0000	0000 0000	0000 0000	uuuu uuuu
PAPU	0000 0000	0000 0000	0000 0000	uuuu uuuu
PA	1111 1111	1111 1111	1111 1111	uuuu uuuu
PAC	1111 1111	1111 1111	1111 1111	uuuu uuuu
PBPU	0000 0000	0000 0000	0000 0000	uuuu uu0
PB	1111 1111	1111 1111	1111 1111	uuuu uuuu
PBC	1111 1111	1111 1111	1111 1111	uuuu uuuu
PCPU	0000 0000	0000 000	0000 0000	uuuu uuuu
PC	1111 1111	1111 1111	1111 1111	uuuu uuuu
PCC	1111 1111	1111 1111	1111 1111	uuuu uuuu
PDPU	--00 0000	--00 0000	--00 0000	--uu uuuu
PD	--11 1111	--11 1111	--11 1111	--uu uuuu
PDC	--11 1111	--11 1111	--11 1111	--uu uuuu
IOHR0	0000 0000	0000 0000	0000 0000	uuuu uuuu
IOHR1	0000 0000	0000 0000	0000 0000	uuuu uuuu
IOHR2	0000 0000	0000 0000	0000 0000	uuuu uuuu
IOHR3	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTRL	0--- -000	0--- -000	0--- -000	u--- -uuu
LVRC	0101 0101	0101 0101	uuuu uuuu	uuuu uuuu
ADRL(ADRFS=0)	xxxx ----	xxxx ----	xxxx ----	uuuu ----
ADRL(ADRFS=1)	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
ADRH(ADRFS=0)	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu

Register	Reset (Power On)	WDT Time-out (Normal Operation)	LVR Reset	WDT Time-out (HALT)
ADRH(ADRFSS=1)	---- xxxx	---- xxxx	---- xxxx	---- uuuu
ADCR0	0110 -000	0110 -000	0110 -000	uuuu -uuu
ADCR1	00--0000	00--0000	00--0000	uu--uuuu
ACERL	1111 1111	1111 1111	1111 1111	uuuu uuuu
TMPC0	-0-0 -0-0	-0-0 -0-0	-0-0 -0-0	-u-u -u-u
EEA	--00 0000	--00 0000	--00 0000	--uu uuuu
EED	0000 0000	0000 0000	0000 0000	uuuu uuuu
EEC	---- 0000	---- 0000	---- 0000	---- uuuu
CTMC0	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMDL	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMDH	---- --00	---- --00	---- --00	---- --uu
CTMAL	0000 0000	0000 0000	0000 0000	uuuu uuuu
CTMAH	---- --00	---- --00	---- --00	---- --uu
STMC0	0000 0---	0000 0---	0000 0---	uuuu u---
STMC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
STMDL	0000 0000	0000 0000	0000 0000	uuuu uuuu
STMDH	0000 0000	0000 0000	0000 0000	uuuu uuuu
STMAL	0000 0000	0000 0000	0000 0000	uuuu uuuu
STMAH	0000 0000	0000 0000	0000 0000	uuuu uuuu
STMRP	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM0C0	0000 0---	0000 0---	0000 0---	uuuu u---
PTM0C1	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM0DL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM0DH	---- --00	---- --00	---- --00	---- --uu
PTM0AL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM0AH	---- --00	---- --00	---- --00	---- --uu
PTM0RPL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM0RPH	---- --00	---- --00	---- --00	---- --uu
PTM1C0	0000 0---	0000 0---	0000 0---	uuuu u---
PTM1C1	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM1DL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM1DH	---- --00	---- --00	---- --00	---- --uu
PTM1AL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM1AH	---- --00	---- --00	---- --00	---- --uu
PTM1RPL	0000 0000	0000 0000	0000 0000	uuuu uuuu
PTM1RPH	---- --00	---- --00	---- --00	---- --uu
SIMTOC	0000 0000	0000 0000	0000 0000	uuuu uuuu
SIMC0	111- 0000	111- 0000	111- 0000	uuu- uuuu
SIMC1	1000 0001	1000 0001	1000 0001	uuuu uuuu
SIMD	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
SIMA	0000 000-	0000 000-	0000 000-	uuuu uuuu
SIMC2	0000 0000	0000 0000	0000 0000	uuuu uuuu
LVDC	--00 -000	--00 -000	--00 -000	--uu -uuu
USR	0000 1011	0000 1011	0000 1011	uuuu uuuu
UCR1	0000 00x0	0000 00x0	0000 00x0	uuuu uuuu
UCR2	0000 0000	0000 0000	0000 0000	uuuu uuuu
BRG	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu

Register	Reset (Power On)	WDT Time-out (Normal Operation)	LVR Reset	WDT Time-out (HALT)
TXR/RXR	xxxx xxxx	xxxx xxxx	xxxx xxxx	uuuu uuuu
SLCDC0	0000 0000	0000 0000	0000 0000	uuuu uuuu
SLCDC1	0000 0000	0000 0000	0000 0000	uuuu uuuu
SLCDC2	0000 0000	0000 0000	0000 0000	uuuu uuuu
SLCDC3	0000 0000	0000 0000	0000 0000	uuuu uuuu
SLCDC4	0000 0000	0000 0000	0000 0000	uuuu uuuu

Note: "-" not implement

"u" stands for "unchanged"

"x" stands for "unknown"

## Input/Output Ports

Holtek microcontrollers offer considerable flexibility on their I/O ports. With the input or output designation of every pin fully under user program control, pull-high selections for all ports and wake-up selections on certain pins, the user is provided with an I/O structure to meet the needs of a wide range of application possibilities.

The device provides bidirectional input/output lines labeled with port names PA~PD. These I/O ports are mapped to the RAM Data Memory with specific addresses as shown in the Special Purpose Data Memory table. All of these I/O ports can be used for input and output operations. For input operation, these ports are non-latching, which means the inputs must be ready at the T2 rising edge of instruction "MOV A, [m]", where m denotes the port address. For output operation, all the data is latched and remains unchanged until the output latch is rewritten.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PAWU	PAWU7	PAWU6	PAWU5	PAWU4	PAWU3	PAWU2	PAWU1	PAWU0
PAPU	PAPU7	PAPU6	PAPU5	PAPU4	PAPU3	PAPU2	PAPU1	PAPU0
PA	PA7	PA6	PA5	PA4	PA3	PA2	PA1	PA0
PAC	PAC7	PAC6	PAC5	PAC4	PAC3	PAC2	PAC1	PAC0
PBPU	PBPU7	PBPU6	PBPU5	PBPU4	PBPU3	PBPU2	PBPU1	PBPU0
PB	PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
PBC	PBC7	PBC6	PBC5	PBC4	PBC3	PBC2	PBC1	PBC0
PCPU	PCPU7	PCPU6	PCPU5	PCPU4	PCPU3	PCPU2	PCPU1	PCPU0
PC	PC7	PC6	PC5	PC4	PC3	PC2	PC1	PC0
PCC	PCC7	PCC6	PCC5	PCC4	PCC3	PCC2	PCC1	PCC0
PDPUn	—	—	PDPUn5	PDPUn4	PDPUn3	PDPUn2	PDPUn1	PDPUn0
PD	—	—	PD5	PD4	PD3	PD2	PD1	PD0
PDC	—	—	PDC5	PDC4	PDC3	PDC2	PDC1	PDC0

### I/O Control Register List

**PAPUn, PBPUUn, PCPUUn, PDPUn:** I/O Port Pull-High Control

0: Disable

1: Enable

**PAWUn:** I/O Port A bit 7 ~ bit 0 Wake Up Control

0: Disable

1: Enable

**PACn, PBCn, PCCn, PDCn:** I/O Port Input/Output Control

0: Output

1: Input

### **Pull-high Resistors**

Many product applications require pull-high resistors for their switch inputs usually requiring the use of an external resistor. To eliminate the need for these external resistors, all I/O pins, when configured as an input have the capability of being connected to an internal pull-high resistor. These pull-high resistors are selected using registers PAPU~PDPU, and are implemented using weak PMOS transistors.

### **Port A Wake-up**

The HALT instruction forces the microcontroller into the SLEEP or IDLE Mode which preserves power, a feature that is important for battery and other low-power applications. Various methods exist to wake-up the microcontroller, one of which is to change the logic condition on one of the Port A pins from high to low. This function is especially suitable for applications that can be woken up via external switches. Each pin on Port A can be selected individually to have this wake-up feature using the PAWU register.

### **I/O Port Control Registers**

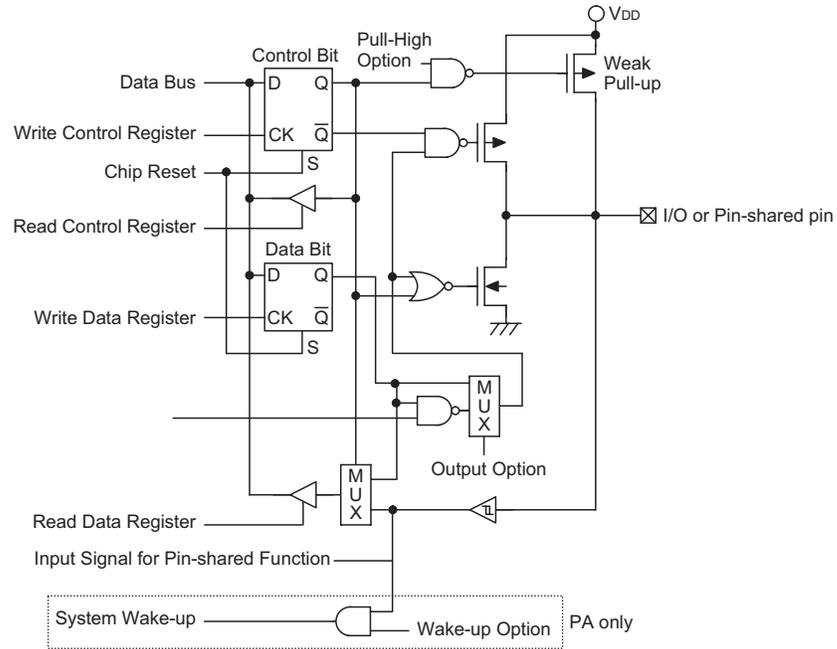
Each I/O port has its own control register known as PAC~PDC, to control the input/output configuration. With this control register, each CMOS output or input can be reconfigured dynamically under software control. Each pin of the I/O ports is directly mapped to a bit in its associated port control register. For the I/O pin to function as an input, the corresponding bit of the control register must be written as a "1". This will then allow the logic state of the input pin to be directly read by instructions. When the corresponding bit of the control register is written as a "0", the I/O pin will be setup as a CMOS output. If the pin is currently setup as an output, instructions can still be used to read the output register. However, it should be noted that the program will in fact only read the status of the output data latch and not the actual logic status of the output pin.

### **Pin-shared Functions**

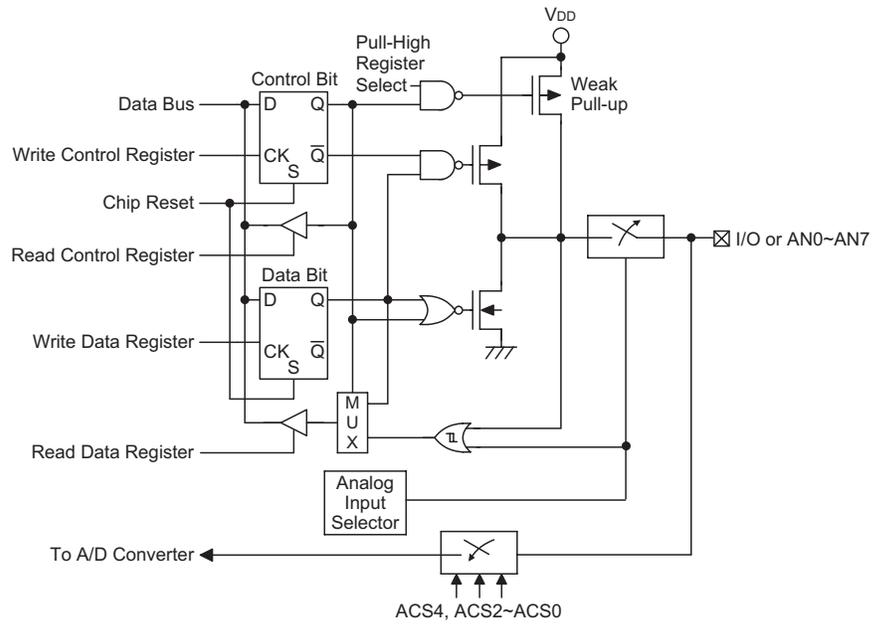
The flexibility of the microcontroller range is greatly enhanced by the use of pins that have more than one function. Limited numbers of pins can force serious design constraints on designers but by supplying pins with multi-functions, many of these difficulties can be overcome. For some pins, the chosen function of the multi-function I/O pins is set by configuration options while for others the function is set by application program control.

### **I/O Pin Structures**

The accompanying diagrams illustrate the internal structures of some generic I/O pin types. As the exact logical construction of the I/O pin will differ from these drawings, they are supplied as a guide only to assist with the functional understanding of the I/O pins. The wide range of pin-shared structures does not permit all types to be shown.



Generic Input/Output Structure



A/D Input/Output Structure

## I/O Port source current selection

The PA~PD input/output ports have programmable source current functions. Four levels of source current can be selected by the registers.

### IOHR0 Register

Bit	7	6	5	4	3	2	1	0
Name	IOHS31	IOHS30	IOHS21	IOHS20	IOHS11	IOHS10	IOHS01	IOHS00
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 **IOHS31~IOHS30**: PA3 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

Bit 5~4 **IOHS21~IOHS20**: PA2 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

Bit 3~2 **IOHS11~IOHS10**: PA1 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

Bit 1~0 **IOHS01~IOHS00**: PA0 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

### IOHR1 Register

Bit	7	6	5	4	3	2	1	0
Name	IOHS71	IOHS70	IOHS61	IOHS60	IOHS51	IOHS50	IOHS41	IOHS40
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~6 **IOHS71~IOHS70**: PA7 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

Bit 5~4 **IOHS61~IOHS60**: PA6 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

Bit 3~2 **IOHS51~IOHS50**: PA5 source current output selection bit

- 00: Full source current output
- 01: 10/22 full source current output
- 10: 7/22 full source current output
- 11: 4/22 full source current output

- Bit 1~0 **IOHS41~IOHS40**: PA4 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output

**IOHR2 Register**

Bit	7	6	5	4	3	2	1	0
Name	IOHSB1	IOHSB0	IOHSA1	IOHSA0	IOHS91	IOHS90	IOHS81	IOHS80
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7~6 **IOHSB1~IOHSB0**: PC4~PC7 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output
- Bit 5~4 **IOHSA1~IOHSA0**: PC0~PC3 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output
- Bit 3~2 **IOHS91~IOHS90**: PB4~PB7 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output
- Bit 1~0 **IOHS81~IOHS80**: PB0~PB3 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output

**IOHR3 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	IOHSD1	IOHSD0	IOHSC1	IOHSC0
R/W	—	—	—	—	R/W	R/W	R/W	R/W
POR	—	—	—	—	0	0	0	0

- Bit 7~4 Unimplemented, read as “0”
- Bit 3~2 **IOHSD1~IOHSD0**: PD4~PD5 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output
- Bit 1~0 **IOHSC1~IOHSC0**: PD0~PD3 source current output selection bit  
 00: Full source current output  
 01: 10/22 full source current output  
 10: 7/22 full source current output  
 11: 4/22 full source current output

## Programming Considerations

Within the user program, one of the first things to consider is port initialisation. After a reset, all of the I/O data and port control registers will be set high. This means that all I/O pins will default to an input state, the level of which depends on the other connected circuitry and whether pull-high selections have been chosen. If the port control registers, PAC~PDC, are then programmed to setup some pins as outputs, these output pins will have an initial high output value unless the associated port data registers, PA~PD, are first programmed. Selecting which pins are inputs and which are outputs can be achieved byte-wide by loading the correct values into the appropriate port control register or by programming individual bits in the port control register using the “SET [m].i” and “CLR [m].i” instructions. Note that when using these bit control instructions, a read-modify-write operation takes place. The microcontroller must first read in the data on the entire port, modify it to the required new bit values and then rewrite this data back to the output ports.

Port A has the additional capability of providing wake-up functions. When the device is in the SLEEP or IDLE Mode, various methods are available to wake the device up. One of these is a high to low transition of any of the Port A pins. Single or multiple pins on Port A can be setup to have this function.

## Timer Modules – TM

One of the most fundamental functions in any microcontroller device is the ability to control and measure time. To implement time related functions the device includes several Timer Modules, abbreviated to the name TM. The TMs are multi-purpose timing units and serve to provide operations such as Timer/Counter, Input Capture, Compare Match Output and Single Pulse Output as well as being the functional unit for the generation of PWM signals. Each of the TMs has two individual interrupts. The addition of input and output pins for each TM ensures that users are provided with timing units with a wide and flexible range of features.

The common features of the different TM types are described here with more detailed information provided in the individual Compact, Standard and Periodic TM sections.

### Introduction

The device contains a 10-bit Compact TM-CTM, a 16-bit Standard TM-STM and two 10-bit Periodic TMs- PTM0 and PTM1. Although similar in nature, the different TM types vary in their feature complexity. The common features to the Compact, Standard and Periodic TMs will be described in this section and the detailed operation will be described in corresponding sections. The main features and differences between the three types of TMs are summarised in the accompanying table.

Function	CTM	STM	PTM
Timer/Counter	√	√	√
I/P Capture	—	√	√
Compare Match Output	√	√	√
PWM Channels	1	1	1
Single Pulse Output	—	1	1
PWM Alignment	Edge	Edge	Edge
PWM Adjustment Period & Duty	Duty or Period	Duty or Period	Duty or Period

**TM Function Summary**

## TM Operation

The three different types of TMs offer a diverse range of functions, from simple timing operations to PWM signal generation. The key to understanding how the TM operates is to see it in terms of a free running counter whose value is then compared with the value of pre-programmed internal comparators. When the free running counter has the same value as the pre-programmed comparator, known as a compare match situation, a TM interrupt signal will be generated which can clear the counter and perhaps also change the condition of the TM output pin. The internal TM counter is driven by a user selectable clock source, which can be an internal clock or an external pin.

## TM Clock Source

The clock source which drives the main counter in each TM can originate from various sources. The selection of the required clock source is implemented using the xTnCK2~xTnCK0 bits in the TM control registers. The clock source can be a ratio of either the system clock  $f_{SYS}$  or the internal high clock  $f_H$ , the  $f_{TBC}$  clock source or the external xTCKn pin. The xTCKn pin clock source is used to allow an external signal to drive the TM as an external clock source or for event counting.

## TM Interrupts

The Compact, Standard and Periodic type TMs each has two internal interrupts, the internal comparator A or comparator P, which generate a TM interrupt when a compare match condition occurs. When a TM interrupt is generated, it can be used to clear the counter and also to change the state of the TM output pin.

## TM External Pins

Each of the TMs, irrespective of what type, has one TM input pin, with the label xTCKn. The TM input pin, is essentially a clock source for the TM and is selected using the xTnCK2~xTnCK0 bits in the xTMnCO register. This external TM input pin allows an external clock source to drive the internal TM. This external TM input pin is shared with other functions but will be connected to the internal TM if selected using the xTnCK2~xTnCK0 bits. The TM input pin can be chosen to have either a rising or falling active edge.

The TMs each have one output pin. When the TM is in the Compare Match Output Mode, these pins can be controlled by the TM to switch to a high or low level or to toggle when a compare match situation occurs. The external xTPn output pin is also the pin where the TM generates the PWM output waveform. As the TM output pins are pin-shared with other function, the TM output function must first be setup using registers. A single bit in one of the registers determines if its associated pin is to be used as an external TM output pin or if it is to have another function. The number of output pins for each TM type is different, the details are provided in the accompanying table.

TM Type	Pin	Input	Output
CTM	CTM	CTCK	CTP
STM	STM	STCK	STP
PTM	PTM0	PTCK0	PTP0
	PTM1	PTCK1	PTP1

**TM Input/Output pin**

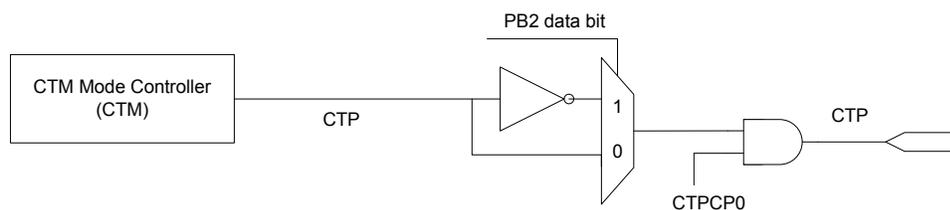
### TM Input/Output Pin Control Register

Selecting to have a TM input/output or whether to retain its other shared function is implemented using one register, with a single bit in the register corresponding to a TM input/output pin. Setting the bit high will setup the corresponding pin as a TM input/output, if reset to zero the pin will retain its original other function.

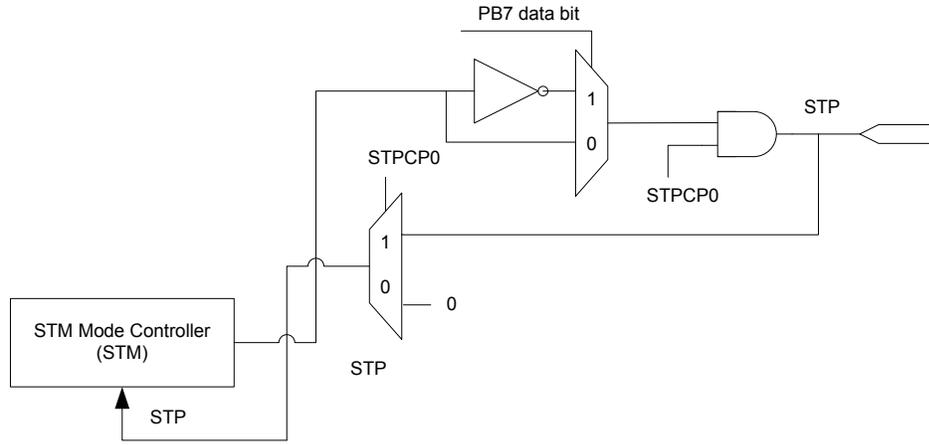
#### TMPC0 Register

Bit	7	6	5	4	3	2	1	0
Name	—	PTP1CP0	—	PTP0CP0	—	STPCP0	—	CTPCP0
R/W	—	R/W	—	R/W	—	R/W	—	R/W
POR	—	0	—	0	—	0	—	0

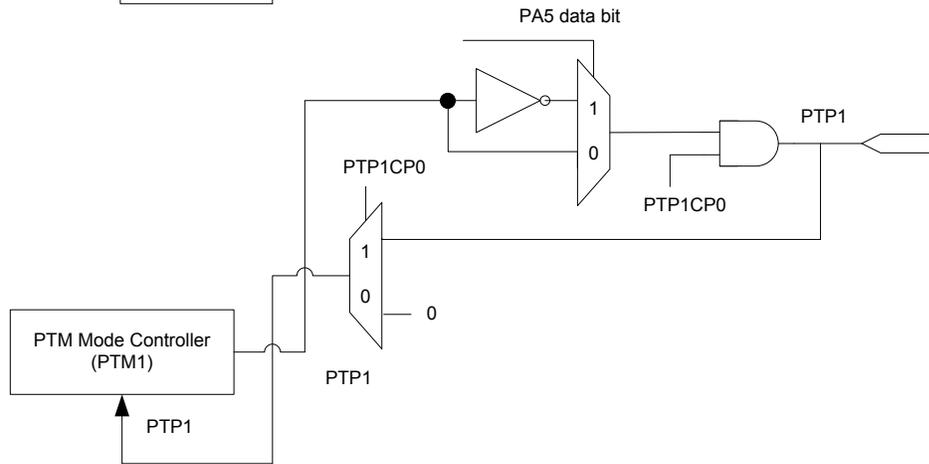
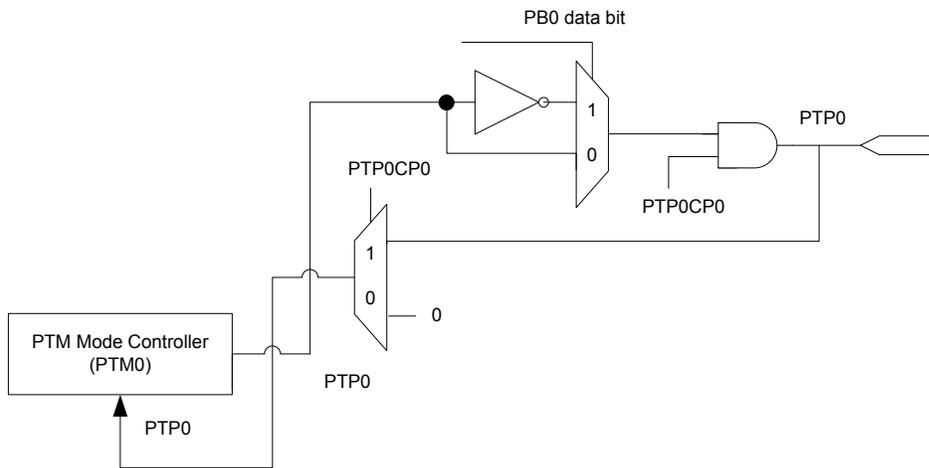
- Bit 7 Unimplemented, read as
- Bit 6 **PTP1CP0**: PTP1 pin Control  
 0: Normal I/O  
 1: PTP1 function
- Bit 5 Unimplemented, read as
- Bit 4 **PTP0CP0**: PTP0 pin Control  
 0: Normal I/O  
 1: PTP0 function
- Bit 3 Unimplemented, read as
- Bit 2 **STPCP0**: STP pin Control  
 0: Normal I/O  
 1: STP function
- Bit 1 Unimplemented, read as
- Bit 0 **CTPCP0**: CTP pin Control  
 0: Normal I/O  
 1: CTP function



**CTM Function Pin Control Block Diagram**



STM Function Pin Control Block Diagram

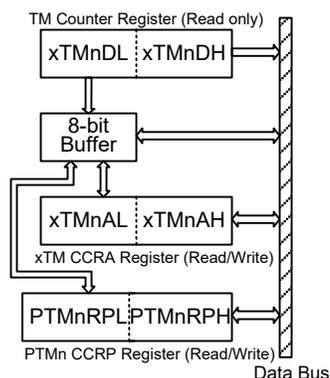


PTM0, PTM1 Function Pin Control Block Diagram

## Programming Considerations

The TM Counter Registers and the Capture/Compare CCRA or CCRP register, being either 10-bit or 16-bit, all have a low and high byte structure. The high bytes can be directly accessed, but as the low bytes can only be accessed via an internal 8-bit buffer, reading or writing to these register pairs must be carried out in a specific way. The important point to note is that data transfer to and from the 8-bit buffer and its related low byte only takes place when a write or read operation to its corresponding high byte is executed.

As the CCRA and CCRP registers are implemented in the way shown in the following diagram and accessing this register is carried out in a specific way described above, it is recommended to use the “MOV” instruction to access the CCRA or CCRP low byte register, named xTMnAL or PTMnRPL, in the following access procedures. Accessing the CCRA or CCRP low byte register without following these access procedures will result in unpredictable values.

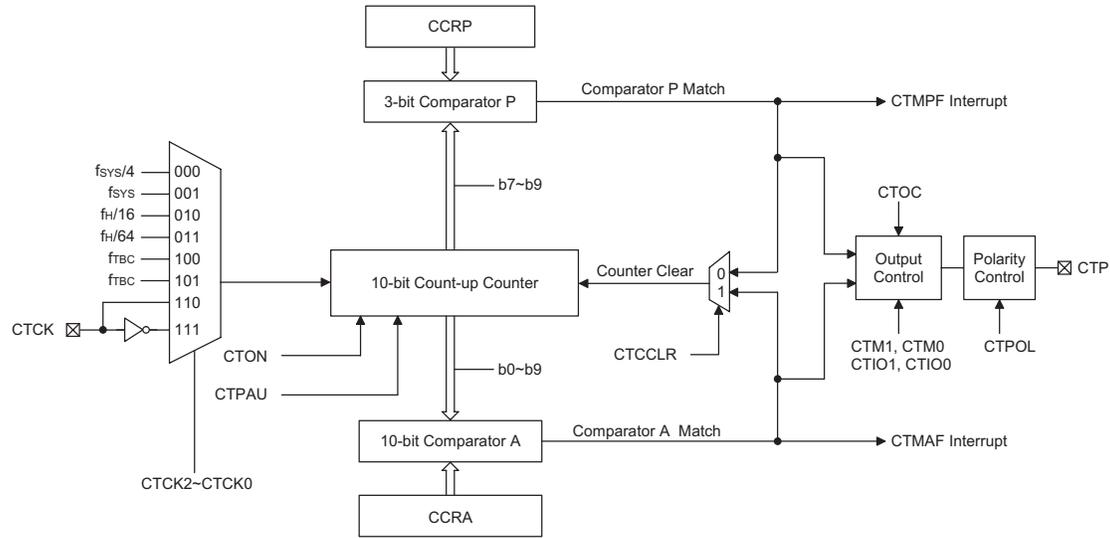


The following steps show the read and write procedures:

- Writing Data to CCRA or CCRP
  - ♦ Step 1. Write data to Low Byte xTMnAL or PTMnRPL
    - Note that here data is only written to the 8-bit buffer.
  - ♦ Step 2. Write data to High Byte xTMnAH or PTMnRPH
    - Here data is written directly to the high byte registers and simultaneously data is latched from the 8-bit buffer to the Low Byte registers.
- Reading Data from the Counter Registers and CCRA or CCRP
  - ♦ Step 1. Read data from the High Byte xTMnDH, xTMnAH or PTMnRPH
    - Here data is read directly from the High Byte registers and simultaneously data is latched from the Low Byte register into the 8-bit buffer.
  - ♦ Step 2. Read data from the Low Byte xTMnDL, xTMnAL or PTMnRPL
    - This step reads data from the 8-bit buffer.

## Compact Type TM – CTM

Although the simplest form of the three TM types, the Compact TM type still contains three operating modes, which are Compare Match Output, Timer/Event Counter and PWM Output modes. The Compact TM can also be controlled with an external input pin and can drive one external output pin.



**Compact Type TM Block Diagram**

### Compact TM Operation

At its core is a 10-bit count-up counter which is driven by a user selectable internal or external clock source. There are also two internal comparators with the names, Comparator A and Comparator P. These comparators will compare the value in the counter with CCRP and CCRA registers. The CCRP is three bits wide whose value is compared with the highest three bits in the counter while the CCRA is the ten bits and therefore compares with all counter bits.

The only way of changing the value of the 10-bit counter using the application program, is to clear the counter by changing the CTON bit from low to high. The counter will also be cleared automatically by a counter overflow or a compare match with one of its associated comparators. When these conditions occur, a TM interrupt signal will also usually be generated. The Compact Type TM can operate in a number of different operational modes, can be driven by different clock sources including an input pin and can also control two output pins. All operating setup conditions are selected using relevant internal registers.

## Compact Type TM Register Description

Overall operation of each Compact TM is controlled using several registers. A read only register pair exists to store the internal counter 10-bit value, while a read/write register pair exists to store the internal 10-bit CCRA value. The remaining two registers are control registers which setup the different operating and control modes as well as the three CCRP bits.

Register Name	Bit							
	7	6	5	4	3	2	1	0
CTMC0	CTPAU	CTCK2	CTCK1	CTCK0	CTON	CTRP2	CTRP1	CTRP0
CTMC1	CTM1	CTM0	CTIO1	CTIO0	CTOC	CTPOL	CTDPX	CTCCLR
CTMDL	D7	D6	D5	D4	D3	D2	D1	D0
CTMDH	—	—	—	—	—	—	D9	D8
CTMAL	D7	D6	D5	D4	D3	D2	D1	D0
CTMAH	—	—	—	—	—	—	D9	D8

**Compact TM Register List**

### CTMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	CTPAU	CTCK2	CTCK1	CTCK0	CTON	CTRP2	CTRP1	CTRP0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **CTPAU**: CTM Counter Pause Control

0: Run  
 1: Pause

The counter can be paused by setting this bit high. Clearing the bit to zero restores normal counter operation. When in a Pause condition the TM will remain powered up and continue to consume power. The counter will retain its residual value when this bit changes from low to high and resume counting from this value when the bit changes to a low value again.

Bit 6 ~ 4 **CTCK2 ~ CTCK0**: Select CTM Counter clock

000:  $f_{SYS}/4$   
 001:  $f_{SYS}$   
 010:  $f_H/16$   
 011:  $f_H/64$   
 100:  $f_{TBC}$   
 101:  $f_{TBC}$   
 110: CTCK rising edge clock  
 111: CTCK falling edge clock

These three bits are used to select the clock source for the CTM. The external pin clock source can be chosen to be active on the rising or falling edge. The clock source  $f_{SYS}$  is the system clock, while  $f_H$  and  $f_{TBC}$  are other internal clocks, the details of which can be found in the oscillator section.

Bit 3 **CTON**: CTM Counter On/Off Control

0: Off  
 1: On

This bit controls the overall on/off function of the TM. Setting the bit high enables the counter to run, clearing the bit disables the TM. Clearing this bit to zero will stop the counter from counting and turn off the TM which will reduce its power consumption. When the bit changes state from low to high the internal counter value will be reset to zero, however when the bit changes from high to low, the internal counter will retain its residual value until the bit returns high again.

If the TM is in the Compare Match Output Mode then the TM output pin will be reset to its initial condition, as specified by the CTOC bit, when the CTON bit changes from low to high.

Bit 2 ~ 0 **CTRP2~CTRP0**: CTM CCRP 3-bit register, compared with the CTM Counter bit9~bit7  
 Comparator P Match Period  
 000: 1024 CTM clocks  
 001: 128 CTM clocks  
 010: 256 CTM clocks  
 011: 384 CTM clocks  
 100: 512 CTM clocks  
 101: 640 CTM clocks  
 110: 768 CTM clocks  
 111: 896 CTM clocks

These three bits are used to setup the value on the internal CCRP 3-bit register, which are then compared with the internal counter's highest three bits. The result of this comparison can be selected to clear the internal counter if the CTCCLR bit is set to zero. Setting the CTCCLR bit to zero ensures that a compare match with the CCRP values will reset the internal counter. As the CCRP bits are only compared with the highest three counter bits, the compare values exist in 128 clock cycle multiples. Clearing all three bits to zero is in effect allowing the counter to overflow at its maximum value.

**CTMC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	CTM1	CTM0	CTIO1	CTIO0	CTOC	CTPOL	CTDPX	CTCCLR
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 6 **CTM1~CTM0**: Select CTM Operating Mode  
 00: Compare Match Output Mode  
 01: Compare Match Output Mode  
 10: PWM Mode  
 11: Timer/Counter Mode

These bits setup the required operating mode for the TM. To ensure reliable operation the TM should be switched off before any changes are made to the bits. In the Timer/Counter Mode, the TM output pin control must be disabled.

Bit 5 ~ 4 **CTIO1~CTIO0**: Select CTM output function  
 Compare Match Output Mode  
 00: No change  
 01: Output low  
 10: Output high  
 11: Toggle output  
 PWM Mode  
 00: PWM output inactive state  
 01: PWM output active state  
 10: PWM output  
 11: Undefined  
 Timer/counter Mode  
 Unused

These two bits are used to determine how the TM output pin changes state when a certain condition is reached. The function that these bits select depends upon in which mode the TM is running.

In the Compare Match Output Mode, the CTIO1~CTIO0 bits determine how the TM output pin changes state when a compare match occurs from the Comparator A. The TM output pin can be setup to switch high, switch low or to toggle its present state when a compare match occurs from the Comparator A. When the CTIO1~CTIO0 bits are both zero, then no change will take place on the output. The initial value of the TM output pin should be setup using the CTOC bit. Note that the output level requested by the CTIO1~CTIO0 bits must be different from the initial value setup using the CTOC bit otherwise no change will occur on the TM output pin when a compare match occurs. After the TM output pin changes state it can be reset to its initial level by changing the level of the CTON bit from low to high.

In the PWM Mode, the CTIO1 and CTIO0 bits determine how the TM output pin changes state when a certain compare match condition occurs. The PWM output function is modified by changing these two bits. It is necessary to change the values of the CTIO1 and CTIO0 bits only after the TM has been switched off. Unpredictable PWM outputs will occur if the CTIO1 and CTIO0 bits are changed when the TM is running.

Bit 3 **CTOC**: CTM Output control bit

Compare Match Output Mode

0: Initial low

1: Initial high

PWM Mode

0: Active low

1: Active high

This is the output control bit for the TM output pin. Its operation depends upon whether TM is being used in the Compare Match Output Mode or in the PWM Mode. It has no effect if the TM is in the Timer/Counter Mode. In the Compare Match Output Mode it determines the logic level of the TM output pin before a compare match occurs. In the PWM Mode it determines if the PWM signal is active high or active low.

Bit 2 **CTPOL**: CTM Output polarity Control

0: Non-invert

1: Invert

This bit controls the polarity of the TM output pin. When the bit is set high the TM output pin will be inverted and not inverted when the bit is zero. It has no effect if the TM is in the Timer/Counter Mode.

Bit 1 **CTDPX**: CTM PWM period/duty Control

0: CCRP - period; CCRA - duty

1: CCRP - duty; CCRA - period

This bit, determines which of the CCRA and CCRP registers are used for period and duty control of the PWM waveform.

Bit 0 **CTCCLR**: Select CTM Counter clear condition

0: CTM Comparatror P match

1: CTM Comparatror A match

This bit is used to select the method which clears the counter. Remember that the Compact TM contains two comparators, Comparator A and Comparator P, either of which can be selected to clear the internal counter. With the CTCCLR bit set high, the counter will be cleared when a compare match occurs from the Comparator A. When the bit is low, the counter will be cleared when a compare match occurs from the Comparator P or with a counter overflow. A counter overflow clearing method can only be implemented if the CCRP bits are all cleared to zero. The CTCCLR bit is not used in the PWM Mode.

**CTMDL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0     **D7~D0**: CTM Counter Low Byte Register bit 7 ~ bit 0  
 CTM 10-bit Counter bit 7 ~ bit 0

**CTMDH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R	R
POR	—	—	—	—	—	—	0	0

Bit 7 ~ 2     Unimplemented, read as “0”  
 Bit 1 ~ 0     **D9~D8**: CTM Counter High Byte Register bit 1 ~ bit 0  
 CTM 10-bit Counter bit 9 ~ bit 8

**CTMAL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0     **D7~D0**: CTM CCRA Low Byte Register bit 7 ~ bit 0  
 CTM 10-bit CCRA bit 7 ~ bit 0

**CTMAH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7 ~ 2     Unimplemented, read as “0”  
 Bit 1 ~ 0     **D9~D8**: CTM CCRA High Byte Register bit 1 ~ bit 0  
 CTM 10-bit CCRA bit 9 ~ bit 8

## **Compact Type TM Operating Modes**

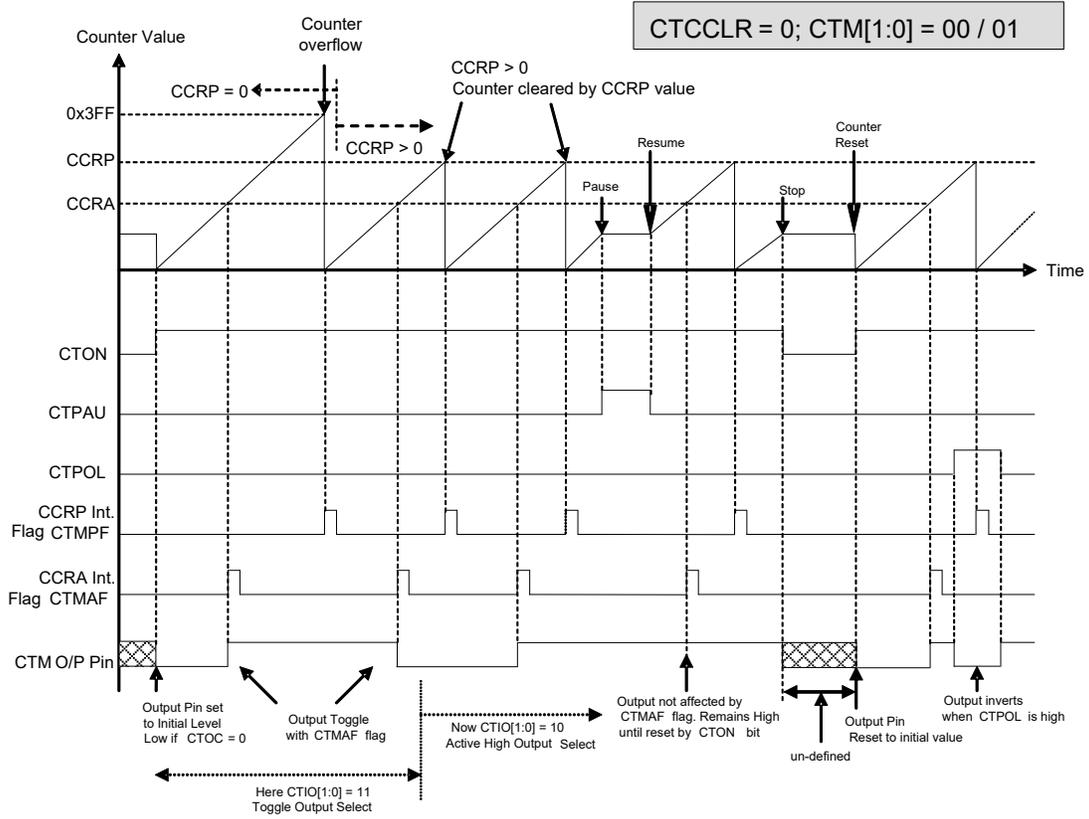
The Compact Type TM can operate in one of three operating modes, Compare Match Output Mode, PWM Output Mode or Timer/Counter Mode. The operating mode is selected using the CTM1 and CTM0 bits in the CTMC1 register.

### **Compare Match Output Mode**

To select this mode, bits CTM1 and CTM0 in the CTMC1 register, should be set to 00 or 01 respectively. In this mode once the counter is enabled and running it can be cleared by three methods. These are a counter overflow, a compare match from Comparator A and a compare match from Comparator P. When the CTCCLR bit is low, there are two ways in which the counter can be cleared. One is when a compare match from Comparator P, the other is when the CCRP bits are all zero which allows the counter to overflow. Here both CTMAF and CTMPF interrupt request flags for Comparator A and Comparator P respectively, will both be generated.

If the CTCCLR bit in the CTMC1 register is high then the counter will be cleared when a compare match occurs from Comparator A. However, here only the CTMAF interrupt request flag will be generated even if the value of the CCRP bits is less than that of the CCRA registers. Therefore when CTCCLR is high no CTMPF interrupt request flag will be generated. If the CCRA bits are all zero, the counter will overflow when it reaches its maximum 10-bit, 3FF Hex, value, however here the CTMAF interrupt request flag will not be generated.

As the name of the mode suggests, after a comparison is made, the TM output pin, will change state. The TM output pin condition however only changes state when a CTMAF interrupt request flag is generated after a compare match occurs from Comparator A. The CTMPF interrupt request flag, generated from a compare match occurs from Comparator P, will have no effect on the TM output pin. The way in which the TM output pin changes state are determined by the condition of the CTIO1 and CTIO0 bits in the CTMC1 register. The TM output pin can be selected using the CTIO1 and CTIO0 bits to go high, to go low or to toggle from its present condition when a compare match occurs from Comparator A. The initial condition of the TM output pin, which is setup after the CTON bit changes from low to high, is setup using the CTOC bit. Note that if the CTIO1 and CTIO0 bits are zero then no pin change will take place.



**Compare Match Output Mode – CTCCLR = 0**

- Note: 1. With CTCCLR = 0, a Comparator P match will clear the counter  
 2. The TM output pin controlled only by the CTMAF flag  
 3. The output pin reset to initial state by a CTON bit rising edge



**Timer/Counter Mode**

To select this mode, bits CTM1 and CTM0 in the CTMC1 register should be set to 11 respectively. The Timer/Counter Mode operates in an identical way to the Compare Match Output Mode generating the same interrupt flags. The exception is that in the Timer/Counter Mode the TM output pin is not used. Therefore the above description and Timing Diagrams for the Compare Match Output Mode can be used to understand its function. As the TM output pin is not used in this mode, the pin can be used as a normal I/O pin or other pin-shared function.

**PWM Output Mode**

To select this mode, bits CTM1 and CTM0 in the CTMC1 register should be set to 10 respectively. The PWM function within the TM is useful for applications which require functions such as motor control, heating control, illumination control etc. By providing a signal of fixed frequency but of varying duty cycle on the TM output pin, a square wave AC waveform can be generated with varying equivalent DC RMS values.

As both the period and duty cycle of the PWM waveform can be controlled, the choice of generated waveform is extremely flexible. In the PWM mode, the CTCCLR bit has no effect as the PWM period. Both of the CCRA and CCRP registers are used to generate the PWM waveform, one register is used to clear the internal counter and thus control the PWM waveform frequency, while the other one is used to control the duty cycle. Which register is used to control either frequency or duty cycle is determined using the CTDPX bit in the CTMC1 register. The PWM waveform frequency and duty cycle can therefore be controlled by the values in the CCRA and CCRP registers.

An interrupt flag, one for each of the CCRA and CCRP, will be generated when a compare match occurs from either Comparator A or Comparator P. The CTOC bit In the CTMC1 register is used to select the required polarity of the PWM waveform while the two CTIO1 and CTIO0 bits are used to enable the PWM output or to force the TM output pin to a fixed high or low level. The CTPOL bit is used to reverse the polarity of the PWM output waveform.

• **CTM, PWM Mode, Edge-aligned Mode, CTDPX=0**

CCRP	001b	010b	011b	100b	101b	110b	111b	000b
Period	128	256	384	512	640	768	896	1024
Duty	CCRA							

If  $f_{SYS} = 16\text{MHz}$ , TM clock source is  $f_{SYS}/4$ , CCRP = 100b, CCRA =128,

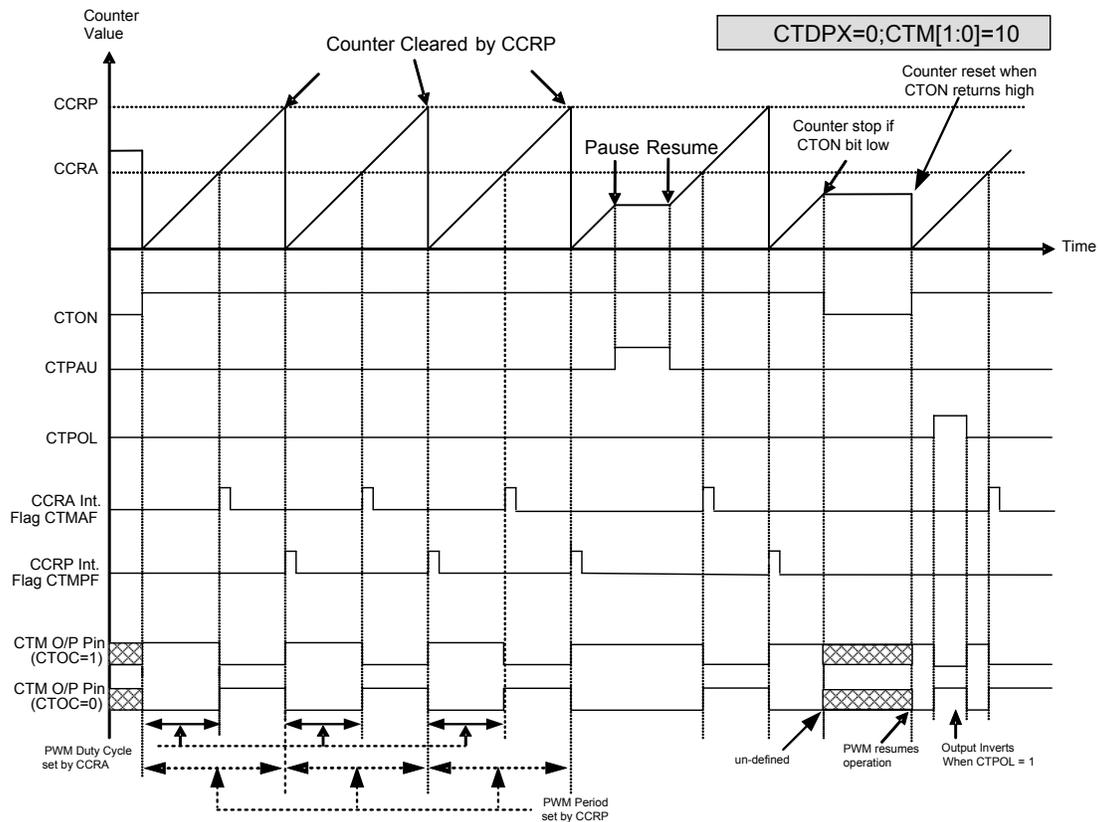
The CTM PWM output frequency =  $(f_{SYS}/4) / 512 = f_{SYS}/2048 = 7.8125 \text{ kHz}$ , duty =  $128/512 = 25\%$ .

If the Duty value defined by the CCRA register is equal to or greater than the Period value, then the PWM output duty is 100%.

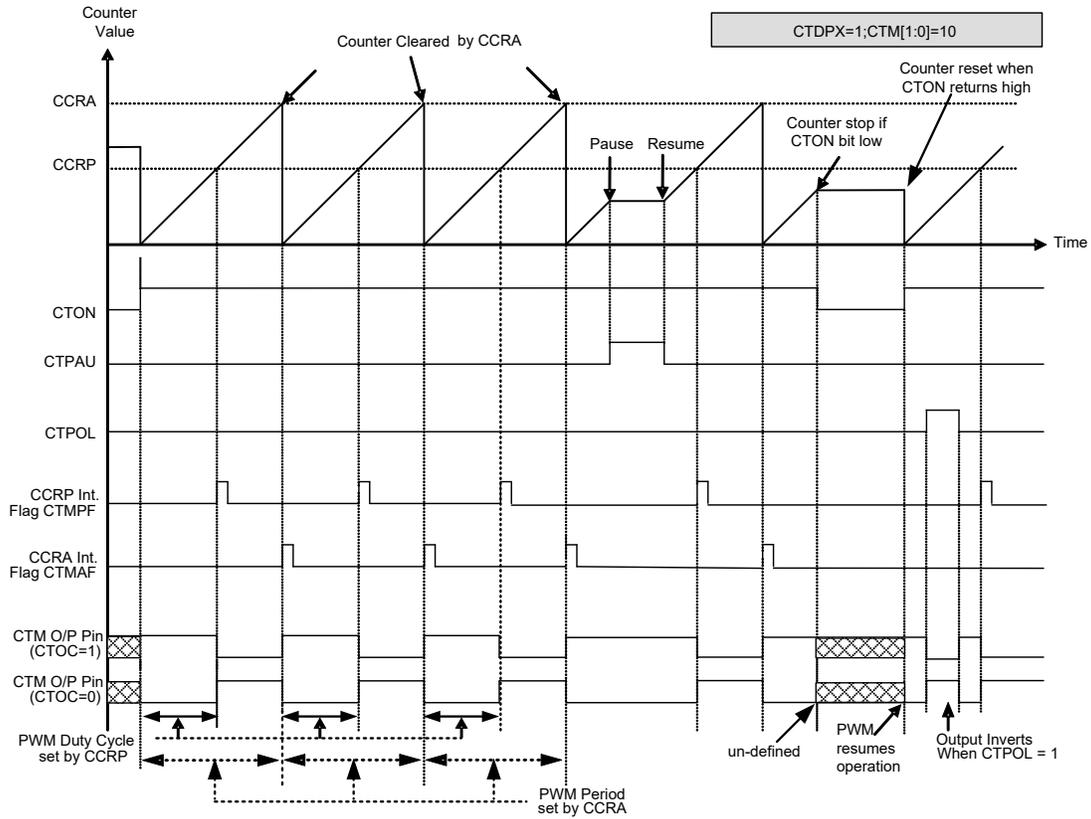
• **CTM, PWM Mode, Edge-aligned Mode, CTDPX=1**

CCRP	001b	010b	011b	100b	101b	110b	111b	000b
Period	CCRA							
Duty	128	256	384	512	640	768	896	1024

The PWM output period is determined by the CCRA register value together with the TM clock while the PWM duty cycle is defined by the CCRP register value.



- Note: 1. Here CTD PX = 0 - Counter cleared by CCRP  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues running even when CTIO[1:0] = 00 or 01  
 4. The CTCCLR bit has no influence on PWM operation



**PWM Mode – CTDPX = 1**

- Note: 1. Here CTDPX = 1 - Counter cleared by CCRA  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues even when CTIO[1:0] = 00 or 01  
 4. The CTCCLR bit has no influence on PWM operation



## Standard Type TM Register Description

Overall operation of the Standard TM is controlled using series of registers. A read only register pair exists to store the internal counter 16-bit value, while a read/write register pair exists to store the internal 16-bit CCRA value. The remaining two registers are control registers which setup the different operating and control modes as well as eight CCRP bits.

Register Name	Bit							
	7	6	5	4	3	2	1	0
STMC0	STPAU	STCK2	STCK1	STCK0	STON	—	—	—
STMC1	STM1	STM0	STIO1	STIO0	STOC	STPOL	STDPX	STCCLR
STMDL	D7	D6	D5	D4	D3	D2	D1	D0
STMDH	D15	D14	D13	D12	D11	D10	D9	D8
STMAL	D7	D6	D5	D4	D3	D2	D1	D0
STMAH	D15	D14	D13	D12	D11	D10	D9	D8
STMRP	D7	D6	D5	D4	D3	D2	D1	D0

16-bit Standard TM Register List

### STMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	STPAU	STCK2	STCK1	STCK0	STON	—	—	—
R/W	R/W	R/W	R/W	R/W	R/W	—	—	—
POR	0	0	0	0	0	—	—	—

Bit 7 **STPAU**: STM Counter Pause Control  
 0: Run  
 1: Pause

The counter can be paused by setting this bit high. Clearing the bit to zero restores normal counter operation. When in a Pause condition the TM will remain powered up and continue to consume power. The counter will retain its residual value when this bit changes from low to high and resume counting from this value when the bit changes to a low value again.

Bit 6 ~ 4 **STCK2 ~ STCK0**: Select STM Counter clock  
 000:  $f_{SYS}/4$   
 001:  $f_{SYS}$   
 010:  $f_H/16$   
 011:  $f_H/64$   
 100:  $f_{TBC}$   
 101:  $f_{TBC}$   
 110: STCK rising edge clock  
 111: STCK falling edge clock

These three bits are used to select the clock source for the TM. The external pin clock source can be chosen to be active on the rising or falling edge. The clock source  $f_{SYS}$  is the system clock, while  $f_H$  and  $f_{TBC}$  are other internal clocks, the details of which can be found in the oscillator section.

Bit 3 **STON**: STM Counter On/Off Control  
 0: Off  
 1: On

This bit controls the overall on/off function of the TM. Setting the bit high enables the counter to run, clearing the bit disables the TM. Clearing this bit to zero will stop the counter from counting and turn off the TM which will reduce its power consumption. When the bit changes state from low to high the internal counter value will be reset to zero, however when the bit changes from high to low, the internal counter will retain its residual value until the bit returns high again.

If the TM is in the Compare Match Output Mode then the TM output pin will be reset to its initial condition, as specified by the STOC bit, when the STON bit changes from low to high.

Bit 2 ~ 0 Unimplemented, read as “0”

**STMC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	STM1	STM0	STIO1	STIO0	STOC	STPOL	STDPX	STCCLR
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 6 **STM1~STM0**: Select STM Operating Mode  
 00: Compare Match Output Mode  
 01: Capture Input Mode  
 10: PWM Mode or Single Pulse Output Mode  
 11: Timer/Counter Mode

These bits setup the required operating mode for the TM. To ensure reliable operation the TM should be switched off before any changes are made to the bits. In the Timer/Counter Mode, the TM output pin control must be disabled.

Bit 5 ~ 4 **STIO1~STIO0**: Select STM output function

Compare Match Output Mode

- 00: No change
- 01: Output low
- 10: Output high
- 11: Toggle output

PWM Mode/Single Pulse Output Mode

- 00: PWM output inactive state
- 01: PWM output active state
- 10: PWM output
- 11: Single pulse output

Capture Input Mode

- 00: Input capture at rising edge of STP
- 01: Input capture at falling edge of STP
- 10: Input capture at falling/rising edge of STP
- 11: Input capture disabled

Timer/counter Mode

Unused

These two bits are used to determine how the TM output pin changes state when a certain condition is reached. The function that these bits select depends upon in which mode the TM is running.

In the Compare Match Output Mode, the STIO1~STIO0 bits determine how the TM output pin changes state when a compare match occurs from the Comparator A. The TM output pin can be setup to switch high, switch low or to toggle its present state when a compare match occurs from the Comparator A. When the STIO1~STIO0 bits are both zero, then no change will take place on the output. The initial value of the TM output pin should be setup using the STOC bit. Note that the output level requested by the STIO1~STIO0 bits must be different from the initial value setup using the STOC bit otherwise no change will occur on the TM output pin when a compare match occurs. After the TM output pin changes state it can be reset to its initial level by changing the level of the STON bit from low to high.

In the PWM Mode, the STIO1 and STIO0 bits determine how the TM output pin changes state when a certain compare match condition occurs. The PWM output function is modified by changing these two bits. It is necessary to change the values of the STIO1 and STIO0 bits only after the TM has been switched off. Unpredictable PWM outputs will occur if the STIO1 and STIO0 bits are changed when the TM is running.

Bit 3 **STOC**: STM Output control bit

Compare Match Output Mode

0: Initial low

1: Initial high

PWM Mode/Single Pulse Output Mode

0: Active low

1: Active high

This is the output control bit for the TM output pin. Its operation depends upon whether TM is being used in the Compare Match Output Mode or in the PWM Mode/Single Pulse Output Mode. It has no effect if the TM is in the Timer/Counter Mode. In the Compare Match Output Mode it determines the logic level of the TM output pin before a compare match occurs. In the PWM Mode it determines if the PWM signal is active high or active low.

Bit 2 **STPOL**: STM Output polarity Control

0: Non-invert

1: Invert

This bit controls the polarity of the TM output pin. When the bit is set high the TM output pin will be inverted and not inverted when the bit is zero. It has no effect if the TM is in the Timer/Counter Mode.

Bit 1 **STDPX**: STM PWM period/duty Control

0: CCRP - period; CCRA - duty

1: CCRP - duty; CCRA - period

This bit, determines which of the CCRA and CCRP registers are used for period and duty control of the PWM waveform.

Bit 0 **STCCLR**: Select STM Counter clear condition

0: TM Comparator P match

1: TM Comparator A match

This bit is used to select the method which clears the counter. Remember that the Standard TM contains two comparators, Comparator A and Comparator P, either of which can be selected to clear the internal counter. With the STCCLR bit set high, the counter will be cleared when a compare match occurs from the Comparator A. When the bit is low, the counter will be cleared when a compare match occurs from the Comparator P or with a counter overflow. A counter overflow clearing method can only be implemented if the CCRP bits are all cleared to zero. The STCCLR bit is not used in the PWM, Single Pulse or Input Capture Mode.

**STMDL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0    **D7~D0:** STM Counter Low Byte Register bit 7 ~ bit 0  
 STM 16-bit Counter bit 7 ~ bit 0

**STMDH Register**

Bit	7	6	5	4	3	2	1	0
Name	D15	D14	D13	D12	D11	D10	D9	D8
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0    **D15~D8:** STM Counter High Byte Register bit 7 ~ bit 0  
 STM 16-bit Counter bit 15 ~ bit 8

**STMAL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0    **D7~D0:** STM CCRA Low Byte Register bit 7 ~ bit 0  
 STM 16-bit CCRA bit 7 ~ bit 0

**STMAH Register**

Bit	7	6	5	4	3	2	1	0
Name	D15	D14	D13	D12	D11	D10	D9	D8
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0    **D15~D8:** STM CCRA High Byte Register bit 7 ~ bit 0  
 STM 16-bit CCRA bit 15 ~ bit 8

**STM RP Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 ~ 0    **STM RP:** STM CCRP High Byte Register bit 7 ~ bit 0  
 STM CCRP 8-bit register, compared with the STM Counter bit 15 ~ bit 8. Comparator P Match Period  
 0: 65536 STM clocks  
 1~255: 256 × (1~255) STM clocks

These eight bits are used to setup the value on the internal CCRP 8-bit register, which are then compared with the internal counter's highest eight bits. The result of this comparison can be selected to clear the internal counter if the STCCLR bit is set to zero. Setting the STCCLR bit to zero ensures that a compare match with the CCRP values will reset the internal counter. As the CCRP bits are only compared with the highest eight counter bits, the compare values exist in 256 clock cycle multiples. Clearing all eight bits to zero is in effect allowing the counter to overflow at its maximum value.

## Standard Type TM Operating Modes

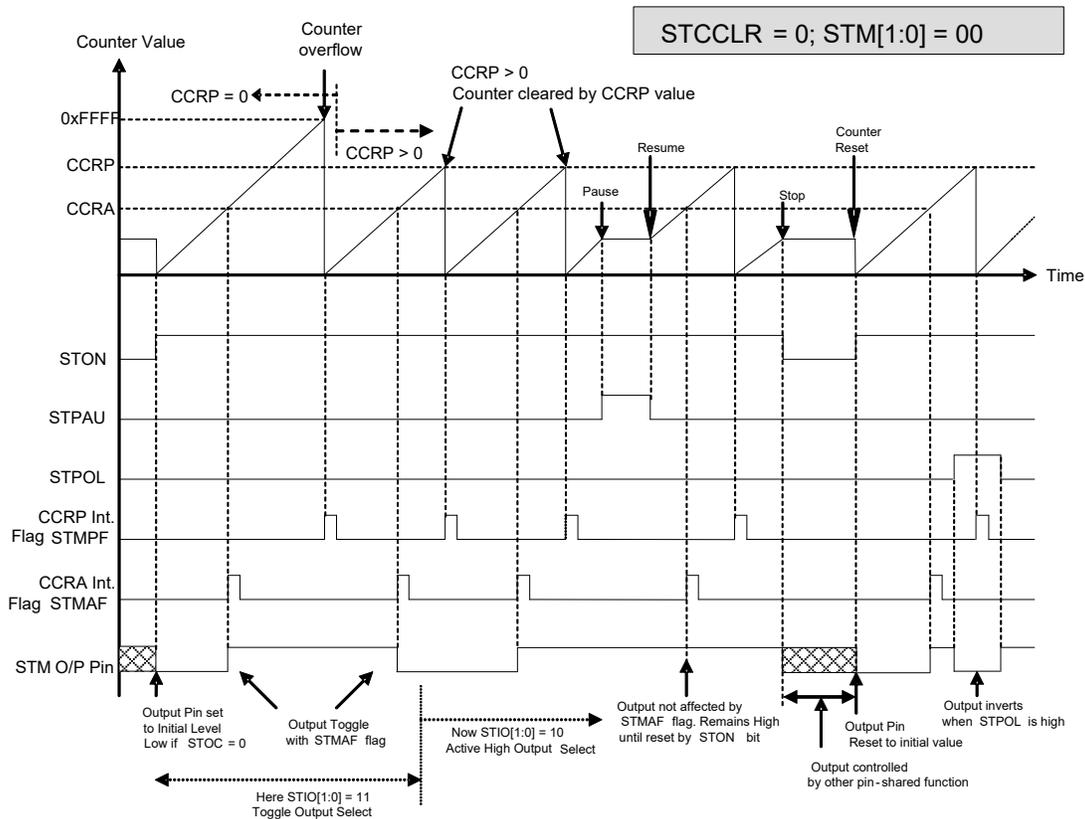
The Standard Type TM can operate in one of five operating modes, Compare Match Output Mode, PWM Output Mode, Single Pulse Output Mode, Capture Input Mode or Timer/Counter Mode. The operating mode is selected using the STM1 and STM0 bits in the STMC1 register.

### Compare Match Output Mode

To select this mode, bits STM1 and STM0 in the STMC1 register, should be set to 00 respectively. In this mode once the counter is enabled and running it can be cleared by three methods. These are a counter overflow, a compare match from Comparator A and a compare match from Comparator P. When the STCCLR bit is low, there are two ways in which the counter can be cleared. One is when a compare match from Comparator P, the other is when the CCRP bits are all zero which allows the counter to overflow. Here both STMAF and STMPF interrupt request flags for Comparator A and Comparator P respectively, will both be generated.

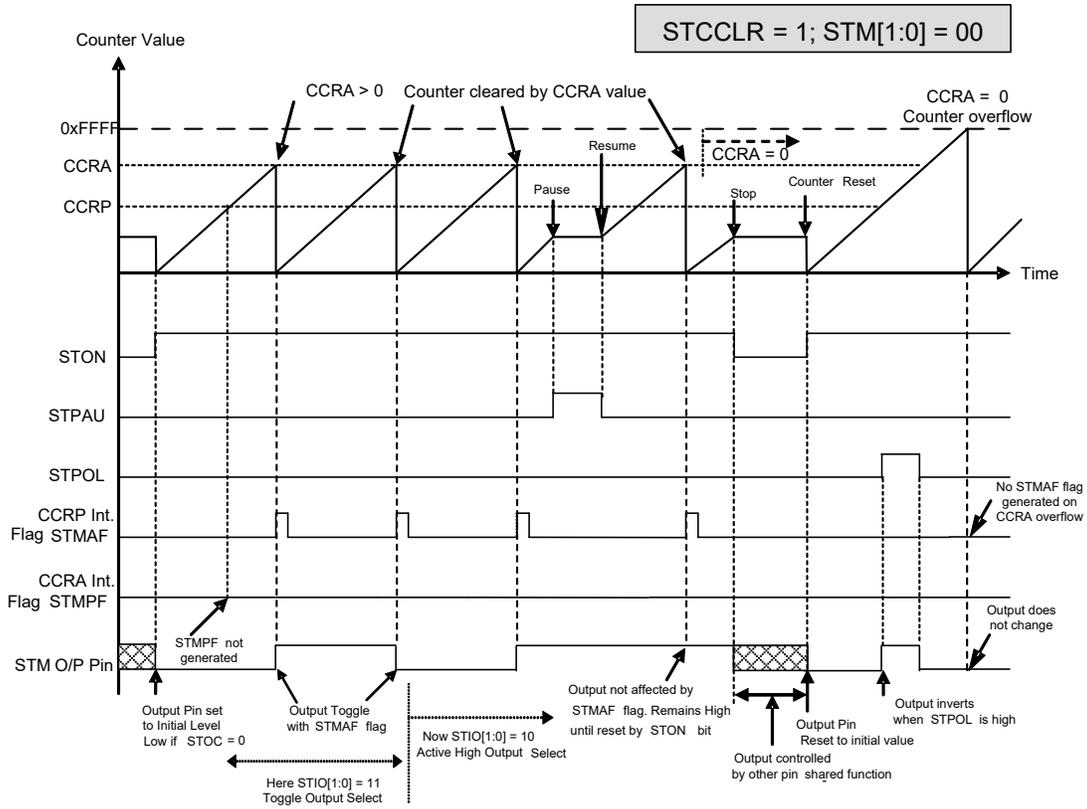
If the STCCLR bit in the STMC1 register is high then the counter will be cleared when a compare match occurs from Comparator A. However, here only the STMAF interrupt request flag will be generated even if the value of the CCRP bits is less than that of the CCRA registers. Therefore when STCCLR is high no STMPF interrupt request flag will be generated. In the Compare Match Output Mode, the CCRA can not be set to “0”.

As the name of the mode suggests, after a comparison is made, the TM output pin, will change state. The TM output pin condition however only changes state when a STMAF interrupt request flag is generated after a compare match occurs from Comparator A. The STMPF interrupt request flag, generated from a compare match occurs from Comparator P, will have no effect on the TM output pin. The way in which the TM output pin changes state are determined by the condition of the STIO1 and STIO0 bits in the STMC1 register. The TM output pin can be selected using the STIO1 and STIO0 bits to go high, to go low or to toggle from its present condition when a compare match occurs from Comparator A. The initial condition of the TM output pin, which is setup after the STON bit changes from low to high, is setup using the STOC bit. Note that if the STIO1 and STIO0 bits are zero then no pin change will take place.



**Compare Match Output Mode – STCCLR = 0**

- Note: 1. With STCCLR = 0 a Comparator P match will clear the counter  
 2. The TM output pin controlled only by the STMAF flag  
 3. The output pin reset to initial state by a STON bit rising edge



**Compare Match Output Mode – STCCLR = 1**

- Note: 1. With STCCLR = 1 a Comparator A match will clear the counter  
 2. The TM output pin controlled only by the STMAF flag  
 3. The output pin reset to initial state by a STON rising edge  
 4. The STMPF flags is not generated when STCCLR = 1

### Timer/Counter Mode

To select this mode, bits STM1 and STM0 in the STMC1 register should be set to 11 respectively. The Timer/Counter Mode operates in an identical way to the Compare Match Output Mode generating the same interrupt flags. The exception is that in the Timer/Counter Mode the TM output pin is not used. Therefore the above description and Timing Diagrams for the Compare Match Output Mode can be used to understand its function. As the TM output pin is not used in this mode, the pin can be used as a normal I/O pin or other pin-shared function.

### PWM Output Mode

To select this mode, bits STM1 and STM0 in the STMC1 register should be set to 10 respectively and also the STIO1 and STIO0 bits should be set to 10 respectively. The PWM function within the TM is useful for applications which require functions such as motor control, heating control, illumination control etc. By providing a signal of fixed frequency but of varying duty cycle on the TM output pin, a square wave AC waveform can be generated with varying equivalent DC RMS values.

As both the period and duty cycle of the PWM waveform can be controlled, the choice of generated waveform is extremely flexible. In the PWM mode, the STCCLR bit has no effect as the PWM period. Both of the CCRA and CCRP registers are used to generate the PWM waveform, one register is used to clear the internal counter and thus control the PWM waveform frequency, while the other one is used to control the duty cycle. Which register is used to control either frequency or duty cycle is determined using the STDPX bit in the STMC1 register.

The PWM waveform frequency and duty cycle can therefore be controlled by the values in the CCRA and CCRP registers. An interrupt flag, one for each of the CCRA and CCRP, will be generated when a compare match occurs from either Comparator A or Comparator P. The STOC bit in the STMC1 register is used to select the required polarity of the PWM waveform while the two STIO1 and STIO0 bits are used to enable the PWM output or to force the TM output pin to a fixed high or low level. The STPOL bit is used to reverse the polarity of the PWM output waveform.

- **16-bit STM, PWM Mode, Edge-aligned Mode, STDPX=0**

CCRP	1~255	0
Period	CCRP×256	65536
Duty	CCRA	

If  $f_{SYS} = 16\text{MHz}$ , TM clock source is  $f_{SYS}/4$ , CCRP = 2 and CCRA = 128,

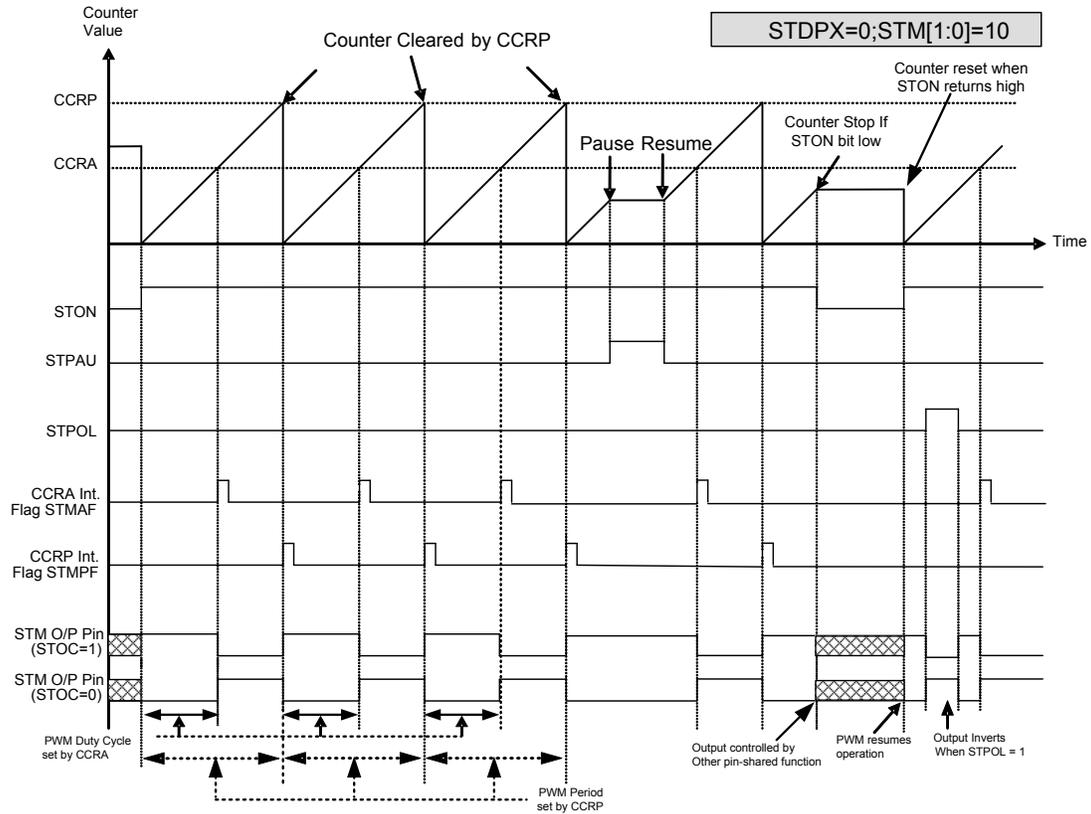
The STM PWM output frequency =  $(f_{SYS}/4) / 512 = f_{SYS}/2048 = 7.8125\text{ kHz}$ , duty =  $128/512 = 25\%$ .

If the Duty value defined by the CCRA register is equal to or greater than the Period value, then the PWM output duty is 100%.

- **16-bit STM, PWM Mode, Edge-aligned Mode, STDPX=1**

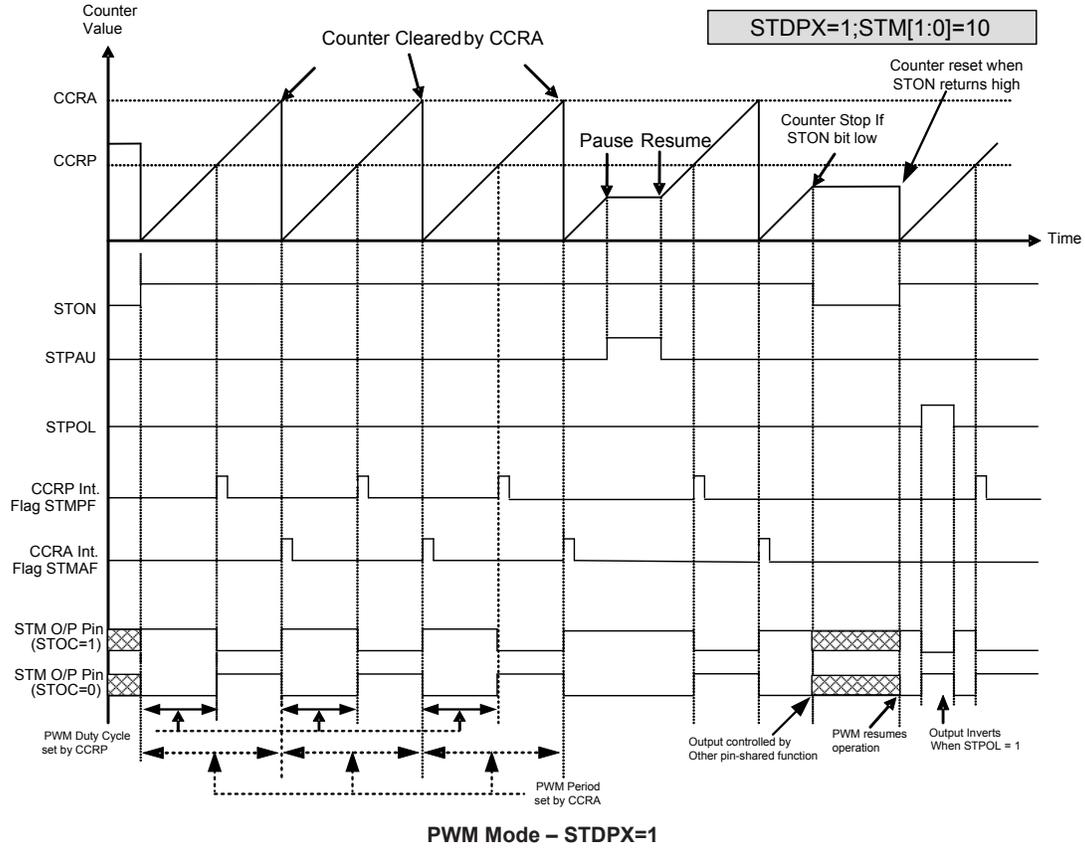
CCRP	1~255	0
Period	CCRA	
Duty	CCRP×256	65536

The PWM output period is determined by the CCRA register value together with the TM clock while the PWM duty cycle is defined by the (CCRP×256) except when the CCRP value is equal to 0.



**PWM Mode – STDPX=0**

- Note: 1. Here STDPX = 0 - Counter cleared by CCRP  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues running even when STIO[1:0] = 00 or 01  
 4. The STCCLR bit has no influence on PWM operation



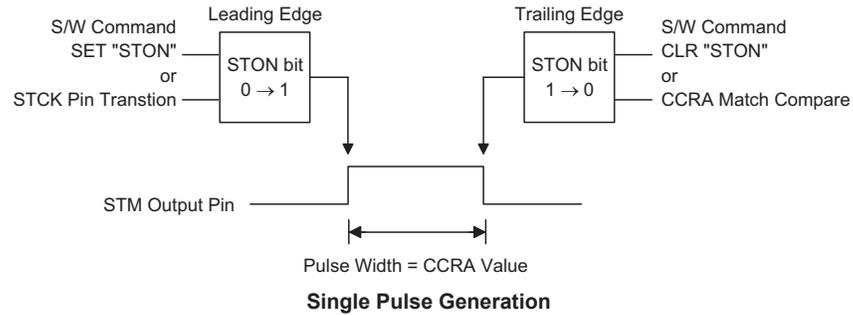
- Note: 1. Here STDPX = 1 - Counter cleared by CCRA  
 2. A counter clear sets PWM Period  
 3. The internal PWM function continues even when STIO[1:0] = 00 or 01  
 4. The STCCLR bit has no influence on PWM operation

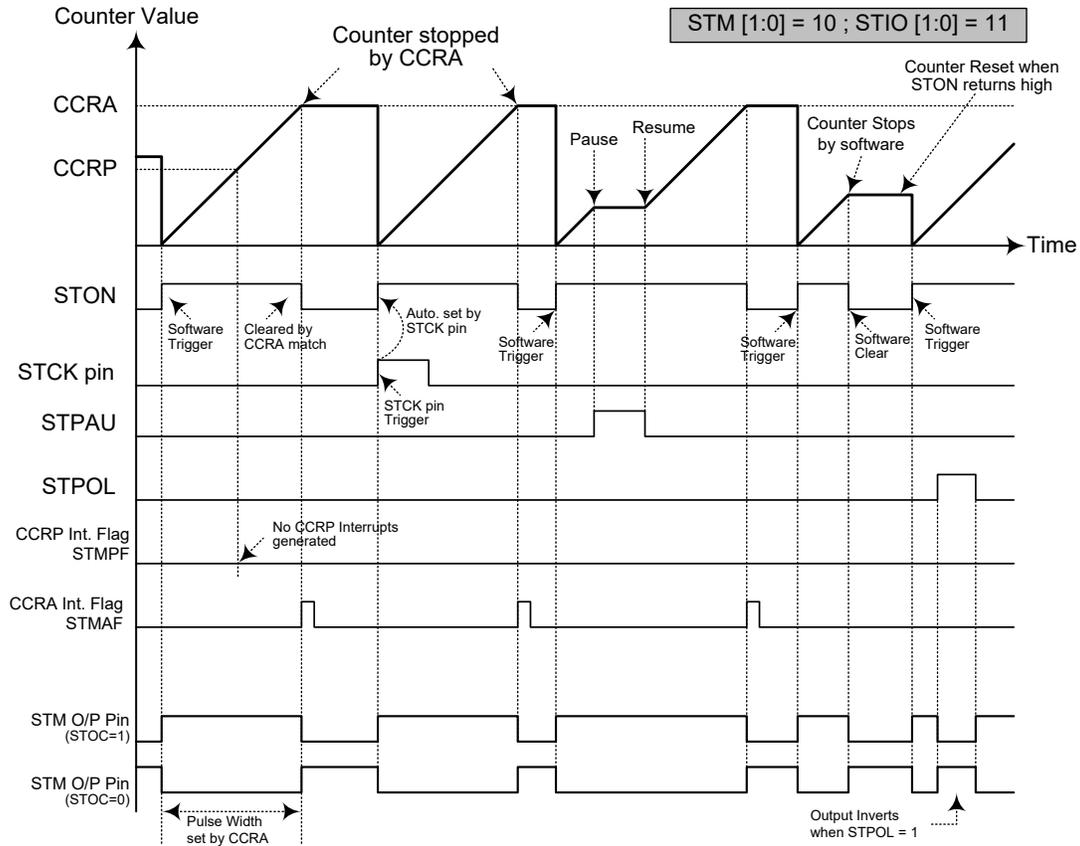
**Single Pulse Mode**

To select this mode, bits STM1 and STM0 in the STMC1 register should be set to 10 respectively and also the STIO1 and STIO0 bits should be set to 11 respectively. The Single Pulse Output Mode, as the name suggests, will generate a single shot pulse on the TM output pin.

The trigger for the pulse output leading edge is a low to high transition of the STON bit, which can be implemented using the application program. However in the Single Pulse Mode, the STON bit can also be made to automatically change from low to high using the external STCK pin, which will in turn initiate the Single Pulse output.

When the STON bit transitions to a high level, the counter will start running and the pulse leading edge will be generated. The STON bit should remain high when the pulse is in its active state. The generated pulse trailing edge will be generated when the STON bit is cleared to zero, which can be implemented using the application program or when a compare match occurs from Comparator A.





**Single Pulse Mode**

- Note:
1. Counter stopped by CCRA match
  2. CCRP is not used
  3. The pulse is triggered by setting the STON bit high or STCK pin
  4. In the Single Pulse Mode, STIO [1:0] must be set to "11" and can not be changed.

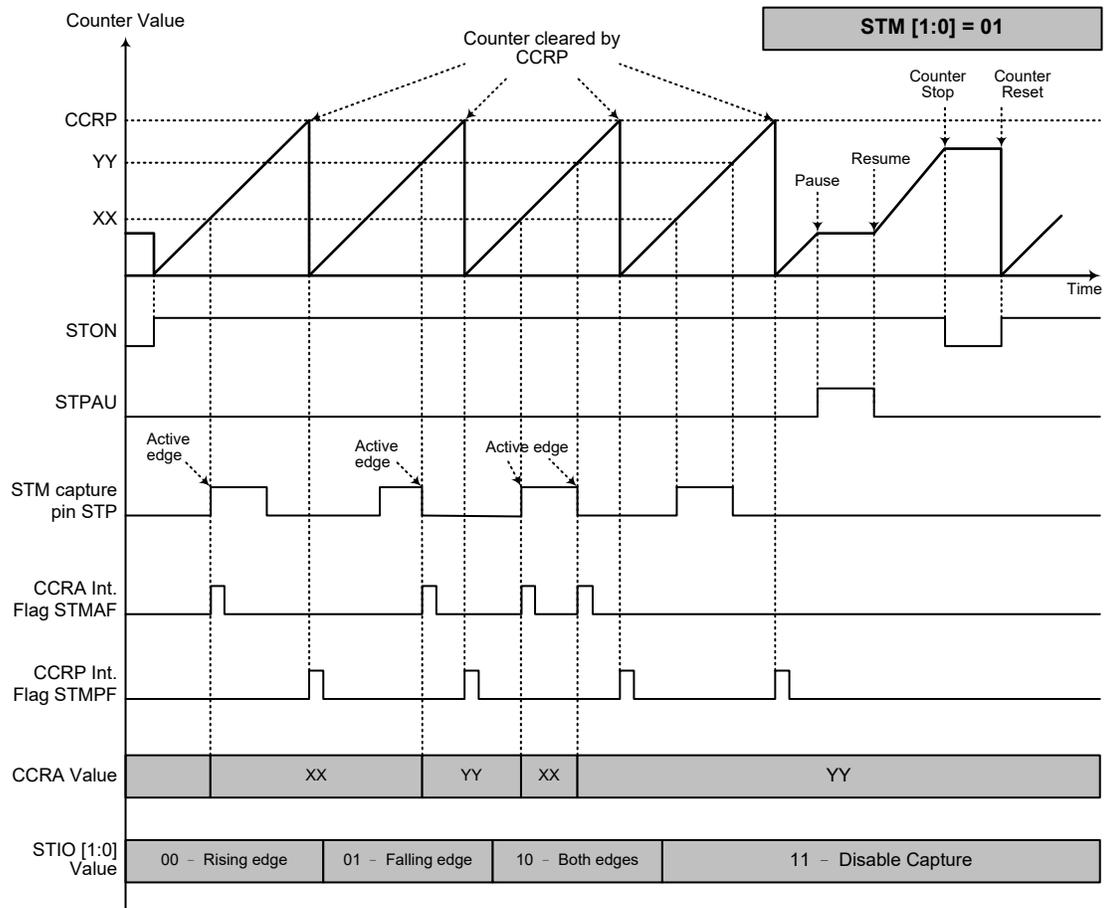
However a compare match from Comparator A will also automatically clear the STON bit and thus generate the Single Pulse output trailing edge. In this way the CCRA value can be used to control the pulse width. A compare match from Comparator A will also generate a TM interrupt. The counter can only be reset back to zero when the STON bit changes from low to high when the counter restarts. In the Single Pulse Mode CCRP is not used. The STCCLR and STDPX bits are not used in this Mode.

### **Capture Input Mode**

To select this mode bits STM1 and STM0 in the STMC1 register should be set to 01 respectively. This mode enables external signals to capture and store the present value of the internal counter and can therefore be used for applications such as pulse width measurements. The external signal is supplied on the STP pin, whose active edge can be either a rising edge, a falling edge or both rising and falling edges; the active edge transition type is selected using the STIO1 and STIO0 bits in the STMC1 register. The counter is started when the STON bit changes from low to high which is initiated using the application program.

When the required edge transition appears on the STP pin the present value in the counter will be latched into the CCRA registers and a TM interrupt generated. Irrespective of what events occur on the STP pin the counter will continue to free run until the STON bit changes from high to low. When a CCRP compare match occurs the counter will reset back to zero; in this way the CCRP value can be used to control the maximum counter value. When a CCRP compare match occurs from Comparator P, a TM interrupt will also be generated. Counting the number of overflow interrupt signals from the CCRP can be a useful method in measuring long pulse widths. The STIO1 and STIO0 bits can select the active trigger edge on the STP pin to be a rising edge, falling edge or both edge types. If the STIO1 and STIO0 bits are both set high, then no capture operation will take place irrespective of what happens on the STP pin, however it must be noted that the counter will continue to run.

As the STP pin is pin shared with other functions, care must be taken if the TM is in the Input Capture Mode. This is because if the pin is setup as an output, then any transitions on this pin may cause an input capture operation to be executed. The STCCLR and STDPX bits are not used in this Mode.



**Capture Input Mode**

- Note: 1. STM[1:0] = 01 and active edge set by the STIO[1:0] bits  
 2. A TM Capture input pin active edge transfers the counter value to CCRA  
 3. The STCCLR bit is not used  
 4. No output function - STOC and STPOL bits are not used  
 5. CCRP determines the counter value and the counter has a maximum count value when CCRP is equal to zero.

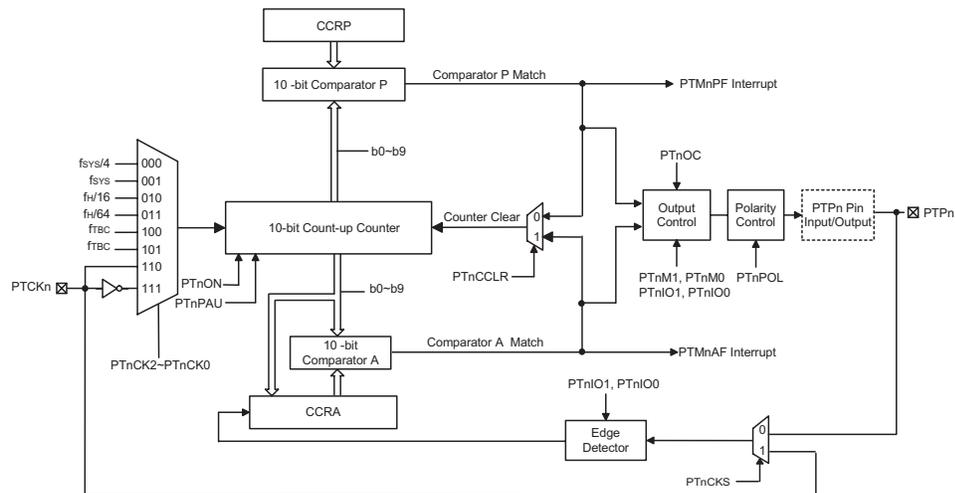
## Periodic Type TM – PTM0, PTM1

The PTM contains five operating modes, which are Compare Match Output, Timer/Event Counter, Capture Input, Single Pulse Output and PWM Output modes. The P-type TM can also be controlled with an external input pin and can drive one output pins.

### Periodic TM Operation

There are two P-type TMs, both are 10-bit wide. At the core is a 10 count-up counter which is driven by a user selectable internal or external clock source. There are also two internal comparators with the names, Comparator A and Comparator P. These comparators will compare the value in the counter with CCRP and CCRA registers. The CCRP comparator is 10-bits wide.

The only way of changing the value of the 10 counter using the application program, is to clear the counter by changing the PTnON bit from low to high. The counter will also be cleared automatically by a counter overflow or a compare match with one of its associated comparators. When these conditions occur, a TM interrupt signal will also usually be generated. The Periodic Type TM can operate in a number of different operational modes, can be driven by different clock sources including an input pin and can also control an output pin. All operating setup conditions are selected using relevant internal registers.



**Periodic Type TM Block Diagram (n=0, 1)**

## PTM register description

Overall operation of the P-type TM is controlled using a series of registers. A read only register pair exists to store the internal counter 10-bit value, while two read/write register pairs exist to store the internal 10-bit CCRA/CCRP value. The remaining two registers are control registers which setup the different operating and control modes.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PTMnC0	PTnPAU	PTnCK2	PTnCK1	PTnCK0	PTnON	—	—	—
PTMnC1	PTnM1	PTnM0	PTnIO1	PTnIO0	PTnOC	PTnPOL	PTnCKS	PTnCCLR
PTMnDL	D7	D6	D5	D4	D3	D2	D1	D0
PTMnDH	—	—	—	—	—	—	D9	D8
PTMnAL	D7	D6	D5	D4	D3	D2	D1	D0
PTMnAH	—	—	—	—	—	—	D9	D8
PTMnRPL	D7	D6	D5	D4	D3	D2	D1	D0
PTMnRPH	—	—	—	—	—	—	D9	D8

10-bit Periodic TM Register List (n=0, 1)

### PTMnC0 Register

Bit	7	6	5	4	3	2	1	0
Name	PTnPAU	PTnCK2	PTnCK1	PTnCK0	PTnON	—	—	—
R/W	R/W	R/W	R/W	R/W	R/W	—	—	—
POR	0	0	0	0	0	—	—	—

Bit 7 **PTnPAU**: PTMn Counter Pause Control

0: Run  
1: Pause

The counter can be paused by setting this bit high. Clearing the bit to zero restores normal counter operation. When in a Pause condition the TM will remain powered up and continue to consume power. The counter will retain its residual value when this bit changes from low to high and resume counting from this value when the bit changes to a low value again.

Bit 6 ~4 **PTnCK2 ~ PTnCK0**: Select PTMn Counter clock

000:  $f_{SYS}/4$   
001:  $f_{SYS}$   
010:  $f_H/16$   
011:  $f_H/64$   
100:  $f_{TBC}$   
101:  $f_{TBC}$   
110: PTCKn rising edge clock  
111: PTCKn falling edge clock

These three bits are used to select the clock source for the TM. The external pin clock source can be chosen to be active on the rising or falling edge. The clock source  $f_{SYS}$  is the system clock, while  $f_H$  and  $f_{TBC}$  are other internal clocks, the details of which can be found in the oscillator section.

- Bit 3     **PTnON**: PTMn Counter On/Off Control  
           0: Off  
           1: On
- This bit controls the overall on/off function of the TM. Setting the bit high enables the counter to run, clearing the bit disables the TM. Clearing this bit to zero will stop the counter from counting and turn off the TM which will reduce its power consumption. When the bit changes state from low to high the internal counter value will be reset to zero, however when the bit changes from high to low, the internal counter will retain its residual value until the bit returns high again. If the TM is in the Compare Match Output Mode then the TM output pin will be reset to its initial condition, as specified by the PTnOC bit, when the PTnON bit changes from low to high.
- Bit 2~0    Unimplemented, read as“0”

**PTMnC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	PTnM1	PTnM0	PTnIO1	PTnIO0	PTnOC	PTnPOL	PTnCKS	PTnCCLR
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7~6     **PTnM1~PTnM0**: Select PTMn Operating Mode  
           00: Compare Match Output Mode  
           01: Capture Input Mode  
           10: PWM Mode or Single Pulse Output Mode  
           11: Timer/Counter Mode
- These bits setup the required operating mode for the TM. To ensure reliable operation the TM should be switched off before any changes are made to the PTnM1 and PTnM0 bits. In the Timer/Counter Mode, the TM output pin control must be disabled.
- Bit 5~4     **PTnIO1~PTnIO0**: Select PTPn output function
- Compare Match Output Mode  
           00: No change  
           01: Output low  
           10: Output high  
           11: Toggle output
- PWM Mode/Single Pulse Output Mode  
           00: Force inactive state  
           01: Force active state  
           10: PWM output  
           11: Single pulse output
- Capture Input Mode  
           00: Input capture at rising edge of PTPn or PTCKn  
           01: Input capture at falling edge of PTPn or PTCKn  
           10: Input capture at falling/rising edge of PTPn or PTCKn  
           11: Input capture disabled
- Timer/counter Mode  
           Unused
- These two bits are used to determine how the TM output pin changes state when a certain condition is reached. The function that these bits select depends upon in which mode the TM is running. In the Compare Match Output Mode, the PTnIO1 and PTnIO0 bits determine how the TM output pin changes state when a compare match occurs from the Comparator A. The TM output pin can be setup to switch high, switch low or to toggle its present state when a compare match occurs from the Comparator A. When the bits are both zero, then no change will take place on the output. The initial value of the TM output pin should be setup using the PTnOC bit in the PTMnC1 register. Note that the output level requested by the PTnIO1 and PTnIO0 bits must be

different from the initial value setup using the PTnOC bit otherwise no change will occur on the TM output pin when a compare match occurs. After the TM output pin changes state it can be reset to its initial level by changing the level of the PTnON bit from low to high.

Bit 3 **PTnOC**: PTPn Output control bit

Compare Match Output Mode

0: Initial low

1: Initial high

PWM Mode/Single Pulse Output Mode

0: Active low

1: Active high

This is the output control bit for the TM output pin. Its operation depends upon whether TM is being used in the Compare Match Output Mode or in the PWM Mode/Single Pulse Output Mode. It has no effect if the TM is in the Timer/Counter Mode. In the Compare Match Output Mode it determines the logic level of the TM output pin before a compare match occurs. In the PWM Mode it determines if the PWM signal is active high or active low.

Bit 2 **PTnPOL**: PTPn Output polarity control

0: Non-invert

1: Invert

This bit controls the polarity of the PTPn output pin. When the bit is set high the TM output pin will be inverted and not inverted when the bit is zero. It has no effect if the TM is in the Timer/Counter Mode.

Bit 1 **PTnCKS**: Input Capture trigger source selection

0: External Clock source of Capture Input Mode comes from PTPn

1: External Clock source of Capture Input Mode comes from PTCKn

Bit 0 **PTnCCLR**: Select PTMn Counter clear condition

0: PTMn Comparator P match

1: PTMn Comparator A match

This bit is used to select the method which clears the counter. Remember that the P-type TM contains two comparators, Comparator A and Comparator P, either of which can be selected to clear the internal counter. With the PTnCCLR bit set high, the counter will be cleared when a compare match occurs from the Comparator A. When the bit is low, the counter will be cleared when a compare match occurs from the Comparator P or with a counter overflow. A counter overflow clearing method can only be implemented if the CCRP bits are all cleared to zero. The PTnCCLR bit is not used in the PWM, Single Pulse or Input Capture Mode.

**PTMnDL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: PTMn Counter Low Byte Register Bit 7~Bit 0  
PTMn 10-bit Counter bit 7 ~ bit 0

**PTMnDH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R	R
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as“0”  
Bit 1~0 **D9~D8**: PTMn Counter High Byte Register Bit 1~Bit 0  
PTMn 10-bit Counter bit 9 ~ bit 8

**PTMnAL Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: PTMn CCRA Low Byte Register bit 7~bit 0  
PTMn 10-bit CCRA bit 7~bit 0

**PTMnAH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as“0”  
Bit 1~0 **D9~D8**: PTMn CCRA High Byte Register Bit 1~Bit 0  
PTMn 10-bit CCRA bit 9 ~ bit 8

**PTMnRPL Register (n=0, 1)**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7~0 **D7~D0**: PTMn CCRP Register bit 7 ~ bit 0  
PTMn 10-bit CCRP bit 7 ~ bit 0

### PTMnRPH Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	D9	D8
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit 7~2 Unimplemented, read as“0”

Bit 1~0 **D9~D8**: PTMn CCRP Register bit 1 ~ bit 0  
 PTMn 10-bit CCRP bit 9 ~ bit 8

### Periodic Type TM Operating Modes

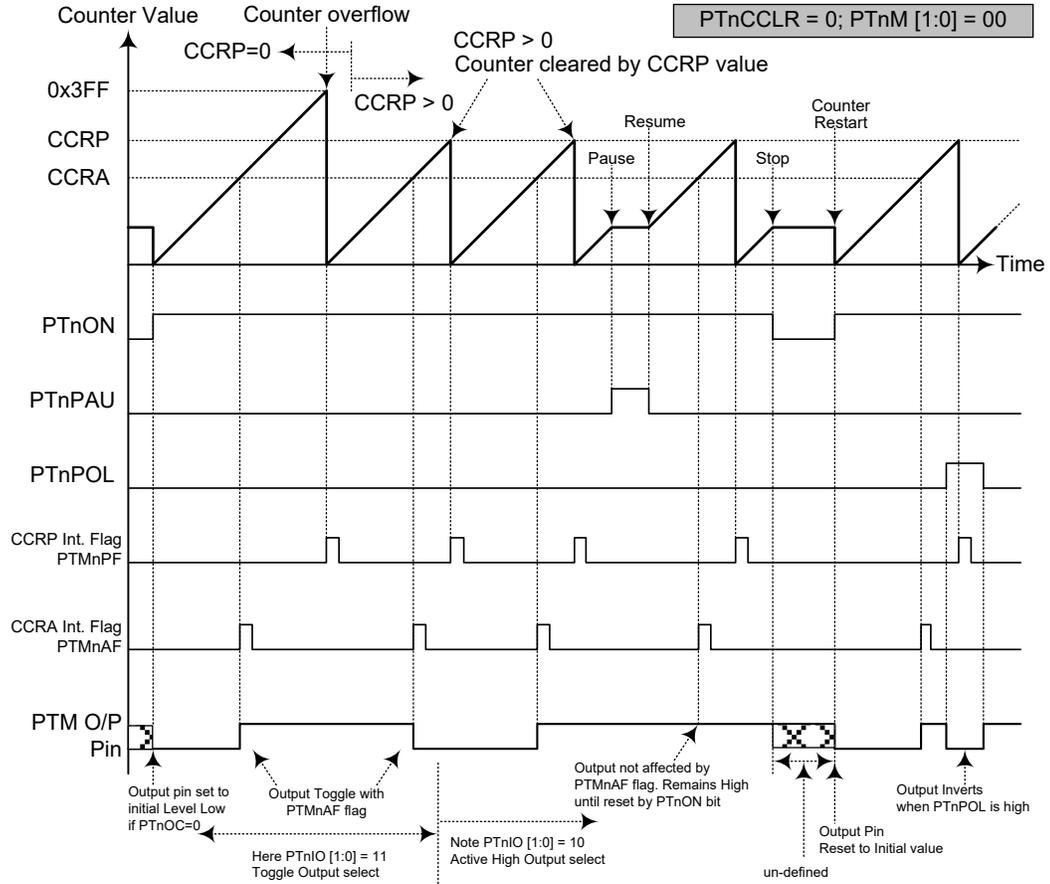
The P-type TM can operate in one of five operating modes, Compare Match Output Mode, PWM Output Mode, Single Pulse Output Mode, Capture Input Mode or Timer/Counter Mode. The operating mode is selected using the PTnM1 and PTnM0 bits in the PTMnC1 register.

#### Compare Output Mode

To select this mode, bits PTnM1 and PTnM0 in the PTMnC1 register, should be set to 00 respectively. In this mode once the counter is enabled and running it can be cleared by three methods. These are a counter overflow, a compare match from Comparator A and a compare match from Comparator P. When the PTnCCLR bit is low, there are two ways in which the counter can be cleared. One is when a compare match from Comparator P, the other is when the CCRP bits are all zero which allows the counter to overflow. Here both PTMnAF and PTMnPF interrupt request flags for Comparator A and Comparator P respectively, will both be generated.

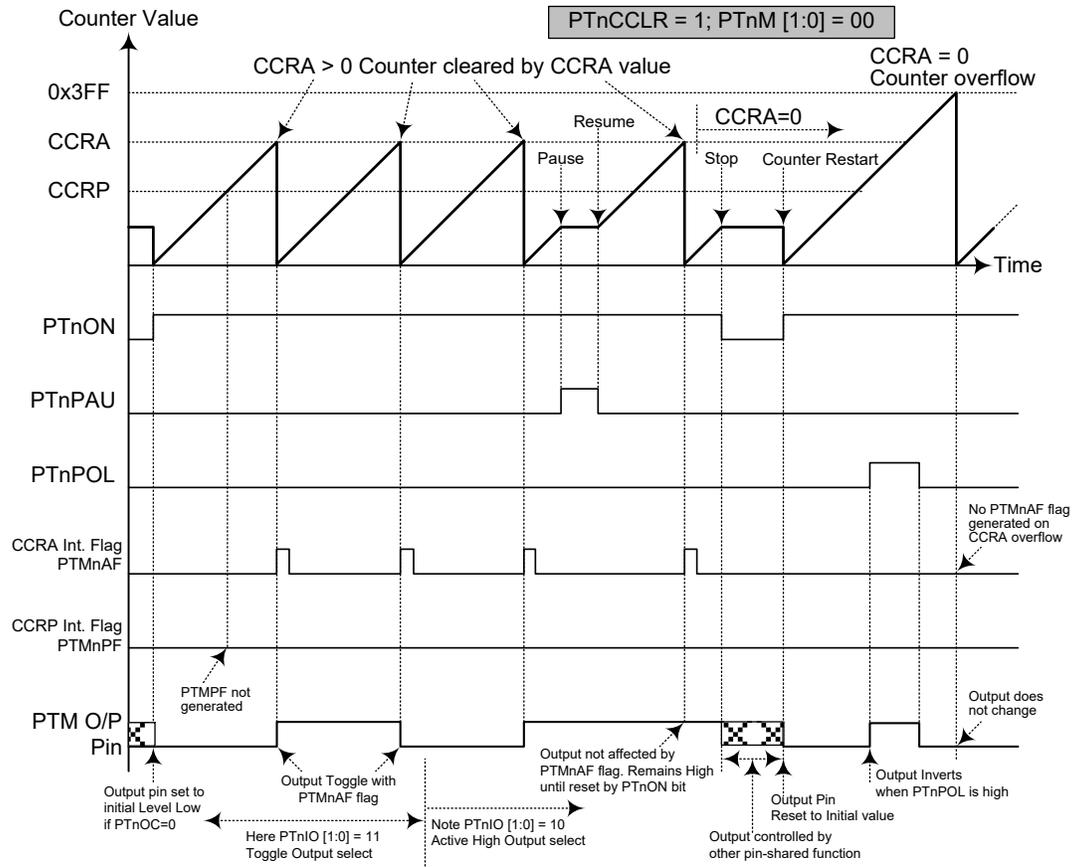
If the PTnCCLR bit in the PTMnC1 register is high then the counter will be cleared when a compare match occurs from Comparator A. However, here only the PTMnAF interrupt request flag will be generated even if the value of the CCRP bits is less than that of the CCRA registers. Therefore when PTnCCLR is high no PTMnPF interrupt request flag will be generated. In the Compare Match Output Mode, the CCRA can not be set to “0”.

As the name of the mode suggests, after a comparison is made, the TM output pin, will change state. The TM output pin condition however only changes state when a PTMnAF interrupt request flag is generated after a compare match occurs from Comparator A. The PTMnPF interrupt request flag, generated from a compare match occurs from Comparator P, will have no effect on the TM output pin. The way in which the TM output pin changes state are determined by the condition of the PTnIO1 and PTnIO0 bits in the PTMnC1 register. The TM output pin can be selected using the PTnIO1 and PTnIO0 bits to go high, to go low or to toggle from its present condition when a compare match occurs from Comparator A. The initial condition of the TM output pin, which is setup after the PTnON bit changes from low to high, is setup using the PTnOC bit. Note that if the PTnIO1 and PTnIO0 bits are zero then no pin change will take place.



**Compare Match Output Mode – PTnCCLR=0 (n=0, 1)**

- Note: 1. With PTnCCLR=0 a Comparator P match will clear the counter  
 2. The TM output pin is controlled only by the PTMnAF flag  
 3. The output pin is reset to its initial state by a PTnON bit rising edge



**Compare Match Output Mode –  $PTnCCLR=1$  ( $n=0, 1$ )**

- Note: 1. With  $PTnCCLR=1$  a Comparator A match will clear the counter  
 2. The TM output pin is controlled only by the PTMnAF flag  
 3. The output pin is reset to its initial state by a PTnON bit rising edge  
 4. A PTMnPF flag is not generated when  $PTnCCLR=1$

### Timer/Counter Mode

To select this mode, bits PTnM1 and PTnM0 in the PTMnC1 register should be set to 11 respectively. The Timer/Counter Mode operates in an identical way to the Compare Match Output Mode generating the same interrupt flags. The exception is that in the Timer/Counter Mode the TM output pin is not used. Therefore the above description and Timing Diagrams for the Compare Match Output Mode can be used to understand its function. As the TM output pin is not used in this mode, the pin can be used as a normal I/O pin or other pin-shared function.

### PWM Output Mode

To select this mode, bits PTnM1 and PTnM0 in the PTMnC1 register should be set to 10 respectively and also the PTnIO1 and PTnIO0 bits should be set to 10 respectively. The PWM function within the TM is useful for applications which require functions such as motor control, heating control, illumination control etc. By providing a signal of fixed frequency but of varying duty cycle on the TM output pin, a square wave AC waveform can be generated with varying equivalent DC RMS values.

As both the period and duty cycle of the PWM waveform can be controlled, the choice of generated waveform is extremely flexible. In the PWM mode, the PTnCCLR bit has no effect as the PWM period. Both of the CCRA and CCRP registers are used to generate the PWM waveform, one register is used to clear the internal counter and thus control the PWM waveform frequency, while the other one is used to control the duty cycle. The PWM waveform frequency and duty cycle can therefore be controlled by the values in the CCRA and CCRP registers.

An interrupt flag, one for each of the CCRA and CCRP, will be generated when a compare match occurs from either Comparator A or Comparator P. The PTnOC bit in the PTMnC1 register is used to select the required polarity of the PWM waveform while the two PTnIO1 and PTnIO0 bits are used to enable the PWM output or to force the TM output pin to a fixed high or low level. The PTnPOL bit is used to reverse the polarity of the PWM output waveform.

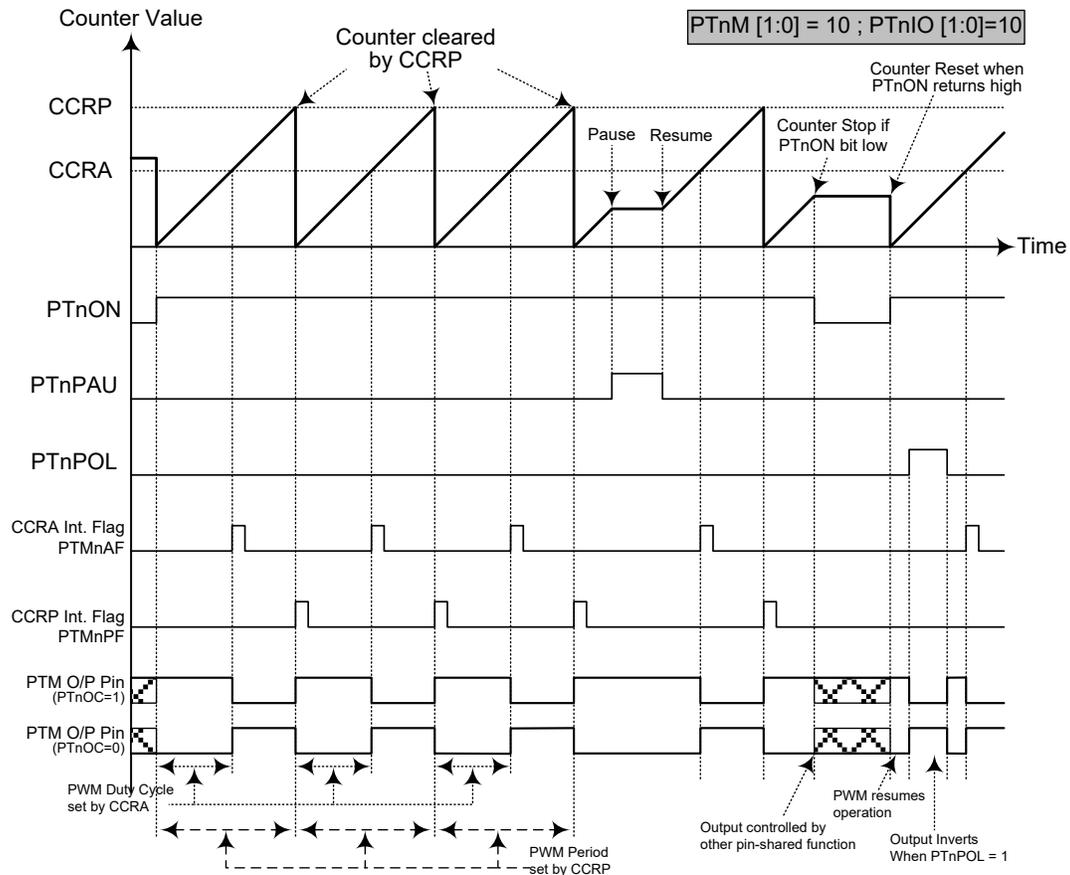
#### • 10-bit PWM Mode, Edge-aligned Mode

CCRP	0	1~1023
Period	1024	1~1023
Duty	CCRA	

If  $f_{SYS} = 16\text{MHz}$ , TM clock source select  $f_{SYS}/4$ , CCRP = 512 and CCRA = 128,

The PTM PWM output frequency =  $(f_{SYS}/4) / 512 = f_{SYS}/2048 = 7.8125\text{kHz}$ , duty =  $128/512 = 25\%$ .

If the Duty value defined by the CCRA register is equal to or greater than the Period value, then the PWM output duty is 100%.



**PWM mode (n=0, 1)**

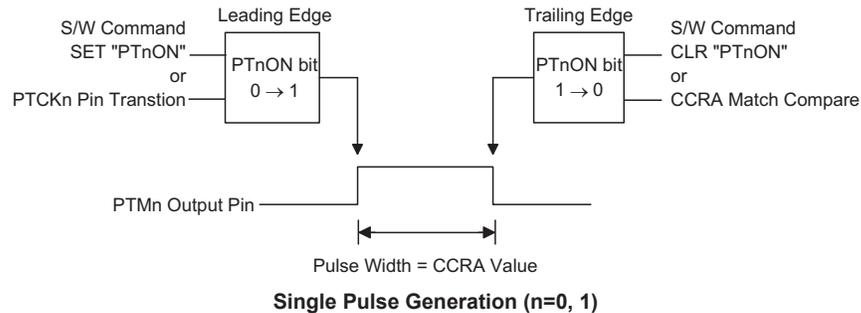
- Note: 1. A counter clear sets the PWM Period  
 2. The internal PWM function continues running even when PTnIO [1:0] = 00 or 01  
 3. The PTnCCLR bit has no influence on PWM operation

### Single Pulse Mode

To select this mode, bits PTnM1 and PTnM0 in the PTMnC1 register should be set to 10 respectively and also the PTnIO1 and PTnIO0 bits should be set to 11 respectively. The Single Pulse Output Mode, as the name suggests, will generate a single shot pulse on the TM output pin.

The trigger for the pulse output leading edge is a low to high transition of the PTnON bit, which can be implemented using the application program. However in the Single Pulse Mode, the PTnON bit can also be made to automatically change from low to high using the external PTCKn pin, which will in turn initiate the Single Pulse output. When the PTnON bit transitions to a high level, the counter will start running and the pulse leading edge will be generated. The PTnON bit should remain high when the pulse is in its active state. The generated pulse trailing edge will be generated when the PTnON bit is cleared to zero, which can be implemented using the application program or when a compare match occurs from Comparator A.

However a compare match from Comparator A will also automatically clear the PTnON bit and thus generate the Single Pulse output trailing edge. In this way the CCRA value can be used to control the pulse width. A compare match from Comparator A will also generate a TM interrupt. The counter can only be reset back to zero when the PTnON bit changes from low to high when the counter restarts. In the Single Pulse Mode CCRP is not used. The PTnCCLR bit is not used in this Mode.



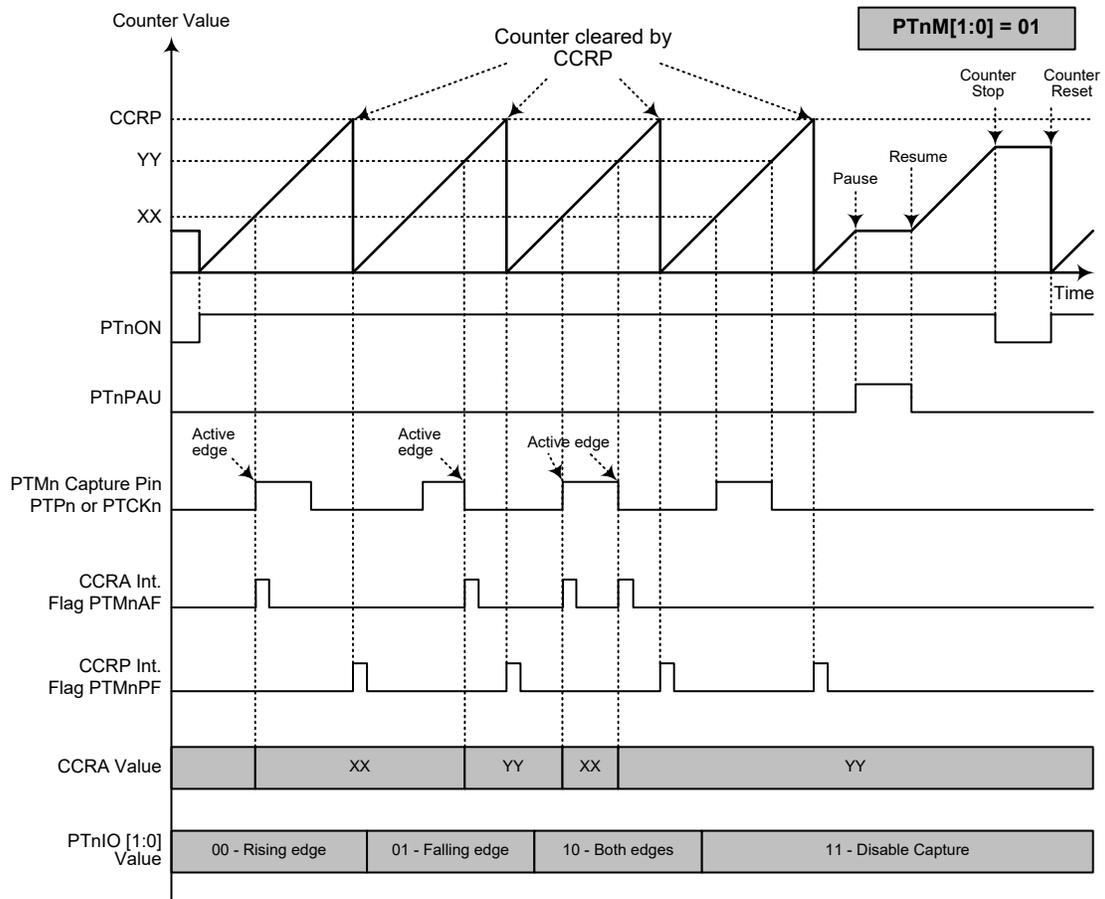


### **Capture Input Mode**

To select this mode bits PTnM1 and PTnM0 in the PTMnC1 register should be set to 01 respectively. This mode enables external signals to capture and store the present value of the internal counter and can therefore be used for applications such as pulse width measurements. The external signal is supplied on the PTPn or PTCKn pin, selected by the PTnCKS bit in the PTMnC1 register. The input pin active edge can be a rising edge, a falling edge or both rising and falling edges; the active edge transition type is selected using the PTnIO1 and PTnIO0 bits in the PTMnC1 register. The counter is started when the PTnON bit changes from low to high which is initiated using the application program.

When the required edge transition appears on the PTPn or PTCKn pin the present value in the counter will be latched into the CCRA registers and a TM interrupt generated. Irrespective of what events occur on the PTPn or PTCKn pin the counter will continue to free run until the PTnON bit changes from high to low. When a CCRP compare match occurs the counter will reset back to zero; in this way the CCRP value can be used to control the maximum counter value. When a CCRP compare match occurs from Comparator P, a TM interrupt will also be generated. Counting the number of overflow interrupt signals from the CCRP can be a useful method in measuring long pulse widths. The PTnIO1 and PTnIO0 bits can select the active trigger edge on the PTPn or PTCKn pin to be a rising edge, falling edge or both edge types. If the PTnIO1 and PTnIO0 bits are both set high, then no capture operation will take place irrespective of what happens on the PTPn or PTCKn pin, however it must be noted that the counter will continue to run.

As the PTPn pin is pin or PTCKn pin shared with other functions, care must be taken if the PTMn is in the Input Capture Mode. This is because if the pin is setup as an output, then any transitions on this pin may cause an input capture operation to be executed. The PTnCCLR, PTnOC and PTnPOL bits are not used in this Mode.



**Capture Input Mode (n=0, 1)**

- Note: 1.  $PTnM[1:0] = 01$  and active edge set by the  $PTnIO[1:0]$  bits  
 2. A TM Capture input pin active edge transfers the counter value to CCRA  
 3.  $PTnCCCLR$  bit not used  
 4. No output function –  $PTnOC$  and  $PTnPOL$  bits are not used  
 5. CCRP determines the counter value and the counter has a maximum count value when CCRP is equal to zero.

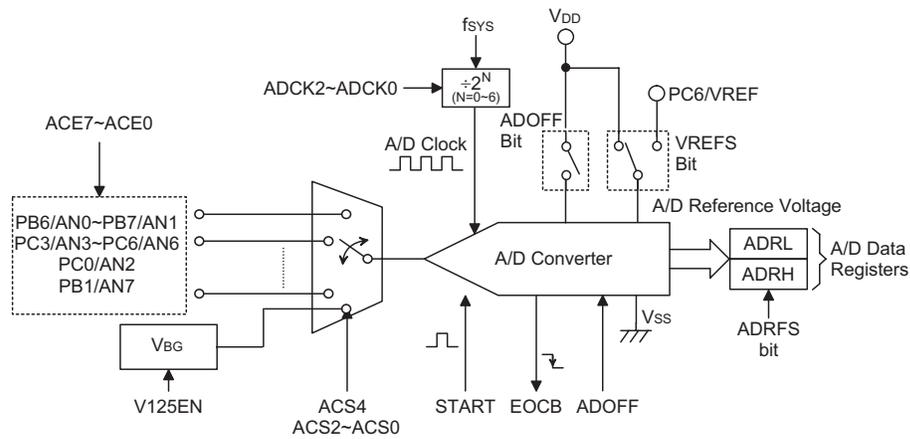
## Analog to Digital Converter

The need to interface to real world analog signals is a common requirement for many electronic systems. However, to properly process these signals by a microcontroller, they must first be converted into digital signals by A/D converters. By integrating the A/D conversion electronic circuitry into the microcontroller, the need for external components is reduced significantly with the corresponding follow-on benefits of lower costs and reduced component space requirements.

### A/D Overview

The device contains a 8-channel analog to digital converter which can directly interface to external analog signals, such as that from sensors or other control signals and convert these signals directly into a 12-bit digital value.

The accompanying block diagram shows the overall internal structure of the A/D converter, together with its associated registers.



**A/D Converter Structure**

### A/D Converter Register Description

Overall operation of the A/D converter is controlled using five registers. A read only register pair exists to store the ADC data 12-bit value. The remaining three registers are control registers which setup the operating and control function of the A/D converter.

Register Name	Bit							
	7	6	5	4	3	2	1	0
ADRL(ADRFS=0)	D3	D2	D1	D0	—	—	—	—
ADRH(ADRFS=0)	D11	D10	D9	D8	D7	D6	D5	D4
ADRL(ADRFS=1)	D7	D6	D5	D4	D3	D2	D1	D0
ADRH(ADRFS=1)	—	—	—	—	D11	D10	D9	D8
ADCR0	START	EOCB	ADOFF	ADRFS	—	ACS2	ACS1	ACS0
ADCR1	ACS4	V125EN	—	VREFS1	VREFS0	ADCK2	ADCK1	ADCK0
ACERL	ACE7	ACE6	ACE5	ACE4	ACE3	ACE2	ACE1	ACE0

**A/D Converter Register List**

### A/D Converter Data Registers – ADRL, ADRH

As the device contains an internal 12-bit A/D converter, it requires two data registers to store the converted value. These are a high byte register, known as ADRH, and a low byte register, known as ADRL. After the conversion process takes place, these registers can be directly read by the microcontroller to obtain the digitised conversion value. As only 12 bits of the 16-bit register space is utilised, the format in which the data is stored is controlled by the ADRFS bit in the ADCR0 register as shown in the accompanying table. D0~D11 are the A/D conversion result data bits. Any unused bits will be read as zero.

ADRFS	ADRH								ADRL							
	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	D11	D10	D9	D8	D7	D6	D5	D4	D3	D2	D1	D0	0	0	0	0
1	0	0	0	0	D11	D10	D9	D8	D7	D6	D5	D4	D3	D2	D1	D0

**A/D Data Registers**

### A/D Converter Control Registers – ADCR0, ADCR1, ACERL

To control the function and operation of the A/D converter, three control registers known as ADCR0, ADCR1, ACERL are provided. These 8-bit registers define functions such as the selection of which analog channel is connected to the internal A/D converter, the digitised data format, the A/D clock source as well as controlling the start function and monitoring the A/D converter end of conversion status. The ACS4 in the ADCR1 register and ACS2~ACS0 bits in the ADCR0 register define the ADC input channel number. As the device contains only one actual analog to digital converter hardware circuit, each of the individual 8 analog inputs must be routed to the converter. It is the function of the ACS4, ACS2 ~ ACS0 bits to determine which analog channel input signals or internal  $V_{BG}$  is actually connected to the internal A/D converter.

The ACERL control register contains the ACE7~ACE0 bits which determine which pins on Port C are used as analog inputs for the A/D converter input and which pins are not to be used as the A/D converter input. Setting the corresponding bit high will select the A/D input function, clearing the bit to zero will select either the I/O or other pin-shared function. When the pin is selected to be an A/D input, its original function whether it is an I/O or other pin-shared function will be removed. In addition, any internal pull-high resistors connected to these pins will be automatically removed if the pin is selected to be an A/D input.

**ADCR0 Register**

Bit	7	6	5	4	3	2	1	0
Name	START	EOCB	ADOFF	ADRF5	—	ACS2	ACS1	ACS0
R/W	R/W	R	R/W	R/W	—	R/W	R/W	R/W
POR	0	1	1	0	—	0	0	0

- Bit 7**     **START:** Start the A/D conversion  
0→1→0: start  
0→1: reset the A/D converter and set EOCB to "1"  
This bit is used to initiate an A/D conversion process. The bit is normally low but if set high and then cleared low again, the A/D converter will initiate a conversion process. When the bit is set high the A/D converter will be reset.
- Bit 6**     **EOCB:** End of A/D conversion flag  
0: A/D conversion ended  
1: A/D conversion in progress  
This read only flag is used to indicate when an A/D conversion process has completed. When the conversion process is running the bit will be high.
- Bit 5**     **ADOFF :** ADC module power on/off control bit  
0: ADC module power on  
1: ADC module power off  
This bit controls the power to the A/D internal function. This bit should be cleared to zero to enable the A/D converter. If the bit is set high then the A/D converter will be switched off reducing the device power consumption. As the A/D converter will consume a limited amount of power, even when not executing a conversion, this may be an important consideration in power sensitive battery powered applications.  
Note: 1. it is recommended to set ADOFF=1 before entering IDLE/SLEEP Mode for saving power.  
2. ADOFF=1 will power down the ADC module.
- Bit 4**     **ADRF5:** ADC Data Format Control  
0: ADC Data MSB is ADRH bit 7, LSB is ADRL bit 4  
1: ADC Data MSB is ADRH bit 3, LSB is ADRL bit 0  
This bit controls the format of the 12-bit converted A/D value in the two A/D data registers. Details are provided in the A/D data register section.
- Bit 3**     Unimplemented, read as "0"
- Bit 2 ~ 0**   **ACS2 ~ ACS0:** Select A/D channel (when ACS4 is "0")  
000: AN0  
001: AN1  
010: AN2  
011: AN3  
100: AN4  
101: AN5  
110: AN6  
111: AN7  
These are the A/D channel select control bits. As there is only one internal hardware A/D converter each of the eight A/D inputs must be routed to the internal converter using these bits. If bit ACS4 in the ADCR1 register is set high then the internal  $V_{BG}$  will be routed to the A/D Converter.

**ADCR1 Register**

Bit	7	6	5	4	3	2	1	0
Name	ACS4	V125EN	—	—	VREFS	ADCK2	ADCK1	ADCK0
R/W	R/W	R/W	—	—	R/W	R/W	R/W	R/W
POR	0	0	—	—	0	0	0	0

- Bit 7 ACS4:** Selecte Internal  $V_{BG}$  as ADC input Control  
 0: Disable  
 1: Enable  
 This bit enables  $V_{BG}$  to be connected to the A/D converter. The V125EN bit must first have been set to enable the bandgap circuit  $V_{BG}$  voltage to be used by the A/D converter. When the ACS4 bit is set high, the bandgap voltage will be routed to the A/D converter and the other A/D input channels disconnected.
- Bit 6 V125EN:** Internal  $V_{BG}$  Control  
 0: Disable  
 1: Enable  
 This bit controls the internal Bandgap circuit on/off function to the A/D converter. When the bit is set high the bandgap voltage  $V_{BG}$  can be used by the A/D converter. If  $V_{BG}$  is not used by the A/D converter and the LVR function is disabled then the bandgap reference circuit will be automatically switched off to conserve power. When  $V_{BG}$  is switched on for use by the A/D converter, a time  $t_{BG}$  should be allowed for the bandgap circuit to stabilise before implementing an A/D conversion.
- Bit 5~4** Unimplemented, read as "0"
- Bit 3 VREFS:** Selecte ADC reference voltage  
 0: Internal ADC power  $AV_{DD}$   
 1: VREF pin  
 This bit is used to select the reference voltage for the A/D converter. If the bit is 1 then the A/D converter reference voltage is supplied on the external pin VREF. If the pin is 0 then the internal reference is used which is taken from the power supply pin  $AV_{DD}$ . When the A/D converter reference voltage is supplied on the external pin VREF which is pin-shared with other functions, all of the other pin-shared functions on this pin are disabled.
- Bit 2 ~ 0 ADCK2 ~ ADCK0:** Select ADC clock source  
 000:  $f_{SYS}$   
 001:  $f_{SYS}/2$   
 010:  $f_{SYS}/4$   
 011:  $f_{SYS}/8$   
 100:  $f_{SYS}/16$   
 101:  $f_{SYS}/32$   
 110:  $f_{SYS}/64$   
 111: Undefined  
 These three bits are used to select the clock source for the A/D converter.

**ACERL Register**

Bit	7	6	5	4	3	2	1	0
Name	ACE7	ACE6	ACE5	ACE4	ACE3	ACE2	ACE1	ACE0
R/W								
POR	1	1	1	1	1	1	1	1

- Bit 7      **ACE7:** Define PB1 is A/D input or not  
0: Not A/D input  
1: A/D input, AN7
- Bit 6      **ACE6:** Define PC6 is A/D input or not  
0: Not A/D input  
1: A/D input, AN6
- Bit 5      **ACE5:** Define PC5 is A/D input or not  
0: Not A/D input  
1: A/D input, AN5
- Bit 4      **ACE4:** Define PC4 is A/D input or not  
0: Not A/D input  
1: A/D input, AN4
- Bit 3      **ACE3:** Define PC3 is A/D input or not  
0: Not A/D input  
1: A/D input, AN3
- Bit 2      **ACE2:** Define PC0 is A/D input or not  
0: Not A/D input  
1: A/D input, AN2
- Bit 1      **ACE1:** Define PB7 is A/D input or not  
0: Not A/D input  
1: A/D input, AN1
- Bit 0      **ACE0:** Define PB6 is A/D input or not  
0: Not A/D input  
1: A/D input, AN0

**A/D Operation**

The START bit in the ADCR0 register is used to start and reset the A/D converter. When the microcontroller sets this bit from low to high and then low again, an analog to digital conversion cycle will be initiated. When the START bit is brought from low to high but not low again, the EOCB bit in the ADCR0 register will be set high and the analog to digital converter will be reset. It is the START bit that is used to control the overall start operation of the internal analog to digital converter.

The EOCB bit in the ADCR0 register is used to indicate when the analog to digital conversion process is complete. This bit will be automatically set to “0” by the microcontroller after a conversion cycle has ended. In addition, the corresponding A/D interrupt request flag will be set in the interrupt control register, and if the interrupts are enabled, an appropriate internal interrupt signal will be generated. This A/D internal interrupt signal will direct the program flow to the associated A/D internal interrupt address for processing. If the A/D internal interrupt is disabled, the microcontroller can be used to poll the EOCB bit in the ADCR0 register to check whether it has been cleared as an alternative method of detecting the end of an A/D conversion cycle.

The clock source for the A/D converter, which originates from the system clock  $f_{SYS}$ , can be chosen to be either  $f_{SYS}$  or a subdivided version of  $f_{SYS}$ . The division ratio value is determined by the ADCK2~ADCK0 bits in the ADCR1 register.

Although the A/D clock source is determined by the system clock  $f_{SYS}$ , and by bits ADCK2~ADCK0, there are some limitations on the A/D clock source speed range that can be selected. As the recommended range of permissible A/D clock period,  $t_{AD}$ , is from 0.5 $\mu$ s to 10 $\mu$ s, care must be taken for system clock frequencies. For example, if the system clock operates at a frequency of 4MHz, the ADCK2~ADCK0 bits should not be set to 000B or 110B. Doing so will give A/D clock periods that are less than the minimum A/D clock period or greater than the maximum A/D clock period which may result in inaccurate A/D conversion values.

Refer to the following table for examples, where values marked with an asterisk \* show where, depending upon the device, special care must be taken, as the values may be less than the specified minimum A/D Clock Period.

$f_{SYS}$	A/D Clock Period ( $t_{AD}$ )							
	ADCK2, ADCK1, ADCK0 =000 ( $f_{SYS}$ )	ADCK2, ADCK1, ADCK0 =001 ( $f_{SYS}/2$ )	ADCK2, ADCK1, ADCK0 =010 ( $f_{SYS}/4$ )	ADCK2, ADCK1, ADCK0 =011 ( $f_{SYS}/8$ )	ADCK2, ADCK1, ADCK0 =100 ( $f_{SYS}/16$ )	ADCK2, ADCK1, ADCK0 =101 ( $f_{SYS}/32$ )	ADCK2, ADCK1, ADCK0 =110 ( $f_{SYS}/64$ )	ADCK2, ADCK1, ADCK0 =111
1MHz	1 $\mu$ s	2 $\mu$ s	4 $\mu$ s	8 $\mu$ s	16 $\mu$ s*	32 $\mu$ s*	64 $\mu$ s*	Undefined
2MHz	500ns	1 $\mu$ s	2 $\mu$ s	4 $\mu$ s	8 $\mu$ s	16 $\mu$ s*	32 $\mu$ s*	Undefined
4MHz	250ns*	500ns	1 $\mu$ s	2 $\mu$ s	4 $\mu$ s	8 $\mu$ s	16 $\mu$ s*	Undefined
8MHz	125ns*	250ns*	500ns	1 $\mu$ s	2 $\mu$ s	4 $\mu$ s	8 $\mu$ s	Undefined
12MHz	83ns*	167ns*	333ns*	667ns	1.33 $\mu$ s	2.67 $\mu$ s	5.33 $\mu$ s	Undefined
16MHz	62ns*	125ns*	250ns*	500ns	1 $\mu$ s	2 $\mu$ s	4 $\mu$ s	Undefined

#### A/D Clock Period Examples

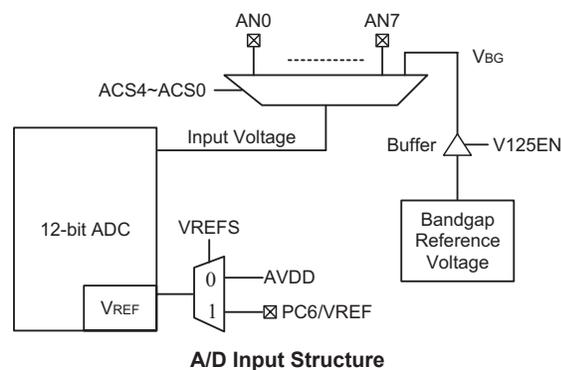
Controlling the power on/off function of the A/D converter circuitry is implemented using the ADOFF bit in the ADCR0 register. This bit must be zero to power on the A/D converter. When the ADOFF bit is cleared to zero to power on the A/D converter internal circuitry a certain delay, as indicated in the timing diagram, must be allowed before an A/D conversion is initiated. Even if no pins are selected for use as A/D inputs by clearing the ACE7~ACE0 bits in the ACERL registers, if the ADOFF bit is zero then some power will still be consumed. In power conscious applications it is therefore recommended that the ADOFF is set high to reduce power consumption when the A/D converter function is not being used.

The reference voltage supply to the A/D Converter can be supplied from either the positive power supply pin, AVDD, or from an external reference source supplied on pin VREF. The desired selection is made using the VREFS bit. As the pin is pin-shared with other function, when the VREF pin function is selected, then the other pin functions will be disabled automatically.

## A/D Input Pins

All of the A/D analog input pins are pin-shared with the I/O pins on Port C as well as other functions. The ACE7~ACE0 bits in the ACERL registers, determine whether the input pins are setup as A/D converter analog inputs or whether they have other functions. If the ACE7~ACE0 bits for its corresponding pin is set high then the pin will be setup to be an A/D converter input and the original pin functions disabled. In this way, pins can be changed under program control to change their function between A/D inputs and other functions. All pull-high resistors, which are setup through register programming, will be automatically disconnected if the pins are setup as A/D inputs. Note that it is not necessary to first setup the A/D pin as an input in the PCC port control register to enable the A/D input as when the ACE7~ACE0 bits enable an A/D input, the status of the port control register will be overridden.

The A/D converter has its own reference voltage pins, VREF, however the reference voltage can also be supplied from the power supply pin, a choice which is made through the VREFS bit in the ADCR1 register.



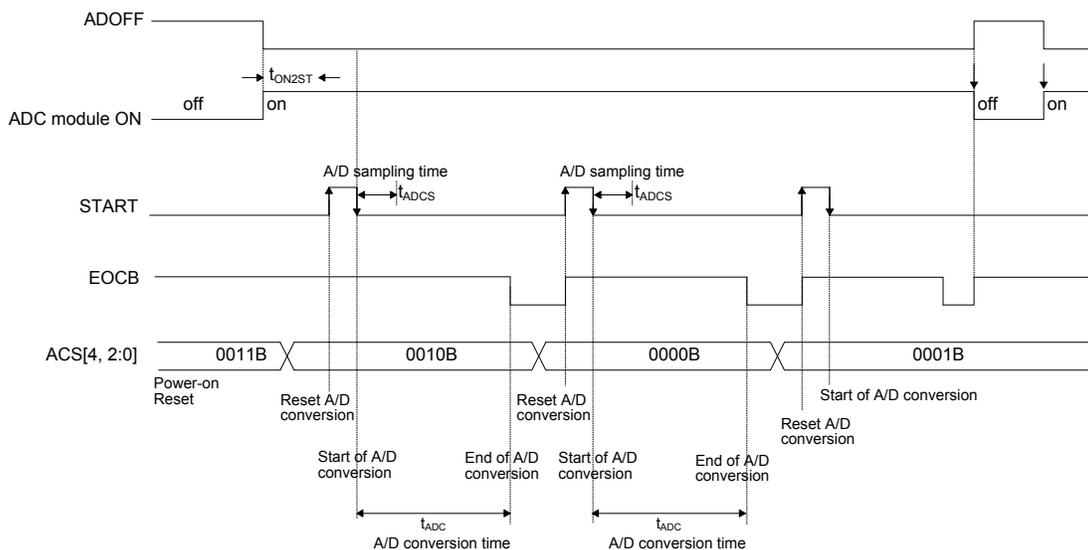
## Summary of A/D Conversion Steps

The following summarises the individual steps that should be executed in order to implement an A/D conversion process.

- Step 1  
Select the required A/D conversion clock by correctly programming bits ADCK2~ADCK0 in the ADCR1 register.
- Step 2  
Enable the A/D by clearing the ADOFF bit in the ADCR0 register to zero.
- Step 3  
Select which channel is to be connected to the internal A/D converter by correctly programming the ACS4, ACS2~ACS0 bits which are also contained in the ADCR1 and ADCR0 register.
- Step 4  
Select which pins are to be used as A/D inputs and configure them by correctly programming the ACE7~ACE0 bits in the ACERL register.
- Step 5  
If the interrupts are to be used, the interrupt control registers must be correctly configured to ensure the A/D converter interrupt function is active. The master interrupt control bit, EMI, and the A/D converter interrupt bit, ADE, must both be set high to do this.

- Step 6  
 The analog to digital conversion process can now be initialised by setting the START bit in the ADCR0 register from low to high and then low again. Note that this bit should have been originally cleared to zero.
- Step 7  
 To check when the analog to digital conversion process is complete, the EOCB bit in the ADCR0 register can be polled. The conversion process is complete when this bit goes low. When this occurs the A/D data registers ADRL and ADRH can be read to obtain the conversion value. As an alternative method, if the interrupts are enabled and the stack is not full, the program can wait for an A/D interrupt to occur.  
 Note: When checking for the end of the conversion process, if the method of polling the EOCB bit in the ADCR0 register is used, the interrupt enable step above can be omitted.

The accompanying diagram shows graphically the various stages involved in an analog to digital conversion process and its associated timing. After an A/D conversion process has been initiated by the application program, the microcontroller internal hardware will begin to carry out the conversion, during which time the program can continue with other functions. The time taken for the A/D conversion is  $16t_{AD}$  where  $t_{AD}$  is equal to the A/D clock period.



Note: A/D clock must be  $f_{sys}$ ,  $f_{sys}/2$ ,  $f_{sys}/4$ ,  $f_{sys}/8$ ,  $f_{sys}/16$ ,  $f_{sys}/32$ ,  $f_{sys}/64$ ,  
 $t_{ADCS} = 4t_{AD}$   
 $t_{ADC} = t_{ADCS} + n \cdot t_{AD}$ ; n=bit count of ADC resolution  
 $t_{ON2ST}$ : at least one instruction cycle ( $f_{sys} = 12\text{MHz}$ )

**A/D Conversion Timing**

**Programming Considerations**

During microcontroller operations where the A/D converter is not being used, the A/D internal circuitry can be switched off to reduce power consumption, by setting bit ADOFF high in the ADCR0 register. When this happens, the internal A/D converter circuits will not consume power irrespective of what analog voltage is applied to their input lines. If the A/D converter input lines are used as normal I/Os, then care must be taken as if the input voltage is not at a valid logic level, then this may lead to some increase in power consumption.

**A/D Transfer Function**

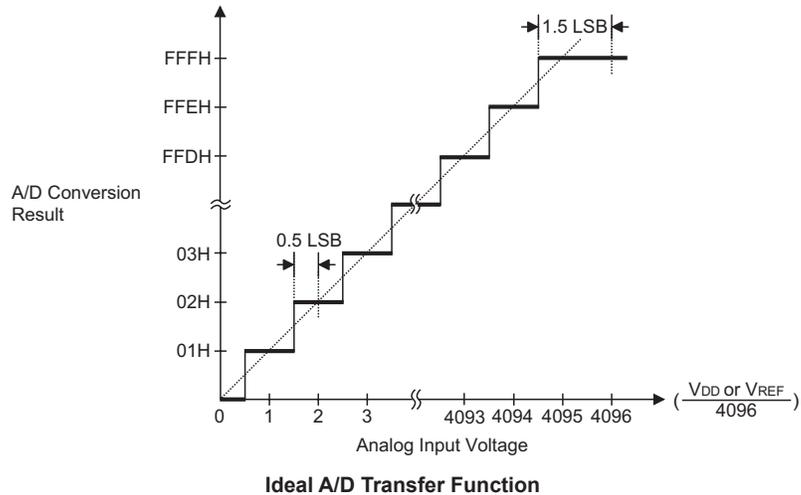
As the device contains a 12-bit A/D converter, its full-scale converted digitised value is equal to FFFH. Since the full-scale analog input value is equal to the  $AV_{DD}$  or  $V_{REF}$  voltage, this gives a single bit analog input value of  $AV_{DD}$  or  $V_{REF}$  divided by 4096.

$$1 \text{ LSB} = (AV_{DD} \text{ or } V_{REF}) / 4096$$

The A/D Converter input voltage value can be calculated using the following equation:

$$\text{A/D input voltage} = \text{A/D output digital value} \times (AV_{DD} \text{ or } V_{REF}) / 4096$$

The diagram shows the ideal transfer function between the analog input value and the digitised output value for the A/D converter. Except for the digitised zero value, the subsequent digitised values will change at a point 0.5 LSB below where they would change without the offset, and the last full scale digitised value will change at a point 1.5 LSB below the  $AV_{DD}$  or  $V_{REF}$  level.



## A/D Programming Examples

The following two programming examples illustrate how to setup and implement an A/D conversion. In the first example, the method of polling the EOCB bit in the ADCR0 register is used to detect when the conversion cycle is complete, whereas in the second example, the A/D interrupt is used to determine when the conversion is complete.

### Example: using an EOCB polling method to detect the end of conversion

```
clr ADE                ; disable ADC interrupt
mov a,03H
mov ADCR1,a            ; select fsys/8 as A/D clock and switch off VBG
clr ADOFF
mov a,0Fh              ; setup ACERL to configure pins AN0~AN3
mov ACERL,a
mov a,00h
mov ADCR0,a            ; enable and connect AN0 channel to A/D converter
:
start_conversion:
clr START              ; high pulse on start bit to initiate conversion
set START              ; reset A/D
clr START              ; start A/D
polling_EOC:
sz EOCB                ; poll the ADCR0 register EOCB bit to detect end of A/D conversion
jmp polling_EOC        ; continue polling
mov a,ADRL              ; read low byte conversion result value
mov ADRL_buffer,a      ; save result to user defined register
mov a,ADRH              ; read high byte conversion result value
mov ADRH_buffer,a      ; save result to user defined register
:
:
jmp start_conversion   ; start next a/d conversion
```

**Example: using the interrupt method to detect the end of conversion**

```
clr ADE          ; disable ADC interrupt
mov a,03H
mov ADCR1,a      ; select fsys/8 as A/D clock and switch off VBG
Clr ADOFF
mov a,0Fh        ; setup ACERL to configure pins AN0~AN3
mov ACERL,a
mov a,00h
mov ADCR0,a      ; enable and connect AN0 channel to A/D converter
Start_conversion:
clr START        ; high pulse on START bit to initiate conversion
set START        ; reset A/D
clr START        ; start A/D
clr ADF          ; clear ADC interrupt request flag
set ADE          ; enable ADC interrupt
set EMI          ; enable global interrupt
:
:
; ADC interrupt service routine
ADC_ISR:
mov acc_stack,a  ; save ACC to user defined memory
mov a,STATUS
mov status_stack,a ; save STATUS to user defined memory
:
:
mov a,ADRL        ; read low byte conversion result value
mov adr1_buffer,a ; save result to user defined register
mov a,ADRH        ; read high byte conversion result value
mov adrh_buffer,a ; save result to user defined register
:
:
EXIT_INT_ISR:
mov a,status_stack
mov STATUS,a     ; restore STATUS from user defined memory
mov a,acc_stack  ; restore ACC from user defined memory
reti
```

## Serial Interface Module – SIM

The device contains a Serial Interface Module, which includes both the four-line SPI interface or two-line I<sup>2</sup>C interface types, to allow an easy method of communication with external peripheral hardware. Having relatively simple communication protocols, these serial interface types allow the microcontroller to interface to external SPI or I<sup>2</sup>C based hardware such as sensors, Flash or EEPROM memory, etc. The SIM interface pins are pin-shared with other I/O pins and therefore the SIM interface functional pins must first be selected using the corresponding pin-shared function selection bits. As both interface types share the same pins and registers, the choice of whether the SPI or I<sup>2</sup>C type is used is made using the SIM operating mode control bits, named SIM2~SIM0, in the SIMC0 register. These pull-high resistors of the SIM pin-shared I/O pins are selected using pull-high control registers when the SIM function is enabled and the corresponding pins are used as SIM input pins.

### SPI Interface

The SPI interface is often used to communicate with external peripheral devices such as sensors, Flash Memory or EEPROM memory devices etc. Originally developed by Motorola, the four line SPI interface is a synchronous serial data interface that has a relatively simple communication protocol simplifying the programming requirements when communicating with external hardware devices.

The communication is full duplex and operates as a slave/master type, where the device can be either master or slave. Although the SPI interface specification can control multiple slave devices from a single master, but the device provides only one  $\overline{SCS}$  pin. If the master needs to control multiple slave devices from a single master, the master can use I/O pin to select the slave devices.

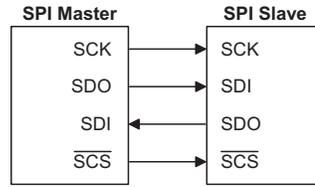
### SPI Interface Operation

The SPI interface is a full duplex synchronous serial data link. It is a four line interface with pin names SDI, SDO, SCK and  $\overline{SCS}$ . Pins SDI and SDO are the Serial Data Input and Serial Data Output lines, SCK is the Serial Clock line and  $\overline{SCS}$  is the Slave Select line. As the SPI interface pins are pin-shared with normal I/O pins and with the I<sup>2</sup>C function pins, the SPI interface pins must first be selected by configuring the pin-shared function selection bits and setting the correct bits in the SIMC0 and SIMC2 registers. After the desired SPI configuration has been set it can also be disabled or enabled using the SIMEN bit in the SIMC0 register. Communication between devices connected to the SPI interface is carried out in a slave/master mode with all data transfer initiations being implemented by the master. The Master also controls the clock signal. As the device only contains a single  $\overline{SCS}$  pin only one slave device can be utilized. The  $\overline{SCS}$  pin is controlled by software, set CSEN bit to “1” to enable  $\overline{SCS}$  pin function, set CSEN bit to “0” the  $\overline{SCS}$  pin will be floating state.

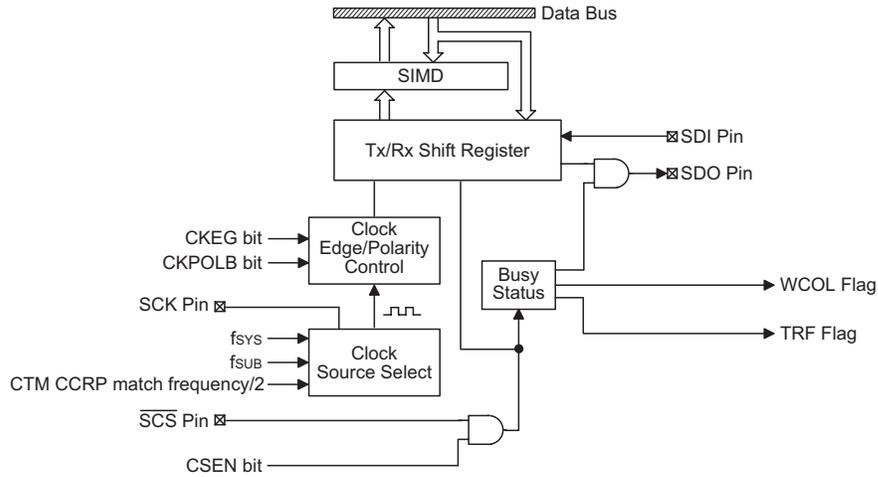
The SPI function in this device offers the following features:

- Full duplex synchronous data transfer
- Both Master and Slave modes
- LSB first or MSB first data transmission modes
- Transmission complete flag
- Rising or falling active clock edge

The status of the SPI interface pins is determined by a number of factors such as whether the device is in the master or slave mode and upon the condition of certain control bits such as CSEN and SIMEN.



**SPI Master/Slave Connection**



**SPI Block Diagram**

**SPI Register**

There are three internal registers which control the overall operation of the SPI interface. These are the SIMD data register and two control registers SIMC0 and SIMC2.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SIMC0	SIM2	SIM1	SIM0	—	SIMDBNC1	SIMDBNC0	SIMEN	SPIICF
SIMD	D7	D6	D5	D4	D3	D2	D1	D0
SIMC2	D7	D6	CKPOLB	CKEG	MLS	CSEN	WCOL	TRF

**SIM Registers List**

The SIMD register is used to store the data being transmitted and received. The same register is used by both the SPI and I<sup>2</sup>C functions. Before the devices write data to the SPI bus, the actual data to be transmitted must be placed in the SIMD register. After the data is received from the SPI bus, the device can read it from the SIMD register. Any transmission or reception of data from the SPI bus must be made via the SIMD register.

### SIMD Register

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

"x": unknown

There are also two control registers for the SPI interface, SIMC0 and SIMC2. Note that the SIMC2 register also has the name SIMA which is used by the I<sup>2</sup>C function. The SIMC1 register is not used by the SPI function, only by the I<sup>2</sup>C function. Register SIMC0 is used to control the enable/disable function and to set the data transmission clock frequency. Although not connected with the SPI function, the SIMC0 register is also used to control the Peripheral Clock Prescaler. Register SIMC2 is used for other control functions such as LSB/MSB selection, write collision flag etc.

### SIMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	SIM2	SIM1	SIM0	—	SIMDBNC1	SIMDBNC0	SIMEN	SPIICF
R/W	R/W	R/W	R/W	—	R/W	R/W	R/W	R/W
POR	1	1	1	—	0	0	0	0

Bit 7 ~ 5 **SIM2~SIM0**: SIM Operating Mode Control  
 000: SPI master mode; SPI clock is  $f_{SYS}/4$   
 001: SPI master mode; SPI clock is  $f_{SYS}/16$   
 010: SPI master mode; SPI clock is  $f_{SYS}/64$   
 011: SPI master mode; SPI clock is  $f_{SUB}$   
 100: SPI master mode; SPI clock is CTM CCRP match frequency/2  
 101: SPI slave mode  
 110: I<sup>2</sup>C slave mode  
 111: Unused mode

These bits setup the overall operating mode of the SIM function. As well as selecting if the I<sup>2</sup>C or SPI function, they are used to control the SPI Master/Slave selection and the SPI Master clock frequency. The SPI clock is a function of the system clock but can also be chosen to be sourced from the CTM. If the SPI Slave Mode is selected then the clock will be supplied by an external Master device.

Bit 4 Unimplemented, read as "0"

Bit 3 ~ 2 **SIMDBNC1~SIMDBNC0**: I<sup>2</sup>C Debounce Time Selection  
 00: No debounce  
 01: 2 system clocks debounce  
 1x: 4 system clocks debounce

Bit 1 **SIMEN**: SIM Control  
 0: Disable  
 1: Enable

The bit is the overall on/off control for the SIM interface. When the SIMEN bit is cleared to zero to disable the SIM interface, the SDI, SDO, SCK and  $\overline{SCS}$ , or SDA and SCL lines will lose their SPI or I<sup>2</sup>C function and the SIM operating current will be reduced to a minimum value. When the bit is high the SIM interface is enabled. If the SIM is configured to operate as an SPI interface via the SIM2~SIM0 bits, the contents of the SPI control registers will remain at the previous settings when the SIMEN bit changes from low to high and should therefore be first initialised by the application program. If the SIM is configured to operate as an I<sup>2</sup>C interface via the SIM2~SIM0 bits and the SIMEN bit changes from low to high, the contents of the I<sup>2</sup>C control bits such as HTX and TXAK will remain at the previous settings and should therefore be first initialised by the application program while the relevant I<sup>2</sup>C flags such as HCF, HAAS, HBB, SRW and RXAK will be set to their default states.

**Bit 0 SPIICF: SPI Incompleted Flag**  
 0: SPI incompleted is not occurred  
 1: SPI incompleted is occurred  
 The SPIICF bit is determined by  $\overline{SCS}$  pin. When  $\overline{SCS}$  pin is set to “1”, it will clear the SPI counter. Meanwhile, the interrupt is occurred, if slave device didn’t complete data received, then the incompleted flag, SPIICF, is set to “1”.

**SIMC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	CKPOLB	CKEG	MLS	CSEN	WCOL	TRF
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

**Bit 7 ~ 6** Undefined bits  
 These bits can be read or written by user software program.

**Bit 5 CKPOLB: Determines the base condition of the clock line**  
 0: The SCK line will be high when the clock is inactive  
 1: The SCK line will be low when the clock is inactive  
 The CKPOLB bit determines the base condition of the clock line, if the bit is high, then the SCK line will be low when the clock is inactive. When the CKPOLB bit is low, then the SCK line will be high when the clock is inactive.

**Bit 4 CKEG: Determines SPI SCK active clock edge type**  
 CKPOLB=0  
 0: SCK is high base level and data capture at SCK rising edge  
 1: SCK is high base level and data capture at SCK falling edge  
 CKPOLB=1  
 0: SCK is low base level and data capture at SCK falling edge  
 1: SCK is low base level and data capture at SCK rising edge  
 The CKEG and CKPOLB bits are used to setup the way that the clock signal outputs and inputs data on the SPI bus. These two bits must be configured before data transfer is executed otherwise an erroneous clock edge may be generated. The CKPOLB bit determines the base condition of the clock line, if the bit is high, then the SCK line will be low when the clock is inactive. When the CKPOLB bit is low, then the SCK line will be high when the clock is inactive. The CKEG bit determines active clock edge type which depends upon the condition of CKPOLB bit.

**Bit 3 MLS: SPI Data shift order**  
 0: LSB  
 1: MSB  
 This is the data shift select bit and is used to select how the data is transferred, either MSB or LSB first. Setting the bit high will select MSB first and low for LSB first.

**Bit 2 CSEN: SPI  $\overline{SCS}$  pin Control**  
 0: Disable  
 1: Enable  
 The CSEN bit is used as an enable/disable for the  $\overline{SCS}$  pin. If this bit is low, then the  $\overline{SCS}$  pin will be disabled and placed into I/O pin or the other functions. If the bit is high the  $\overline{SCS}$  pin will be enabled and used as a select pin.

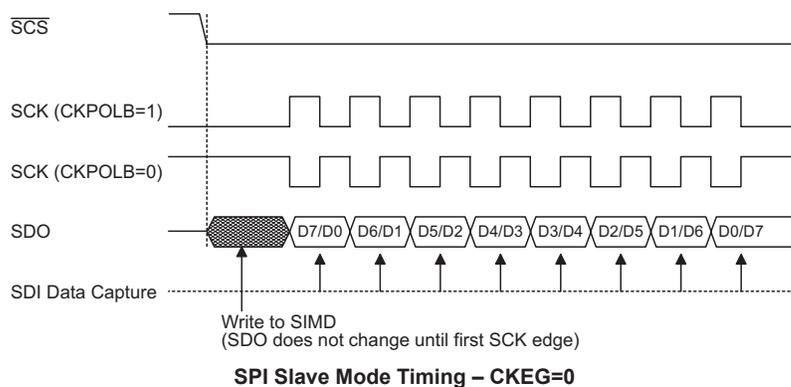
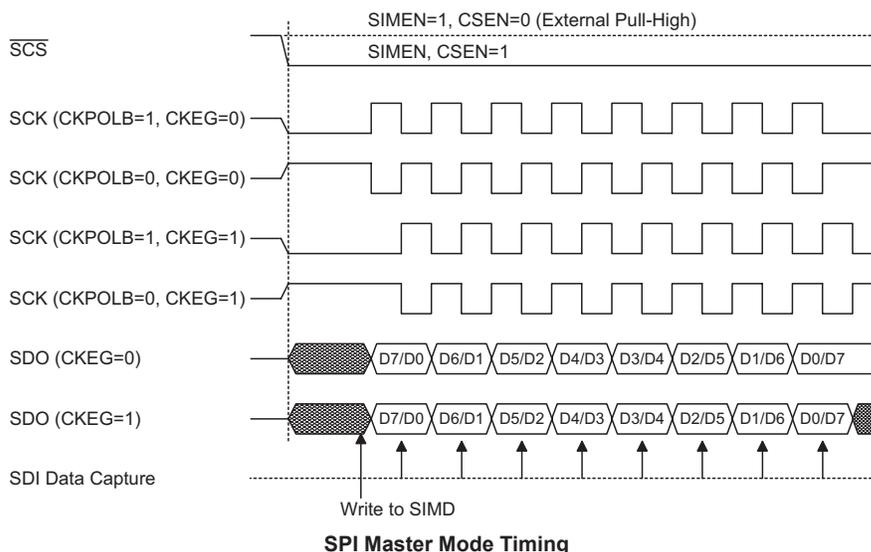
**Bit 1 WCOL: SPI Write Collision flag**  
 0: No collision  
 1: Collision  
 The WCOL flag is used to detect if a data collision has occurred. If this bit is high it means that data has been attempted to be written to the SIMD register during a data transfer operation. This writing operation will be ignored if data is being transferred. The bit can be cleared by the application program.

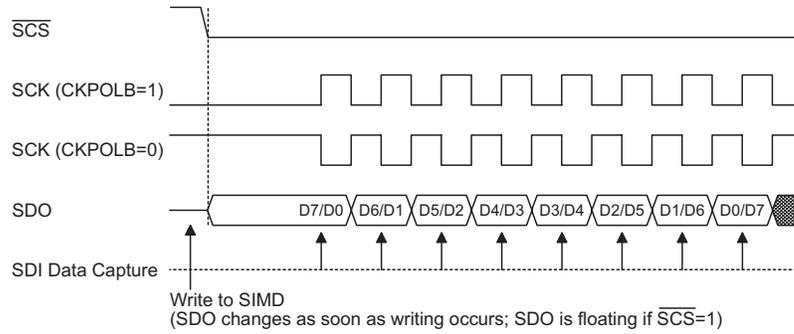
Bit 0 **TRF**: SPI Transmit/Receive Complete flag  
 0: Data is being transferred  
 1: SPI data transmission is completed

The TRF bit is the Transmit/Receive Complete flag and is set “1” automatically when an SPI data transmission is completed, but must set to “0” by the application program. It can be used to generate an interrupt.

### SPI Communication

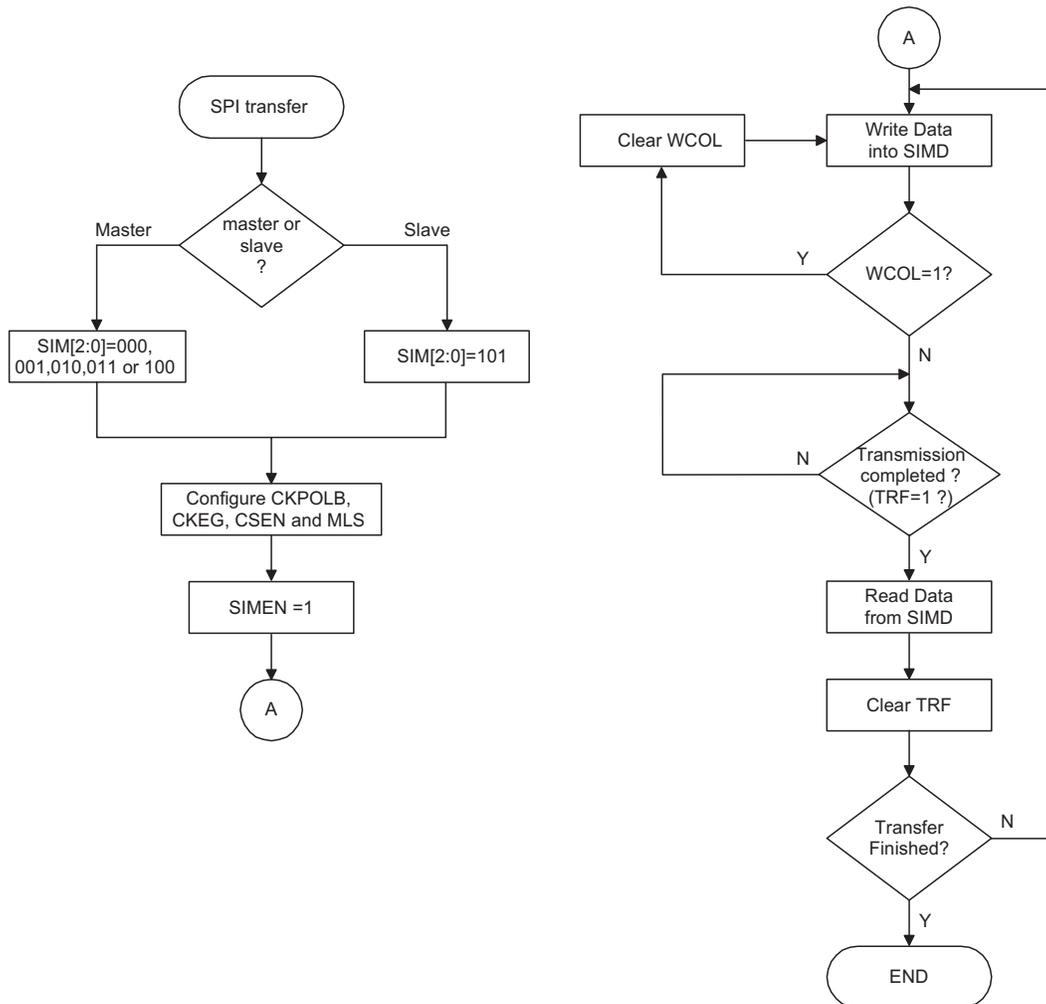
After the SPI interface is enabled by setting the SIMEN bit high, then in the Master Mode, when data is written to the SIMD register, transmission/reception will begin simultaneously. When the data transfer is complete, the TRF flag will be set automatically, but must be cleared using the application program. In the Slave Mode, when the clock signal from the master has been received, any data in the SIMD register will be transmitted and any data on the SDI pin will be shifted into the SIMD register. The master should output a  $\overline{SCS}$  signal to enable the slave device before a clock signal is provided. The slave data to be transferred should be well prepared at the appropriate moment relative to the  $\overline{SCS}$  signal depending upon the configurations of the CKPOLB bit and CKEG bit. The accompanying timing diagram shows the relationship between the slave data and  $\overline{SCS}$  signal for various configurations of the CKPOLB and CKEG bits. The SPI will continue to function even in the IDLE Mode.





**Note:** For SPI slave mode, if SIMEN=1 and CSEN=0, SPI is always enabled and ignores the  $\overline{SCS}$  level.

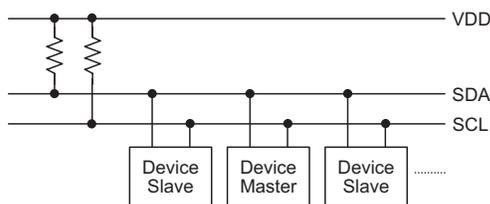
**SPI Slave Mode Timing – CKEG=1**



**SPI Transfer Control Flowchart**

## I<sup>2</sup>C Interface

The I<sup>2</sup>C interface is used to communicate with external peripheral devices such as sensors, EEPROM memory etc. Originally developed by Philips, it is a two line low speed serial interface for synchronous serial data transfer. The advantage of only two lines for communication, relatively simple communication protocol and the ability to accommodate multiple devices on the same bus has made it an extremely popular interface type for many applications.

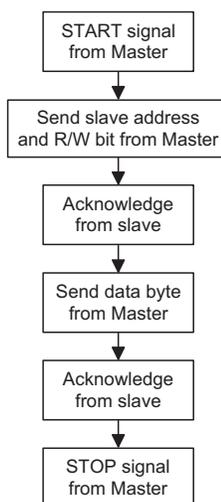


**I<sup>2</sup>C Master Slave Bus Connection**

## I<sup>2</sup>C Interface Operation

The I<sup>2</sup>C serial interface is a two line interface, a serial data line, SDA, and serial clock line, SCL. As many devices may be connected together on the same bus, their outputs are both open drain types. For this reason it is necessary that external pull-high resistors are connected to these outputs. Note that no chip select line exists, as each device on the I<sup>2</sup>C bus is identified by a unique address which will be transmitted and received on the I<sup>2</sup>C bus.

When two devices communicate with each other on the bidirectional I<sup>2</sup>C bus, one is known as the master device and one as the slave device. Both master and slave can transmit and receive data, however, it is the master device that has overall control of the bus. For these devices, which only operates in slave mode, there are two methods of transferring data on the I<sup>2</sup>C bus, the slave transmit mode and the slave receive mode. The pull-up control function pin-shared with SCL/SDA pin is still applicable even if I<sup>2</sup>C device is activated and the related internal pull-up register could be controlled by its corresponding pull-up control register.



## I<sup>2</sup>C Register

There are four control registers associated with the I<sup>2</sup>C bus, SIMC0, SIMC1, SIMA and SIMTOC and one data register, SIMD. The SIMD register, which is shown in the above SPI section, is used to store the data being transmitted and received on the I<sup>2</sup>C bus. Before the microcontroller writes data to the I<sup>2</sup>C bus, the actual data to be transmitted must be placed in the SIMD register. After the data is received from the I<sup>2</sup>C bus, the microcontroller can read it from the SIMD register. Any transmission or reception of data from the I<sup>2</sup>C bus must be made via the SIMD register.

Note that the SIMA register also has the name SIMC2 which is used by the SPI function. Bit SIMEN and bits SIM2~SIM0 in register SIMC0 are used by the I<sup>2</sup>C interface. The SIMTOC register is used for I<sup>2</sup>C time-out control function.

Register Name	Bit							
	7	6	5	4	3	2	1	0
SIMC0	SIM2	SIM1	SIM0	—	SIMDEB1	SIMDEB0	SIMEN	SPIICF
SIMC1	HCF	HAAS	HBB	HTX	TXAK	SRW	RCIN	RXAK
SIMD	D7	D6	D5	D4	D3	D2	D1	D0
SIMA	IICA6	IICA5	IICA4	IICA3	IICA2	IICA1	IICA0	D0
SIMTOC	SIMTOEN	SIMTOF	SIMTOS5	SIMTOS4	SIMTOS3	SIMTOS2	SIMTOS1	SIMTOS0

I<sup>2</sup>C Register List

### SIMC0 Register

Bit	7	6	5	4	3	2	1	0
Name	SIM2	SIM1	SIM0	—	SIMDBNC1	SIMDBNC0	SIMEN	SPIICF
R/W	R/W	R/W	R/W	—	R/W	R/W	R/W	R/W
POR	1	1	1	—	0	0	0	0

Bit 7 ~ 5 **SIM2~SIM0**: SIM Operating Mode Control  
 000: SPI master mode; SPI clock is  $f_{SYS}/4$   
 001: SPI master mode; SPI clock is  $f_{SYS}/16$   
 010: SPI master mode; SPI clock is  $f_{SYS}/64$   
 011: SPI master mode; SPI clock is  $f_{SUB}$   
 100: SPI master mode; SPI clock is CTM CCRP match frequency/2  
 101: SPI slave mode  
 110: I<sup>2</sup>C slave mode  
 111: Unused mode

These bits setup the overall operating mode of the SIM function. As well as selecting if the I<sup>2</sup>C or SPI function, they are used to control the SPI Master/Slave selection and the SPI Master clock frequency. The SPI clock is a function of the system clock but can also be chosen to be sourced from the CTM. If the SPI Slave Mode is selected then the clock will be supplied by an external Master device.

Bit 4 Unimplemented, read as “0”

Bit 3 ~ 2 **SIMDBNC1~SIMDBNC0**: I<sup>2</sup>C Debounce Time Selection  
 00: No debounce  
 01: 2 system clocks debounce  
 1x: 4 system clocks debounce

- Bit 1 SIMEN: SIM Control**  
 0: Disable  
 1: Enable
- The bit is the overall on/off control for the SIM interface. When the SIMEN bit is cleared to zero to disable the SIM interface, the SDI, SDO, SCK and  $\overline{SCS}$ , or SDA and SCL lines will lose their SPI or I<sup>2</sup>C function and the SIM operating current will be reduced to a minimum value. When the bit is high the SIM interface is enabled. If the SIM is configured to operate as an SPI interface via the SIM2~SIM0 bits, the contents of the SPI control registers will remain at the previous settings when the SIMEN bit changes from low to high and should therefore be first initialised by the application program. If the SIM is configured to operate as an I<sup>2</sup>C interface via the SIM2~SIM0 bits and the SIMEN bit changes from low to high, the contents of the I<sup>2</sup>C control bits such as HTX and TXAK will remain at the previous settings and should therefore be first initialised by the application program while the relevant I<sup>2</sup>C flags such as HCF, HAAS, HBB, SRW and RXAK will be set to their default states.
- Bit 0 SPIICF: SPI Incompleted Flag**  
 0: SPI incompleted is not occurred  
 1: SPI incompleted is occurred
- The SPIICF bit is determined by  $\overline{SCS}$  pin. When  $\overline{SCS}$  pin is set to “1”, it will clear the SPI counter. Meanwhile, the interrupt is occurred, if slave device didn't complete data received, then the incompleted flag, SPIICF, is set to”1”.

**SIMC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	HCF	HAAS	HBB	HTX	TXAK	SRW	IAMWU	RXAK
R/W	R	R	R	R/W	R/W	R	R/W	R
POR	1	0	0	0	0	0	0	1

- Bit 7 HCF: I<sup>2</sup>C Bus data transfer completion flag**  
 0: Data is being transferred  
 1: Completion of an 8-bit data transfer
- The HCF flag is the data transfer flag. This flag will be zero when data is being transferred. Upon completion of an 8-bit data transfer the flag will go high and an interrupt will be generated.
- Bit 6 HAAS: I<sup>2</sup>C Bus address match flag**  
 0: Not address match  
 1: Address match
- The HAAS flag is the address match flag. This flag is used to determine if the slave device address is the same as the master transmit address. If the addresses match then this bit will be high, if there is no match then the flag will be low.
- Bit 5 HBB: I<sup>2</sup>C Bus busy flag**  
 0: I<sup>2</sup>C Bus is not busy  
 1: I<sup>2</sup>C Bus is busy
- The HBB flag is the I<sup>2</sup>C busy flag. This flag will be “1” when the I<sup>2</sup>C bus is busy which will occur when a START signal is detected. The flag will be set to “0” when the bus is free which will occur when a STOP signal is detected.
- Bit 4 HTX: Select I<sup>2</sup>C slave device is transmitter or receiver**  
 0: Slave device is the receiver  
 1: Slave device is the transmitter

- Bit 3      **TXAK:** I<sup>2</sup>C Bus transmit acknowledge flag  
             0: Slave send acknowledge flag  
             1: Slave do not send acknowledge flag  
 The TXAK bit is the transmit acknowledge flag. After the slave device receipt of 8-bits of data, this bit will be transmitted to the bus on the 9<sup>th</sup> clock from the slave device. The slave device must always set TXAK bit to “0” before further data is received.
- Bit 2      **SRW:** I<sup>2</sup>C Slave Read/Write flag  
             0: Slave device should be in receive mode  
             1: Slave device should be in transmit mode  
 The SRW flag is the I<sup>2</sup>C Slave Read/Write flag. This flag determines whether the master device wishes to transmit or receive data from the I<sup>2</sup>C bus. When the transmitted address and slave address is match, that is when the HAAS flag is set high, the slave device will check the SRW flag to determine whether it should be in transmit mode or receive mode. If the SRW flag is high, the master is requesting to read data from the bus, so the slave device should be in transmit mode. When the SRW flag is zero, the master will write data to the bus, therefore the slave device should be in receive mode to read this data.
- Bit 1      **IAMWU:** I<sup>2</sup>C Address Match Wake-up Control  
             0: Disable  
             1: Enable – must be cleared by the application program after wake-up  
 This bit should be set to 1 to enable the I<sup>2</sup>C address match wake up from the SLEEP or IDLE Mode. If the IAMWU bit has been set before entering the SLEEP or IDLE mode to enable the I<sup>2</sup>C address match wake up, then this bit must be cleared by the application program after wake-up to ensure correction device operation.
- Bit 0      **RXAK:** I<sup>2</sup>C Bus Receive acknowledge flag  
             0: Slave receive acknowledge flag  
             1: Slave do not receive acknowledge flag  
 The RXAK flag is the receiver acknowledge flag. When the RXAK flag is “0”, it means that a acknowledge signal has been received at the 9<sup>th</sup> clock, after 8 bits of data have been transmitted. When the slave device in the transmit mode, the slave device checks the RXAK flag to determine if the master receiver wishes to receive the next byte. The slave transmitter will therefore continue sending out data until the RXAK flag is “1”. When this occurs, the slave transmitter will release the SDA line to allow the master to send a STOP signal to release the I<sup>2</sup>C Bus.  
 The SIMD register is used to store the data being transmitted and received. The same register is used by both the SPI and I<sup>2</sup>C functions. Before the device writes data to the I<sup>2</sup>C bus, the actual data to be transmitted must be placed in the SIMD register. After the data is received from the I<sup>2</sup>C bus, the device can read it from the SIMD register. Any transmission or reception of data from the I<sup>2</sup>C bus must be made via the SIMD register.

### SIMD Register

Bit	7	6	5	4	3	2	1	0
Name	D7	D6	D5	D4	D3	D2	D1	D0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

"x": unknown

### SIMA Register

Bit	7	6	5	4	3	2	1	0
Name	A6	A5	A4	A3	A2	A1	A0	—
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	—
POR	0	0	0	0	0	0	0	—

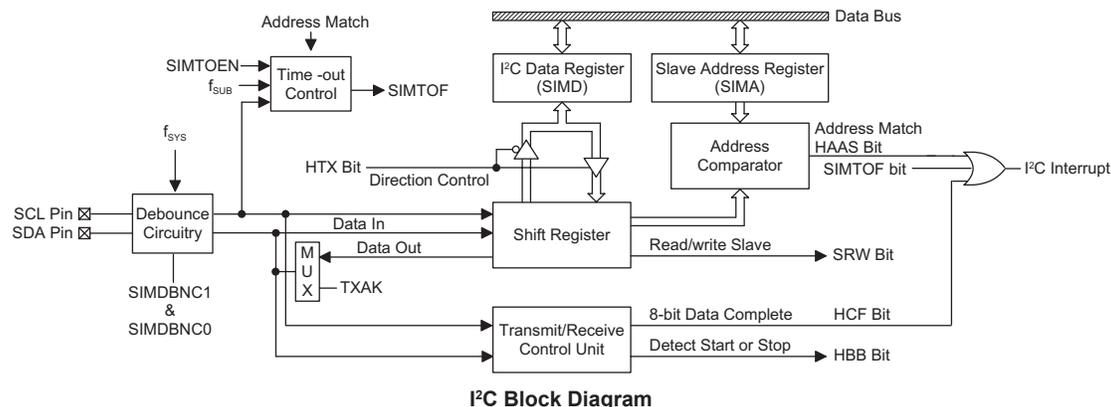
bit 7 ~ 1 **A6~A0: I<sup>2</sup>C slave address**

A6~A0 is the I<sup>2</sup>C slave address bit 6 ~ bit 0.

The SIMA register is also used by the SPI interface but has the name SIMC2. The SIMA register is the location where the 7-bit slave address of the slave device is stored. Bits 7~1 of the SIMA register define the device slave address. Bit 0 is not defined. When a master device, which is connected to the I<sup>2</sup>C bus, sends out an address, which matches the slave address in the SIMA register, the slave device will be selected. Note that the SIMA register is the same register address as SIMC2 which is used by the SPI interface.

Bit 0 **Undefined bit**

This bit can be read or written by user software program.



**I<sup>2</sup>C Bus Communication**

Communication on the I<sup>2</sup>C bus requires four separate steps, a START signal, a slave device address transmission, a data transmission and finally a STOP signal. When a START signal is placed on the I<sup>2</sup>C bus, all devices on the bus will receive this signal and be notified of the imminent arrival of data on the bus. The first seven bits of the data will be the slave address with the first bit being the MSB. If the address of the slave device matches that of the transmitted address, the HAAS bit in the SIMC1 register will be set and an I<sup>2</sup>C interrupt will be generated. After entering the interrupt service routine, the slave device must first check the condition of the HAAS bit to determine whether the interrupt source originates from an address match or from the completion of an 8-bit data transfer. During a data transfer, note that after the 7-bit slave address has been transmitted, the following bit, which is the 8<sup>th</sup> bit, is the read/write bit whose value will be placed in the SRW bit. This bit will be checked by the slave device to determine whether to go into transmit or receive mode. Before any transfer of data to or from the I<sup>2</sup>C bus, the microcontroller must initialize the bus, the following are steps to achieve this:

**Step 1**

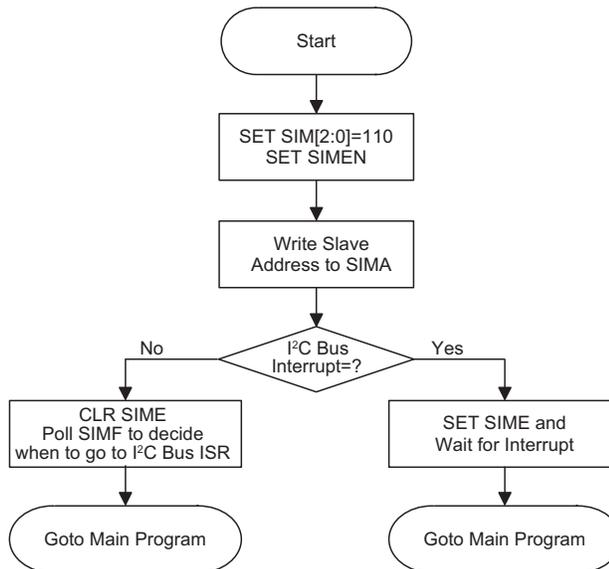
Set the SIM2~SIM0 and SIMEN bits in the SIMC0 register to “110” and “1” to enable the I<sup>2</sup>C bus.

**Step 2**

Write the slave address of the device to the I<sup>2</sup>C bus address register SIMA.

**Step 3**

Set the SIME bit of the interrupt control register to enable the SIM interrupt.



**I<sup>2</sup>C Bus Initialisation Flow Chart**

## **I<sup>2</sup>C Bus Start Signal**

The START signal can only be generated by the master device connected to the I<sup>2</sup>C bus and not by the slave device. This START signal will be detected by all devices connected to the I<sup>2</sup>C bus. When detected, this indicates that the I<sup>2</sup>C bus is busy and therefore the HBB bit will be set. A START condition occurs when a high to low transition on the SDA line takes place when the SCL line remains high.

## **Slave Address**

The transmission of a START signal by the master will be detected by all devices on the I<sup>2</sup>C bus. To determine which slave device the master wishes to communicate with, the address of the slave device will be sent out immediately following the START signal. All slave devices, after receiving this 7-bit address data, will compare it with their own 7-bit slave address. If the address sent out by the master matches the internal address of the microcontroller slave device, then an internal I<sup>2</sup>C bus interrupt signal will be generated. The next bit following the address, which is the 8<sup>th</sup> bit, defines the read/write status and will be saved to the SRW bit of the SIMC1 register. The slave device will then transmit an acknowledge bit, which is a low level, as the 9<sup>th</sup> bit. The slave device will also set the status flag HAAS when the addresses match.

As an I<sup>2</sup>C bus interrupt can come from two sources, when the program enters the interrupt subroutine, the HAAS bit should be examined to see whether the interrupt source has come from a matching slave address or from the completion of a data byte transfer. When a slave address is matched, the device must be placed in either the transmit mode and then write data to the SIMD register, or in the receive mode where it must implement a dummy read from the SIMD register to release the SCL line.

## **I<sup>2</sup>C Bus Read/Write Signal**

The SRW bit in the SIMC1 register defines whether the slave device wishes to read data from the I<sup>2</sup>C bus or write data to the I<sup>2</sup>C bus. The slave device should examine this bit to determine if it is to be a transmitter or a receiver. If the SRW flag is “1” then this indicates that the master device wishes to read data from the I<sup>2</sup>C bus, therefore the slave device must be setup to send data to the I<sup>2</sup>C bus as a transmitter. If the SRW flag is “0” then this indicates that the master wishes to send data to the I<sup>2</sup>C bus, therefore the slave device must be setup to read data from the I<sup>2</sup>C bus as a receiver.

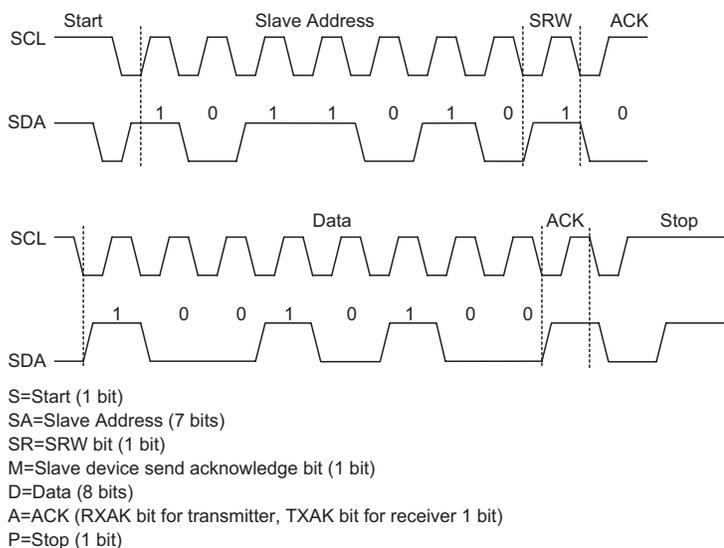
## **I<sup>2</sup>C Bus Slave Address Acknowledge Signal**

After the master has transmitted a calling address, any slave device on the I<sup>2</sup>C bus, whose own internal address matches the calling address, must generate an acknowledge signal. The acknowledge signal will inform the master that a slave device has accepted its calling address. If no acknowledge signal is received by the master then a STOP signal must be transmitted by the master to end the communication. When the HAAS flag is high, the addresses have matched and the slave device must check the SRW flag to determine if it is to be a transmitter or a receiver. If the SRW flag is high, the slave device should be setup to be a transmitter so the HTX bit in the SIMC1 register should be set to “1”. If the SRW flag is low, then the microcontroller slave device should be setup as a receiver and the HTX bit in the SIMC1 register should be set to “0”.

**I<sup>2</sup>C Bus Data and Acknowledge Signal**

The transmitted data is 8-bits wide and is transmitted after the slave device has acknowledged receipt of its slave address. The order of serial bit transmission is the MSB first and the LSB last. After receipt of 8-bits of data, the receiver must transmit an acknowledge signal, level “0”, before it can receive the next data byte. If the slave transmitter does not receive an acknowledge bit signal from the master receiver, then the slave transmitter will release the SDA line to allow the master to send a STOP signal to release the I<sup>2</sup>C Bus. The corresponding data will be stored in the SIMD register. If setup as a transmitter, the slave device must first write the data to be transmitted into the SIMD register. If setup as a receiver, the slave device must read the transmitted data from the SIMD register.

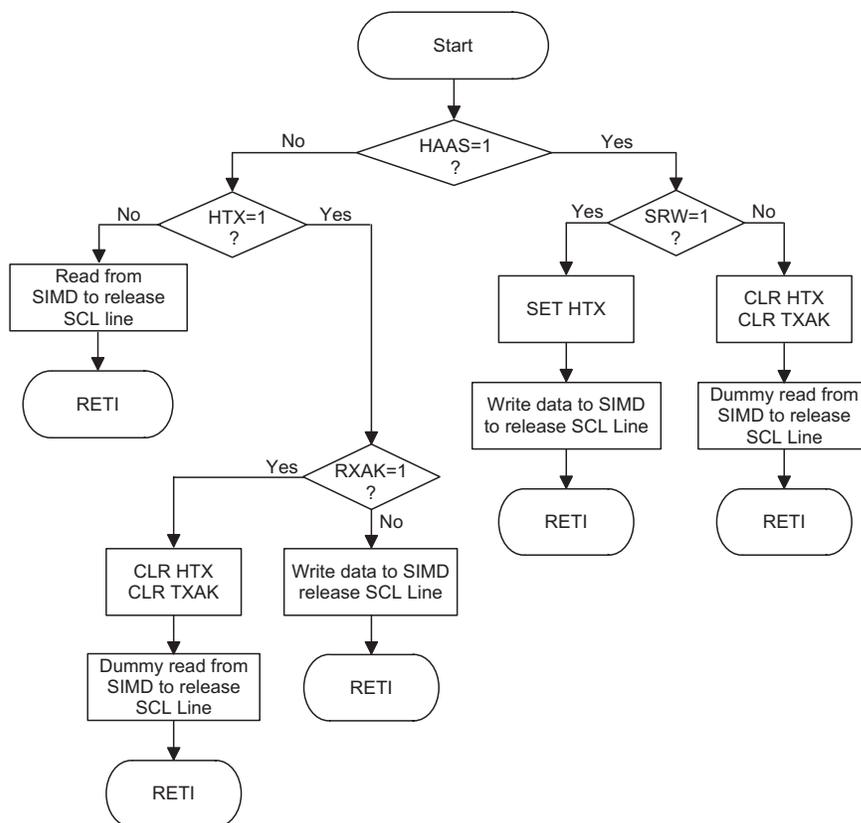
When the slave receiver receives the data byte, it must generate an acknowledge bit, known as TXAK, on the 9<sup>th</sup> clock. The slave device, which is setup as a transmitter will check the RXAK bit in the SIMC1 register to determine if it is to send another data byte, if not then it will release the SDA line and await the receipt of a STOP signal from the master.



S	SA	SR	M	D	A	D	A	.....	S	SA	SR	M	D	A	D	A	.....	P
---	----	----	---	---	---	---	---	-------	---	----	----	---	---	---	---	---	-------	---

Note: \*When a slave address is matched, the devices must be placed in either the transmit mode and then write data to the SIMD register, or in the receive mode where it must implement a dummy read from the SIMD register to release the SCL line.

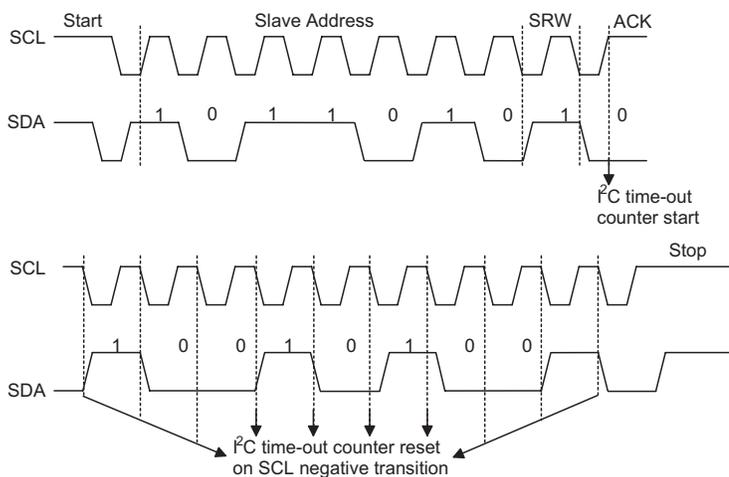
**I<sup>2</sup>C Communication Timing Diagram**



I<sup>2</sup>C Bus ISR Flow Chart

### I<sup>2</sup>C Time-out Control

In order to reduce the I<sup>2</sup>C lockup problem due to reception of erroneous clock sources, a time-out function is provided. If the clock source connected to the I<sup>2</sup>C bus is not received for a while, then the I<sup>2</sup>C circuitry and registers will be reset after a certain time-out period. The time-out counter starts to count on an I<sup>2</sup>C bus “START” & “address match” condition, and is cleared by an SCL falling edge. Before the next SCL falling edge arrives, if the time elapsed is greater than the time-out period specified by the SIMTOC register, then a time-out condition will occur. The time-out function will stop when an I<sup>2</sup>C “STOP” condition occurs.



**I<sup>2</sup>C Time-out**

When an I<sup>2</sup>C time-out counter overflow occurs, the counter will stop and the SIMTOEN bit will be cleared to zero and the SIMTOF bit will be set high to indicate that a time-out condition has occurred. The time-out condition will also generate an interrupt which uses the I<sup>2</sup>C interrupt vector. When an I<sup>2</sup>C time-out occurs, the I<sup>2</sup>C internal circuitry will be reset and the registers will be reset into the following condition:

Register	After I <sup>2</sup> C Time-out
SIMD, SIMA, SIMC0	No change
SIMC1	Reset to POR condition

**I<sup>2</sup>C Registers after Time-out**

The SIMTOF flag can be cleared by the application program. There are 64 time-out period selections which can be selected using the SIMTOS bits in the SIMTOC register. The time-out duration is calculated by the formula:  $((1-64) \times (32/f_{SUB}))$ . This gives a time-out period which ranges from about 1ms to 64ms.

**SIMTOC Register**

Bit	7	6	5	4	3	2	1	0
Name	SIMTOEN	SIMTOF	SIMTOS5	SIMTOS4	SIMTOS3	SIMTOS2	SIMTOS1	SIMTOS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **SIMTOEN**: I<sup>2</sup>C Time-out Control

- 0: Disable
- 1: Enable

Bit 6 **SIMTOF**: I<sup>2</sup>C Time-out flag

- 0: No time-out occurred
- 1: Time-out occurred

Bit 5~0 **SIMTOS5~SIMTOS0**: I<sup>2</sup>C Time-out Time Selection

I<sup>2</sup>C Time-out clock source is  $f_{SUB}/32$

I<sup>2</sup>C Time-out time is given by:  $(SIMTOS [5:0] + 1) \times (32/f_{SUB})$

## UART Module Serial Interface

The device contains an integrated full-duplex asynchronous serial communications UART interface that enables communication with external devices that contain a serial interface. The UART function has many features and can transmit and receive data serially by transferring a frame of data with eight or nine data bits per transmission as well as being able to detect errors when the data is overwritten or incorrectly framed. The UART function possesses its own internal interrupt which can be used to indicate when a reception occurs or when a transmission terminates.

The integrated UART function contains the following features:

- Full-duplex, Universal Asynchronous Receiver and Transmitter (UART) communication
- 8 or 9 bits character length
- Even, odd or no parity options
- One or two stop bits
- Baud rate generator with 8-bit prescaler
- Parity, framing, noise and overrun error detection
- Support for interrupt on address detect (last character bit=1)
- Transmitter and receiver enabled independently
- 2-byte Deep FIFO Receive Data Buffer
- Transmit and receive interrupts
- Transmit and Receive Multiple Interrupt Generation Sources:
  - ♦ Transmitter Empty
  - ♦ Transmitter Idle
  - ♦ Receiver Full
  - ♦ Receiver Overrun
  - ♦ Address Mode Detect
  - ♦ RX pin Wake-up

### UART External Interface

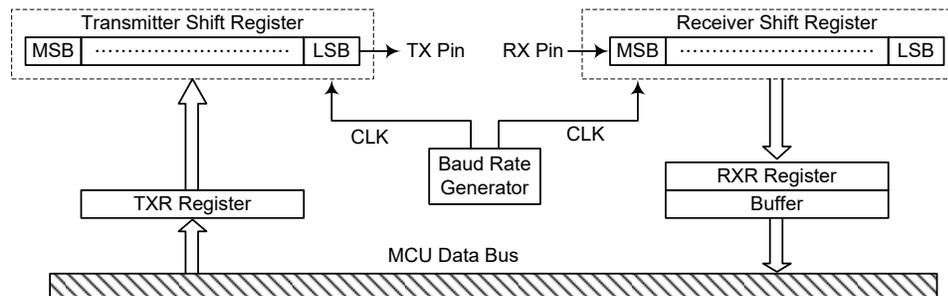
To communicate with an external serial interface, the internal UART has two external pins known as TX and RX. The TX pin is the UART transmitter pin, which can be used as a general purpose I/O or other pin-shared functional pin if the pin is not configured as a UART transmitter, which occurs when the TXEN bit in the UCR2 control register is equal to zero. Similarly, the RX pin is the UART receiver pin, which can also be used as a general purpose I/O or other pin-shared functional pin, if the pin is not configured as a receiver, which occurs if the RXEN bit in the UCR2 register is equal to zero. Along with the UARTEN bit, the TXEN and RXEN bits, if set, will automatically setup these I/O or other pin-shared functional pins to their respective TX output and RX input conditions and disable any pull-high resistor option which may exist on the RX pin.

### UART Data Transfer Scheme

The block diagram shows the overall data transfer structure arrangement for the UART. The actual data to be transmitted from the MCU is first transferred to the TXR register by the application program. The data will then be transferred to the Transmit Shift Register from where it will be shifted out, LSB first, onto the TX pin at a rate controlled by the Baud Rate Generator. Only the TXR register is mapped onto the MCU Data Memory, the Transmit Shift Register is not mapped and is therefore inaccessible to the application program.

Data to be received by the UART is accepted on the external RX pin, from where it is shifted in, LSB first, to the Receiver Shift Register at a rate controlled by the Baud Rate Generator. When the shift register is full, the data will then be transferred from the shift register to the internal RXR register, where it is buffered and can be manipulated by the application program. Only the RXR register is mapped onto the MCU Data Memory, the Receiver Shift Register is not mapped and is therefore inaccessible to the application program.

It should be noted that the actual register for data transmission and reception, although referred to in the text, and in application programs, as separate TXR and RXR registers, only exists as a single shared register in the Data Memory. This shared register known as the TXR/RXR register is used for both data transmission and data reception.



**UART Data Transfer Scheme**

### UART Status and Control Registers

There are five control registers associated with the UART function. The USR, UCR1 and UCR2 registers control the overall function of the UART, while the BRG register controls the Baud rate. The actual data to be transmitted and received on the serial interface is managed through the TXR/RXR data registers.

Register Name	Bit							
	7	6	5	4	3	2	1	0
USR	PERR	NF	FERR	OERR	RIDLE	RXIF	TIDLE	TXIF
UCR1	UARTEN	BNO	PREN	PRT	STOPS	TXBRK	RX8	TX8
UCR2	TXEN	RXEN	BRGH	ADDEN	WAKE	RIE	TIIE	TEIE
TXR/RXR	TXRX7	TXRX6	TXRX5	TXRX4	TXRX3	TXRX2	TXRX1	TXRX0
BRG	BRG7	BRG6	BRG5	BRG4	BRG3	BRG2	BRG1	BRG0

**UART Register Summary**

### USR register

The USR register is the status register for the UART, which can be read by the program to determine the present status of the UART. All flags within the USR register are read only. Further explanation on each of the flags is given below:

Bit	7	6	5	4	3	2	1	0
Name	PERR	NF	FERR	OERR	RIDLE	RXIF	TIDLE	TXIF
R/W	R	R	R	R	R	R	R	R
POR	0	0	0	0	1	0	1	1

- Bit 7 PERR:** Parity error flag  
 0: No parity error is detected  
 1: Parity error is detected
- The PERR flag is the parity error flag. When this read only flag is “0”, it indicates a parity error has not been detected. When the flag is “1”, it indicates that the parity of the received word is incorrect. This error flag is applicable only if Parity mode (odd or even) is selected. The flag can also be cleared by a software sequence which involves a read to the status register USR followed by an access to the RXR data register.
- Bit 6 NF:** Noise flag  
 0: No noise is detected  
 1: Noise is detected
- The NF flag is the noise flag. When this read only flag is "0", it indicates no noise condition. When the flag is "1", it indicates that the UART has detected noise on the receiver input. The NF flag is set during the same cycle as the RXIF flag but will not be set in the case of an overrun. The NF flag can be cleared by a software sequence which will involve a read to the status register USR followed by an access to the RXR data register.
- Bit 5 FERR:** Framing error flag  
 0: No framing error is detected  
 1: Framing error is detected
- The FERR flag is the framing error flag. When this read only flag is “0”, it indicates that there is no framing error. When the flag is “1”, it indicates that a framing error has been detected for the current character. The flag can also be cleared by a software sequence which will involve a read to the status register USR followed by an access to the RXR data register.
- Bit 4 OERR:** Overrun error flag  
 0: No overrun error is detected  
 1: Overrun error is detected
- The OERR flag is the overrun error flag which indicates when the receiver buffer has overflowed. When this read only flag is “0”, it indicates that there is no overrun error. When the flag is “1”, it indicates that an overrun error occurs which will inhibit further transfers to the RXR receive data register. The flag is cleared by a software sequence, which is a read to the status register USR followed by an access to the RXR data register.
- Bit 3 RIDLE:** Receiver status  
 0: Data reception is in progress (data being received)  
 1: No data reception is in progress (receiver is idle)
- The RIDLE flag is the receiver status flag. When this read only flag is “0”, it indicates that the receiver is between the initial detection of the start bit and the completion of the stop bit. When the flag is “1”, it indicates that the receiver is idle. Between the completion of the stop bit and the detection of the next start bit, the RIDLE bit is “1” indicating that the UART receiver is idle and the RX pin stays in logic high condition.

- Bit 2**     **RXIF:** Receive RXR data register status  
               0: RXR data register is empty  
               1: RXR data register has available data  
 The RXIF flag is the receive data register status flag. When this read only flag is “0”, it indicates that the RXR read data register is empty. When the flag is “1”, it indicates that the RXR read data register contains new data. When the contents of the shift register are transferred to the RXR register, an interrupt is generated if RIE=1 in the UCR2 register. If one or more errors are detected in the received word, the appropriate receive-related flags NF, FERR, and/or PERR are set within the same clock cycle. The RXIF flag is cleared when the USR register is read with RXIF set, followed by a read from the RXR register, and if the RXR register has no data available.
- Bit 1**     **TIDLE:** Transmission idle  
               0: Data transmission is in progress (data being transmitted)  
               1: No data transmission is in progress (transmitter is idle)  
 The TIDLE flag is known as the transmission complete flag. When this read only flag is “0”, it indicates that a transmission is in progress. This flag will be set to “1” when the TXIF flag is “1” and when there is no transmit data or break character being transmitted. When TIDLE is equal to “1”, the TX pin becomes idle with the pin state in logic high condition. The TIDLE flag is cleared by reading the USR register with TIDLE set and then writing to the TXR register. The flag is not generated when a data character or a break is queued and ready to be sent.
- Bit 0**     **TXIF:** Transmit TXR data register status  
               0: Character is not transferred to the transmit shift register  
               1: Character has transferred to the transmit shift register (TXR data register is empty)  
 The TXIF flag is the transmit data register empty flag. When this read only flag is “0”, it indicates that the character is not transferred to the transmitter shift register. When the flag is “1”, it indicates that the transmitter shift register has received a character from the TXR data register. The TXIF flag is cleared by reading the UART status register (USR) with TXIF set and then writing to the TXR data register. Note that when the TXEN bit is set, the TXIF flag bit will also be set since the transmit data register is not yet full.

**UCR1 register**

The UCR1 register together with the UCR2 register are the two UART control registers that are used to set the various options for the UART function, such as overall on/off control, parity control, data transfer bit length etc. Further explanation on each of the bits is given below:

Bit	7	6	5	4	3	2	1	0
Name	UARTEN	BNO	PREN	PRT	STOPS	TXBRK	RX8	TX8
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R	W
POR	0	0	0	0	0	0	x	0

“x” unknown

- Bit 7**     **UARTEN:** UART function enable control  
               0: Disable UART. TX and RX pins are as I/O or other pin-shared functional pins  
               1: Enable UART. TX and RX pins function as UART pins  
 The UARTEN bit is the UART enable bit. When this bit is equal to “0”, the UART will be disabled and the RX pin as well as the TX pin will be as General Purpose I/O or other pin-shared functional pins. When the bit is equal to “1”, the UART will be enabled and the TX and RX pins will function as defined by the TXEN and RXEN enable control bits.

When the UART is disabled, it will empty the buffer so any character remaining in the buffer will be discarded. In addition, the value of the baud rate counter will be reset. If the UART is disabled, all error and status flags will be reset. Also the TXEN, RXEN, TXBRK, RXIF, OERR, FERR, PERR and NF bits will be cleared, while the TIDLE, TXIF and RIDLE bits will be set. Other control bits in UCR1, UCR2 and BRG registers will remain unaffected. If the UART is active and the UARTEN bit is cleared, all pending transmissions and receptions will be terminated and the module will be reset as defined above. When the UART is re-enabled, it will restart in the same configuration.

- Bit 6 **BNO**: Number of data transfer bits selection  
0: 8-bit data transfer  
1: 9-bit data transfer  
This bit is used to select the data length format, which can have a choice of either 8-bit or 9-bit format. When this bit is equal to “1”, a 9-bit data length format will be selected. If the bit is equal to “0”, then an 8-bit data length format will be selected. If 9-bit data length format is selected, then bits RX8 and TX8 will be used to store the 9<sup>th</sup> bit of the received and transmitted data respectively.
- Bit 5 **PREN**: Parity function enable control  
0: Parity function is disabled  
1: Parity function is enabled  
This is the parity enable bit. When this bit is equal to “1”, the parity function will be enabled. If the bit is equal to “0”, then the parity function will be disabled. Replace the most significant bit position with a parity bit.
- Bit 4 **PRT**: Parity type selection bit  
0: Even parity for parity generator  
1: Odd parity for parity generator  
This bit is the parity type selection bit. When this bit is equal to “1”, odd parity type will be selected. If the bit is equal to “0”, then even parity type will be selected.
- Bit 3 **STOPS**: Number of Stop bits selection  
0: One stop bit format is used  
1: Two stop bits format is used  
This bit determines if one or two stop bits are to be used. When this bit is equal to “1”, two stop bits are used. If this bit is equal to “0”, then only one stop bit is used.
- Bit 2 **TXBRK**: Transmit break character  
0: No break character is transmitted  
1: Break characters transmit  
The TXBRK bit is the Transmit Break Character bit. When this bit is “0”, there are no break characters and the TX pin operates normally. When the bit is “1”, there are transmit break characters and the transmitter will send logic zeros. When this bit is equal to “1”, after the buffered data has been transmitted, the transmitter output is held low for a minimum of a 13-bit length and until the TXBRK bit is reset.
- Bit 1 **RX8**: Receive data bit 8 for 9-bit data transfer format (read only)  
This bit is only used if 9-bit data transfers are used, in which case this bit location will store the 9<sup>th</sup> bit of the received data known as RX8. The BNO bit is used to determine whether data transfers are in 8-bit or 9-bit format.
- Bit 0 **TX8**: Transmit data bit 8 for 9-bit data transfer format (write only)  
This bit is only used if 9-bit data transfers are used, in which case this bit location will store the 9<sup>th</sup> bit of the transmitted data known as TX8. The BNO bit is used to determine whether data transfers are in 8-bit or 9-bit format.

### UCR2 register

The UCR2 register is the second of the two UART control registers and serves several purposes. One of its main functions is to control the basic enable/disable operation of the UART Transmitter and Receiver as well as enabling the various UART interrupt sources. The register also serves to control the baud rate speed, receiver wake-up enable and the address detect enable. Further explanation on each of the bits is given below:

Bit	7	6	5	4	3	2	1	0
Name	TXEN	RXEN	BRGH	ADDEN	WAKE	RIE	TIIE	TEIE
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **TXEN**: UART Transmitter enabled control

0: UART transmitter is disabled

1: UART transmitter is enabled

The bit named TXEN is the Transmitter Enable Bit. When this bit is equal to “0”, the transmitter will be disabled with any pending data transmissions being aborted. In addition the buffers will be reset. In this situation the TX pin will be used as an I/O or other pin-shared functional pin.

If the TXEN bit is equal to “1” and the UARTEN bit is also equal to “1”, the transmitter will be enabled and the TX pin will be controlled by the UART. Clearing the TXEN bit during a transmission will cause the data transmission to be aborted and will reset the transmitter. If this situation occurs, the TX pin will be used as an I/O or other pin-shared functional pin.

Bit 6 **RXEN**: UART Receiver enabled control

0: UART receiver is disabled

1: UART receiver is enabled

The bit named RXEN is the Receiver Enable Bit. When this bit is equal to “0”, the receiver will be disabled with any pending data receptions being aborted. In addition the receive buffers will be reset. In this situation the RX pin will be used as an I/O or other pin-shared functional pin. If the RXEN bit is equal to “1” and the UARTEN bit is also equal to “1”, the receiver will be enabled and the RX pin will be controlled by the UART. Clearing the RXEN bit during a reception will cause the data reception to be aborted and will reset the receiver. If this situation occurs, the RX pin will be used as an I/O or other pin-shared functional pin.

Bit 5 **BRGH**: Baud Rate speed selection

0: Low speed baud rate

1: High speed baud rate

The bit named BRGH selects the high or low speed mode of the Baud Rate Generator. This bit, together with the value placed in the baud rate register BRG, controls the Baud Rate of the UART. If this bit is equal to “1”, the high speed mode is selected. If the bit is equal to “0”, the low speed mode is selected.

Bit 4 **ADDEN**: Address detect function enable control

0: Address detect function is disabled

1: Address detect function is enabled

The bit named ADDEN is the address detect function enable control bit. When this bit is equal to “1”, the address detect function is enabled. When it occurs, if the 8<sup>th</sup> bit, which corresponds to RX7 if BNO=0 or the 9<sup>th</sup> bit, which corresponds to RX8 if BNO=1, has a value of “1”, then the received word will be identified as an address, rather than data. If the corresponding interrupt is enabled, an interrupt request will be generated each time the received word has the address bit set, which is the 8<sup>th</sup> or 9<sup>th</sup> bit depending on the value of BNO. If the address bit known as the 8<sup>th</sup> or 9<sup>th</sup> bit of the received word is “0” with the address detect function being enabled, an interrupt will not be generated and the received data will be discarded.

- Bit 3 WAKE:** RX pin falling edge wake-up function enable control  
 0: RX pin wake-up function is disabled  
 1: RX pin wake-up function is enabled  
 This bit enables or disables the receiver wake-up function. If this bit is equal to “1” and the MCU is in IDLE or SLEEP mode, a falling edge on the RX input pin will wake-up the device. Please reference the UART RX pin wake-up functions in different operating mode for the detail. If this bit is equal to “0” and the MCU is in IDLE or SLEEP mode, any edge transitions on the RX pin will not wake-up the device.
- Bit 2 RIE:** Receiver interrupt enable control  
 0: Receiver related interrupt is disabled  
 1: Receiver related interrupt is enabled  
 This bit enables or disables the receiver interrupt. If this bit is equal to “1” and when the receiver overrun flag OERR or receive data available flag RXIF is set, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt request flag will not be influenced by the condition of the OERR or RXIF flags.
- Bit 1 TIE:** Transmitter Idle interrupt enable control  
 0: Transmitter idle interrupt is disabled  
 1: Transmitter idle interrupt is enabled  
 This bit enables or disables the transmitter idle interrupt. If this bit is equal to “1” and when the transmitter idle flag TIDLE is set, due to a transmitter idle condition, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt request flag will not be influenced by the condition of the TIDLE flag.
- Bit 0 TEIE:** Transmitter Empty interrupt enable control  
 0: Transmitter empty interrupt is disabled  
 1: Transmitter empty interrupt is enabled  
 This bit enables or disables the transmitter empty interrupt. If this bit is equal to “1” and when the transmitter empty flag TXIF is set, due to a transmitter empty condition, the UART interrupt request flag will be set. If this bit is equal to “0”, the UART interrupt request flag will not be influenced by the condition of the TXIF flag.

**TXR/RXR register**

Bit	7	6	5	4	3	2	1	0
Name	TXRX7	TXRX6	TXRX5	TXRX4	TXRX3	TXRX2	TXRX1	TXRX0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

“x” unknown

Bit 7~0 **TXRX7~TXRX0:** UART Transmit/Receive Data bit 7 ~ bit 0

**BRG Register**

Bit	7	6	5	4	3	2	1	0
Name	BRG7	BRG6	BRG5	BRG4	BRG3	BRG2	BRG1	BRG0
R/W								
POR	x	x	x	x	x	x	x	x

“x” unknown

Bit 7~0 **BRG7~BRG0:** Baud Rate values  
 By programming the BRGH bit in UCR2 Register which allows selection of the related formula described above and programming the required value in the BRG register, the required baud rate can be setup.  
 Note: Baud rate=  $f_{SYS}/[64 \times (N+1)]$  if BRGH=0.  
 Baud rate=  $f_{SYS}/[16 \times (N+1)]$  if BRGH=1.

## Baud Rate Generator

To setup the speed of the serial data communication, the UART function contains its own dedicated baud rate generator. The baud rate is controlled by its own internal free running 8-bit timer, the period of which is determined by two factors. The first of these is the value placed in the baud rate register BRG and the second is the value of the BRGH bit with the control register UCR2. The BRGH bit decides if the baud rate generator is to be used in a high speed mode or low speed mode, which in turn determines the formula that is used to calculate the baud rate. The value N in the BRG register which is used in the following baud rate calculation formula determines the division factor. Note that N is the decimal value placed in the BRG register and has a range of between 0 and 255.

UCR2 BRGH Bit	0	1
Baud Rate (BR)	$f_{\text{SYS}} / [64 (N+1)]$	$f_{\text{SYS}} / [16 (N+1)]$

By programming the BRGH bit which allows selection of the related formula and programming the required value in the BRG register, the required baud rate can be setup. Note that because the actual baud rate is determined using a discrete value, N, placed in the BRG register, there will be an error associated between the actual and requested value. The following example shows how the BRG register value N and the error value can be calculated.

### Calculating the register and error values

For a clock frequency of 4MHz, and with BRGH set to “0” determine the BRG register value N, the actual baud rate and the error value for a desired baud rate of 4800.

From the above table the desired baud rate  $BR = f_{\text{SYS}} / [64 (N+1)]$

Re-arranging this equation gives  $N = [f_{\text{SYS}} / (BR \times 64)] - 1$

Giving a value for  $N = [4000000 / (4800 \times 64)] - 1 = 12.0208$

To obtain the closest value, a decimal value of 12 should be placed into the BRG register. This gives an actual or calculated baud rate value of  $BR = 4000000 / [64 \times (12 + 1)] = 4808$

Therefore the error is equal to  $(4808 - 4800) / 4800 = 0.16\%$

The following tables show actual values of baud rate and error values for the two values of BRGH.

Baud Rate K/BPS	Baud Rates for BRGH=0					
	$f_{\text{SYS}} = 16 \text{ MHz}$			$f_{\text{SYS}} = 20 \text{ MHz}$		
	BRG	Kbaud	Error (%)	BRG	Kbaud	Error (%)
0.3	—	—	—	—	—	—
1.2	207	1.202	0.16	259	1.202	0.16
2.4	103	2.404	0.16	129	2.404	0.16
4.8	51	4.808	0.16	64	4.808	0.16
9.6	25	9.615	0.16	32	9.470	-1.36
19.2	12	19.231	0.16	15	19.531	1.73
38.4	6	35.714	-6.99	7	39.063	1.73
57.6	3	62.5	8.51	4	62.5	8.51
115.2	1	125	8.51	2	104.17	-9.58
250	0	250	0	0	312.5	25

**Baud Rates and Error Values for BRGH = 0**

Baud Rate K/BPS	Baud Rates for BRGH=1					
	f <sub>sys</sub> = 16 MHz			f <sub>sys</sub> = 20 MHz		
	BRG	Kbaud	Error (%)	BRG	Kbaud	Error (%)
0.3	—	—	—	—	—	—
1.2	—	—	—	—	—	—
2.4	—	—	—	—	—	—
4.8	207	4.808	0.16	—	—	—
9.6	103	9.615	0.16	129	9.615	0.16
19.2	51	19.231	0.16	64	19.231	0.16
38.4	25	38.462	0.16	32	37.879	-1.36
57.6	16	58.824	2.12	21	56.818	-1.35
115.2	8	111.11	-3.55	10	113.636	-1.36
250	3	250	0	4	250	0

**Baud Rates and Error Values for BRGH = 1**

### UART Setup and Control

For data transfer, the UART function utilizes a non-return-to-zero, more commonly known as NRZ, format. This is composed of one start bit, eight or nine data bits, and one or two stop bits. Parity is supported by the UART hardware, and can be setup to be even, odd or no parity. For the most common data format, 8 data bits along with no parity and one stop bit, denoted as 8, N, 1, is used as the default setting, which is the setting at power-on. The number of data bits and stop bits, along with the parity, are setup by programming the corresponding BNO, PRT, PREN, and STOPS bits in the UCR1 register. The baud rate used to transmit and receive data is setup using the internal 8-bit baud rate generator, while the data is transmitted and received LSB first. Although the UART transmitter and receiver are functionally independent, they both use the same data format and baud rate. In all cases stop bits will be used for data transmission.

#### Enabling/disabling the UART

The basic on/off function of the internal UART function is controlled using the URTEN bit in the UCR1 register. If the URTEN, TXEN and RXEN bits are set, then these two UART pins will act as normal TX output pin and RX input pin respectively. If no data is being transmitted on the TX pin, then it will default to a logic high value.

Clearing the URTEN bit will disable the TX and RX pins and allow these two pins to be used as normal I/O or other pin-shared functional pins. When the UART function is disabled the buffer will be reset to an empty condition, at the same time discarding any remaining residual data. Disabling the UART will also reset the error and status flags with bits TXEN, RXEN, TXBRK, RXIF, OERR, FERR, PERR and NF being cleared while bits TIDLE, TXIF and RIDLE will be set. The remaining control bits in the UCR1, UCR2 and BRG registers will remain unaffected. If the URTEN bit in the UCR1 register is cleared while the UART is active, then all pending transmissions and receptions will be immediately suspended and the UART will be reset to a condition as defined above. If the UART is then subsequently re-enabled, it will restart again in the same configuration.

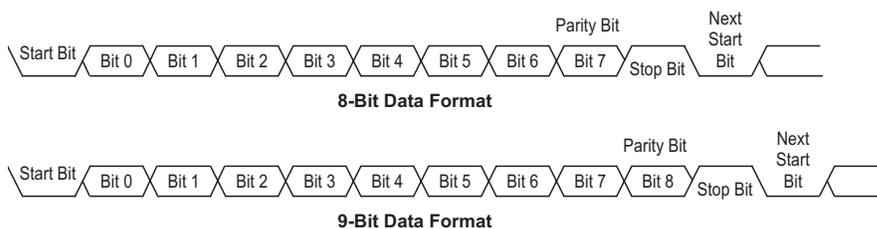
### Data, parity and stop bit selection

The format of the data to be transferred is composed of various factors such as data bit length, parity on/off, parity type, address bits and the number of stop bits. These factors are determined by the setup of various bits within the UCR1 register. The BNO bit controls the number of data bits which can be set to either 8 or 9, the PRT bit controls the choice of odd or even parity, the PREN bit controls the parity on/off function and the STOPS bit decides whether one or two stop bits are to be used. The following table shows various formats for data transmission. The address bit identifies the frame as an address character. The number of stop bits, which can be either one or two, is independent of the data length.

Start Bit	Data Bits	Address Bits	Parity Bits	Stop Bit
<b>Example of 8-bit Data Formats</b>				
1	8	0	0	1
1	7	0	1	1
1	7	1	0	1
<b>Example of 9-bit Data Formats</b>				
1	9	0	0	1
1	8	0	1	1
1	8	1	0	1

#### Transmitter Receiver Data Format

The following diagram shows the transmit and receive waveforms for both 8-bit and 9-bit data formats.



### UART transmitter

Data word lengths of either 8 or 9 bits, can be selected by programming the BNO bit in the UCR1 register. When BNO bit is set, the word length will be set to 9 bits. In this case the 9<sup>th</sup> bit, which is the MSB, needs to be stored in the TX8 bit in the UCR1 register. At the transmitter core lies the Transmitter Shift Register, more commonly known as the TSR, whose data is obtained from the transmit data register, which is known as the TXR register. The data to be transmitted is loaded into this TXR register by the application program. The TSR register is not written to with new data until the stop bit from the previous transmission has been sent out. As soon as this stop bit has been transmitted, the TSR can then be loaded with new data from the TXR register, if it is available. It should be noted that the TSR register, unlike many other registers, is not directly mapped into the Data Memory area and as such is not available to the application program for direct read/write operations. An actual transmission of data will normally be enabled when the TXEN bit is set, but the data will not be transmitted until the TXR register has been loaded with data and the baud rate generator has defined a shift clock source. However, the transmission can also be initiated by first loading data into the TXR register, after which the TXEN bit can be set. When a transmission of data begins, the TSR is normally empty, in which case a transfer to the TXR register will result in an immediate transfer to the TSR. If during a transmission the TXEN bit is cleared, the transmission will immediately cease and the transmitter will be reset. The TX output pin will then return to the I/O or other pin-shared function.

### Transmitting data

When the UART is transmitting data, the data is shifted on the TX pin from the shift register, with the least significant bit first. In the transmit mode, the TXR register forms a buffer between the internal bus and the transmitter shift register. It should be noted that if 9-bit data format has been selected, then the MSB will be taken from the TX8 bit in the UCR1 register. The steps to initiate a data transfer can be summarized as follows:

- Make the correct selection of the BNO, PRT, PREN and STOPS bits to define the required word length, parity type and number of stop bits.
- Setup the BRG register to select the desired baud rate.
- Set the TXEN bit to ensure that the TX pin is used as a UART transmitter pin.
- Access the USR register and write the data that is to be transmitted into the TXR register. Note that this step will clear the TXIF bit.
- This sequence of events can now be repeated to send additional data.

It should be noted that when TXIF=0, data will be inhibited from being written to the TXR register. Clearing the TXIF flag is always achieved using the following software sequence:

1. A USR register access
2. A TXR register write execution

The read-only TXIF flag is set by the UART hardware and if set indicates that the TXR register is empty and that other data can now be written into the TXR register without overwriting the previous data. If the TEIE bit is set then the TXIF flag will generate an interrupt.

During a data transmission, a write instruction to the TXR register will place the data into the TXR register, which will be copied to the shift register at the end of the present transmission. When there is no data transmission in progress, a write instruction to the TXR register will place the data directly into the shift register, resulting in the commencement of data transmission, and the TXIF bit being immediately set. When a frame transmission is complete, which happens after stop bits are sent or after the break frame, the TIDLE bit will be set. To clear the TIDLE bit the following software sequence is used:

1. A USR register access
2. A TXR register write execution

Note that both the TXIF and TIDLE bits are cleared by the same software sequence.

### Transmit break

If the TXBRK bit is set then break characters will be sent on the next transmission. Break character transmission consists of a start bit, followed by  $13 \times N$  '0' bits and stop bits, where  $N=1, 2$ , etc. If a break character is to be transmitted then the TXBRK bit must be first set by the application program, then cleared to generate the stop bits. Transmitting a break character will not generate a transmit interrupt. Note that a break condition length is at least 13 bits long. If the TXBRK bit is continually kept at a logic high level then the transmitter circuitry will transmit continuous break characters. After the application program has cleared the TXBRK bit, the transmitter will finish transmitting the last break character and subsequently send out one or two stop bits. The automatic logic highs at the end of the last break character will ensure that the start bit of the next frame is recognized.

### UART receiver

The UART is capable of receiving word lengths of either 8 or 9 bits. If the BNO bit is set, the word length will be set to 9 bits with the MSB being stored in the RX8 bit of the UCR1 register. At the receiver core lies the Receive Serial Shift Register, commonly known as the RSR. The data which is received on the RX external input pin, is sent to the data recovery block. The data recovery block operating speed is 16 times that of the baud rate, while the main receive serial shifter operates at the baud rate. After the RX pin is sampled for the stop bit, the received data in RSR is transferred to the receive data register, if the register is empty. The data which is received on the external RX input pin is sampled three times by a majority detect circuit to determine the logic level that has been placed onto the RX pin. It should be noted that the RSR register, unlike many other registers, is not directly mapped into the Data Memory area and as such is not available to the application program for direct read/write operations.

### Receiving data

When the UART receiver is receiving data, the data is serially shifted in on the external RX input pin, LSB first. In the read mode, the RXR register forms a buffer between the internal bus and the receiver shift register. The RXR register is a two byte deep FIFO data buffer, where two bytes can be held in the FIFO while a third byte can continue to be received. Note that the application program must ensure that the data is read from RXR before the third byte has been completely shifted in, otherwise this third byte will be discarded and an overrun error OERR will be subsequently indicated. The steps to initiate a data transfer can be summarized as follows:

- Make the correct selection of BNO, PRT, PREN and STOPS bits to define the word length, parity type and number of stop bits.
- Setup the BRG register to select the desired baud rate.
- Set the RXEN bit to ensure that the RX pin is used as a UART receiver pin.

At this point the receiver will be enabled which will begin to look for a start bit.

When a character is received the following sequence of events will occur:

- The RXIF bit in the USR register will be set when RXR register has data available, at least one character can be read.
- When the contents of the shift register have been transferred to the RXR register, then if the RIE bit is set, an interrupt will be generated.
- If during reception, a frame error, noise error, parity error, or an overrun error has been detected, then the error flags can be set.

The RXIF bit can be cleared using the following software sequence:

1. A USR register access
2. An RXR register read execution

### **Receive break**

Any break character received by the UART will be managed as a framing error. The receiver will count and expect a certain number of bit times as specified by the values programmed into the BNO and STOPS bits. If the break is much longer than 13 bit times, the reception will be considered as complete after the number of bit times specified by BNO and STOPS. The RXIF bit is set, FERR is set, zeros are loaded into the receive data register, interrupts are generated if appropriate and the RIDLE bit is set. If a long break signal has been detected and the receiver has received a start bit, the data bits and the invalid stop bit, which sets the FERR flag, the receiver must wait for a valid stop bit before looking for the next start bit. The receiver will not make the assumption that the break condition on the line is the next start bit. A break is regarded as a character that contains only zeros with the FERR flag set. The break character will be loaded into the buffer and no further data will be received until stop bits are received. It should be noted that the RIDLE read only flag will go high when the stop bits have not yet been received. The reception of a break character on the UART registers will result in the following:

- The framing error flag, FERR, will be set.
- The receive data register, RXR, will be cleared.
- The OERR, NF, PERR, RIDLE or RXIF flags will possibly be set.

### **Idle status**

When the receiver is reading data, which means it will be in between the detection of a start bit and the reading of a stop bit, the receiver status flag in the USR register, otherwise known as the RIDLE flag, will have a zero value. In between the reception of a stop bit and the detection of the next start bit, the RIDLE flag will have a high value, which indicates the receiver is in an idle condition.

### **Receiver interrupt**

The read only receive interrupt flag RXIF in the USR register is set by an edge generated by the receiver. An interrupt is generated if RIE=1, when a word is transferred from the Receive Shift Register, RSR, to the Receive Data Register, RXR. An overrun error can also generate an interrupt if RIE=1.

## **Managing receiver errors**

Several types of reception errors can occur within the UART module, the following section describes the various types and how they are managed by the UART.

### **Overrun Error – OERR flag**

The RXR register is composed of a two byte deep FIFO data buffer, where two bytes can be held in the FIFO register, while a third byte can continue to be received. Before the third byte has been entirely shifted in, the data should be read from the RXR register. If this is not done, the overrun error flag OERR will be consequently indicated.

In the event of an overrun error occurring, the following will happen:

- The OERR flag in the USR register will be set.
- The RXR contents will not be lost.
- The shift register will be overwritten.
- An interrupt will be generated if the RIE bit is set.

The OERR flag can be cleared by an access to the USR register followed by a read to the RXR register.

### Noise Error – NF Flag

Over-sampling is used for data recovery to identify valid incoming data and noise. If noise is detected within a frame the following will occur:

- The read only noise flag, NF, in the USR register will be set on the rising edge of the RXIF bit.
- Data will be transferred from the Shift register to the RXR register.
- No interrupt will be generated. However this bit rises at the same time as the RXIF bit which itself generates an interrupt.

Note that the NF flag is reset by a USR register read operation followed by an RXR register read operation.

### Framing Error – FERR Flag

The read only framing error flag, FERR, in the USR register, is set if a zero is detected instead of stop bits. If two stop bits are selected, both stop bits must be high, otherwise the FERR flag will be set. The FERR flag is buffered along with the received data and is cleared on any reset.

### Parity Error – PERR Flag

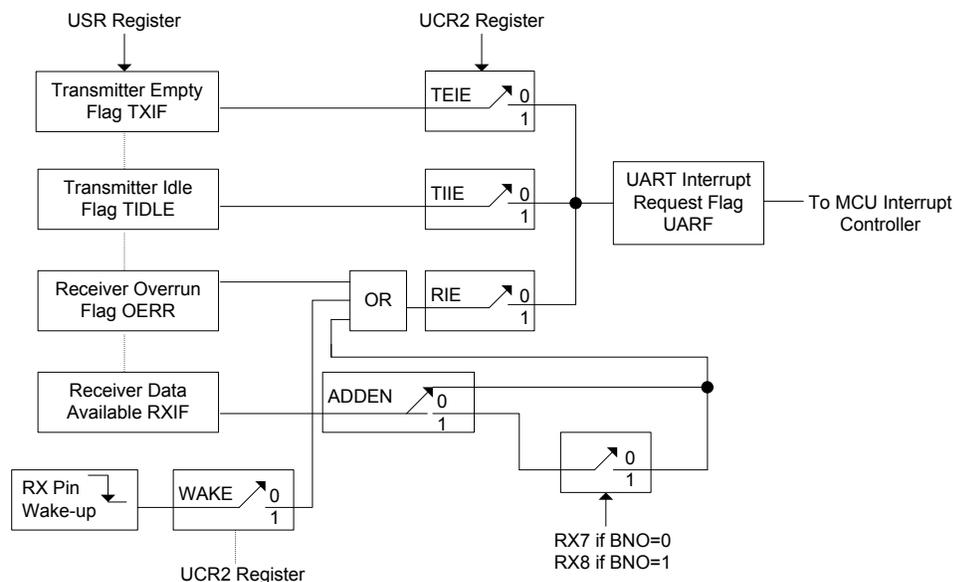
The read only parity error flag, PERR, in the USR register, is set if the parity of the received word is incorrect. This error flag is only applicable if the parity is enabled, PREN = 1, and if the parity type, odd or even is selected. The read only PERR flag is buffered along with the received data bytes. It is cleared on any reset. It should be noted that the FERR and PERR flags are buffered along with the corresponding word and should be read before reading the data word.

## UART Module Interrupt Structure

Several individual UART conditions can generate a UART interrupt. When these conditions exist, a low pulse will be generated to get the attention of the microcontroller. These conditions are a transmitter data register empty, transmitter idle, receiver data available, receiver overrun, address detect and an RX pin wake-up. When any of these conditions are created, if its corresponding interrupt control is enabled and the stack is not full, the program will jump to its corresponding interrupt vector where it can be serviced before returning to the main program. Four of these conditions have the corresponding USR register flags which will generate a UART interrupt if its associated interrupt enable control bit in the UCR2 register is set. The two transmitter interrupt conditions have their own corresponding enable control bits, while the two receiver interrupt conditions have a shared enable control bit. These enable bits can be used to mask out individual UART interrupt sources.

The address detect condition, which is also a UART interrupt source, does not have an associated flag, but will generate a UART interrupt when an address detect condition occurs if its function is enabled by setting the ADDEN bit in the UCR2 register. An RX pin wake-up, which is also a UART interrupt source, does not have an associated flag, but will generate a UART interrupt if the microcontroller is woken up by a falling edge on the RX pin, if the WAKE and RIE bits in the UCR register are set. Note that in the event of an RX wake-up interrupt occurring, there will be a certain period of delay, commonly known as the System Start-up Time, for the oscillator to restart and stabilize before the system resumes normal operation.

Note that the USR register flags are read only and cannot be cleared or set by the application program, neither will they be cleared when the program jumps to the corresponding interrupt servicing routine, as is the case for some of the other interrupts. The flags will be cleared automatically when certain actions are taken by the UART, the details of which are given in the UART register section. The overall UART interrupt can be disabled or enabled by the related interrupt enable control bits in the interrupt control registers of the microcontroller to decide whether the interrupt requested by the UART module is masked out or allowed.



UART Interrupt Scheme

### Address detect mode

Setting the Address Detect Mode bit, ADDEN, in the UCR2 register, enables this special mode. If this bit is enabled then an additional qualifier will be placed on the generation of a Receiver Data Available interrupt, which is requested by the RXIF flag. If the ADDEN bit is enabled, then when data is available, an interrupt will only be generated, if the highest received bit has a high value. Note that the MFE, URE and EMI interrupt enable bits must also be enabled for correct interrupt generation. This highest address bit is the 9<sup>th</sup> bit if BNO=1 or the 8<sup>th</sup> bit if BNO=0. If this bit is high, then the received word will be defined as an address rather than data. A Data Available interrupt will be generated every time the last bit of the received word is set. If the ADDEN bit is not enabled, then a Receiver Data Available interrupt will be generated each time the RXIF flag is set, irrespective of the data last bit status. The address detect mode and parity enable are mutually exclusive functions. Therefore if the address detect mode is enabled, then to ensure correct operation, the parity function should be disabled by resetting the parity enable bit to zero.

ADDEN	Bit 9 if BNO=1, Bit 8 if BNO=0	UART Interrupt Generated
0	0	√
	1	√
1	0	×
	1	√

ADDEN Bit Function

## UART Module Power Down and Wake-up

When the  $f_{SYS}$  is off, the UART will cease to function. All clock sources to the module are shutdown. If the  $f_{SYS}$  is off while a transmission is still in progress, then the transmission will be paused until the UART clock source derived from the microcontroller is activated. In a similar way, if the MCU enters the Power Down Mode while receiving data, then the reception of data will likewise be paused. When the MCU enters the Power Down Mode, note that the USR, UCR1, UCR2, transmit and receive registers, as well as the BRG register will not be affected. It is recommended to make sure first that the UART data transmission or reception has been finished before the microcontroller enters the Power Down mode.

The UART function contains a receiver RX pin wake-up function, which is enabled or disabled by the WAKE bit in the UCR2 register. If this bit, along with the UART enable bit, UARTEN, the receiver enable bit, RXEN and the receiver interrupt bit, RIE, are all set before the MCU enters the Power Down Mode, then a falling edge on the RX pin will wake up the MCU from the Power Down Mode. Note that as it takes certain system clock cycles after a wake-up, before normal microcontroller operation resumes, any data received during this time on the RX pin will be ignored. For a UART wake-up interrupt to occur, in addition to the bits for the wake-up being set, the global interrupt enable bit, EMI, and the UART interrupt enable bit, UARE, must also be set. If these two bits are not set then only a wake up event will occur and no interrupt will be generated. Note also that as it takes certain system clock cycles after a wake-up before normal microcontroller resumes, the UART interrupt will not be generated until after this time has elapsed.

## Interrupts

Interrupts are an important part of any microcontroller system. When an external event or an internal function such as a Timer Module or an A/D converter requires microcontroller attention, their corresponding interrupt will enforce a temporary suspension of the main program allowing the microcontroller to direct attention to their respective needs. The device contains an external interrupt and internal interrupts functions. The external interrupt is generated by the action of the external INTn pins, while the internal interrupts are generated by various internal functions such as the TMs, Time Base, EEPROM, SIM, UART and the A/D converter.

### Interrupt Registers

Overall interrupt control, which basically means the setting of request flags when certain microcontroller conditions occur and the setting of interrupt enable bits by the application program, is controlled by a series of registers, located in the Special Purpose Data Memory, as shown in the accompanying table. The number of registers depends upon the device chosen but fall into three categories. The first is the INTC0~INTC3 registers which setup the primary interrupts, the second is the MFI0~MFI2 registers which setup the Multi-function interrupts. Finally there are INTEG0~INTEG1 registers to setup the external interrupt trigger edge type.

Each register contains a number of enable bits to enable or disable individual registers as well as interrupt flags to indicate the presence of an interrupt request. The naming convention of these follows a specific pattern. First is listed an abbreviated interrupt type, then the (optional) number of that interrupt followed by either an "E" for enable/disable bit or "F" for request flag.

Function	Enable Bit	Request Flag	Notes
Global	EMI	—	—
INTn Pin	INTnE	INTnF	n=0~5
UART	UARE	UARF	—
Time Base	TBnE	TBnF	n=0 or 1
A/D Converter	ADE	ADF	—
Multi-function	MFnE	MFnF	n=0~2
EEPROM	DEE	DEF	—
SIM	SIME	SIMF	—
CTM	CTMPE	CTMPF	—
	CTMAE	CTMAF	
STM	STMPE	STMPF	—
	STMAE	STMAF	
PTM	PTMnPE	PTMnPF	n=0 or 1
	PTMnAE	PTMnAF	

#### Interrupt Register Bit Naming Conventions

Register Name	Bit							
	7	6	5	4	3	2	1	0
INTEG0	INT3S1	INT3S0	INT2S1	INT2S0	INT1S1	INT1S0	INT0S1	INT0S0
INTEG1	—	—	—	—	INT5S1	INT5S0	INT4S1	INT4S0
INTC0	—	INT2F	INT1F	INT0F	INT2E	INT1E	INT0E	EMI
INTC1	UARF	INT5F	INT4F	INT3F	UARE	INT5E	INT4E	INT3E
INTC2	ADF	MF2F	MF1F	MF0F	ADE	MF2E	MF1E	MF0E
INTC3	SIMF	DEF	TB1F	TB0F	SIME	DEE	TB1E	TB0E
MF10	—	—	CTMAF	CTMPF	—	—	CTMAE	CTMPE
MF11	LVF	—	STMAF	STMPF	LVE	—	STMAE	STMPE
MF12	PTM1AF	PTM1PF	PTM0AF	PTM0PF	PTM1AE	PTM1PE	PTM0AE	PTM0PE

#### Interrupt Register Contents

##### INTEG0 Register

Bit	7	6	5	4	3	2	1	0
Name	INT3S1	INT3S0	INT2S1	INT2S0	INT1S1	INT1S0	INT0S1	INT0S0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7 ~ 6    **INT3S1, INT3S0:** Defines INT3 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt
- Bit 5 ~ 4    **INT2S1, INT2S0:** Defines INT2 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt
- Bit 3 ~ 2    **INT1S1, INT1S0:** Defines INT1 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt

Bit 1 ~ 0 **INT0S1, INT0S0**: Defines INT0 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt

**INTEG1 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	INT5S1	INT5S0	INT4S1	INT4S0
R/W	—	—	—	—	R/W	R/W	R/W	R/W
POR	—	—	—	—	0	0	0	0

Bit 7 ~ 4 Unimplemented, read as "0"

Bit 3 ~ 2 **INT5S1, INT5S0**: Defines INT5 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt

Bit 1 ~ 0 **INT4S1, INT4S0**: Defines INT4 interrupt active edge  
 00: Disabled Interrupt  
 01: Rising Edge Interrupt  
 10: Falling Edge Interrupt  
 11: Dual Edge Interrupt

**INTC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	INT2F	INT1F	INT0F	INT2E	INT1E	INT0E	EMI
R/W	—	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	0	0	0	0	0	0	0

Bit 7 Unimplemented, read as "0"

Bit 6 **INT2F**: INT2 Interrupt Request Flag  
 0: No request  
 1: Interrupt request

Bit 5 **INT1F**: INT1 Interrupt Request Flag  
 0: No request  
 1: Interrupt request

Bit 4 **INT0F**: INT0 Interrupt Request Flag  
 0: No request  
 1: Interrupt request

Bit 3 **INT2E**: INT2 Interrupt Control  
 0: Disable  
 1: Enable

Bit 2 **INT1E**: INT1 Interrupt Control  
 0: Disable  
 1: Enable

Bit 1 **INT0E**: INT0 Interrupt Control  
 0: Disable  
 1: Enable

Bit 0 **EMI**: Global Interrupt Control  
 0: Disable  
 1: Enable

**INTC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	UARF	INT5F	INT4F	INT3F	UARE	INT5E	INT4E	INT3E
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **UARF**: UART Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 6      **INT5F**: INT5 Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 5      **INT4F**: INT4 Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 4      **INT3F**: INT3 Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 3      **UARE**: UART Interrupt Control  
             0: Disable  
             1: Enable
- Bit 2      **INT5E**: INT5 Interrupt Control  
             0: Disable  
             1: Enable
- Bit 1      **INT4E**: INT4 Interrupt Control  
             0: Disable  
             1: Enable
- Bit 0      **INT3E**: INT3 Interrupt Control  
             0: Disable  
             1: Enable

**INTC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	ADF	MF2F	MF1F	MF0F	ADE	MF2E	MF1E	MF0E
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **ADF**: A/D Converter Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 6      **MF2F**: Multi-function 2 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 5      **MF1F**: Multi-function 1 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 4      **MF0F**: Multi-function 0 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 3      **ADE**: A/D Converter Interrupt Control  
0: Disable  
1: Enable
- Bit 2      **MF2E**: Multi-function 2 Interrupt Control  
0: Disable  
1: Enable
- Bit 1      **MF1E**: Multi-function 1 Interrupt Control  
0: Disable  
1: Enable
- Bit 0      **MF0E**: Multi-function 0 Interrupt Control  
0: Disable  
1: Enable

**INTC3 Register**

Bit	7	6	5	4	3	2	1	0
Name	SIMF	DEF	TB1F	TB0F	SIME	DEE	TB1E	TB0E
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7 **SIMF**: Serial Interface Module Interrupt Request Flag

0: No request

1: Interrupt request

Bit 6 **DEF**: Data EEPROM Interrupt Request Flag

0: No request

1: Interrupt request

Bit 5 **TB1F** : Time Base 1 Interrupt Request Flag

0: No request

1: Interrupt request

Bit 4 **TB0F**: Time Base 0 Interrupt Request Flag

0: No request

1: Interrupt request

Bit 3 **SIME**: Serial Interface Module Interrupt Control

0: Disable

1: Enable

Bit 2 **DEE**: Data EEPROM Interrupt Control

0: Disable

1: Enable

Bit 1 **TB1E** : Time Base 1 Interrupt Control

0: Disable

1: Enable

Bit 0 **TB0E**: Time Base 0 Interrupt Control

0: Disable

1: Enable

**MFIO Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	CTMAF	CTMPF	—	—	CTMAE	CTMPE
R/W	—	—	R/W	R/W	—	—	R/W	R/W
POR	—	—	0	0	—	—	0	0

- Bit 7 ~ 6     Unimplemented, read as "0"
- Bit 5        **CTMAF**: CTM Comparator A match interrupt request flag  
               0: No request  
               1: Interrupt request
- Bit 4        **CTMPF**: CTM Comparator P match interrupt request flag  
               0: No request  
               1: Interrupt request
- Bit 3 ~ 2     Unimplemented, read as "0"
- Bit 1        **CTMAE**: CTM Comparator A match interrupt control  
               0: Disable  
               1: Enable
- Bit 0        **CTMPE**: CTM Comparator P match interrupt control  
               0: Disable  
               1: Enable

**MFII Register**

Bit	7	6	5	4	3	2	1	0
Name	LVF	—	STMAF	STMPF	LVE	—	STMAE	STMPE
R/W	R/W	—	R/W	R/W	R/W	—	R/W	R/W
POR	0	—	0	0	0	—	0	0

- Bit 7        **LVF**: LVD interrupt request flag  
               0: No request  
               1: Interrupt request
- Bit 6        Unimplemented, read as "0"
- Bit 5        **STMAF**: STM Comparator A match interrupt request flag  
               0: No request  
               1: Interrupt request
- Bit 4        **STMPF**: STM Comparator P match interrupt request flag  
               0: No request  
               1: Interrupt request
- Bit 3        **LVE**: LVD interrupt control  
               0: Disable  
               1: Enable
- Bit 2        Unimplemented, read as "0"
- Bit 1        **STMAE**: STM Comparator A match interrupt control  
               0: Disable  
               1: Enable
- Bit 0        **STMPE**: STM Comparator P match interrupt control  
               0: Disable  
               1: Enable

**MFI2 Register**

Bit	7	6	5	4	3	2	1	0
Name	PTM1AF	PTM1PF	PTM0AF	PTM0PF	PTM1AE	PTM1PE	PTM0AE	PTM0PE
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **PTM1AF**: PTM1 Comparator A match interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 6      **PTM1PF**: PTM1 Comparator P match interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 5      **PTM0AF**: PTM0 Comparator A match interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 4      **PTM0PF**: PTM0 Comparator P match interrupt request flag  
             0: No request  
             1: Interrupt request
- Bit 3      **PTM1AE**: PTM1 Comparator A match interrupt control  
             0: Disable  
             1: Enable
- Bit 2      **PTM1PE**: PTM1 Comparator P match interrupt control  
             0: Disable  
             1: Enable
- Bit 1      **PTM0AE**: PTM0 Comparator A match interrupt control  
             0: Disable  
             1: Enable
- Bit 0      **PTM0PE**: PTM0 Comparator P match interrupt control  
             0: Disable  
             1: Enable

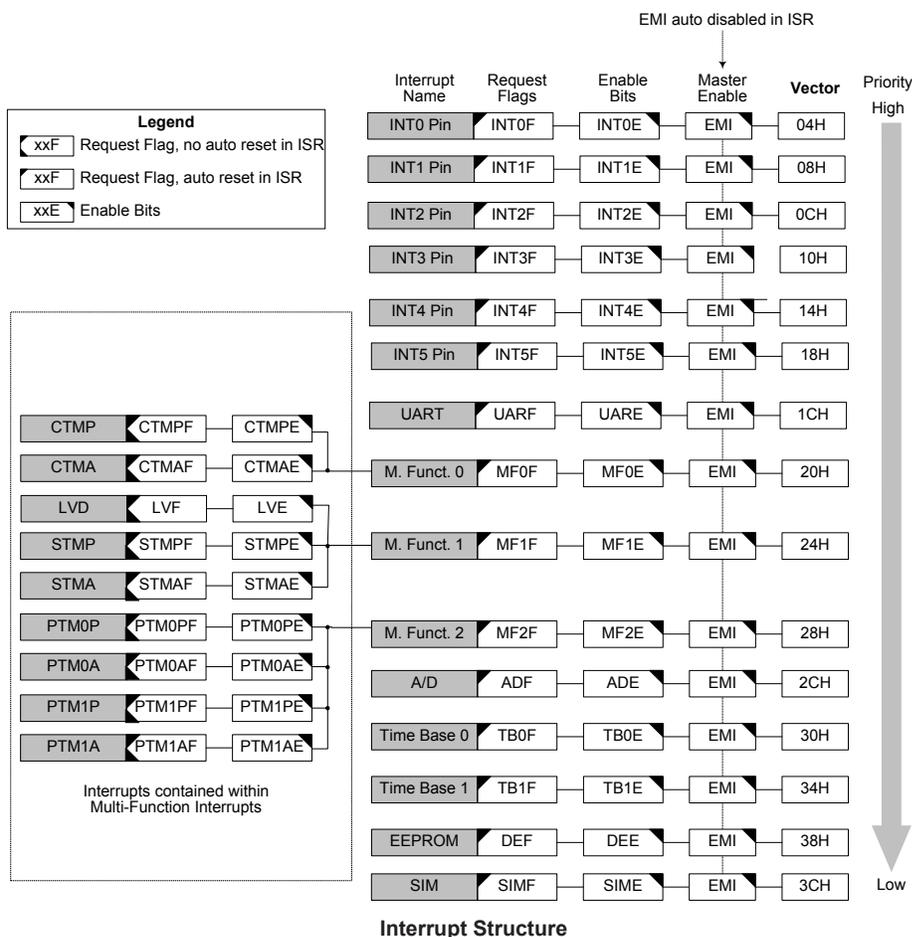
## Interrupt Operation

When the conditions for an interrupt event occur, such as a TM Comparator P or Comparator A match or A/D conversion completion etc, the relevant interrupt request flag will be set. Whether the request flag actually generates a program jump to the relevant interrupt vector is determined by the condition of the interrupt enable bit. If the enable bit is set high then the program will jump to its relevant vector; if the enable bit is zero then although the interrupt request flag is set an actual interrupt will not be generated and the program will not jump to the relevant interrupt vector. The global interrupt enable bit, if cleared to zero, will disable all interrupts.

When an interrupt is generated, the Program Counter, which stores the address of the next instruction to be executed, will be transferred onto the stack. The Program Counter will then be loaded with a new address which will be the value of the corresponding interrupt vector. The microcontroller will then fetch its next instruction from this interrupt vector. The instruction at this vector will usually be a “JMP” which will jump to another section of program which is known as the interrupt service routine. Here is located the code to control the appropriate interrupt. The interrupt service routine must be terminated with a “RETI”, which retrieves the original Program Counter address from the stack and allows the microcontroller to continue with normal execution at the point where the interrupt occurred.

The various interrupt enable bits, together with their associated request flags, are shown in the accompanying diagrams with their order of priority. Some interrupt sources have their own individual vector while others share the same multi-function interrupt vector. Once an interrupt subroutine is serviced, all the other interrupts will be blocked, as the global interrupt enable bit, EMI bit will be cleared automatically. This will prevent any further interrupt nesting from occurring. However, if other interrupt requests occur during this interval, although the interrupt will not be immediately serviced, the request flag will still be recorded.

If an interrupt requires immediate servicing while the program is already in another interrupt service routine, the EMI bit should be set after entering the routine, to allow interrupt nesting. If the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the Stack Pointer is decremented. If immediate service is desired, the stack must be prevented from becoming full. In case of simultaneous requests, the accompanying diagram shows the priority that is applied. All of the interrupt request flags when set will wake-up the device if it is in SLEEP or IDLE Mode, however to prevent a wake-up from occurring the corresponding flag should be set before the device is in SLEEP or IDLE Mode.



## External Interrupt

The external interrupt is controlled by signal transitions on the pins INT0~INT5. An external interrupt request will take place when the external interrupt request flag, INTnF, is set, which will occur when a transition, whose type is chosen by the edge select bits, appears on the external interrupt pin. To allow the program to branch to the interrupt vector address, the global interrupt enable bit, EMI, and the external interrupt enable bit, INTnE, must first be set. Additionally the correct interrupt edge type must be selected using the INTEGn register to enable the external interrupt function and to choose the trigger edge type. As the external interrupt pin is pin-shared with an I/O pin, it can only be configured as external interrupt pin if its external interrupt enable bit in the corresponding interrupt register has been set. The pin must also be setup as an input by setting the corresponding bit in the port control register. When the interrupt is enabled, the stack is not full and the correct transition type appears on the external interrupt pin, a subroutine call to the external interrupt vector, will take place. When the interrupt is serviced, the external interrupt request flag, INTnF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts. Note that the pull-high resistor selection on the external interrupt pin will remain valid even if the pin is used as an external interrupt input.

The INTEG0~INTEG1 registers are used to select the type of active edge that will trigger the external interrupt. A choice of either rising or falling or both edge types can be chosen to trigger an external interrupt. Note that the INTEG0~INTEG1 registers can also be used to disable the external interrupt function.

### Multi-function Interrupt

Within this device there are up to three Multi-function interrupts. Unlike the other independent interrupts, these interrupts have no independent source, but rather are formed from other existing interrupt sources, namely the TM Interrupts and LVD interrupt.

A Multi-function interrupt request will take place when any of the Multi-function interrupt request flags, MF0F~MF2F are set. The Multi-function interrupt flags will be set when any of their included functions generate an interrupt request flag. To allow the program to branch to its respective interrupt vector address, when the Multi-function interrupt is enabled and the stack is not full, and either one of the interrupts contained within each of Multi-function interrupt occurs, a subroutine call to one of the Multi-function interrupt vectors will take place. When the interrupt is serviced, the related Multi-Function request flag will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

However, it must be noted that, although the Multi-function Interrupt flags will be automatically reset when the interrupt is serviced, the request flags from the original source of the Multi-function interrupts, namely the TM Interrupts and LVD interrupt, will not be automatically reset and must be manually reset by the application program.

### A/D Converter Interrupt

The device contains an A/D converter which has its own independent interrupt. The A/D Converter Interrupt is controlled by the termination of an A/D conversion process. An A/D Converter Interrupt request will take place when the A/D Converter Interrupt request flag, ADF, is set, which occurs when the A/D conversion process finishes. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and A/D Interrupt enable bit, ADE, must first be set. When the interrupt is enabled, the stack is not full and the A/D conversion process has ended, a subroutine call to the A/D Converter Interrupt vector, will take place. When the interrupt is serviced, the A/D Converter Interrupt flag, ADF, will be automatically cleared. The EMI bit will also be automatically cleared to disable other interrupts.

### Time Base Interrupts

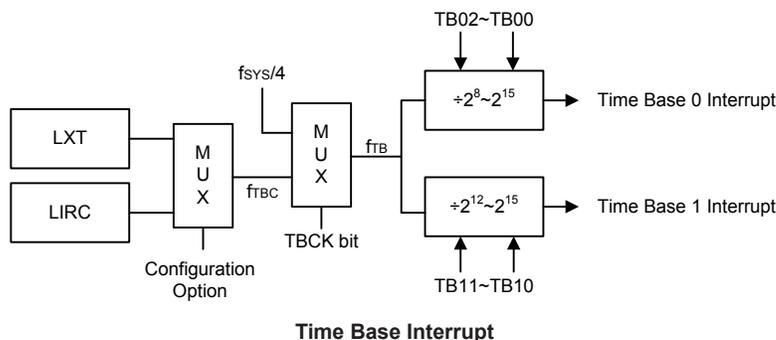
The function of the Time Base Interrupts is to provide regular time signal in the form of an internal interrupt. They are controlled by the overflow signals from their respective timer functions. When these happens their respective interrupt request flags, TB0F or TB1F will be set. To allow the program to branch to their respective interrupt vector addresses, the global interrupt enable bit, EMI and Time Base enable bits, TB0E or TB1E, must first be set. When the interrupt is enabled, the stack is not full and the Time Base overflows, a subroutine call to their respective vector locations will take place. When the interrupt is serviced, the respective interrupt request flag, TB0F or TB1F, will be automatically reset and the EMI bit will be cleared to disable other interrupts.

The purpose of the Time Base Interrupt is to provide an interrupt signal at fixed time periods. Their clock sources originate from the internal clock source  $f_{TB}$ . This  $f_{TB}$  input clock passes through a divider, the division ratio of which is selected by programming the appropriate bits in the TBC register to obtain longer interrupt periods whose value ranges. The clock source that generates  $f_{TB}$ , which in turn controls the Time Base interrupt period, can originate from several different sources, as shown in the System Operating Mode section.

### TBC Register

Bit	7	6	5	4	3	2	1	0
Name	TBON	TBCK	TB11	TB10	LXTLP	TB02	TB01	TB00
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	1	1	0	1	1	1

- Bit 7 **TBON**: TB0 and TB1 Control bit  
0: Disable  
1: Enable
- Bit 6 **TBCK**: Select  $f_{TB}$  Clock  
0:  $f_{TB}$   
1:  $f_{SYS}/4$
- Bit 5 ~ 4 **TB11 ~ TB10**: Select Time Base 1 Time-out Period  
00:  $4096/f_{TB}$   
01:  $8192/f_{TB}$   
10:  $16384/f_{TB}$   
11:  $32768/f_{TB}$
- Bit 3 **LXTLP**: LXT Low Power Mode control  
0: Disable (Quick Start Mode)  
1: Enable (Low Power Mode)
- Bit 2 ~ 0 **TB02 ~ TB00**: Select Time Base 0 Time-out Period  
000:  $2^8/f_{TB}$   
001:  $2^9/f_{TB}$   
010:  $2^{10}/f_{TB}$   
011:  $2^{11}/f_{TB}$   
100:  $2^{12}/f_{TB}$   
101:  $2^{13}/f_{TB}$   
110:  $2^{14}/f_{TB}$   
111:  $2^{15}/f_{TB}$



### Serial Interface Module Interrupts

A SIM Interrupt request will take place when the SIM Interrupt request flag, SIMF, is set, which occurs when a byte of data has been received or transmitted by the SIM interface. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and the Serial Interface Interrupt enable bit, SIME, must first be set. When the interrupt is enabled, the stack is not full and a byte of data has been transmitted or received by the SIM interface, a subroutine call to the respective Interrupt vector, will take place. When the interrupt is serviced, the respective interrupt request flag, SIMF, will be automatically reset and the EMI bit will be cleared to disable other interrupts.

## EEPROM Interrupt

An EEPROM Interrupt request will take place when the EEPROM Interrupt request flag, DEF, is set, which occurs when an EEPROM Write cycle ends. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and EEPROM Interrupt enable bit, DEE, must first be set. When the interrupt is enabled, the stack is not full and an EEPROM Write cycle ends, a subroutine call to the respective EEPROM Interrupt vector, will take place. When the EEPROM Interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts, and the EEPROM interrupt request flag, DEF, will also be automatically cleared.

## TM Interrupts

The Compact, Standard and Periodic Type TMs each has two interrupts. All of the TM interrupts are contained within the Multi-function Interrupts. For each of the Compact, Standard and Periodic Type TMs there are two interrupt request flags xTMnPF and xTMnAF and two enable bits xTMnPE and xTMnAE. A TM interrupt request will take place when any of the TM request flags are set, a situation which occurs when a TM comparator P or comparator A match situation happens.

To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and the respective TM Interrupt enable bit, and associated Multi-function interrupt enable bit, MFnF, must first be set. When the interrupt is enabled, the stack is not full and a TM comparator match situation occurs, a subroutine call to the relevant TM Interrupt vector locations, will take place. When the TM interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts, however only the related MFnF flag will be automatically cleared. As the TM interrupt request flags will not be automatically cleared, they have to be cleared by the application program.

## LVD Interrupt

The Low Voltage Detector Interrupt is contained within the Multi-function Interrupt. An LVD Interrupt request will take place when the LVD Interrupt request flag, LVF, is set, which occurs when the Low Voltage Detector function detects a low power supply voltage. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, Low Voltage Interrupt enable bit, LVE, and associated Multi-function interrupt enable bit, must first be set. When the interrupt is enabled, the stack is not full and a low voltage condition occurs, a subroutine call to the Multi-function Interrupt vector will take place. When the Low Voltage Interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts, however only the Multi-function interrupt request flag will be also automatically cleared. As the LVF flag will not be automatically cleared, it has to be cleared by the application program.

## UART Interrupt

Several individual UART conditions can generate a UART interrupt. When these conditions exist, a low pulse will be generated to get the attention of the microcontroller. These conditions are a transmitter data register empty, transmitter idle, receiver data available, receiver overrun, address detect and an RX pin wake-up. To allow the program to branch to the respective interrupt vector addresses, the global interrupt enable bit, EMI, and UART interrupt enable bit, UARE, must first be set. When the interrupt is enabled, the stack is not full and any of these conditions are created, a subroutine call to the UART Interrupt vector will take place. When the interrupt is serviced, the UART Interrupt flag, UARF, will be automatically cleared. The EMI bit will also be automatically cleared to disable other interrupts. However, the USR register flags will be cleared automatically when certain actions are taken by the UART, the details of which are given in the UART section.

## **Interrupt Wake-up Function**

Each of the interrupt functions has the capability of waking up the microcontroller when in the SLEEP or IDLE Mode. A wake-up is generated when an interrupt request flag changes from low to high and is independent of whether the interrupt is enabled or not. Therefore, even though the device is in the SLEEP or IDLE Mode and its system oscillator stopped, situations such as external edge transitions on the external interrupt pin, a low power supply voltage or comparator input change may cause their respective interrupt flag to be set high and consequently generate an interrupt. Care must therefore be taken if spurious wake-up situations are to be avoided. If an interrupt wake-up function is to be disabled then the corresponding interrupt request flag should be set high before the device enters the SLEEP or IDLE Mode. The interrupt enable bits have no effect on the interrupt wake-up function.

## **Programming Considerations**

By disabling the relevant interrupt enable bits, a requested interrupt can be prevented from being serviced, however, once an interrupt request flag is set, it will remain in this condition in the interrupt register until the corresponding interrupt is serviced or until the request flag is cleared by the application program.

Where a certain interrupt is contained within a Multi-function interrupt, then when the interrupt service routine is executed, as only the Multi-function interrupt request flags, MF0F~MF2F, will be automatically cleared, the individual request flag for the function needs to be cleared by the application program.

It is recommended that programs do not use the “CALL” instruction within the interrupt service subroutine. Interrupts often occur in an unpredictable manner or need to be serviced immediately. If only one stack is left and the interrupt is not well controlled, the original control sequence will be damaged once a CALL subroutine is executed in the interrupt subroutine.

Every interrupt has the capability of waking up the microcontroller when it is in SLEEP or IDLE Mode, the wake up being generated when the interrupt request flag changes from low to high. If it is required to prevent a certain interrupt from waking up the microcontroller then its respective request flag should be first set high before enter SLEEP or IDLE Mode.

As only the Program Counter is pushed onto the stack, then when the interrupt is serviced, if the contents of the accumulator, status register or other registers are altered by the interrupt service program, their contents should be saved to the memory at the beginning of the interrupt service routine.

To return from an interrupt subroutine, either a RET or RETI instruction may be executed. The RETI instruction in addition to executing a return to the main program also automatically sets the EMI bit high to allow further interrupts. The RET instruction however only executes a return to the main program leaving the EMI bit in its present zero state and therefore disabling the execution of further interrupts.

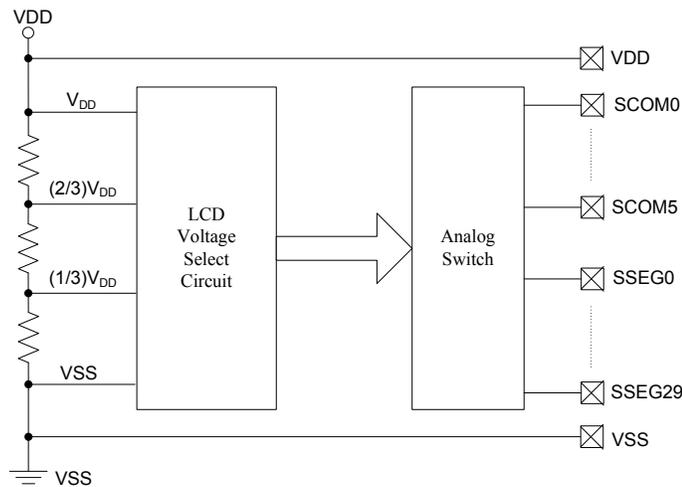
## Software LCD Driver

The devices have the capability of driving external LCD panels. The common pins for LCD driving, SCOM0~SCOM5, and segment pins, SSEG0~SSEG29, are pin shared with certain pin on the I/O ports. The LCD signals (COM and SEG) are generated using the application program.

### LCD operation

An external LCD panel can be driven using this device by configuring the I/O pins as common pins and configuring the I/O pins as segment pins. The LCD driver function is controlled using the SLCDC0~SLCDC4 registers which in addition to controlling the overall on/off function, LCD bias current setup function also controls the pin function selection. This enables the LCD COM and SEG driver to generate the necessary  $V_{SS}$ ,  $(1/3)V_{DD}$ ,  $(2/3)V_{DD}$  voltage and  $V_{DD}$  levels for LCD 1/3 bias operation.

The LCDEN bit in the SLCDC0 register is the overall master control for the LCD driver. The LCD SCOMn pin is selected to be used for LCD driving by the corresponding pin-shared function selection bits. Note that the Port Control register does not need to first setup the pins as outputs to enable the LCD driver operation.



**LCD Driver structure**

### LCD Control Registers

The LCD COM driver enables a range of selections to be provided to suit the requirement of the LCD panel which are being used. The bias current choice is implemented using the ISEL1 and ISEL0 bits in the SLCDC0 register, the FRAME bit is used to select the output frame.

**SLCDC0 Register**

Bit	7	6	5	4	3	2	1	0
Name	FRAME	ISEL1	ISEL0	LCDEN	COM3EN	COM2EN	COM1EN	COM0EN
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **FRAME**: Output frame0 or frame1  
0: Frame 0  
1: Frame 1
- Bit 6~5    **ISEL1~ISEL0**: LCD bias current selection ( $V_{DD}=5V$ )  
00: 8.3 $\mu$ A  
01: 16.7 $\mu$ A  
10: 50 $\mu$ A  
11: 100 $\mu$ A
- Bit 4      **LCDEN**: SCOM and SSEG module on/off control  
0: Off  
1: On  
SCOMn and SSEGm can be enable by COMnEN and SEGmEN if LCDEN=1  
If LCDEN=0, SCOMn and SSEGm output  $V_{SS}$   
When LCDEN is set, it will turn on the DC path of resistor to generate LCD Bias voltage.
- Bit 3      **COM3EN**: LCD or other function selection  
0: Other function  
1: SCOM3/SSEG3
- Bit 2      **COM2EN**: LCD or other function selection  
0: Other function  
1: SCOM2/SSEG2
- Bit 1      **COM1EN**: LCD or other function selection  
0: Other function  
1: SCOM1/SSEG1
- Bit 0      **COM0EN**: LCD or other function selection  
0: Other function  
1: SCOM0/SSEG0

**SLCDC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	COM5EN	COM4EN	COMSEGS5	COMSEGS4	COMSEGS3	COMSEGS2	COMSEGS1	COMSEGS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **COM5EN**: LCD or other function selection  
             0: Other function  
             1: SCOM5/SSEG5
- Bit 6      **COM4EN**: LCD or other function selection  
             0: Other function  
             1: SCOM4/SSEG4
- Bit 5      **COMSEGS5**: SCOM5 or SSEG 5 selection  
             0: SCOM5  
             1: SSEG5
- Bit 4      **COMSEGS4**: SCOM4 or SSEG 4 selection  
             0: SCOM4  
             1: SSEG4
- Bit 3      **COMSEGS3**: SCOM3 or SSEG 3 selection  
             0: SCOM3  
             1: SSEG3
- Bit 2      **COMSEGS2**: SCOM2 or SSEG 2 selection  
             0: SCOM2  
             1: SSEG2
- Bit 1      **COMSEGS1**: SCOM1 or SSEG 1 selection  
             0: SCOM1  
             1: SSEG1

**SLCDC2 Register**

Bit	7	6	5	4	3	2	1	0
Name	SEG13EN	SEG12EN	SEG11EN	SEG10EN	SEG9EN	SEG8EN	SEG7EN	SEG6EN
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **SEG13EN**: SSEG13 function control  
             0: Disable  
             1: Enable
  
- Bit 6      **SEG12EN**: SSEG12 function control  
             0: Disable  
             1: Enable
  
- Bit 5      **SEG11EN**: SSEG11 function control  
             0: Disable  
             1: Enable
  
- Bit 4      **SEG10EN**: SSEG10 function control  
             0: Disable  
             1: Enable
  
- Bit 3      **SEG9EN**: SSEG9 function control  
             0: Disable  
             1: Enable
  
- Bit 2      **SEG8EN**: SSEG8 function control  
             0: Disable  
             1: Enable
  
- Bit 1      **SEG7EN**: SSEG7 function control  
             0: Disable  
             1: Enable
  
- Bit 0      **SEG6EN**: SSEG6 function control  
             0: Disable  
             1: Enable

**SLCDC3 Register**

Bit	7	6	5	4	3	2	1	0
Name	SEG21EN	SEG20EN	SEG19EN	SEG18EN	SEG17EN	SEG16EN	SEG15EN	SEG14EN
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7      **SEG21EN**: SSEG21 function control

0: Disable  
1: Enable

Bit 6      **SEG20EN**: SSEG20 function control

0: Disable  
1: Enable

Bit 5      **SEG19EN**: SSEG19 function control

0: Disable  
1: Enable

Bit 4      **SEG18EN**: SSEG18 function control

0: Disable  
1: Enable

Bit 3      **SEG17EN**: SSEG17 function control

0: Disable  
1: Enable

Bit 2      **SEG16EN**: SSEG16 function control

0: Disable  
1: Enable

Bit 1      **SEG15EN**: SSEG15 function control

0: Disable  
1: Enable

Bit 0      **SEG14EN**: SSEG14 function control

0: Disable  
1: Enable

Note: SSEG14, SSEG15 share the pins with OSC1/XT1, OSC2/XT2 functions, if select the OSC2/OSC1 or XT2/XT1pin function, then the SEG14, SEG15 have no actions even if the functions are enabled.

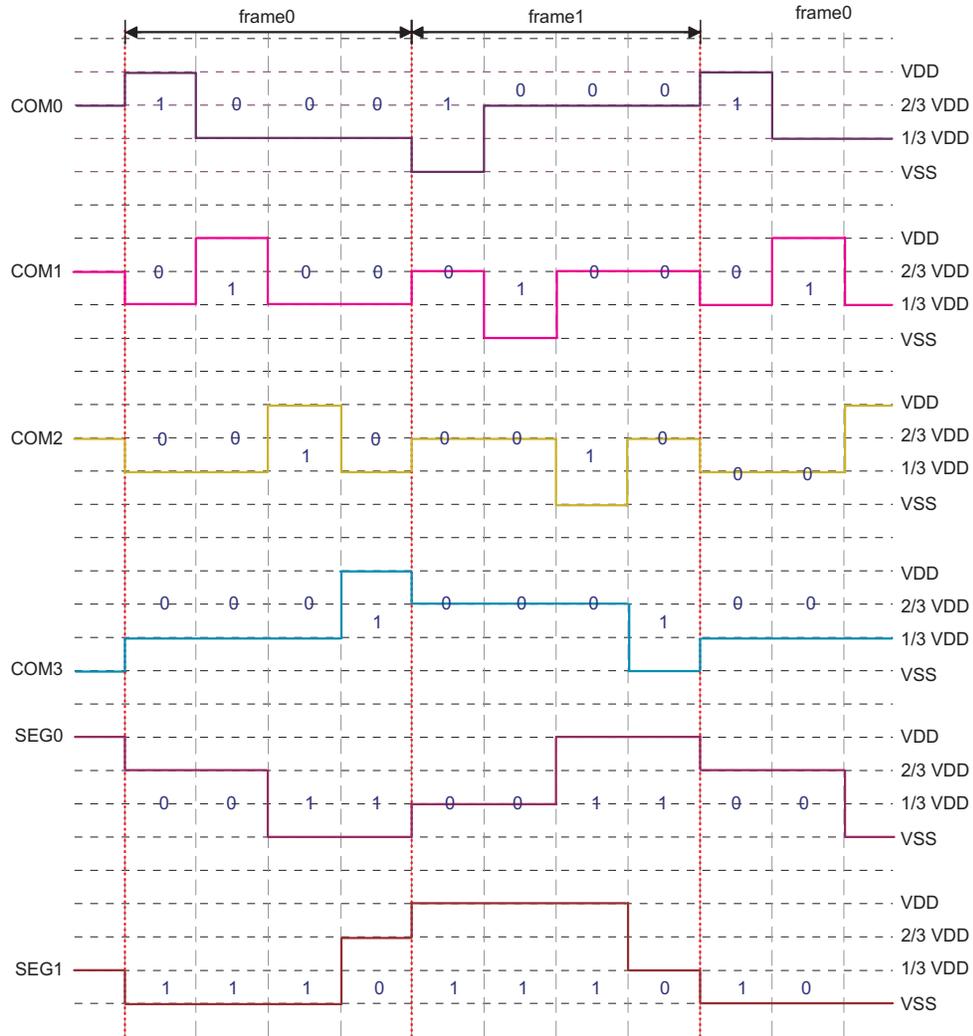
**SLCDC4 Register**

Bit	7	6	5	4	3	2	1	0
Name	SEG29EN	SEG28EN	SEG27EN	SEG26EN	SEG25EN	SEG24EN	SEG23EN	SEG22EN
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **SEG29EN**: SSEG29 function control  
             0: Disable  
             1: Enable
  
- Bit 6      **SEG28EN**: SSEG28 function control  
             0: Disable  
             1: Enable
  
- Bit 5      **SEG27EN**: SSEG27 function control  
             0: Disable  
             1: Enable
  
- Bit 4      **SEG26EN**: SSEG26 function control  
             0: Disable  
             1: Enable
  
- Bit 3      **SEG25EN**: SSEG25 function control  
             0: Disable  
             1: Enable
  
- Bit 2      **SEG24EN**: SSEG24 function control  
             0: Disable  
             1: Enable
  
- Bit 1      **SEG23EN**: SSEG23 function control  
             0: Disable  
             1: Enable
  
- Bit 0      **SEG22EN**: SSEG22 function control  
             0: Disable  
             1: Enable

**LCD waveform**

The accompanying waveform diagram shows a typical 1/3 Bias LCD waveform generated using the application program. Note that the depiction of a "1" in the diagram illustrates an illuminated LCD pixel. The COM signal polarity generated on pins SCOMn, whether 0 or 1, are generated using the corresponding I/O data registers.



1/3 Bias LCD Waveform

A cyclic LCD waveform includes two frames, known as Frame 0 and Frame 1 which is selected by the FRAME bit in the SLCDC0 register. The following offers a functional explanation.

To select Frame 0 clear the FRAME bit to 0.

In frame 0, the COM signal output can have a value of  $V_{DD}$ , or have a Vbias value of  $1/3V_{DD}$  (SEG\_L). The SEG signal can have a value of  $V_{SS}$ , or have a Vbias value of  $2/3V_{DD}$  (SEG\_H).

In frame 1, the COM signal output can have a value of  $V_{SS}$ , or have a Vbias value of  $2/3V_{DD}$  (SEG\_H). The SEG signal can have a value of  $V_{DD}$  have a Vbias value of  $1/3V_{DD}$  (SEG\_L).

The SCOM0~SCOMn waveform is controlled by the application program using the FRAME bit, and the corresponding I/O data register for the respective SCOM pin to determine whether the SCOM0~SCOMn output has a value of either  $V_{DD}$ ,  $V_{SS}$  or Vbias. The SSEG0~SSEGm waveform is controlled in a similar way using the FRAME bit and the corresponding I/O data register for the respective SSEG pin to determine whether the SSEG0~SSEGn output has a value of either  $V_{DD}$ ,  $V_{SS}$  or Vbias.

## Low Voltage Detector – LVD

Each device has a Low Voltage Detector function, also known as LVD. This enabled the device to monitor the power supply voltage,  $V_{DD}$ , and provide a warning signal should it fall below a certain level. This function may be especially useful in battery applications where the supply voltage will gradually reduce as the battery ages, as it allows an early warning battery low signal to be generated. The Low Voltage Detector also has the capability of generating an interrupt signal.

### LVD Register

The Low Voltage Detector function is controlled using a single register with the name LVDC. Three bits in this register, VLVD2~VLVD0, are used to select one of eight fixed voltages below which a low voltage condition will be determined. A low voltage condition is indicated when the LVDO bit is set. If the LVDO bit is low, this indicates that the  $V_{DD}$  voltage is above the preset low voltage value. The LVDEN bit is used to control the overall on/off function of the low voltage detector. Setting the bit high will enable the low voltage detector. Clearing the bit to zero will switch off the internal low voltage detector circuits. As the low voltage detector will consume a certain amount of power, it may be desirable to switch off the circuit when not in use, an important consideration in power sensitive battery powered applications.

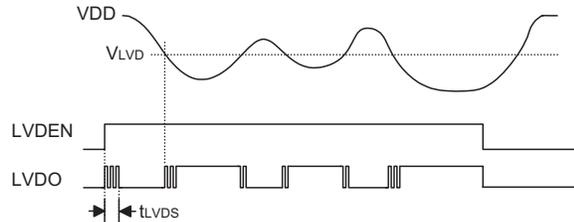
### LVDC Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	LVDO	LVDEN	—	VLVD2	VLVD1	VLVD0
R/W	—	—	R	R/W	—	R/W	R/W	R/W
POR	—	—	0	0	—	0	0	0

- Bit 7~6 Unimplemented, read as "0"
- Bit 5 **LVDO**: LVD Output Flag  
 0: No Low Voltage Detect  
 1: Low Voltage Detect
- Bit 4 **LVDEN**: Low Voltage Detector Control  
 0: Disable  
 1: Enable
- Bit 3 Unimplemented, read as "0"
- Bit 2~0 **VLVD2~VLVD0**: Select LVD Voltage  
 000: 2.0V  
 001: 2.2V  
 010: 2.4V  
 011: 2.7V  
 100: 3.0V  
 101: 3.3V  
 110: 3.6V  
 111: 4.0V

**LVD Operation**

The Low Voltage Detector function operates by comparing the power supply voltage,  $V_{DD}$ , with a pre-specified voltage level stored in the LVDC register. This has a range of between 2.0V and 4.0V. When the power supply voltage,  $V_{DD}$ , falls below this pre-determined value, the LVDO bit will be set high indicating a low power supply voltage condition. The Low Voltage Detector function is supplied by a reference voltage which will be automatically enabled. When the device is powered down the low voltage detector will remain active if the LVDEN bit is high. After enabling the Low Voltage Detector, a time delay  $t_{LVDS}$  should be allowed for the circuitry to stabilise before reading the LVDO bit. Note also that as the  $V_{DD}$  voltage may rise and fall rather slowly, at the voltage nears that of  $V_{LVD}$ , there may be multiple bit LVDO transitions.



**LVD Operation**

The Low Voltage Detector also has its own interrupt which is contained within one of the Multi-function interrupts, providing an alternative means of low voltage detection, in addition to polling the LVDO bit. The interrupt will only be generated after a delay of  $t_{LVD}$  after the LVDO bit has been set high by a low voltage condition. When the device is powered down the Low Voltage Detector will remain active if the LVDEN bit is high. In this case, the LVF interrupt request flag will be set, causing an interrupt to be generated if  $V_{DD}$  falls below the preset LVD voltage. This will cause the device to wake-up from the SLEEP Mode, however if the Low Voltage Detector wake up function is not required then the LVF flag should be first set high before the device enters the SLEEP Mode.

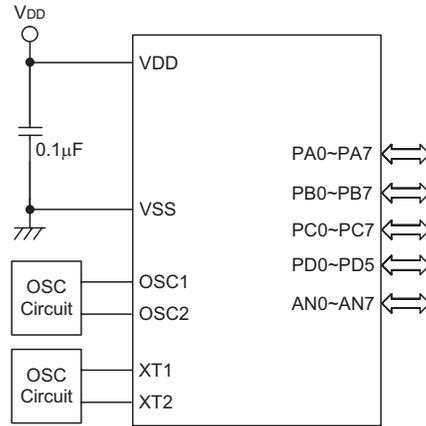
When LVD function is enabled, it is recommended to clear LVD flag first, and then enables interrupt function to avoid mistake action.

## Configuration Option

Configuration options refer to certain options within the MCU that are programmed into the device during the programming process. During the development process, these options are selected using the HT-IDE software development tools. As these options are programmed into the device using the hardware programming tools, once they are selected they cannot be changed later using the application program. All options must be defined for proper system function, the details of which are shown in the table.

No.	Options
<b>Oscillator Option</b>	
1	High Speed/Low Speed System Oscillator Selection – $f_{osc}$ : HIRC+LIRC HIRC+LXT HXT+LIRC

## Application Circuit



## Instruction Set

### Introduction

Central to the successful operation of any microcontroller is its instruction set, which is a set of program instruction codes that directs the microcontroller to perform certain operations. In the case of Holtek microcontroller, a comprehensive and flexible set of over 60 instructions is provided to enable programmers to implement their application with the minimum of programming overheads.

For easier understanding of the various instruction codes, they have been subdivided into several functional groupings.

### Instruction Timing

Most instructions are implemented within one instruction cycle. The exceptions to this are branch, call, or table read instructions where two instruction cycles are required. One instruction cycle is equal to 4 system clock cycles, therefore in the case of an 8MHz system oscillator, most instructions would be implemented within 0.5 $\mu$ s and branch or call instructions would be implemented within 1 $\mu$ s. Although instructions which require one more cycle to implement are generally limited to the JMP, CALL, RET, RETI and table read instructions, it is important to realize that any other instructions which involve manipulation of the Program Counter Low register or PCL will also take one more cycle to implement. As instructions which change the contents of the PCL will imply a direct jump to that new address, one more cycle will be required. Examples of such instructions would be “CLR PCL” or “MOV PCL, A”. For the case of skip instructions, it must be noted that if the result of the comparison involves a skip operation then this will also take one more cycle, if no skip is involved then only one cycle is required.

### Moving and Transferring Data

The transfer of data within the microcontroller program is one of the most frequently used operations. Making use of several kinds of MOV instructions, data can be transferred from registers to the Accumulator and vice-versa as well as being able to move specific immediate data directly into the Accumulator. One of the most important data transfer applications is to receive data from the input ports and transfer data to the output ports.

### Arithmetic Operations

The ability to perform certain arithmetic operations and data manipulation is a necessary feature of most microcontroller applications. Within the Holtek microcontroller instruction set are a range of add and subtract instruction mnemonics to enable the necessary arithmetic to be carried out. Care must be taken to ensure correct handling of carry and borrow data when results exceed 255 for addition and less than 0 for subtraction. The increment and decrement instructions such as INC, INCA, DEC and DECA provide a simple means of increasing or decreasing by a value of one of the values in the destination specified.

## Logical and Rotate Operation

The standard logical operations such as AND, OR, XOR and CPL all have their own instruction within the Holtek microcontroller instruction set. As with the case of most instructions involving data manipulation, data must pass through the Accumulator which may involve additional programming steps. In all logical data operations, the zero flag may be set if the result of the operation is zero. Another form of logical data manipulation comes from the rotate instructions such as RR, RL, RRC and RLC which provide a simple means of rotating one bit right or left. Different rotate instructions exist depending on program requirements. Rotate instructions are useful for serial port programming applications where data can be rotated from an internal register into the Carry bit from where it can be examined and the necessary serial bit set high or low. Another application which rotate data operations are used is to implement multiplication and division calculations.

## Branches and Control Transfer

Program branching takes the form of either jumps to specified locations using the JMP instruction or to a subroutine using the CALL instruction. They differ in the sense that in the case of a subroutine call, the program must return to the instruction immediately when the subroutine has been carried out. This is done by placing a return instruction “RET” in the subroutine which will cause the program to jump back to the address right after the CALL instruction. In the case of a JMP instruction, the program simply jumps to the desired location. There is no requirement to jump back to the original jumping off point as in the case of the CALL instruction. One special and extremely useful set of branch instructions are the conditional branches. Here a decision is first made regarding the condition of a certain data memory or individual bits. Depending upon the conditions, the program will continue with the next instruction or skip over it and jump to the following instruction. These instructions are the key to decision making and branching within the program perhaps determined by the condition of certain input switches or by the condition of internal data bits.

## Bit Operations

The ability to provide single bit operations on Data Memory is an extremely flexible feature of all Holtek microcontrollers. This feature is especially useful for output port bit programming where individual bits or port pins can be directly set high or low using either the “SET [m].i” or “CLR [m].i” instructions respectively. The feature removes the need for programmers to first read the 8-bit output port, manipulate the input data to ensure that other bits are not changed and then output the port with the correct new data. This read-modify-write process is taken care of automatically when these bit operation instructions are used.

## Table Read Operations

Data storage is normally implemented by using registers. However, when working with large amounts of fixed data, the volume involved often makes it inconvenient to store the fixed data in the Data Memory. To overcome this problem, Holtek microcontrollers allow an area of Program Memory to be setup as a table where data can be directly stored. A set of easy to use instructions provides the means by which this fixed data can be referenced and retrieved from the Program Memory.

## Other Operations

In addition to the above functional instructions, a range of other instructions also exist such as the “HALT” instruction for Power-down operations and instructions to control the operation of the Watchdog Timer for reliable program operations under extreme electric or electromagnetic environments. For their relevant operations, refer to the functional related sections.

## Instruction Set Summary

The instructions related to the data memory access in the following table can be used when the desired data memory is located in Data Memory sector 0.

### Table Conventions

x: Bits immediate data  
m: Data Memory address  
A: Accumulator  
i: 0~7 number of bits  
addr: Program memory address

Mnemonic	Description	Cycles	Flag Affected
<b>Arithmetic</b>			
ADD A,[m]	Add Data Memory to ACC	1	Z, C, AC, OV, SC
ADDM A,[m]	Add ACC to Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC
ADD A,x	Add immediate data to ACC	1	Z, C, AC, OV, SC
ADC A,[m]	Add Data Memory to ACC with Carry	1	Z, C, AC, OV, SC
ADCM A,[m]	Add ACC to Data memory with Carry	1 <sup>Note</sup>	Z, C, AC, OV, SC
SUB A,x	Subtract immediate data from the ACC	1	Z, C, AC, OV, SC, CZ
SUB A,[m]	Subtract Data Memory from ACC	1	Z, C, AC, OV, SC, CZ
SUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
SBC A,x	Subtract immediate data from ACC with Carry	1	Z, C, AC, OV, SC, CZ
SBC A,[m]	Subtract Data Memory from ACC with Carry	1	Z, C, AC, OV, SC, CZ
SBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
DAA [m]	Decimal adjust ACC for Addition with result in Data Memory	1 <sup>Note</sup>	C
<b>Logic Operation</b>			
AND A,[m]	Logical AND Data Memory to ACC	1	Z
OR A,[m]	Logical OR Data Memory to ACC	1	Z
XOR A,[m]	Logical XOR Data Memory to ACC	1	Z
ANDM A,[m]	Logical AND ACC to Data Memory	1 <sup>Note</sup>	Z
ORM A,[m]	Logical OR ACC to Data Memory	1 <sup>Note</sup>	Z
XORM A,[m]	Logical XOR ACC to Data Memory	1 <sup>Note</sup>	Z
AND A,x	Logical AND immediate Data to ACC	1	Z
OR A,x	Logical OR immediate Data to ACC	1	Z
XOR A,x	Logical XOR immediate Data to ACC	1	Z
CPL [m]	Complement Data Memory	1 <sup>Note</sup>	Z
CPLA [m]	Complement Data Memory with result in ACC	1	Z
<b>Increment &amp; Decrement</b>			
INCA [m]	Increment Data Memory with result in ACC	1	Z
INC [m]	Increment Data Memory	1 <sup>Note</sup>	Z
DECA [m]	Decrement Data Memory with result in ACC	1	Z
DEC [m]	Decrement Data Memory	1 <sup>Note</sup>	Z
<b>Rotate</b>			
RRA [m]	Rotate Data Memory right with result in ACC	1	None
RR [m]	Rotate Data Memory right	1 <sup>Note</sup>	None
RRCA [m]	Rotate Data Memory right through Carry with result in ACC	1	C
RRC [m]	Rotate Data Memory right through Carry	1 <sup>Note</sup>	C
RLA [m]	Rotate Data Memory left with result in ACC	1	None
RL [m]	Rotate Data Memory left	1 <sup>Note</sup>	None
RLCA [m]	Rotate Data Memory left through Carry with result in ACC	1	C
RLC [m]	Rotate Data Memory left through Carry	1 <sup>Note</sup>	C

Mnemonic	Description	Cycles	Flag Affected
<b>Data Move</b>			
MOV A,[m]	Move Data Memory to ACC	1	None
MOV [m],A	Move ACC to Data Memory	1 <sup>Note</sup>	None
MOV A,x	Move immediate data to ACC	1	None
<b>Bit Operation</b>			
CLR [m].i	Clear bit of Data Memory	1 <sup>Note</sup>	None
SET [m].i	Set bit of Data Memory	1 <sup>Note</sup>	None
<b>Branch Operation</b>			
JMP addr	Jump unconditionally	2	None
SZ [m]	Skip if Data Memory is zero	1 <sup>Note</sup>	None
SZA [m]	Skip if Data Memory is zero with data movement to ACC	1 <sup>Note</sup>	None
SZ [m].i	Skip if bit i of Data Memory is zero	1 <sup>Note</sup>	None
SNZ [m]	Skip if Data Memory is not zero	1 <sup>Note</sup>	None
SNZ [m].i	Skip if bit i of Data Memory is not zero	1 <sup>Note</sup>	None
SIZ [m]	Skip if increment Data Memory is zero	1 <sup>Note</sup>	None
SDZ [m]	Skip if decrement Data Memory is zero	1 <sup>Note</sup>	None
SIZA [m]	Skip if increment Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
SDZA [m]	Skip if decrement Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
CALL addr	Subroutine call	2	None
RET	Return from subroutine	2	None
RET A,x	Return from subroutine and load immediate data to ACC	2	None
RETI	Return from interrupt	2	None
<b>Table Read Operation</b>			
TABRD [m]	Read table to TBLH and Data Memory	2 <sup>Note</sup>	None
TABRDL [m]	Read table (last page) to TBLH and Data Memory	2 <sup>Note</sup>	None
ITABRD [m]	Increment table pointer TBLP first and Read table to TBLH and Data Memory	2 <sup>Note</sup>	None
ITABRDL [m]	Increment table pointer TBLP first and Read table (last page) to TBLH and Data Memory	2 <sup>Note</sup>	None
<b>Miscellaneous</b>			
NOP	No operation	1	None
CLR [m]	Clear Data Memory	1 <sup>Note</sup>	None
SET [m]	Set Data Memory	1 <sup>Note</sup>	None
CLR WDT	Clear Watchdog Timer	1	TO, PDF
SWAP [m]	Swap nibbles of Data Memory	1 <sup>Note</sup>	None
SWAPA [m]	Swap nibbles of Data Memory with result in ACC	1	None
HALT	Enter power down mode	1	TO, PDF

Note: 1. For skip instructions, if the result of the comparison involves a skip then up to three cycles are required, if no skip takes place only one cycle is required.

2. Any instruction which changes the contents of the PCL will also require 2 cycles for execution.

3. For the “CLR WDT” instruction the TO and PDF flags may be affected by the execution status. The TO and PDF flags are cleared after the “CLR WDT” instructions is executed. Otherwise the TO and PDF flags remain unchanged.

### Extended Instruction Set

The extended instructions are used to support the full range address access for the data memory. When the accessed data memory is located in any data memory sections except sector 0, the extended instruction can be used to access the data memory instead of using the indirect addressing access to improve the CPU firmware performance.

Mnemonic	Description	Cycles	Flag Affected
<b>Arithmetic</b>			
LADD A,[m]	Add Data Memory to ACC	2	Z, C, AC, OV, SC
LADDM A,[m]	Add ACC to Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC
LADC A,[m]	Add Data Memory to ACC with Carry	2	Z, C, AC, OV, SC
LADCM A,[m]	Add ACC to Data memory with Carry	2 <sup>Note</sup>	Z, C, AC, OV, SC
LSUB A,[m]	Subtract Data Memory from ACC	2	Z, C, AC, OV, SC, CZ
LSUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
LSBC A,[m]	Subtract Data Memory from ACC with Carry	2	Z, C, AC, OV, SC, CZ
LSBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	2 <sup>Note</sup>	Z, C, AC, OV, SC, CZ
LDAA [m]	Decimal adjust ACC for Addition with result in Data Memory	2 <sup>Note</sup>	C
<b>Logic Operation</b>			
LAND A,[m]	Logical AND Data Memory to ACC	2	Z
LOR A,[m]	Logical OR Data Memory to ACC	2	Z
LXOR A,[m]	Logical XOR Data Memory to ACC	2	Z
LANDM A,[m]	Logical AND ACC to Data Memory	2 <sup>Note</sup>	Z
LORM A,[m]	Logical OR ACC to Data Memory	2 <sup>Note</sup>	Z
LXORM A,[m]	Logical XOR ACC to Data Memory	2 <sup>Note</sup>	Z
LCPL [m]	Complement Data Memory	2 <sup>Note</sup>	Z
LCPLA [m]	Complement Data Memory with result in ACC	2	Z
<b>Increment &amp; Decrement</b>			
LINCA [m]	Increment Data Memory with result in ACC	2	Z
LINC [m]	Increment Data Memory	2 <sup>Note</sup>	Z
LDECA [m]	Decrement Data Memory with result in ACC	2	Z
LDEC [m]	Decrement Data Memory	2 <sup>Note</sup>	Z
<b>Rotate</b>			
LRRRA [m]	Rotate Data Memory right with result in ACC	2	None
LRR [m]	Rotate Data Memory right	2 <sup>Note</sup>	None
LRRCA [m]	Rotate Data Memory right through Carry with result in ACC	2	C
LRRC [m]	Rotate Data Memory right through Carry	2 <sup>Note</sup>	C
LRLA [m]	Rotate Data Memory left with result in ACC	2	None
LRL [m]	Rotate Data Memory left	2 <sup>Note</sup>	None
LRLCA [m]	Rotate Data Memory left through Carry with result in ACC	2	C
LRLC [m]	Rotate Data Memory left through Carry	2 <sup>Note</sup>	C
<b>Data Move</b>			
LMOV A,[m]	Move Data Memory to ACC	2	None
LMOV [m],A	Move ACC to Data Memory	2 <sup>Note</sup>	None
<b>Bit Operation</b>			
LCLR [m].i	Clear bit of Data Memory	2 <sup>Note</sup>	None
LSET [m].i	Set bit of Data Memory	2 <sup>Note</sup>	None

Mnemonic	Description	Cycles	Flag Affected
<b>Branch</b>			
LSZ [m]	Skip if Data Memory is zero	2 <sup>Note</sup>	None
LSZA [m]	Skip if Data Memory is zero with data movement to ACC	2 <sup>Note</sup>	None
LSNZ [m]	Skip if Data Memory is not zero	2 <sup>Note</sup>	None
LSZ [m].i	Skip if bit i of Data Memory is zero	2 <sup>Note</sup>	None
LSNZ [m].i	Skip if bit i of Data Memory is not zero	2 <sup>Note</sup>	None
LSIZ [m]	Skip if increment Data Memory is zero	2 <sup>Note</sup>	None
LSDZ [m]	Skip if decrement Data Memory is zero	2 <sup>Note</sup>	None
LSIZA [m]	Skip if increment Data Memory is zero with result in ACC	2 <sup>Note</sup>	None
LSDZA [m]	Skip if decrement Data Memory is zero with result in ACC	2 <sup>Note</sup>	None
<b>Table Read</b>			
LTABRD [m]	Read table to TBLH and Data Memory	3 <sup>Note</sup>	None
LTABRDL [m]	Read table (last page) to TBLH and Data Memory	3 <sup>Note</sup>	None
LITABRD [m]	Increment table pointer TBLP first and Read table to TBLH and Data Memory	3 <sup>Note</sup>	None
LITABRDL [m]	Increment table pointer TBLP first and Read table (last page) to TBLH and Data Memory	3 <sup>Note</sup>	None
<b>Miscellaneous</b>			
LCLR [m]	Clear Data Memory	2 <sup>Note</sup>	None
LSET [m]	Set Data Memory	2 <sup>Note</sup>	None
LSWAP [m]	Swap nibbles of Data Memory	2 <sup>Note</sup>	None
LSWAPA [m]	Swap nibbles of Data Memory with result in ACC	2	None

Note: 1. For these extended skip instructions, if the result of the comparison involves a skip then up to four cycles are required, if no skip takes place two cycles is required.

- Any extended instruction which changes the contents of the PCL register will also require three cycles for execution.

## Instruction Definition

<b>ADC A,[m]</b>	Add Data Memory to ACC with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADCM A,[m]</b>	Add ACC to Data Memory with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADD A,[m]</b>	Add Data Memory to ACC
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADD A,x</b>	Add immediate data to ACC
Description	The contents of the Accumulator and the specified immediate data are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + x$
Affected flag(s)	OV, Z, AC, C, SC
<b>ADDM A,[m]</b>	Add ACC to Data Memory
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>AND A,[m]</b>	Logical AND Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
<b>AND A,x</b>	Logical AND immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bit wise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } x$
Affected flag(s)	Z
<b>ANDM A,[m]</b>	Logical AND ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z

<b>CALL addr</b>	Subroutine call
Description	Unconditionally calls a subroutine at the specified address. The Program Counter then increments by 1 to obtain the address of the next instruction which is then pushed onto the stack. The specified address is then loaded and the program continues execution from this new address. As this instruction requires an additional operation, it is a two cycle instruction.
Operation	Stack ← Program Counter + 1 Program Counter ← addr
Affected flag(s)	None
<b>CLR [m]</b>	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	[m] ← 00H
Affected flag(s)	None
<b>CLR [m].i</b>	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	[m].i ← 0
Affected flag(s)	None
<b>CLR WDT</b>	Clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared.
Operation	WDT cleared TO ← 0 PDF ← 0
Affected flag(s)	TO, PDF
<b>CPL [m]</b>	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	[m] ← $\overline{[m]}$
Affected flag(s)	Z
<b>CPLA [m]</b>	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC ← $\overline{[m]}$
Affected flag(s)	Z
<b>DAA [m]</b>	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	[m] ← ACC + 00H or [m] ← ACC + 06H or [m] ← ACC + 60H or [m] ← ACC + 66H
Affected flag(s)	C

<b>DEC [m]</b>	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	$[m] \leftarrow [m] - 1$
Affected flag(s)	Z
<b>DECA [m]</b>	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Z
<b>HALT</b>	Enter power down mode
Description	This instruction stops the program execution and turns off the system clock. The contents of the Data Memory and registers are retained. The WDT and prescaler are cleared. The power down flag PDF is set and the WDT time-out flag TO is cleared.
Operation	$TO \leftarrow 0$ $PDF \leftarrow 1$
Affected flag(s)	TO, PDF
<b>INC [m]</b>	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	$[m] \leftarrow [m] + 1$
Affected flag(s)	Z
<b>INCA [m]</b>	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] + 1$
Affected flag(s)	Z
<b>JMP addr</b>	Jump unconditionally
Description	The contents of the Program Counter are replaced with the specified address. Program execution then continues from this new address. As this requires the insertion of a dummy instruction while the new address is loaded, it is a two cycle instruction.
Operation	Program Counter $\leftarrow$ addr
Affected flag(s)	None
<b>MOV A,[m]</b>	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	$ACC \leftarrow [m]$
Affected flag(s)	None
<b>MOV A,x</b>	Move immediate data to ACC
Description	The immediate data specified is loaded into the Accumulator.
Operation	$ACC \leftarrow x$
Affected flag(s)	None
<b>MOV [m],A</b>	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	$[m] \leftarrow ACC$
Affected flag(s)	None

<b>NOP</b>	No operation
Description	No operation is performed. Execution continues with the next instruction.
Operation	No operation
Affected flag(s)	None
<b>OR A,[m]</b>	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "OR" [m]
Affected flag(s)	Z
<b>OR A,x</b>	Logical OR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "OR" x
Affected flag(s)	Z
<b>ORM A,[m]</b>	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "OR" [m]
Affected flag(s)	Z
<b>RET</b>	Return from subroutine
Description	The Program Counter is restored from the stack. Program execution continues at the restored address.
Operation	Program Counter ← Stack
Affected flag(s)	None
<b>RET A,x</b>	Return from subroutine and load immediate data to ACC
Description	The Program Counter is restored from the stack and the Accumulator loaded with the specified immediate data. Program execution continues at the restored address.
Operation	Program Counter ← Stack ACC ← x
Affected flag(s)	None
<b>RETI</b>	Return from interrupt
Description	The Program Counter is restored from the stack and the interrupts are re-enabled by setting the EMI bit. EMI is the master interrupt global enable bit. If an interrupt was pending when the RETI instruction is executed, the pending Interrupt routine will be processed before returning to the main program.
Operation	Program Counter ← Stack EMI ← 1
Affected flag(s)	None
<b>RL [m]</b>	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	[m].(i+1) ← [m].i; (i=0~6) [m].0 ← [m].7
Affected flag(s)	None

<b>RLA [m]</b>	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim6)$ $ACC.0 \leftarrow [m].7$
Affected flag(s)	None
<b>RLC [m]</b>	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	$[m].(i+1) \leftarrow [m].i; (i=0\sim6)$ $[m].0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>RLCA [m]</b>	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim6)$ $ACC.0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>RR [m]</b>	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim6)$ $[m].7 \leftarrow [m].0$
Affected flag(s)	None
<b>RRA [m]</b>	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim6)$ $ACC.7 \leftarrow [m].0$
Affected flag(s)	None
<b>RRC [m]</b>	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim6)$ $[m].7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C

<b>RRCA [m]</b>	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.i ← [m].(i+1); (i=0~6) ACC.7 ← C C ← [m].0
Affected flag(s)	C
<b>SBC A,[m]</b>	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	ACC ← ACC - [m] - $\bar{C}$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SBC A, x</b>	Subtract immediate data from ACC with Carry
Description	The immediate data and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	ACC ← ACC - [m] - $\bar{C}$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SBCM A,[m]</b>	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	[m] ← ACC - [m] - $\bar{C}$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SDZ [m]</b>	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	[m] ← [m] - 1 Skip if [m]=0
Affected flag(s)	None
<b>SDZA [m]</b>	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	ACC ← [m] - 1 Skip if ACC=0
Affected flag(s)	None

<b>SET [m]</b>	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
<b>SET [m].i</b>	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None
<b>SIZ [m]</b>	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>SIZA [m]</b>	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>SNZ [m].i</b>	Skip if Data Memory is not 0
Description	If the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None
<b>SNZ [m]</b>	Skip if Data Memory is not 0
Description	If the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m] \neq 0$
Affected flag(s)	None
<b>SUB A,[m]</b>	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>SUBM A,[m]</b>	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SUB A,x</b>	Subtract immediate data from ACC
Description	The immediate data specified by the code is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - x$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>SWAP [m]</b>	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	$[m].3 \sim [m].0 \leftrightarrow [m].7 \sim [m].4$
Affected flag(s)	None
<b>SWAPA [m]</b>	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC.3 \sim ACC.0 \leftarrow [m].7 \sim [m].4$ $ACC.7 \sim ACC.4 \leftarrow [m].3 \sim [m].0$
Affected flag(s)	None
<b>SZ [m]</b>	Skip if Data Memory is 0
Description	If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if $[m]=0$
Affected flag(s)	None
<b>SZA [m]</b>	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m]$ Skip if $[m]=0$
Affected flag(s)	None
<b>SZ [m].i</b>	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if $[m].i=0$
Affected flag(s)	None

<b>TABRD [m]</b>	Read table (current page) to TBLH and Data Memory
Description	The low byte of the program code (current page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>TABRDL [m]</b>	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>ITABRD [m]</b>	Increment table pointer low byte first and read table to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the program code addressed by the table pointer (TBLP and TBLH) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>ITABRDL [m]</b>	Increment table pointer low byte first and read table (last page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>XOR A,[m]</b>	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" [m]
Affected flag(s)	Z
<b>XORM A,[m]</b>	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "XOR" [m]
Affected flag(s)	Z
<b>XOR A,x</b>	Logical XOR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" x
Affected flag(s)	Z

## Extended Instruction Definition

The extended instructions are used to directly access the data stored in any data memory sections.

<b>LADC A,[m]</b>	Add Data Memory to ACC with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>LADCM A,[m]</b>	Add ACC to Data Memory with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C, SC
<b>LADD A,[m]</b>	Add Data Memory to ACC
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>LADDM A,[m]</b>	Add ACC to Data Memory
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C, SC
<b>LAND A,[m]</b>	Logical AND Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
<b>LANDM A,[m]</b>	Logical AND ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
<b>LCLR [m]</b>	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	$[m] \leftarrow 00H$
Affected flag(s)	None
<b>LCLR [m].i</b>	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	$[m].i \leftarrow 0$
Affected flag(s)	None

<b>LCPL [m]</b>	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	$[m] \leftarrow \overline{[m]}$
Affected flag(s)	Z
<b>LCPLA [m]</b>	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow \overline{[m]}$
Affected flag(s)	Z
<b>LDAA [m]</b>	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	$[m] \leftarrow ACC + 00H$ or $[m] \leftarrow ACC + 06H$ or $[m] \leftarrow ACC + 60H$ or $[m] \leftarrow ACC + 66H$
Affected flag(s)	C
<b>LDEC [m]</b>	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	$[m] \leftarrow [m] - 1$
Affected flag(s)	Z
<b>LDECA [m]</b>	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Z
<b>LINC [m]</b>	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	$[m] \leftarrow [m] + 1$
Affected flag(s)	Z
<b>LINCA [m]</b>	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] + 1$
Affected flag(s)	Z

<b>LMOV A,[m]</b>	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	$ACC \leftarrow [m]$
Affected flag(s)	None
<b>LMOV [m],A</b>	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	$[m] \leftarrow ACC$
Affected flag(s)	None
<b>LOR A,[m]</b>	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "OR" } [m]$
Affected flag(s)	Z
<b>LORM A,[m]</b>	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "OR" } [m]$
Affected flag(s)	Z
<b>LRL [m]</b>	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	$[m].(i+1) \leftarrow [m].i; (i=0\sim 6)$ $[m].0 \leftarrow [m].7$
Affected flag(s)	None
<b>LRLA [m]</b>	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim 6)$ $ACC.0 \leftarrow [m].7$
Affected flag(s)	None
<b>LRLC [m]</b>	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	$[m].(i+1) \leftarrow [m].i; (i=0\sim 6)$ $[m].0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C
<b>LRLCA [m]</b>	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.(i+1) \leftarrow [m].i; (i=0\sim 6)$ $ACC.0 \leftarrow C$ $C \leftarrow [m].7$
Affected flag(s)	C

<b>LRR [m]</b>	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim 6)$ $[m].7 \leftarrow [m].0$
Affected flag(s)	None
<b>LRRRA [m]</b>	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim 6)$ $ACC.7 \leftarrow [m].0$
Affected flag(s)	None
<b>LRRRC [m]</b>	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim 6)$ $[m].7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>LRRCA [m]</b>	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim 6)$ $ACC.7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>LSBC A,[m]</b>	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>LSBCM A,[m]</b>	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m] - C$
Affected flag(s)	OV, Z, AC, C, SC, CZ

<b>LSDZ [m]</b>	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] - 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>LSDZA [m]</b>	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] - 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>LSET [m]</b>	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
<b>LSET [m].i</b>	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None
<b>LSIZ [m]</b>	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if $[m]=0$
Affected flag(s)	None
<b>LSIZA [m]</b>	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if $ACC=0$
Affected flag(s)	None
<b>LSNZ [m].i</b>	Skip if Data Memory is not 0
Description	If the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None

<b>LSNZ [m]</b>	Skip if Data Memory is not 0
Description	If the content of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if [m] ≠ 0
Affected flag(s)	None
<b>LSUB A,[m]</b>	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	ACC ← ACC – [m]
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>LSUBM A,[m]</b>	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	[m] ← ACC – [m]
Affected flag(s)	OV, Z, AC, C, SC, CZ
<b>LSWAP [m]</b>	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	[m].3~[m].0 ↔ [m].7~[m].4
Affected flag(s)	None
<b>LSWAPA [m]</b>	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	ACC.3~ACC.0 ← [m].7~[m].4 ACC.7~ACC.4 ← [m].3~[m].0
Affected flag(s)	None
<b>LSZ [m]</b>	Skip if Data Memory is 0
Description	If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if [m]=0
Affected flag(s)	None
<b>LSZA [m]</b>	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	ACC ← [m] Skip if [m]=0
Affected flag(s)	None

<b>LSZ [m].i</b>	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if [m].i=0
Affected flag(s)	None
<b>LTABRD [m]</b>	Read table (current page) to TBLH and Data Memory
Description	The low byte of the program code (current page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LTABRDL [m]</b>	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LITABRD [m]</b>	Increment table pointer low byte first and read table to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the program code addressed by the table pointer (TBLP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LITABRDL [m]</b>	Increment table pointer low byte first and read table (last page) to TBLH and Data Memory
Description	Increment table pointer low byte, TBLP, first and then the low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>LXOR A,[m]</b>	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" [m]
Affected flag(s)	Z
<b>LXORM A,[m]</b>	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "XOR" [m]
Affected flag(s)	Z

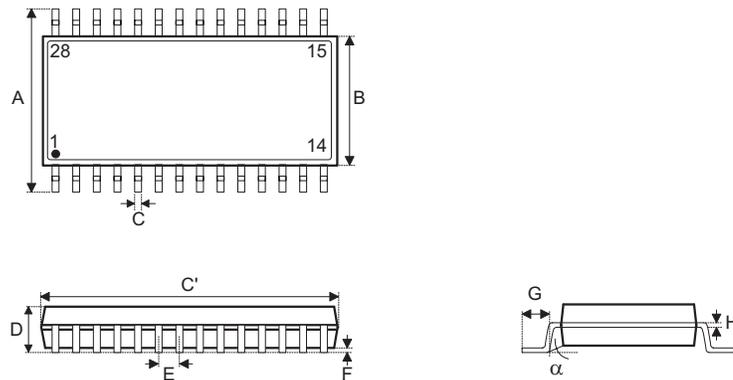
## Package Information

Note that the package information provided here is for consultation purposes only. As this information may be updated at regular intervals users are reminded to consult the [Holtek website](#) for the latest version of the [Package/Carton Information](#).

Additional supplementary information with regard to packaging is listed below. Click on the relevant section to be transferred to the relevant website page.

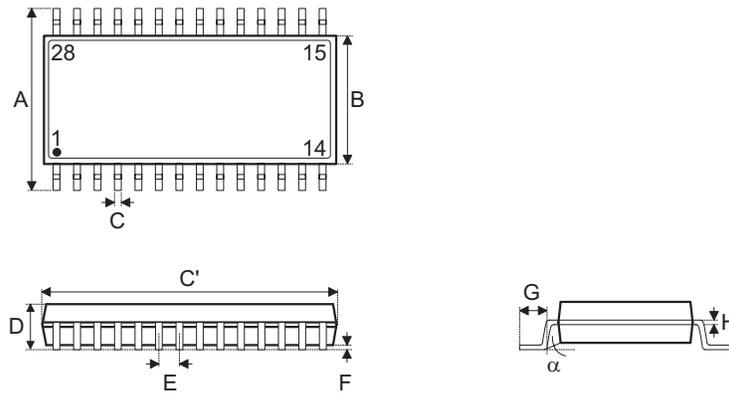
- Package Information (include Outline Dimensions, Product Tape and Reel Specifications)
- The Operation Instruction of Packing Materials
- Carton information

28-pin SOP (300mil) Outline Dimensions



Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	—	0.406 BSC	—
B	—	0.295 BSC	—
C	0.012	—	0.020
C'	—	0.705 BSC	—
D	—	—	0.104
E	—	0.050 BSC	—
F	0.004	—	0.012
G	0.016	—	0.050
H	0.008	—	0.013
α	0°	—	8°

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	—	10.30 BSC	—
B	—	7.50 BSC	—
C	0.31	—	0.51
C'	—	17.9 BSC	—
D	—	—	2.65
E	—	1.27 BSC	—
F	0.10	—	0.30
G	0.40	—	1.27
H	0.20	—	0.33
α	0°	—	8°

**28-pin SSOP (150mil) Outline Dimensions**


Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	—	0.236 BSC	—
B	—	0.154 BSC	—
C	0.008	—	0.012
C'	—	0.390 BSC	—
D	—	—	0.069
E	—	0.025 BSC	—
F	0.004	—	0.010
G	0.016	—	0.050
H	0.004	—	0.010
$\alpha$	0°	—	8°

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	—	6.0 BSC	—
B	—	3.9 BSC	—
C	0.20	—	0.30
C'	—	9.9 BSC	—
D	—	—	1.75
E	—	0.635 BSC	—
F	0.10	—	0.25
G	0.41	—	1.27
H	0.10	—	0.25
$\alpha$	0°	—	8°

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