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**Induction Cooker Flash MCU**

**HT45F0057**

Revision: V1.30 Date: May 17, 2022

[www.holtek.com](http://www.holtek.com)

## Features

### CPU Features

- Operating Voltage
  - ♦  $f_{SYS} = 8\text{MHz}$ : 2.2V~5.5V
- Up to 0.5 $\mu\text{s}$  instruction cycle with 8MHz system clock at  $V_{DD}=5\text{V}$
- Power down and wake-up functions to reduce power consumption
- Two Oscillators
  - ♦ Internal High Speed RC – HIRC
  - ♦ Internal Low Speed 32kHz RC – LIRC
- Fully integrated internal 8MHz oscillator requires no external components
- All instructions executed in one or two instruction cycles
- Table read instructions
- 63 powerful instructions
- 6-level subroutine nesting
- Bit manipulation instruction

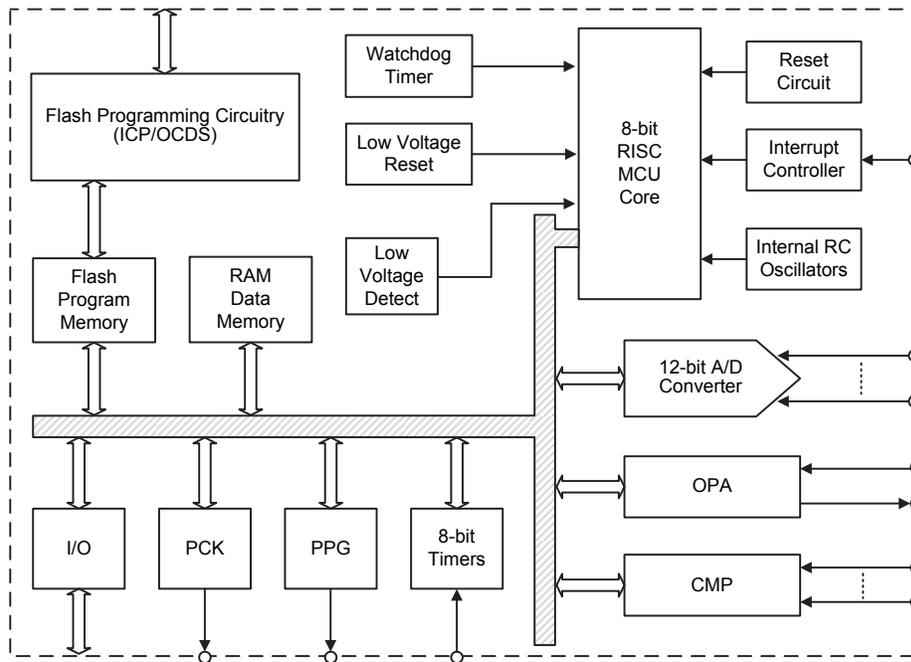
### Peripheral Features

- Program Memory: 4K $\times$ 16
- Data Memory: 208 $\times$ 8
- Watchdog Timer function
- Up to 13 bidirectional I/O lines
- 9-channel 12-bit resolution A/D converter
- 9-bit programmable pulse generator
  - ♦ Pulse width limit
  - ♦ Two PPG preload registers
  - ♦ Non-retriggered control
  - ♦ Active high or low output is selected by configuration option
- Single pin-shared external interrupt
  - ♦ Integrated de-bounce circuit for the external interrupt
  - ♦ De-bounce time is selected by software options
- Four comparator functions
- Single Operational Amplifier – input voltage offset adjustable by software
- Peripheral clock output
- Three 8-bit programmable timer/event counters
  - ♦ Timer 0 can be configured to count synchronism pulse number or measure synchronism pulse high or low period
  - ♦ Timer 1 can be configured to implement PPG non-retriggered function
- Low Voltage Reset function
- Low Voltage Detect function
- Package types: 16-pin NSOP

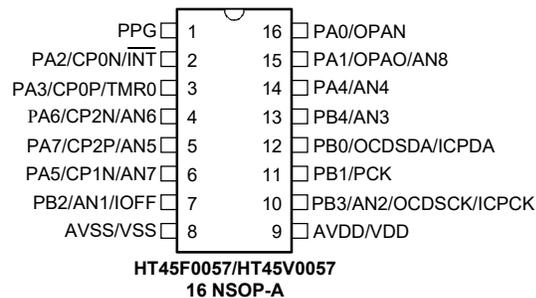
## General Description

The device is a Flash Memory A/D type 8-bit high performance RISC architecture microcontroller specifically designed for induction cooker applications. The device includes a wide range of functions and features such as low power consumption, I/O flexibility, timer functions, multi-channel A/D converter, power down and wake-up functions that enhance the versatility of this device making it especially suitable for induction cooker applications but which may have other application possibilities which require similar functions. The device also provides four comparators, an operational amplifier and a programmable pulse generator.

## Block Diagram



## Pin Assignment



Note: The OCSDA and OCDSCK pins are supplied for the OCDS dedicated pins and as such only available for the HT45V0057 device which is the OCDS EV chip for the HT45F0057 device.

## Pin Description

Pin Name	Function	OPT	I/T	O/T	Description
PA0/OPAN	PA0	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	OPAN	CO	AN	—	OPA inverting input.
PA1/OPAO/AN8	PA1	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	OPAO	CO	—	AN	OPA output pin.
	AN8	ADCR PCRH	AN	—	A/D Converter analog input.
PA2/CP0N/ $\overline{\text{INT}}$	PA2	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	CP0N	CO	AN	—	Comparator 0 inverting input pin.
	$\overline{\text{INT}}$	INTC0	ST	—	External interrupt input.
PA3/CP0P/TMR0	PA3	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	CP0P	CO CMPSC	AN	—	Comparator 0 non-inverting input pin, Comparator 1 inverting input pin or Comparator 3 inverting input pin.
	TMR0	—	ST	—	External Timer 0 clock input.
PA4/AN4	PA4	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	AN4	ADCR PCRL	AN	—	A/D Converter analog input.
PA5/CP1N/AN7	PA5	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	CP1N	CO CMPSC	AN	—	Comparator 1 inverting input pin.
	AN7	ADCR PCRL	AN	—	A/D Converter analog input.
PA6/CP2N/AN6	PA6	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	CP2N	CO	AN	—	Comparator 2 inverting input pin.
	AN6	ADCR PCRL	AN	—	A/D Converter analog input.
PA7/CP2P/AN5	PA7	CO	ST	CMOS	General purpose I/O. Pull-high and wake-up functions are selected by configuration options.
	CP2P	CO	AN	—	Comparator 2 non-inverting input pin.
	AN5	ADCR PCRL	AN	—	A/D Converter analog input.
PB0/OCSDA/ICPDA	PB0	CO	ST	CMOS	General purpose I/O. Pull-high function is selected by configuration options.
	OCSDA	—	ST	CMOS	OCDS Address/Data, for EV chip only.
	ICPDA	—	ST	CMOS	ICP Data/Address.
PB1/PCK	PB1	CO	ST	CMOS	General purpose I/O. Pull-high function is selected by configuration options.
	PCK	CO	—	CMOS	PCK output.
PB2/AN1/IOFF	PB2	CO	ST	CMOS	General purpose I/O. Pull-high function is selected by configuration options.
	IOFF	CO	—	CMOS	IOFF output.
	AN1	ADCR PCRL	AN	—	A/D Converter analog input.



## D.C. Electrical Characteristics

Operating Temperature: Ta = 25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>DD</sub>	Operating Voltage (HIRC)	—	f <sub>SYS</sub> = f <sub>HIRC</sub> = 8MHz	2.2	—	5.5	V
I <sub>DD</sub>	Operating Current (HIRC)	5V	No load, all peripherals off, f <sub>SYS</sub> = f <sub>HIRC</sub> = 8MHz	—	4.0	6.0	mA
I <sub>STB</sub>	Standby Current (SLEEP mode)	3V	No load, all peripherals off, WDT off	—	—	1	μA
		5V		—	—	2	μA
		3V	No load, all peripherals off, WDT on	—	—	5	μA
		5V		—	—	10	μA
V <sub>IL</sub>	Input Low Voltage for I/O Ports, TMR0 and INT	5V	—	0	—	1.5	V
		—	—	0	—	0.2V <sub>DD</sub>	V
V <sub>IH</sub>	Input High Voltage for I/O Ports, TMR0 and INT	5V	—	3.5	—	5	V
		—	—	0.8V <sub>DD</sub>	—	V <sub>DD</sub>	V
I <sub>OL</sub>	Sink Current for PPG, I/O Ports	3V	V <sub>OL</sub> = 0.1V <sub>DD</sub>	4	8	—	mA
		5V		10	20	—	mA
I <sub>OH</sub>	Source Current for PPG, I/O Ports	3V	V <sub>OH</sub> = 0.9V <sub>DD</sub>	-2	-4	—	mA
		5V		-5	-10	—	mA
R <sub>PH</sub>	Pull-high Resistance	3V	—	20	60	100	kΩ
		5V	—	10	30	50	kΩ

## A.C. Electrical Characteristics

Operating Temperature: Ta = 25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
f <sub>SYS</sub>	System Clock (HIRC)	2.2V ~ 5.5V	f <sub>SYS</sub> = f <sub>HIRC</sub> = 8MHz	—	8	—	MHz
f <sub>HIRC</sub>	High Speed Internal RC Oscillator (HIRC)	5V	—	- 2%	8	+ 2%	MHz
		5V ± 0.5V	Ta = 0°C ~ 55°C	- 4%	8	+ 4%	MHz
t <sub>WDTOSC</sub>	Watchdog Oscillator Period	3V	—	45	90	180	μs
		5V	—	32	65	130	μs
t <sub>SST</sub>	System Start-up Timer Period	—	Power on or wake-up from power-down mode	—	1024	—	t <sub>SYS</sub>
t <sub>INT</sub>	External Interrupt Minimum Pulse Width	—	—	1	—	—	μs
t <sub>TC</sub>	TMR0 Input Pin Minimum Pulse Width	—	—	0.1	—	—	μs
t <sub>RSTD</sub>	System Reset Delay Time	—	—	50	100	200	ms

Note: 1. t<sub>SYS</sub> = 1/f<sub>SYS</sub>

2. Minimum pulse width is for guaranteed interrupt. It is possible for a smaller pulse width to be recognized.

## LVD & LVR Electrical Characteristics

Operating Temperature: Ta = 25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>LVR</sub>	Low Voltage Reset Voltage	—	LVR enable, voltage select 2.1V	1.98	2.1	2.22	V
V <sub>LVD</sub>	Low Voltage Detection Voltage	—	LVD enable	4.12	4.4	4.7	V
t <sub>LVR</sub>	Minimum Low Voltage Width to Reset	—	—	0.25	1	2	ms

## A/D Converter Characteristics

Operating Temperature: Ta = 25°C, unless otherwise specify

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>DD</sub>	Operating Voltage	—	—	2.7	—	5.5	V
V <sub>ADI</sub>	Input Voltage	—	—	0	—	V <sub>REF</sub>	V
V <sub>REF</sub>	Reference Voltage	—	—	—	V <sub>DD</sub>	—	V
DNL	Differential Nonlinearity	3V	V <sub>REF</sub> = V <sub>DD</sub> , t <sub>ADCK</sub> = 1.0μs	—	—	±3	LSB
		5V					
		3V	V <sub>REF</sub> = V <sub>DD</sub> , t <sub>ADCK</sub> = 4μs				
		5V					
INL	Integral Nonlinearity	3V	V <sub>REF</sub> = V <sub>DD</sub> , t <sub>ADCK</sub> = 1.0μs	—	—	±4	LSB
		5V					
		3V	V <sub>REF</sub> = V <sub>DD</sub> , t <sub>ADCK</sub> = 4μs				
		5V					
I <sub>ADC</sub>	Additional Current for A/D Converter Enable	3V	No load (t <sub>ADCK</sub> = 1.0μs)	—	1	2	mA
		5V		—	1.5	3	mA
t <sub>ADCK</sub>	Clock Period	—	—	0.5	—	10	μs
t <sub>ADS</sub>	Sampling Time	—	—	—	4	—	t <sub>ADCK</sub>
t <sub>ADC</sub>	Conversion Time (Include A/D Converter Sample And Hold Time)	—	—	—	16	—	t <sub>ADCK</sub>

## Comparator Characteristics

Operating Temperature: Ta = 25°C, unless otherwise specify

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>CP</sub>	Operating Voltage	—	—	3.0	—	5.5	V
I <sub>CMP</sub>	Additional Current for Comparator Enable	3V	—	—	150	225	μA
		5V	—	—	200	300	μA
V <sub>OS</sub>	Input Offset Voltage	5V	Without calibration (CnCOF[4:0]=10000B)	-15	—	15	mV
		5V	With calibration	-2	—	2	mV
V <sub>CM</sub>	Common Mode Voltage Range	—	—	0	—	V <sub>DD</sub> - 1.4	V
V <sub>HYS</sub>	Hysteresis	5V	—	20	40	60	mV
t <sub>C0IRP</sub>	Comparator 0 Interrupt Response Time	5V	Disable Hysteresis, de-bounce disable	—	0.5	1	μs
		5V	Disable Hysteresis, de-bounce enable	—	1	2	μs
t <sub>C1IRP</sub>	Comparator 1 Interrupt Response Time	5V	Disable Hysteresis, de-bounce disable	—	0.5	1	μs
		5V	Disable Hysteresis, de-bounce enable	—	1	2	μs
t <sub>C2IRP</sub>	Comparator 2 Interrupt Response Time	5V	Disable Hysteresis, de-bounce disable	—	0.5	1	μs
		5V	Disable Hysteresis, de-bounce enable	—	1	2	μs
t <sub>C3IRP</sub>	Comparator 3 Interrupt Response Time	5V	Disable Hysteresis, de-bounce disable	—	0.5	1	μs
		5V	Disable Hysteresis, de-bounce enable	—	1	2	μs

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>R1</sub>	Reference Voltage for Comparator 1	5V	Ta=-40°C ~ 85°C	- 5%	0.6V <sub>DD</sub> ~0.775V <sub>DD</sub> , 0.025V <sub>DD</sub> /step	+ 5%	V <sub>DD</sub>
V <sub>R2</sub>	Reference Voltage for Comparator 2	5V	Ta=-40°C ~ 85°C	- 5%	0.6V <sub>DD</sub> ~0.775V <sub>DD</sub> , 0.025V <sub>DD</sub> /step	+ 5%	V <sub>DD</sub>
V <sub>R3</sub>	Reference Voltage for Comparator 2	5V	Ta=-40°C ~ 85°C	- 5%	0.075V <sub>DD</sub> ~0.25V <sub>DD</sub> , 0.025 V <sub>DD</sub> /step	+ 5%	V <sub>DD</sub>
V <sub>R4</sub>	Reference Voltage for Comparator 3	5V	Ta=-40°C ~ 85°C	- 5%	0.6V <sub>DD</sub> ~0.775V <sub>DD</sub> , 0.025V <sub>DD</sub> /step	+ 5%	V <sub>DD</sub>

## Operational Amplifier Characteristics

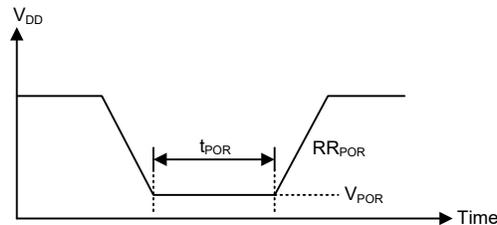
Operating Temperature: Ta = 25°C, unless otherwise specify

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V <sub>DD</sub>	Conditions				
V <sub>OP</sub>	Operating Voltage	—	—	3.0	—	5.5	V
I <sub>OPA</sub>	Additional Current for OPA Enable	3V	no load	—	150	300	μA
		5V		—	300	450	μA
V <sub>OS</sub>	Input Offset Voltage	5V	Without calibration (OPAOF[4:0] = 10000B)	-15	—	15	mV
		5V	With calibration	-2	—	2	mV
V <sub>CM</sub>	Common Mode Voltage Range	—	—	0	—	V <sub>DD</sub> - 1.4	V
V <sub>OR</sub>	Maximum Output Voltage Range	5V	—	V <sub>SS</sub> + 0.1	—	V <sub>DD</sub> - 0.1	V
SR	Slew Rate	5V	No load	0.6	1.8	—	V/μs
GBW	Gain Bandwidth	5V	R <sub>LOAD</sub> = 1MΩ, C <sub>LOAD</sub> = 100pF	600	2200	—	kHz
PSRR	Power Supply Rejection Ratio	5V	—	60	80	—	dB
CMRR	Common Mode Rejection Ratio	5V	—	60	80	—	dB
R <sub>OPAR</sub>	Typical Value of OPAR1	5V	Ta = 25°C	45	60	75	kΩ
	Temperature Drift Percentage Relative to 25°C	5V	Ta = 0°C ~ 50°C	-3	—	+3	%
	Typical Value of OPAR2	5V	Ta = 25°C	0.75	1	1.25	kΩ
	Temperature Drift Percentage Relative to 25°C	5V	Ta = 0°C ~ 50°C	-3	—	+3	%
	Typical Value of OPAR3	5V	Ta = 25°C	0.75	1	1.25	kΩ
	Temperature Drift Percentage Relative to 25°C	5V	Ta = 0°C ~ 50°C	-3	—	+3	%
	Typical Value of OPAR4	5V	Ta = 25°C	—	10	—	kΩ

## Power-on Reset Characteristics

Operating Temperature:  $T_a = 25^\circ\text{C}$ , unless otherwise specify

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		$V_{DD}$	Conditions				
$V_{POR}$	$V_{DD}$ Start Voltage to Ensure Power-on Reset	—	—	—	—	100	mV
$RR_{POR}$	$V_{DD}$ Raising Rate to Ensure Power-on Reset	—	—	0.035	—	—	V/ms
$t_{POR}$	Minimum Time for $V_{DD}$ Stays at $V_{POR}$ to Ensure Power-on Reset	—	—	1	—	—	ms

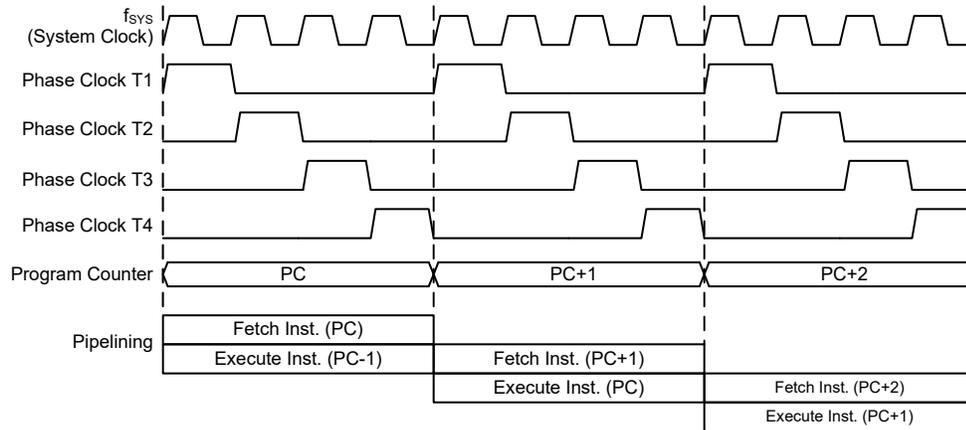


## System Architecture

A key factor in the high-performance features of the Holtek range of microcontrollers is attributed to their internal system architecture. The device takes advantage of the usual features found within RISC microcontrollers providing increased speed of operation and enhanced performance. The pipelining scheme is implemented in such a way that instruction fetching and instruction execution are overlapped, hence instructions are effectively executed in one cycle, with the exception of branch or call instructions. An 8-bit wide ALU is used in practically all instruction set operations, which carries out arithmetic operations, logic operations, rotation, increment, decrement, branch decisions, etc. The internal data path is simplified by moving data through the Accumulator and the ALU. Certain internal registers are implemented in the Data Memory and can be directly or indirectly addressed. The simple addressing methods of these registers along with additional architectural features ensure that a minimum of external components is required to provide a functional I/O and A/D control system with maximum reliability and flexibility. This makes the device suitable for low-cost, high-volume production for controller applications.

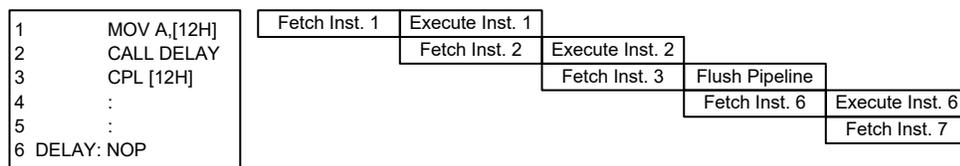
### Clocking and Pipelining

The main system clock, derived from either a HIRC or LIRC oscillator is subdivided into four internally generated non-overlapping clocks, T1~T4. The Program Counter is incremented at the beginning of the T1 clock during which time a new instruction is fetched. The remaining T2~T4 clocks carry out the decoding and execution functions. In this way, one T1~T4 clock cycle forms one instruction cycle. Although the fetching and execution of instructions takes place in consecutive instruction cycles, the pipelining structure of the microcontroller ensures that instructions are effectively executed in one instruction cycle. The exception to this are instructions where the contents of the Program Counter are changed, such as subroutine calls or jumps, in which case the instruction will take one more instruction cycle to execute.



### System Clocking and Pipelining

For instructions involving branches, such as jump or call instructions, two machine cycles are required to complete instruction execution. An extra cycle is required as the program takes one cycle to first obtain the actual jump or call address and then another cycle to actually execute the branch. The requirement for this extra cycle should be taken into account by programmers in timing sensitive applications.



### Instruction Fetching

### Program Counter

During program execution, the Program Counter is used to keep track of the address of the next instruction to be executed. It is automatically incremented by one each time an instruction is executed except for instructions, such as “JMP” or “CALL” that demands a jump to a non-consecutive Program Memory address. Only the lower 8 bits, known as the Program Counter Low Register, are directly addressable by the application program.

When executing instructions requiring jumps to non-consecutive addresses such as a jump instruction, a subroutine call, interrupt or reset, etc., the microcontroller manages program control by loading the required address into the Program Counter. For conditional skip instructions, once the condition has been met, the next instruction, which has already been fetched during the present instruction execution, is discarded and a dummy cycle takes its place while the correct instruction is obtained.

Program Counter	
Program Counter High Byte	PCL Register
PC11~PC8	PCL7~PCL0

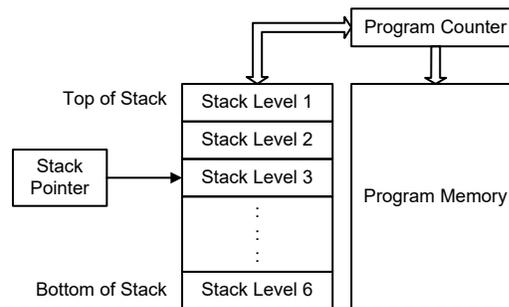
### Program Counter

The lower byte of the Program Counter, known as the Program Counter Low register or PCL, is available for program control and is a readable and writeable register. By transferring data directly into this register, a short program jump can be executed directly. However, as only this low byte is available for manipulation, the jumps are limited to the present page of memory that is 256 locations. When such program jumps are executed it should also be noted that a dummy cycle will be inserted. Manipulating the PCL register may cause program branching, so an extra cycle is needed to pre-fetch.

**Stack**

This is a special part of the memory which is used to save the contents of the Program Counter only. The stack is organized into 6 levels and is neither part of the data nor part of the program space, and is neither readable nor writeable. The activated level is indexed by the Stack Pointer, and is neither readable nor writeable. At a subroutine call or interrupt acknowledge signal, the contents of the Program Counter are pushed onto the stack. At the end of a subroutine or an interrupt routine, signaled by a return instruction, RET or RETI, the Program Counter is restored to its previous value from the stack. After a device reset, the Stack Pointer will point to the top of the stack.

If the stack is full and an enabled interrupt takes place, the interrupt request flag will be recorded but the acknowledge signal will be inhibited. When the Stack Pointer is decremented, by RET or RETI, the interrupt will be serviced. This feature prevents stack overflow allowing the programmer to use the structure more easily. However, when the stack is full, a CALL subroutine instruction can still be executed which will result in a stack overflow. Precautions should be taken to avoid such cases which might cause unpredictable program branching. If the stack is overflow, the first Program Counter save in the stack will be lost.



**Arithmetic and Logic Unit – ALU**

The arithmetic-logic unit or ALU is a critical area of the microcontroller that carries out arithmetic and logic operations of the instruction set. Connected to the main microcontroller data bus, the ALU receives related instruction codes and performs the required arithmetic or logical operations after which the result will be placed in the specified register. As these ALU calculation or operations may result in carry, borrow or other status changes, the status register will be correspondingly updated to reflect these changes. The ALU supports the following functions:

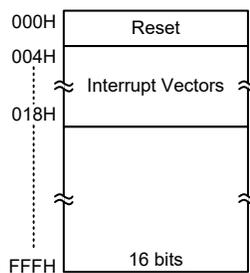
- Arithmetic operations: ADD, ADDM, ADC, ADCM, SUB, SUBM, SBC, SBCM, DAA
- Logic operations: AND, OR, XOR, ANDM, ORM, XORM, CPL, CPLA
- Rotation: RRA, RR, RRCA, RRC, RLA, RL, RLCA, RLC
- Increment and Decrement: INCA, INC, DECA, DEC
- Branch decision: JMP, SZ, SZA, SNZ, SIZ, SDZ, SIZA, SDZA, CALL, RET, RETI

## Flash Program Memory

The Program Memory is the location where the user code or program is stored. For this device the Program Memory is Flash type, which means it can be programmed and re-programmed a large number of times, allowing the user the convenience of code modification on the same device. By using the appropriate programming tools, the Flash device offer users the flexibility to conveniently debug and develop their applications while also offering a means of field programming and updating.

### Structure

The Program Memory has a capacity of 4K×16 bits. The Program Memory is addressed by the Program Counter and also contains data, table information and interrupt entries. Table data, which can be setup in any location within the Program Memory, is addressed by a separate table pointer register.



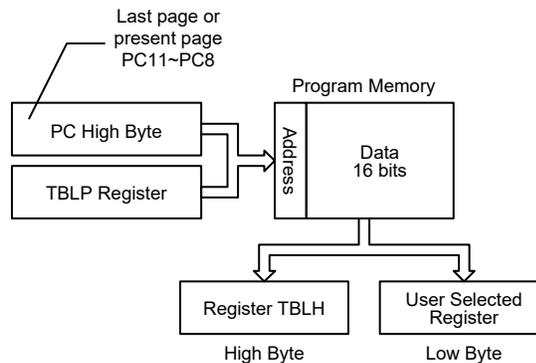
**Program Memory Structure**

### Look-up Table

Any location within the Program Memory can be defined as a look-up table where programmers can store fixed data. To use the look-up table, the table pointer must first be setup by placing the address of the look up data to be retrieved in the table pointer register, TBLP. These registers define the total address of the look-up table.

After setting up the table pointer, the table data can be retrieved from the Program Memory using the “TABRDC [m]” or “TABRDL [m]” instructions, respectively. When the instruction is executed, the lower order table byte from the Program Memory will be transferred to the user defined Data Memory register [m] as specified in the instruction. The higher order table data byte from the Program Memory will be transferred to the TBLH special register. Any unused bits in this transferred higher order byte will be read as 0.

The accompanying diagram illustrates the addressing data flow of the look-up table.



### Table Program Example

The following example shows how the table pointer and table data is defined and retrieved from the microcontroller. This example uses raw table data located in the Program Memory which is stored there using the ORG statement. The value at this ORG statement is “0F00H” which refers to the start address of the last page within the 4K Program Memory of the device. The table pointer is setup here to have an initial value of “06H”. This will ensure that the first data read from the data table will be at the Program Memory address “0F06H” or 6 locations after the start of the last page. Note that the value for the table pointer is referenced to the first address of the present page if the “TABRDC [m]” instruction is being used. The high byte of the table data which in this case is equal to zero will be transferred to the TBLH register automatically when the “TABRDC [m]” instruction is executed.

Because the TBLH register is a read-only register and cannot be restored, care should be taken to ensure its protection if both the main routine and Interrupt Service Routine use table read instructions. If using the table read instructions, the Interrupt Service Routines may change the value of the TBLH and subsequently cause errors if used again by the main routine. As a rule it is recommended that simultaneous use of the table read instructions should be avoided. However, in situations where simultaneous use cannot be avoided, the interrupts should be disabled prior to the execution of any main routine table-read instructions. Note that all table related instructions require two instruction cycles to complete their operation.

#### Table Read Program Example

```
tempreg1 db ?          ; temporary register #1
tempreg2 db ?          ; temporary register #2
:
mov a, 06h             ; initialise low table pointer - note that this address is
                       ; referenced
mov tblp, a           ; to the last page or present page
:
:
tabrdc tempreg1       ; transfers value in table referenced by table pointer data at
                       ; program memory address 0F06H transferred to tempreg1 and TBLH
dec tblp              ; reduce value of table pointer by one
tabrdc tempreg2       ; transfers value in table referenced by table pointer data at
                       ; program memory address 0F05H transferred to tempreg2 and TBLH
                       ; in this example the data 1AH is transferred to tempreg1 and data
                       ; 0FH to register tempreg2
:
org 0F00h             ; sets initial address of program memory
dc 00Ah, 00Bh, 00Ch, 00Dh, 00Eh, 00Fh, 01Ah, 01Bh
:
```

### In Circuit Programming – ICP

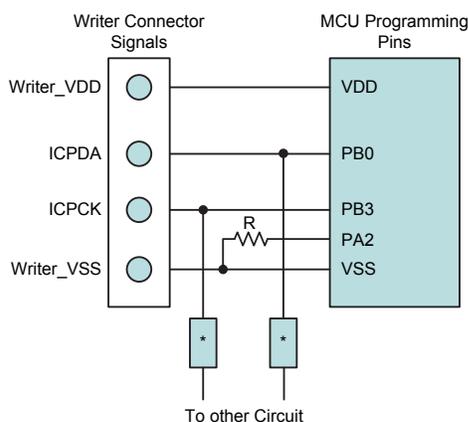
The provision of Flash type Program Memory provides the user with a means of convenient and easy upgrades and modifications to their programs on the same device. As an additional convenience, Holtek has provided a means of programming the microcontroller in-circuit using a 4-pin interface. This provides manufacturers with the possibility of manufacturing their circuit boards complete with a programmed or un-programmed microcontroller, and then programming or upgrading the program at a later stage. This enables product manufacturers to easily keep their manufactured products supplied with the latest program releases without removal and re-insertion of the device.

The Holtek Flash MCU to Writer Programming Pin correspondence table is as follows:

Holtek Writer Pins	MCU Programming Pins	Pin Description
ICPDA	PB0	Programming Serial Data/Address
ICPCK	PB3	Programming Clock
VDD	VDD	Power Supply
VSS	VSS	Ground

The Program Memory can be programmed serially in-circuit using this 4-wire interface. Data is downloaded and uploaded serially on a single pin with an additional line for the clock. Two additional lines are required for the power supply. The technical details regarding the in-circuit programming of the device are beyond the scope of this document and will be supplied in supplementary literature.

During the programming process, taking control of the ICPDA and ICPCK pins for data and clock programming purposes. The user must there take care to ensure that no other outputs are connected to these two pins.



- Note: 1.\* may be resistor or capacitor. The resistance of \* must be greater than 1kΩ or the capacitance of \* must be less than 1nF.
2. A Resistor with a value less than 1MΩ is recommended to be connected between PA2 and VSS pins during programming.

### On-Chip Debug Support – OCDS

There is an EV chip named HT45V0057 which is used to emulate the HT45F0057. The EV chip device provides an “On-Chip Debug” function to debug the corresponding MCU device during the development process. The EV chip and the actual MCU device are almost functionally compatible except for the “On-Chip Debug” function. Users can use the EV chip device to emulate the real chip device behavior by connecting the OCSDA and OCDSCK pins to the Holtek HT-IDE development tools. The OCSDA pin is the OCDS Data/Address input/output pin while the OCDSCK pin is the OCDS clock input pin. When using the EV device during debug, if there are other pin functions which are shared with the OCSDA and OCDSCK pins, then these functions will have no effect in the EV chip. The two OCDS pins, which are pin-shared with the ICP programming pins, are still used as Flash Memory programming pins for ICP. For a more detailed OCDS description, refer to the corresponding document named “Holtek e-Link for 8-bit MCU OCDS User’s Guide”.

Holtek e-Link Pins	EV Chip Pins	Pin Description
OCSDA	OCSDA	On-chip Debug Support Data/Address input/output
OCDSCK	OCDSCK	On-chip Debug Support Clock input
VDD	VDD	Power Supply
VSS	VSS	Ground

## Data Memory

The Data Memory is a volatile area of 8-bit wide RAM internal memory and is the location where temporary information is stored.

### Structure

Divided into two sections, the first of these is an area of RAM, known as the Special Function Data Memory. Here are located registers which are necessary for correct operation of the device. Many of these registers can be read from and written to directly under program control, however, some remain protected from user manipulation. The second area of Data Memory is known as the General Purpose Data Memory, which is reserved for general purpose use. All locations within this area are read and write accessible under program control.

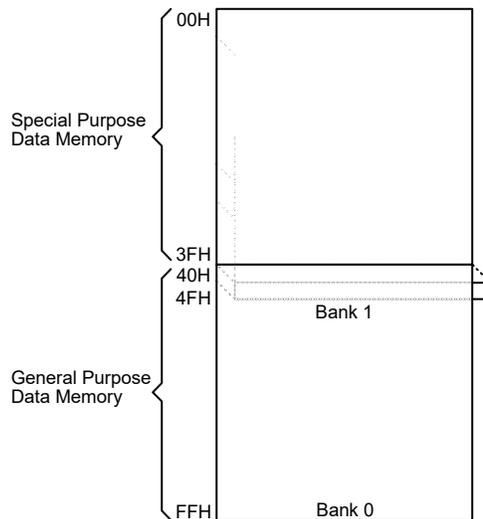
The overall Data Memory is subdivided into two banks. The Special Purpose Data Memory registers are accessible in all banks. Switching between the different Data Memory banks is achieved by setting the Bank Pointer to the correct value. The start address of the Data Memory is the address 00H.

### General Purpose Data Memory

All microcontroller programs require an area of read/write memory where temporary data can be stored and retrieved for use later. It is this area of RAM memory that is known as General Purpose Data Memory. This area of Data Memory is fully accessible by the user program for both reading and writing operations. By using the “SET [m].i” and “CLR [m].i” instructions individual bits can be set or reset under program control giving the user a large range of flexibility for bit manipulation in the Data Memory.

Special Purpose Data Memory		General Purpose Data Memory	
Available Banks	Banks	Capacity	Banks
0,1	Bank 0: 00H~3FH Bank 1: 00H~3FH	208 × 8	Bank 0: 40H~FFH Bank 1: 40H~4FH

**Data Memory Summary**



**Data Memory Structure**

### Special Purpose Data Memory

This area of Data Memory is where registers, necessary for the correct operation of the microcontroller, are stored. Most of the registers are both readable and writeable but some are protected and are readable only, the details of which are located under the relevant Special Function Register section. Note that for locations that are unused, any read instruction to these addresses will return the value “00H”.

Bank 0,1		Bank0;Bank1	
00H	IAR0	20H	PPGC
01H	MP0	21H	PPGTA
02H	IAR1	22H	PPGTB
03H	MP1	23H	PPGTEX
04H	BP	24H	ADRL
05H	ACC	25H	ADRH
06H	PCL	26H	ADCR
07H	TBLP	27H	ACSR
08H	TBLH	28H	CTRL1
09H		29H	OPAC
0AH	STATUS	2AH	CTRL2
0BH	INTC0	2BH	PWLT
0CH		2CH	CTRL3
0DH	TMR0	2DH	TMR2
0EH	TMR0C	2EH	TMR2C
0FH		2FH	CMPSC
10H	TMR1	30H	
11H	TMR1C	31H	
12H	PA	32H	
13H	PAC	33H	
14H	PB	34H	
15H	PBC	35H	
16H		36H	
17H		37H	
18H		38H	
19H	39H		
1AH	CTRL0	3AH	MFI
1BH	CMP0C	3BH	PCRL
1CH	CMP1C	3CH	PCRH
1DH	CMP2C	3DH	
1EH	INTC1	3EH	
1FH		3FH	

 : Unused, read as “00”

### Special Purpose Data Memory

## Special Function Register Description

Most of the Special Function Register details will be described in the relevant functional sections, however several registers require a separate description in this section.

### Indirect Addressing Register – IAR0, IAR1

The Indirect Addressing Registers, IAR0 and IAR1, although having their locations in normal RAM register space, do not actually physically exist as normal registers. The method of indirect addressing for RAM data manipulation uses these Indirect Addressing Registers and Memory Pointers, in contrast to direct memory addressing, where the actual memory address is specified. Actions on the IAR0 and IAR1 registers will result in no actual read or write operation to these registers but rather to the memory location specified by their corresponding Memory Pointers, MP0 or MP1. Acting as a pair, IAR0 and MP0 can together access data from Bank 0 while the IAR1 and MP1 register pair can access data from any bank. As the Indirect Addressing Registers are not physically implemented, reading the Indirect Addressing Registers indirectly will return a result of “00H” and writing to the registers indirectly will result in no operation.

### Memory Pointers – MP0, MP1

Two Memory Pointers, known as MP0 and MP1 are provided. These Memory Pointers are physically implemented in the Data Memory and can be manipulated in the same way as normal registers providing a convenient way with which to address and track data. When any operation to the relevant Indirect Addressing Registers is carried out, the actual address that the microcontroller is directed to, is the address specified by the related Memory Pointer. MP0, together with Indirect Addressing Register, IAR0, are used to access data from Bank 0, while MP1 and IAR1 are used to access data from all banks according to BP register. Direct Addressing can only be used with Bank 0, all other Banks must be addressed indirectly using MP1 and IAR1.

The following example shows how to clear a section of four Data Memory locations already defined as locations `adres1` to `adres4`.

### Indirect Addressing Program Example

```
data .section 'data'
adres1 db ?
adres2 db ?
adres3 db ?
adres4 db ?
block db ?
code .section at 0 code
org 00h
start:
mov a, 04h ; setup size of block
mov block, a
mov a, offset adres1 ; Accumulator loaded with first RAM address
mov mp0, a ; setup memory pointer with first RAM address
loop:
clr IAR0 ; clear the data at address defined by MP0
inc mp0 ; increment memory pointer
sdz block ; check if last memory location has been cleared
jmp loop
continue:
```

The important point to note here is that in the example shown above, no reference is made to specific Data Memory addresses.

### Bank Pointer – BP

For this device, the Data Memory is divided into two banks, Bank0 and Bank1. Selecting the required Data Memory area is achieved using the Bank Pointer. Bit 0 of the Bank Pointer is used to select Data Memory Banks 0~1.

The Data Memory is initialised to Bank 0 after a reset, except for a WDT time-out reset in the Power Down Mode, in which case, the Data Memory bank remains unaffected. It should be noted that the Special Function Data Memory is not affected by the bank selection, which means that the Special Function Registers can be accessed from within any bank. Directly addressing the Data Memory will always result in Bank 0 being accessed irrespective of the value of the Bank Pointer. Accessing data from Bank1 must be implemented using Indirect Addressing.

- **BP Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	—	DMBP
R/W	—	—	—	—	—	—	—	R/W
POR	—	—	—	—	—	—	—	0

Bit 7~1 Unimplemented, read as “0”

Bit 0 **DMBP**: Select Data Memory Banks  
 0: Bank 0  
 1: Bank 1

### Accumulator – ACC

The Accumulator is central to the operation of any microcontroller and is closely related with operations carried out by the ALU. The Accumulator is the place where all intermediate results from the ALU are stored. Without the Accumulator it would be necessary to write the result of each calculation or logical operation such as addition, subtraction, shift, etc., to the Data Memory resulting in higher programming and timing overheads. Data transfer operations usually involve the temporary storage function of the Accumulator; for example, when transferring data between one user defined register and another, it is necessary to do this by passing the data through the Accumulator as no direct transfer between two registers is permitted.

### Program Counter Low Register – PCL

To provide additional program control functions, the low byte of the Program Counter is made accessible to programmers by locating it within the Special Purpose area of the Data Memory. By manipulating this register, direct jumps to other program locations are easily implemented. Loading a value directly into this PCL register will cause a jump to the specified Program Memory location, however, as the register is only 8-bit wide, only jumps within the current Program Memory page are permitted. When such operations are used, note that a dummy cycle will be inserted.

### Look-up Table Registers – TBLP, TBLH

These two special function registers are used to control operation of the look-up table which is stored in the Program Memory. TBLP is the table pointer and indicates the location where the table data is located. Its value must be setup before any table read commands are executed. Its value can be changed, for example using the “INC” or “DEC” instructions, allowing for easy table data pointing and reading. TBLH is the location where the high order byte of the table data is stored after a table read data instruction has been executed. Note that the lower order table data byte is transferred to a user defined location.

## Status Register – STATUS

This 8-bit register contains the zero flag (Z), carry flag (C), auxiliary carry flag (AC), overflow flag (OV), power down flag (PDF), and watchdog time-out flag (TO). These arithmetic/logical operation and system management flags are used to record the status and operation of the microcontroller.

With the exception of the TO and PDF flags, bits in the status register can be altered by instructions like most other registers. Any data written into the status register will not change the TO or PDF flag. In addition, operations related to the status register may give different results due to the different instruction operations. The TO flag can be affected only by a system power-up, a WDT time-out or by executing the “CLR WDT” or “HALT” instruction. The PDF flag is affected only by executing the “HALT” or “CLR WDT” instruction or during a system power-up.

The Z, OV, AC and C flags generally reflect the status of the latest operations.

- C is set if an operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation; otherwise C is cleared. C is also affected by a rotate through carry instruction.
- AC is set if an operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction; otherwise AC is cleared.
- Z is set if the result of an arithmetic or logical operation is zero; otherwise Z is cleared.
- OV is set if an operation results in a carry into the highest-order bit but not a carry out of the highest-order bit, or vice versa; otherwise OV is cleared.
- PDF is cleared by a system power-up or executing the “CLR WDT” instruction. PDF is set by executing the “HALT” instruction.
- TO is cleared by a system power-up or executing the “CLR WDT” or “HALT” instruction. TO is set by a WDT time-out.

In addition, on entering an interrupt sequence or executing a subroutine call, the status register will not be pushed onto the stack automatically. If the contents of the status registers are important and if the subroutine can corrupt the status register, precautions must be taken to correctly save it.

Note: that bits 0~3 of the STATUS register are both readable and writeable bits.

• **STATUS Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	TO	PDF	OV	Z	AC	C
R/W	—	—	R	R	R/W	R/W	R/W	R/W
POR	—	—	0	0	x	x	x	x

“x”: unknown

- Bit 7~6 Unimplemented, read as “0”
- Bit 5 **TO**: Watchdog Time-Out flag  
 0: After power up or executing the “CLR WDT” or “HALT” instruction  
 1: A watchdog time-out occurred.
- Bit 4 **PDF**: Power down flag  
 0: After power up or executing the “CLR WDT” instruction  
 1: By executing the “HALT” instruction
- Bit 3 **OV**: Overflow flag  
 0: No overflow  
 1: An operation results in a carry into the highest-order bit but not a carry out of the highest-order bit or vice versa.
- Bit 2 **Z**: Zero flag  
 0: The result of an arithmetic or logical operation is not zero  
 1: The result of an arithmetic or logical operation is zero
- Bit 1 **AC**: Auxiliary flag  
 0: No auxiliary carry  
 1: An operation results in a carry out of the low nibbles in addition, or no borrow from the high nibble into the low nibble in subtraction
- Bit 0 **C**: Carry flag  
 0: No carry-out  
 1: An operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation  
 C is also affected by a rotate through carry instruction.

### System Control Registers – CTRL0, CTRL1, CTRL2, CTRL3, CTRL4

These registers are used to provide control over various internal functions. Some of these include the inverting control of the comparator output signal, TMR0 external clock source selection, peripheral clock prescaler stage, oscillator stop function control, internal reference voltage selection of the comparator, etc.

#### • CTRL0 Register

Bit	7	6	5	4	3	2	1	0
Name	C2VOINV	TMR0ECS	—	—	PCKPSC	—	LVDO	LVDC
R/W	R/W	R/W	—	—	R/W	—	R	R/W
POR	0	0	—	—	0	—	0	0

Bit 7      **C2VOINV**: Inverting control of the comparator 2 output signal

0: Non-inverted  
1: Inverted

Bit 6      **TMR0ECS**: Select TMR0 external clock source

0: TMR0 pin  
1: INT0

Bit 5~4    Unimplemented, read as “0”

Bit 3      **PCKPSC**: Peripheral clock prescaler stage

$f_{PCK} =$   
0:  $f_{SYS}/2048$   
1:  $f_{SYS}/4096$

Bit 2      Unimplemented, read as “0”

Bit 1      **LVDO**: Low voltage detector output

0: Normal voltage  
1: Low voltage detected

Bit 0      **LVDC**: Low voltage detector control

0: Disable  
1: Enable

#### • CTRL1 Register

Bit	7	6	5	4	3	2	1	0
Name	INTINV	INTS	DBC5	DBC4	DBC3	DBC2	DBC1	DBC0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

Bit 7      **INTINV**: Inverting control of the de-bounced external interrupt input signal

0: Non-inverted  
1: Inverted

Bit 6      **INTS**: Select external interrupt source

0: INT  
1: Comparator 0 output “C0VO”

Bit 5~0    **DBC5~DBC0**: Select external interrupt input de-bounce time ( $f_{SYS}=8\text{MHz}$ )

000000: bypass digital de-bounce circuit  
000001:  $0 \sim 1/f_{SYS}$ , about  $0.125\mu\text{s}$   
000010:  $1/f_{SYS} \sim 2/f_{SYS}$ , about  $0.25\mu\text{s}$   
:  
:  
101111:  $46/f_{SYS} \sim 47/f_{SYS}$ , about  $5.875\mu\text{s}$   
11xxxx:  $47/f_{SYS} \sim 48/f_{SYS}$ , about  $6\mu\text{s}$

• CTRL2 Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	CVREF5	CVREF4	CVREF3	CVREF2	CVREF1	CVREF0
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	0	0	0	0

Bit 7~6 Unimplemented, read as “0”

Bit 5~3 **CVREF5~CVREF3**: Select internal reference voltage  $V_{R2}$  for comparator 2

000:  $0.600V_{DD}$   
 001:  $0.625V_{DD}$   
 010:  $0.650V_{DD}$   
 011:  $0.675V_{DD}$   
 100:  $0.700V_{DD}$   
 101:  $0.725V_{DD}$   
 110:  $0.750V_{DD}$   
 111:  $0.775V_{DD}$

Bit 2~0 **CVREF2~CVREF0**: Select internal reference voltage  $V_{R1}$  for comparator 1

000:  $0.600V_{DD}$   
 001:  $0.625V_{DD}$   
 010:  $0.650V_{DD}$   
 011:  $0.675V_{DD}$   
 100:  $0.700V_{DD}$   
 101:  $0.725V_{DD}$   
 110:  $0.750V_{DD}$   
 111:  $0.775V_{DD}$

• CTRL3 Register

Bit	7	6	5	4	3	2	1	0
Name	—	C3VOINV	CVREF11	CVREF10	CVREF9	CVREF8	CVREF7	CVREF6
R/W	—	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	0	0	0	0	0	0	0

Bit 7 Unimplemented, read as “0”

Bit 6 **C3VOINV**: Inverting control of the comparator 3 output signal

0: Non-inverted  
 1: Inverted

Bit 5~3 **CVREF11~CVREF9**: Select internal reference voltage  $V_{R4}$  for comparator 3

000:  $0.600V_{DD}$   
 001:  $0.625V_{DD}$   
 010:  $0.650V_{DD}$   
 011:  $0.675V_{DD}$   
 100:  $0.700V_{DD}$   
 101:  $0.725V_{DD}$   
 110:  $0.750V_{DD}$   
 111:  $0.775V_{DD}$

Bit 2~0 **CVREF8~CVREF6**: Select internal reference voltage  $V_{R3}$  for comparator 2

000:  $0.075V_{DD}$   
 001:  $0.100V_{DD}$   
 010:  $0.125V_{DD}$   
 011:  $0.150V_{DD}$   
 100:  $0.175V_{DD}$   
 101:  $0.200V_{DD}$   
 110:  $0.225V_{DD}$   
 111:  $0.250V_{DD}$

• **CTRL4 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	PPGDL5	PPGDL4	PPGDL3	PPGDL2	PPGDL1	PPGDL0
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	0	0	0	0

Bit 7~6 Unimplemented, read as “0”

Bit 5~0 **PPGDL5~PPGDL0**: Select trigger delay for PPG ( $f_{SYS}=8\text{MHz}$ )

000000: 0 $\mu\text{s}$   
 000001: 0.125 $\mu\text{s}$   
 000010: 0.25 $\mu\text{s}$   
 :  
 :  
 101111: 5.875 $\mu\text{s}$   
 1xxxxx: 6 $\mu\text{s}$

Note: 1. Trigger delay means the time of the falling edge of INT0S to PPG trigger signal (INT00) being sent.

2. Falling edge of INT0S during trigger delay period will be ignored.

INT0: Inverted or non-inverted de-bounce signal from  $\overline{\text{INT}}$  or comparator 0 output “COVO” by software option;

INT0S: Signal represents single or double falling edge of INT0;

INT00: PPG hardware trigger signal.

## Oscillators

Various oscillator options offer the user a wide range of functions according to their various application requirements. The flexible features of the oscillator functions ensure that the best optimisation can be achieved in terms of speed and power saving. Oscillator selections and operation are selected through configuration options.

### Oscillator Overview

In addition to being the source of the main system clock the oscillators also provide clock sources for the Watchdog Timer. Two fully integrated internal oscillators, requiring no external components, are provided to form a range of both fast and slow system oscillators. The higher frequency oscillator provides higher performance but carry with it the disadvantage of higher power requirements, while the opposite is of course true for the lower frequency oscillator. With the capability of dynamically switching between fast and slow system clock, the device has the flexibility to optimize the performance/power ratio, a feature especially important in power sensitive portable applications.

Type	Name	Frequency
Internal High Speed RC	HIRC	8MHz
Internal Low Speed RC	LIRC	32kHz

### System Clock Configurations

There are two methods of generating the system clock, a high speed oscillator and a low speed oscillator. The high speed oscillator is the internal 8MHz RC oscillator. The low speed oscillator is the internal 32kHz RC oscillator. Selecting whether the low or high speed oscillator is used as the system oscillator is implemented by the configuration options.

### **High Speed Internal RC Oscillator – HIRC**

The high speed internal RC oscillator is a fully integrated system oscillator requiring no external components. The internal RC oscillator has a fixed frequency of 8MHz. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised.

### **Internal 32kHz Oscillator – LIRC**

The internal 32kHz System Oscillator is the low frequency oscillator. It is a fully integrated RC oscillator with a typical frequency of 32kHz at 5V, requiring no external components for its implementation. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised.

## **Power Down Mode and Wake-up**

### **Power Down Mode**

All of the Holtek microcontrollers have the ability to enter a Power Down Mode. When the device enters this mode, the normal operating current, will be reduced to an extremely low standby current level. This occurs because when the device enters the Power Down Mode, the system oscillator is stopped which reduces the power consumption to extremely low levels, however, as the device maintains its present internal condition, it can be woken up at a later stage and continue running, without requiring a full reset. This feature is extremely important in application areas where the MCU must have its power supply constantly maintained to keep the device in a known condition but where the power supply capacity is limited such as in battery applications.

### **Entering the Power Down Mode**

There is only one way for the device to enter the Power Down Mode and that is to execute the “HALT” instruction in the application program. When this instruction is executed, the following will occur:

- The system oscillator will stop running and the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition
- The WDT will be cleared and resume counting if the WDT clock source is selected to come from the WDT oscillator. The WDT will stop if its clock source originates from the system clock.
- The I/O ports will maintain their present condition.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.

### **Standby Current Considerations**

As the main reason for entering the Power Down Mode is to keep the current consumption of the MCU to as low a value as possible, perhaps only in the order of several micro-amps, there are other considerations which must also be taken into account by the circuit designer if the power consumption is to be minimized. Special attention must be made to the I/O pins on the device. All high-impedance input pins must be connected to either a fixed high or low level as any floating input pins could create internal oscillations and result in increased current consumption. This also applies to the device which has different package types, as there may be unconnected pins, which must either be setup as outputs or if setup as inputs must have pull-high resistors connected. Care must also be

taken with the loads, which are connected to I/O pins, which are setup as outputs. These should be placed in a condition in which minimum current is drawn or connected only to external circuits that do not draw current, such as other CMOS inputs. Also note that additional standby current will also be required if the configuration options have enabled the Watchdog Timer internal oscillator.

## Wake-up

After the system enters the Power Down Mode, it can be woken up from one of various sources listed as follows:

- An external falling edge on Port A
- A system interrupt
- A WDT overflow

If the device is woken up by a WDT overflow, a Watchdog Timer reset will be initiated. Although the wake-up method will initiate a reset operation, the actual source of the wake-up can be determined by examining the TO and PDF flags. The PDF flag is cleared by a system power-up or executing the clear Watchdog Timer instructions and is set when executing the “HALT” instruction. The TO flag is set if a WDT time-out occurs, and causes a wake-up that only resets the Program Counter and Stack Pointer, the other flags remain in their original status.

Each pin on Port A can be setup via an individual configuration option to permit a negative transition on the pin to wake-up the system. When a Port A pin wake-up occurs, the program will resume execution at the instruction following the “HALT” instruction.

If the system is woken up by an interrupt, then two possible situations may occur. The first is where the related interrupt is disabled or the interrupt is enabled but the stack is full, in which case the program will resume execution at the instruction following the “HALT” instruction. In this situation, the interrupt which woke-up the device will not be immediately serviced, but will rather be serviced later when the related interrupt is finally enabled or when a stack level becomes free. The other situation is where the related interrupt is enabled and the stack is not full, in which case the regular interrupt response takes place. If an interrupt request flag is set to “1” before entering the Power Down Mode, the wake-up function of the related interrupt will be disabled.

No matter what the source of the wake-up event is, once a wake-up situation occurs, a time period equal to 1024 system clock periods will be required before normal system operation resumes. However, if the wake-up has originated due to an interrupt, the actual interrupt subroutine execution will be delayed by an additional one or more cycles. If the wake-up results in the execution of the next instruction following the “HALT” instruction, this will be executed immediately after the 1024 system clock period delay has ended.

## Watchdog Timer

The Watchdog Timer is provided to prevent program malfunctions or sequences from jumping to unknown locations, due to certain uncontrollable external events such as electrical noise.

### Watchdog Timer Clock Source

The Watchdog Timer clock source is provided by the LIRC or  $f_{SYS}/4$  clock, selected by configuration option. The WDT oscillator has an approximate period of  $65\mu s$  at a supply voltage of 5V. Note that the oscillation frequency can vary with  $V_{DD}$ , temperature and process variations. Another clock source is supplied by  $f_{SYS}/4$ . The Watchdog Timer source clock is then subdivided by a ratio of  $2^{13}$  to  $2^{16}$  to give longer timeouts, the actual value being chosen using the configuration option.

Note that the LIRC is running during MCU normal mode. Once the device entering power down mode, the LIRC will keep running if the WDT is enabled and stop if the WDT is disabled.

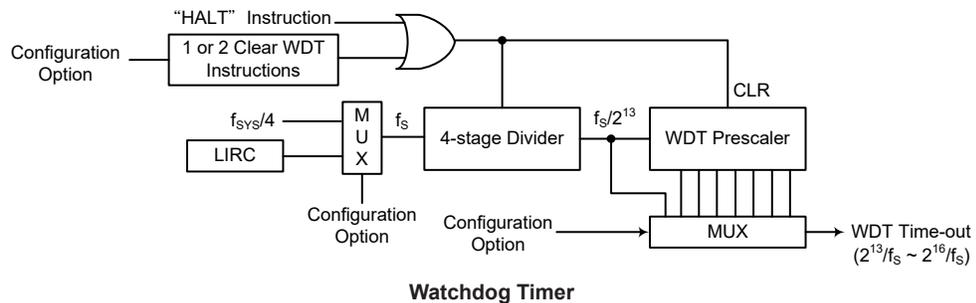
## Watchdog Timer Operation

The Watchdog Timer operates by providing a device reset when its timer overflows. This means that in the application program and during normal operation the user has to strategically clear the Watchdog Timer before it overflows to prevent the Watchdog Timer from executing a reset. This is done using the clear watchdog instructions. If the program malfunctions for whatever reason, jumps to an unknown location, or enters an endless loop, these clear instructions will not be executed in the correct manner, in which case the Watchdog Timer will overflow and reset the device. Most of the Watchdog Timer options, such as enable/disable, Watchdog Timer clock source and clear instruction type are selected using configuration options. There are not any related register. Note that if the Watchdog Timer configuration option has been disabled, then any instruction relating to its operation will result in no operation.

Under normal program operation, a Watchdog Timer time-out will initialise a device reset and set the status bit TO. However, if the system is in the Power Down Mode, when a Watchdog Timer time-out occurs, the TO bit in the status register will be set and only the Program Counter and Stack Pointer will be reset. Two methods can be adopted to clear the contents of the Watchdog Timer. The first is using the Watchdog Timer software clear instructions and the second is via a HALT instruction.

There are two methods of using software instructions to clear the Watchdog Timer, one of which must be chosen by configuration option. The first option is to use the single “CLR WDT” instruction while the second is to use the two commands “CLR WDT1” and “CLR WDT2”. For the first option, a simple execution of “CLR WDT” will clear the WDT while for the second option, both “CLR WDT1” and “CLR WDT2” must both be executed to successfully clear the Watchdog Timer. Note that for this second option, if “CLR WDT1” is used to clear the Watchdog Timer, successive executions of this instruction will have no effect, only the execution of a “CLR WDT2” instruction will clear the Watchdog Timer. Similarly after the “CLR WDT2” instruction has been executed, only a successive “CLR WDT1” instruction can clear the Watchdog Timer.

The maximum time out period is when the  $2^{16}$  division ratio is selected. As an example, with the WDT oscillator as its source clock, this will give a maximum watchdog period of around  $4.7\mu\text{s}$  for the  $2^{16}$  division ratio. If the device operates in a noisy environment, using the WDT oscillator is strongly recommended, since the system clock will be stopped under the power down mode.



## Reset and Initialisation

A reset function is a fundamental part of any microcontroller ensuring that the device can be set to some predetermined condition irrespective of outside parameters. The most important reset condition is after power is first applied to the microcontroller. In this case, internal circuitry will ensure that the microcontroller, after a short delay, will be in a well-defined state and ready to execute the first program instruction. After this power-on reset, certain important internal registers will be set to defined states before the program commences. One of these registers is the Program Counter, which will be reset to zero forcing the microcontroller to begin program execution from the lowest Program Memory address.

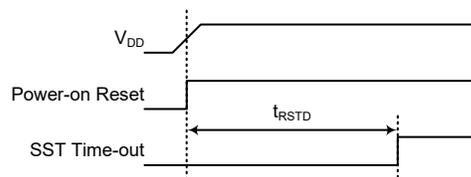
In addition to the power-on reset, another reset exists in the form of a Low Voltage Reset, LVR, where a full reset is implemented in situations where the power supply voltage falls below a certain threshold. Another type of reset is when the Watchdog Timer overflows and resets the microcontroller. All types of reset operations result in different register conditions being setup.

### Reset Functions

There are several ways in which a microcontroller reset can occur, through events occurring internally:

#### Power-on Reset

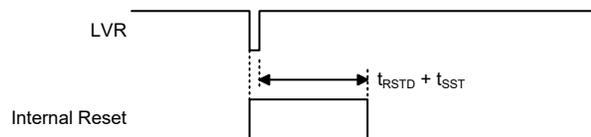
The most fundamental and unavoidable reset is the one that occurs after power is first applied to the microcontroller. As well as ensuring that the Program Memory begins execution from the first memory address, a power-on reset also ensures that certain other registers are preset to known conditions. All the I/O port and port control registers will power up in a high condition ensuring that all pins will be first set to inputs.



**Power-On Reset Timing Chart**

#### Low Voltage Reset – LVR

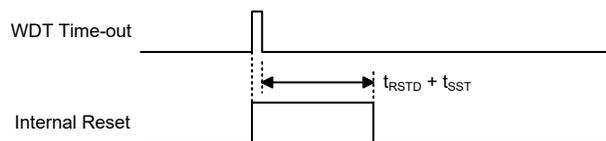
The device contains a low voltage reset circuit in order to monitor the supply voltage of the device. The LVR function is always enabled with a specific LVR voltage  $V_{LVR}$ . If the supply voltage of the device drops to within a range of  $0.9V \sim V_{LVR}$  such as might occur when changing the battery, the LVR will automatically reset the device internally. For a valid LVR signal, a low supply voltage, i.e., a voltage in the range between  $0.9V \sim V_{LVR}$  must exist for a time greater than that specified by  $t_{LVR}$  in the LVD & LVR Electrical Characteristics. If the low supply voltage state does not exceed this value, the LVR will ignore the low supply voltage and will not perform a reset function. The actual VLVR is fixed at a voltage value of 2.1V. Note that the LVR function will be automatically disabled when the device enters the power down mode.



**Low Voltage Reset Timing Chart**

### Watchdog Time-out Reset during Normal Operation

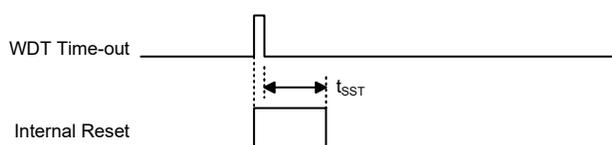
The Watchdog time-out Reset during normal operation is the same as a LVR reset except that the Watchdog time-out flag TO will be set to “1”.



WDT Time-out Reset during Normal Operation Timing Chart

### Watchdog Time-out Reset during Power Down Mode

The Watchdog time-out Reset during Power Down Mode is a little different from other kinds of reset. Most of the conditions remain unchanged except that the Program Counter and the Stack Pointer will be cleared to “0” and the TO flag will be set to “1”. Refer to the A.C. Characteristics for  $t_{SST}$  details.



WDT Time-out Reset during Power Down Timing Chart

### Reset Initial Conditions

The different types of reset described affect the reset flags in different ways. These flags, known as PDF and TO are located in the status register and are controlled by various microcontroller operations, such as the Power Down Mode function or Watchdog Timer. The reset flags are shown in the table:

TO	PDF	Reset Conditions
0	0	Power-on reset
u	u	LVR reset during Normal Mode operation
1	u	WDT time-out reset during Normal Mode operation
1	1	WDT time-out reset during Power Down Mode operation

Note: “u” stands for unchanged

The following table indicates the way in which the various components of the microcontroller are affected after a power-on reset occurs.

Item	Condition after Reset
Program Counter	Reset to zero
Interrupts	All interrupts will be disabled
WDT	Clear after reset, WDT begins counting
Timer/Event Counter	All Timer/Counters will be turned off
PPG Counter	The PPG Counter will be stopped
PPG, IOFF Output	Floating
Input/Output Ports	I/O ports will be setup as inputs
Stack Pointer	Stack Pointer will point to the top of the stack

The different kinds of resets all affect the internal registers of the microcontroller in different ways. To ensure reliable continuation of normal program execution after a reset occurs, it is important to know what condition the microcontroller is in after a particular reset occurs. The following table describes how each type of reset affects each of the microcontroller internal registers.

Register	Power On Reset	LVR Reset	WDT Time-out (Normal Operation)	WDT Time-out (Powe Down Mode)
IAR0	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
MP0	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
IAR1	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
MP1	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
BP	- - - - - - 0	- - - - - - 0	- - - - - - 0	- - - - - - u
ACC	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
PCL	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0
TBLP	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
TBLH	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
STATUS	- - 0 0 x x x x	- - u u u u u u	- - 1 u u u u u	- - 1 1 u u u u
INTC0	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- u u u u u u u u
TMR0	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
TMR0C	0 0 - 0 1 0 0 0 0	0 0 - 0 1 0 0 0 0	0 0 - 0 1 0 0 0 0	u u - u u u u u
TMR1	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
TMR1C	0 0 - 0 - 0 0 0 0 0	0 0 - 0 - 0 0 0 0 0	0 0 - 0 - 0 0 0 0 0	u u - u - u u u
PA	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	u u u u u u u u
PAC	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	u u u u u u u u
PB	- - - 1 1 1 1 1 1	- - - 1 1 1 1 1 1	- - - 1 1 1 1 1 1	- - - u u u u u
PBC	- - - 1 1 1 1 1 1	- - - 1 1 1 1 1 1	- - - 1 1 1 1 1 1	- - - u u u u u
CTRL0	0 0 - - 0 - 0 0	0 0 - - 0 - 0 0	0 0 - - 0 - 0 0	u u - - u - u u
CMP0C	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	u u u u u u u u
CMP1C	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	u u u u u u u u
CMP2C	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	u u u u u u u u
INTC1	- 0 0 - - 0 0 -	- 0 0 - - 0 0 -	- 0 0 - - 0 0 -	- u u - - u u -
PPGC	0 0 0 0 - - 0 0	0 0 0 0 - - 0 0	0 0 0 0 - - 0 0	u u u u - - u u
PPGTA	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
PPGTB	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
PPGTEX	- - - x - - - x	- - - x - - - x	- - - x - - - x	- - - u - - - u
ADRL	x x x x - - - -	x x x x x x x x	x x x x x x x x	u u u u u u u u
ADRH	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
ADCR	0 1 - - 0 0 0 0	0 1 - - 0 0 0 0	0 1 - - 0 0 0 0	u u - - u u u u
ACSR	- - - - - - 0 0	- - - - - - 0 0	- - - - - - 0 0	- - - - - - u u
CTRL1	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	u u u u u u u u
OPAC	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	u u u u u u u u
CTRL2	- - 0 0 0 0 0 0	- - 0 0 0 0 0 0	- - 0 0 0 0 0 0	- - u u u u u u
PWLT	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
CTRL3	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- u u u u u u u
TMR2	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
TMR2C	- - - 0 - 0 0 0	- - - 0 - 0 0 0	- - - 0 - 0 0 0	- - - u - u u u
CMPSC	0 0 0 0 - - 0 0	0 0 0 0 - - 0 0	0 0 0 0 - - 0 0	u u u u - - u u
CMP3C	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	0 0 0 1 0 0 0 0	u u u u u u u u
CTRL4	- - 0 0 0 0 0 0	- - 0 0 0 0 0 0	- - 0 0 0 0 0 0	- - u u u u u u
MFI	- 0 0 0 - 0 0 0	- 0 0 0 - 0 0 0	- 0 0 0 - 0 0 0	- u u u - u u u
PCRL	0 0 0 0 0 0 0 -	0 0 0 0 0 0 0 -	0 0 0 0 0 0 0 -	u u u u u u u -
PCRH	- - - - - - 0 0	- - - - - - 0 0	- - - - - - 0 0	- - - - - - u u

Note: “u” stands for unchanged  
“x” stands for unknown  
“-” stands for unimplemented

## Input/Output Ports

Holtek microcontrollers offer considerable flexibility on their I/O ports. With the input or output designation of every pin fully under user program control, Additionally, as there are pull-high resistors and wake-up configuration options, the user is provided with an I/O structure to meet the needs of a wide range of application possibilities.

The device provides bidirectional input/output lines labeled with port names PA~PB. These I/O ports are mapped to the RAM Data Memory with specific addresses as shown in the Special Purpose Data Memory table. All of these I/O ports can be used for input and output operations. For input operation, these ports are non-latching, which means the inputs must be ready at the T2 rising edge of instruction “MOV A, [m]”, where m denotes the port address. For output operation, all the data is latched and remains unchanged until the output latch is rewritten.

Register Name	Bit							
	7	6	5	4	3	2	1	0
PA	PA7	PA6	PA5	PA4	PA3	PA2	PA1	PA0
PAC	PAC7	PAC6	PAC5	PAC4	PAC3	PAC2	PAC1	PAC0
PB	—	—	—	PB4	PB3	PB2	PB1	PB0
PBC	—	—	—	PBC4	PBC3	PBC2	PBC1	PBC0

“—”: Unimplemented, read as “0”

I/O Logic Function Register List

### Pull-high Resistors

Many product applications require pull-high resistors for their switch inputs usually requiring the use of an external resistor. To eliminate the need for these external resistors, when configured as an input have the capability of being connected to an internal pull-high resistor. These pull-high resistors are selectable via configuration options. The pull-high resistors are implemented using weak PMOS transistors.

### Port A Wake-up

The HALT instruction forces the microcontroller into the Power Down Mode which preserves power, a feature that is important for battery and other low-power applications. Various methods exist to wake-up the microcontroller, one of which is to change the logic condition on one of the Port A pins from high to low. After a HALT instruction forces the microcontroller into entering the Power Down Mode, the processor will remain in a low-power state until the logic condition of the selected wake-up pin on Port A changes from high to low. This function is especially suitable for applications that can be woken up via external switches. Each pin on Port A can be selected individually by configuration options to have this wake-up feature.

### I/O Port Control Registers

Each I/O port has its own control register known as PAC~PBC, to control the input/output configuration. With this control register, each CMOS output or input can be reconfigured dynamically under software control. Each pin of the I/O ports is directly mapped to a bit in its associated port control register. For the I/O pin to function as an input, the corresponding bit of the control register must be written as a “1”. This will then allow the logic state of the input pin to be directly read by instructions. When the corresponding bit of the control register is written as a “0”, the I/O pin will be setup as a CMOS output. If the pin is currently setup as an output, instructions can still be used to read the output register. However, it should be noted that the program will in fact only read the status of the output data latch and not the actual logic status of the output pin.

• **PxC Register**

Bit	7	6	5	4	3	2	1	0
Name	PxC7	PxC6	PxC5	PxC4	PxC3	PxC2	PxC1	PxC0
R/W								
POR	1	1	1	1	1	1	1	1

**PxCn:** I/O Port x Pin type selection

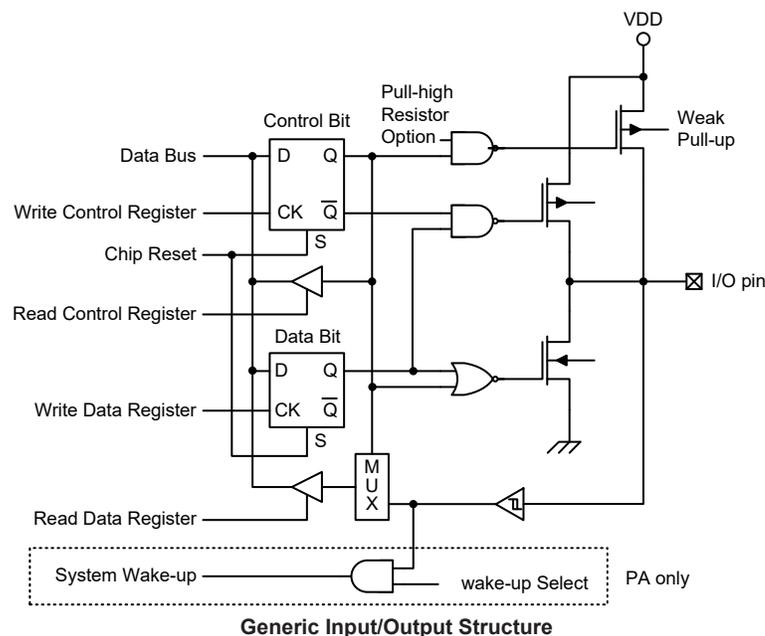
0: Output

1: Input

The PxCn bit is used to control the pin type selection. Here the “x” can be A and B. However, the actual available bits for each I/O Port may be different.

**I/O Pin Structures**

The accompanying diagram illustrates the internal structure of the I/O logic function. As the exact logical construction of the I/O pin will differ from this diagram, it is supplied as a guide only to assist with the functional understanding of the logic function I/O pins. The wide range of pin-shared structures does not permit all types to be shown.



**Programming Considerations**

Within the user program, one of the first things to consider is port initialisation. After a reset, all of the I/O data and port control registers will be set high. This means that all I/O pins will default to an input state, the level of which depends on the other connected circuitry and whether pull-high selections have been chosen. If the port control registers, PAC~PBC, are then programmed to setup some pins as outputs, these output pins will have an initial high output value unless the associated port data registers, PA~PB, are first programmed. Selecting which pins are inputs and which are outputs can be achieved byte-wide by loading the correct values into the appropriate port control register or by programming individual bits in the port control register using the “SET [m].i” and “CLR [m].i” instructions. Note that when using these bit control instructions, a read-modify-write operation takes place. The microcontroller must first read in the data on the entire port, modify it to the required new bit values and then rewrite this data back to the output ports.

Port A has the additional capability of providing wake-up functions. When the device is in the Power Down Mode, various methods are available to wake the device up. One of these is a high to low transition of any of the Port A pins. Single or multiple pins on Port A can be setup to have this function.

## Timer/Event Counters

The provision of timers form an important part of any microcontroller, giving the designer a means of carrying out time related functions. The device contains three count-up timer of 8-bit capacity. The Timer/Event Counter 0 has three different operating modes, it can be configured to operate as a general timer, an external event counter or as a pulse width capture device. The Timer/Event Counter 1 has only one operating mode – the PPG non-retriggered function. The Timer/Event Counter 2 has only one operating mode which is timer mode. The provision of an internal prescaler to the clock circuitry on gives added range to the timers.

There are two types of registers related to the Timer/Event Counters. The first are the registers that contain the actual value of the timer and into which an initial value can be preloaded. Reading from these registers retrieves the contents of the Timer/Event Counter. The second type of associated registers are the Timer Control Registers which define the timer options and determines how the timers are to be used.

### Configuring the Timer/Event Counter Input Clock Source

For the Timer/Event Counter 0, the clock source can be from an external source or an internal clock source, when selecting an external source, the choice of which is determined by the TMR0ECS bit in the CTRL0 register, it can be from the TMR0 pin or the INT0 signal. For the Timer/Event Counter 1/2, there is only one clock source,  $f_{SYS}$ , which is also divided by a prescaler, the division ratio is conditioned by the Timer Control Register bits TnPSC0~TnPSC2. The external clock input allows the user to count external events, measure time internals or pulse widths. While using the internal clock allows the user to generate an accurate time base.

#### • CTRL0 Register

Bit	7	6	5	4	3	2	1	0
Name	C2VOINV	TMR0ECS	—	—	PCKPSC	—	LVDO	LVDC
R/W	R/W	R/W	—	—	R/W	—	R	R/W
POR	0	0	—	—	0	—	0	0

Bit 7      **C2VOINV**: Inverting control of the comparator 2 output signal  
 Described elsewhere

Bit 6      **TMR0ECS**: Select TMR0 external clock source  
 0: TMR0 pin  
 1: INT0

Bit 5~4    Unimplemented, read as “0”

Bit 3      **PCKPSC**: Peripheral clock prescaler stage  
 Described elsewhere

Bit 2      Unimplemented, read as “0”

Bit 1      **LVDO**: Low voltage detector output  
 Described elsewhere

Bit 0      **LVDC**: Low voltage detector control  
 Described elsewhere

### Timer/Event Counter Registers – TMR0, TMR1, TMR2

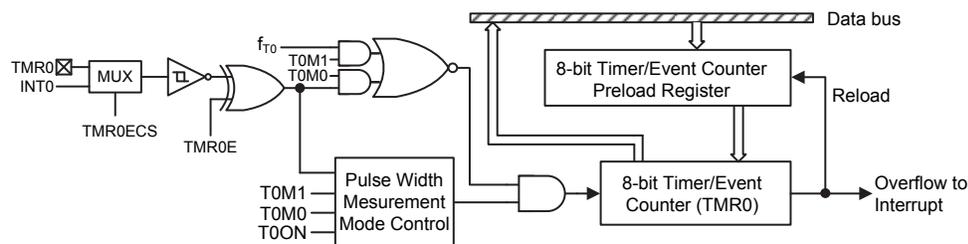
The timer registers are special function registers located in the Special Purpose Data Memory and is the place where the actual timer value is stored. The value in the timer registers increases by one each time an internal clock pulse is received or an external transition occurs on the external timer pin. The timer will count from the initial value loaded by the preload register to the full count of FFH at which point the timer overflows and an internal interrupt signal is generated. Then the timer value will be reset with the initial preload register value and continue counting.

Note that to achieve a maximum full range count of FFH, all the preload registers must first be cleared to zero. It should be noted that after power-on, the preload registers will be in an unknown condition. Note that if the Timer/Event Counter is in an OFF condition and data is written to its preload register, this data will be immediately written into the actual counter. However, if the counter is enabled and counting, any new data written into the preload data register during this period will remain in the preload register and will only be written into the actual counter the next time an overflow occurs.

### Timer Control Registers – TMR0C, TMR1C, TMR2C

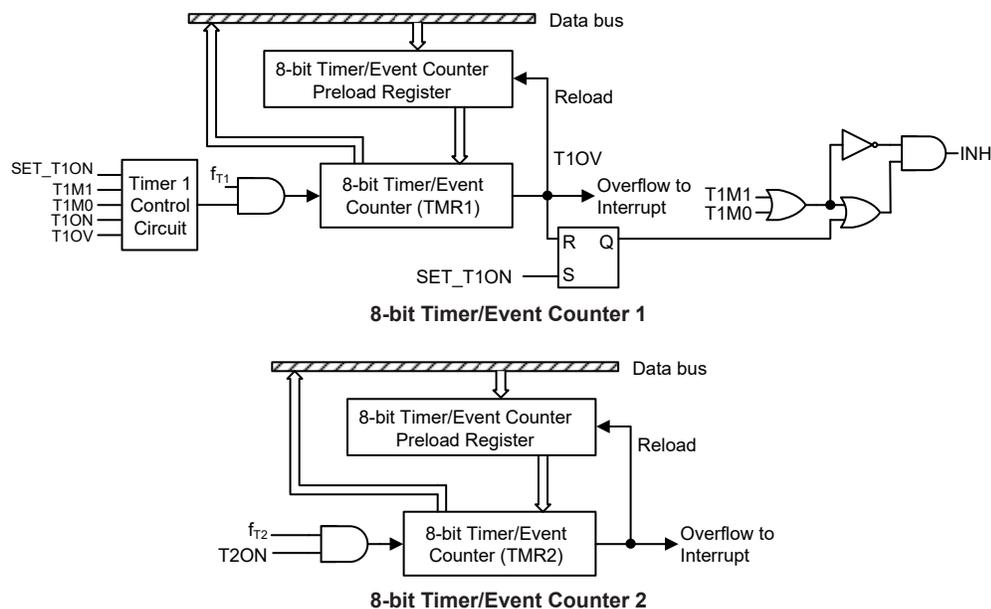
The flexible features of the Holtek microcontroller Timer/Event Counters enable them to operate in several different modes, the options of which are determined by the contents of their respective control register. The Timer Control Register is known as TMRnC. It is the Timer Control Register together with its corresponding timer registers that control the full operation of the Timer/Event Counter. Before the timer can be used, it is essential that the Timer Control Register is fully programmed with the right data to ensure its correct operation, a process that is normally carried out during program initialisation.

To choose which of the several modes the timer is to operate in, either in the timer mode, the event counting mode, the pulse width capture mode or the PPG non-retriggered function mode 0, bits 7 and 6 of the Timer Control Register, which are known as the bit pair TOM1/TOM0 or T1M1/T1M0, must be set to the required logic levels. The timer-on bit, which is bit 4 of the Timer Control Register and known as TnON, provides the basic on/off control of the respective timer. Setting the bit high allows the counter to run. Clearing the bit stops the counter. Bits 0~2 of the Timer Control Register determine the division ratio of the input clock prescaler. The prescaler bit settings have no effect if an external clock source is used. If the timer is in the event count or pulse width capture mode, the active transition edge level type is selected by the logic level of bit 3 of the Timer Control Register TMR0C which is known as TMR0E.



Note: INT0 is inverted or non-inverted de-bounce signal from  $\overline{INT}$  pin or comparator 0 output “C0VO” by software option.

#### 8-bit Timer/Event Counter 0



• **TMR0C Register**

Bit	7	6	5	4	3	2	1	0
Name	T0M1	T0M0	—	T0ON	TMR0E	T0PSC2	T0PSC1	T0PSC0
R/W	R/W	R/W	—	R/W	R/W	R/W	R/W	R/W
POR	0	0	—	0	1	0	0	0

- Bit 7~6    **T0M1~T0M0**: Timer 0 operation mode selection  
 00: No mode available  
 01: Event counter mode  
 10: Timer mode  
 11: Pulse width measurement mode
- Bit 5    Unimplemented, read as “0”
- Bit 4    **T0ON**: Timer/event counter counting enable  
 0: Disable  
 1: Enable
- Bit 3    **TMR0E**: Event counter active edge selection  
 In Event Counter Mode:  
 0: Count on rising edge  
 1: Count on falling edge  
 In Pulse Width measurement mode:  
 0: Start counting on the falling edge, stop on the rising edge  
 1: Start counting on the rising edge, stop on the falling edge
- Bit 2~0    **T0PSC2~T0PSC0**: Timer prescaler rate selection  
 Timer internal clock  $f_{T0}$  =  
 000:  $f_{SYS}$   
 001:  $f_{SYS}/2$   
 010:  $f_{SYS}/4$   
 011:  $f_{SYS}/8$   
 100:  $f_{SYS}/16$   
 101:  $f_{SYS}/32$   
 110:  $f_{SYS}/64$   
 111:  $f_{SYS}/128$

• **TMR1C Register**

Bit	7	6	5	4	3	2	1	0
Name	T1M1	T1M0	—	T1ON	—	T1PSC2	T1PSC1	T1PSC0
R/W	R/W	R/W	—	R/W	—	R/W	R/W	R/W
POR	0	0	—	0	—	0	0	0

Bit 7~6 **T1M1~T1M0**: Timer 1 operation mode selection

00: Mode 0 (PPG non-retriggered function)

01: No mode available

10: No mode available

11: No mode available

Bit 5 Unimplemented, read as “0”

Bit 4 **T1ON**: Timer/event counter counting enable

0: Disable

1: Enable

Bit 3 Unimplemented, read as “0”

Bit 2~0 **T1PSC2~T1PSC0**: Timer prescaler rate selection

Timer internal clock  $f_{T1}$ =

000:  $f_{SYS}$

001:  $f_{SYS}/2$

010:  $f_{SYS}/4$

011:  $f_{SYS}/8$

100:  $f_{SYS}/16$

101:  $f_{SYS}/32$

110:  $f_{SYS}/64$

111:  $f_{SYS}/128$

• **TMR2C Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	T2ON	—	T2PSC2	T2PSC1	T2PSC0
R/W	—	—	—	R/W	—	R/W	R/W	R/W
POR	—	—	—	0	—	0	0	0

Bit 7~5 Unimplemented, read as “0”

Bit 4 **T2ON**: Timer/event counter counting enable

0: Disable

1: Enable

Bit 3 Unimplemented, read as “0”

Bit 2~0 **T2PSC2~T2PSC0**: Timer prescaler rate selection

Timer internal clock  $f_{T2}$ =

000:  $f_{SYS}$

001:  $f_{SYS}/2$

010:  $f_{SYS}/4$

011:  $f_{SYS}/8$

100:  $f_{SYS}/16$

101:  $f_{SYS}/32$

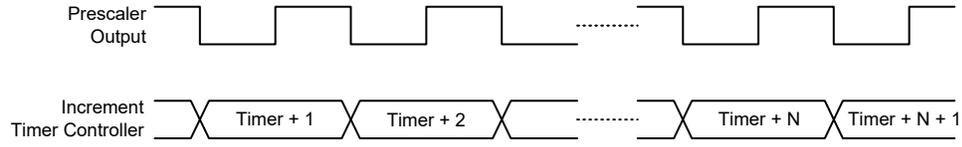
110:  $f_{SYS}/64$

111:  $f_{SYS}/128$

**Timer Mode**

In this mode, the Timer/Event Counter can be utilised to measure fixed time intervals, providing an internal interrupt signal each time the Timer/Event Counter overflows. To operate in this mode, the Operating Mode Select bit pair, T0M1/T0M0 or T2M1/T2M0, in the Timer Control Register must be set to 1/0.

In this mode the internal clock is used as the timer clock. The timer input clock source is  $f_{SYS}$ . However, this timer clock source is further divided by a prescaler, the value of which is determined by the bits T0PSC2~T0PSC0 or T2PSC2~T2PSC0 in the Timer Control Register. The timer-on bit, T0ON or T2ON must be set high to enable the timer to run. Each time an internal clock high to low transition occurs, the timer increments by one. When the timer is full and overflows, an interrupt signal is generated and the timer will reload the value already loaded into the preload register and continue counting. A timer overflow condition and corresponding internal interrupts are two of the wake-up sources. However, the internal interrupts can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bits in the Interrupt control registers are reset to zero.



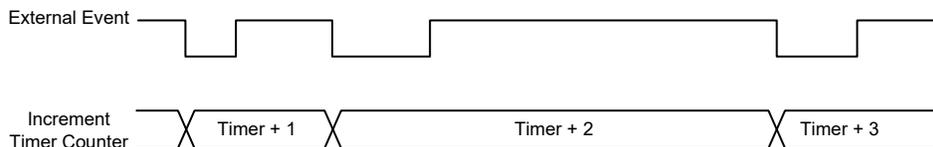
**Timer Mode Timing Chart**

### Event Counter Mode

This mode only exists in the Timer/Event Counter 0. In this mode, a number of externally changing logic events, occurring on the external timer pin TMR0, can be recorded by the Timer/Event Counter 0. To operate in this mode, the Operating Mode Select bit pair, T0M1/T0M0, in the Timer Control Register must be set to 0/1.

In this mode, the external timer pin TMR0 or the INT0 signal can be used as the Timer/Event Counter 0 clock source, however it is not divided by the internal prescaler. After the other bits in the Timer Control Register have been setup, the enable bit T0ON, which is bit 4 of the Timer Control Register, can be set high to enable the Timer/Event Counter to run. If the Active Edge Select bit, TMR0E, which is bit 3 of the Timer Control Register, is low, the Timer/Event Counter will increment each time the external timer pin receives a low to high transition. If the TMR0E is high, the counter will increment each time the external timer pin receives a high to low transition. When it is full and overflows, an interrupt signal is generated and the Timer/Event Counter will reload the value already loaded into the preload register and continue counting. The interrupt can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bit in the corresponding Interrupt Control Register. It is reset to zero.

As the external timer pin is shared with an I/O pin, to ensure that the pin is configured to operate as an event counter input pin, two things have to happen. The first is to ensure that the Operating Mode Select bits in the Timer Control Register place the Timer/Event Counter 0 in the Event Counting Mode. The second is to ensure that the port control register configures the pin as an input. It should be noted that in the event counting mode, even if the microcontroller is in the Power Down Mode, the Timer/Event Counter 0 will continue to record externally changing logic events on the timer input TMR0 pin or INT0 signal. As a result when the timer overflows it will generate a timer interrupt and corresponding wake-up source.



**Event Counter Mode Timing Chart (TMR0E=1)**

**Pulse Width Measurement Mode**

This mode only exists in the Timer/Event Counter 0. In this mode, the Timer/Event Counter can be utilised to measure the width of external pulses applied to the external timer pin. To operate in this mode, the Operating Mode Select bit pair, T0M1/T0M0, in the Timer Control Register must be set to 1/1.

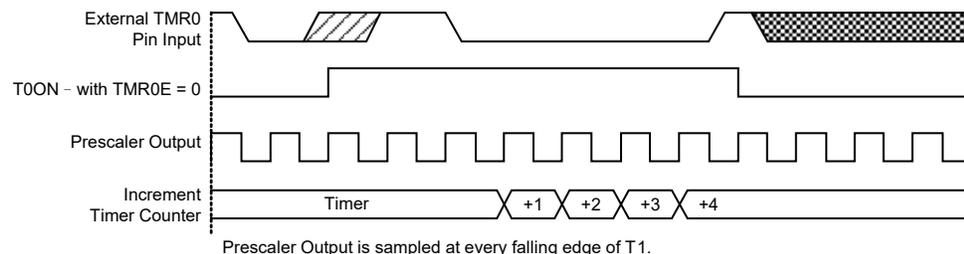
In this mode, the external timer pin TMR0 or the INT0 signal can be used as the Timer/Event Counter clock source, however it is not divided by the internal prescaler. After the other bits in the Timer Control Register have been setup, the enable bit T0ON, which is bit 4 of the Timer Control Register, can be set high to enable the Timer/Event Counter to run. However it will not actually start counting until an active edge is received on the external timer pin.

If the Active Edge Select bit TMR0E, which is bit 3 of the Timer Control Register, is low, once a high to low transition has been received on the external timer pin, the Timer/Event Counter will start counting until the external timer pin returns to its original high level. At this point the enable bit will be automatically reset to zero and the Timer/Event Counter will stop counting. If the Active Edge Select bit is high, the Timer/Event Counter will begin counting once a low to high transition has been received on the external timer pin and stop counting when the external timer pin returns to its original low level. As before, the enable bit will be automatically reset to zero and the Timer/Event Counter will stop counting. It is important to note that in the pulse width capture mode, the enable bit is automatically reset to zero when the external control signal on the external timer pin returns to its original level, whereas in the other two modes the enable bit can only be reset to zero under program control.

The residual value in the Timer/Event Counter, which can now be read by the program, therefore represents the length of the pulse received on the TMR0 pin or INT0 signal. As the enable bit has now been reset, any further transitions on the external timer pin will be ignored. The timer cannot begin further pulse width capture until the enable bit is set high again by the program. In this way, single shot pulse measurements can be easily made.

It should be noted that in this mode the Timer/Event Counter is controlled by logical transitions on the external timer pin and not by the logic level. When the Timer/Event Counter is full and overflows, an interrupt signal is generated and the Timer/Event Counter will reload the value already loaded into the preload register and continue counting. The interrupt can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bit in the corresponding Interrupt Control Register, it is reset to zero.

As the TMR0 pin is shared with an I/O pin, to ensure that the pin is configured to operate as a pulse width capture pin, two things have to be implemented. The first is to ensure that the Operating Mode Select bits in the Timer Control Register place the Timer/Event Counter in the pulse width capture mode, the second is to ensure that the port control register configure the pin as an input.



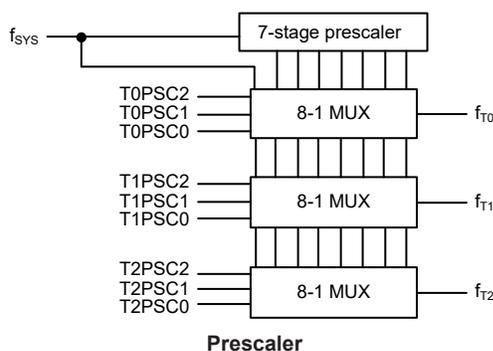
**Pulse Width Capture Mode Timing Chart (TMR0E=0)**

## Mode 0

The Timer/Event Counter 1 has a mode 0 for PPG usage. This mode is used to implement the PPG non-retriggered function. In this mode, the timer starts counting when PPG is stopped and stops when overflow. That means the TION will be set once PPG stopped and cleared when overflow. Once an overflow occurs, the counter is reloaded from the timer/event counter preload register, and generates an interrupt request flag. The interrupt can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bit in the corresponding Interrupt Control Register, it is reset to zero.

## Prescaler

Bits TnPSC0~TnPSC2 of the TMRnC register can be used to define a division ratio for the internal clock source of the Timer/Event Counter enabling longer time out periods to be setup.



## I/O Interfacing

The Timer/Event Counter, when configured to run in the event counter or pulse width capture mode, requires the use of an external timer pin for its operation. As this pin is a shared pin it must be configured correctly to ensure that it is setup for use as a Timer/Event Counter input pin. This is achieved by ensuring that the mode selects bits in the Timer/Event Counter control register, either the event counter or pulse width capture mode. Additionally the corresponding Port Control Register bit must be set high to ensure that the pin is setup as an input. Any pull-high resistor connected to this pin will remain valid even if the pin is used as a Timer/Event Counter input.

## Programming Considerations

When configured to run in the timer mode, the internal system clock is used as the timer clock source and is therefore synchronised with the overall operation of the microcontroller. In this mode when the appropriate timer register is full, the microcontroller will generate an internal interrupt signal directing the program flow to the respective internal interrupt vector. For the pulse width capture mode, the internal system clock is also used as the timer clock source but the timer will only run when the correct logic condition appears on the external timer input pin. As this is an external event and not synchronised with the internal timer clock, the microcontroller will only see this external event when the next timer clock pulse arrives. As a result, there may be small differences in measured values requiring programmers to take this into account during programming. The same applies if the timer is configured to be in the event counting mode, which again is an external event and not synchronised with the internal system or timer clock.

When the Timer/Event Counter is read, or if data is written to the preload register, the clock is inhibited to avoid errors, however as this may result in a counting error, this should be taken into account by the programmer. Care must be taken to ensure that the timers are properly initialised before using them for the first time. The associated timer enable bits in the interrupt control

register must be properly set otherwise the internal interrupt associated with the timer will remain inactive. The edge select, timer mode and clock source control bits in timer control register must also be correctly set to ensure the timer is properly configured for the required application. It is also important to ensure that an initial value is first loaded into the timer registers before the timer is switched on; this is because after power-on the initial values of the timer registers are unknown. After the timer has been initialized the timer can be turned on and off by controlling the enable bit in the timer control register.

When the Timer/Event Counter overflows, its corresponding interrupt request flag in the interrupt control register will be set. If the Timer/Event Counter interrupt is enabled this will in turn generate an interrupt signal. However irrespective of whether the interrupts are enabled or not, a Timer/Event Counter overflow will also generate a wake-up signal if the device is in a Power-down condition. This situation may occur if the Timer/Event Counter is in the Event Counting Mode and if the external signal continues to change state. In such a case, the Timer/Event Counter will continue to count these external events and if an overflow occurs the device will be woken up from its Power-down condition. To prevent such a wake-up from occurring, the timer interrupt request flag should first be set high before issuing the “HALT” instruction to enter the Power down Mode.

### Timer Program Example

This program example shows how the Timer/Event Counter registers are setup, along with how the interrupts are enabled and managed. Note how the Timer/Event Counter is turned on, by setting bit 4 of the Timer Control Register. The Timer/Event Counter can be turned off in a similar way by clearing the same bit. This example program sets the Timer/Event Counters to be in the timer mode, which uses the internal system clock as their clock source.

#### Timer Programming Example

```

org 0ch          ; external interrupt vector
org 14h          ; Timer Counter 0 interrupt vector
jmp tmr0int      ; jump here when Timer 0 overflows
:
org 20h          ; main program
:
                ; internal Timer 0 interrupt routine

tmr0int:
:
                ; Timer 0 main program placed here

begin:
                ; setup Timer 0 registers
mov a, 09bh     ; setup Timer 0 preload value
mov tmr0, a
mov a, 081h     ; setup Timer 0 control register
mov tmr0c, a   ; timer mode and prescaler set to fsys/2 setup interrupt register
mov a, 001h     ; enable master interrupt and both timer interrupts
mov intc0, a
mov a, 020h
mov intc1, a
:
set tmr0c.4     ; start Timer 0
:

```

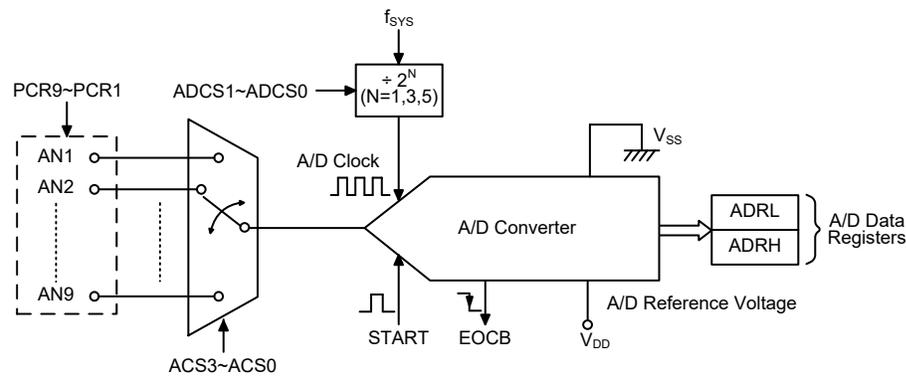
## Analog to Digital Converter

The need to interface to real world analog signals is a common requirement for many electronic systems. However, to properly process these signals by a microcontroller, they must first be converted into digital signals by A/D converters. By integrating the A/D conversion electronic circuitry into the microcontroller, the need for external components is reduced significantly with the corresponding follow-on benefits of lower costs and reduced component space requirements.

### A/D Overview

The device contains a 9-channel analog to digital converter which can directly interface to external analog signals, such as that from sensors or other control signals and convert these signals directly into a 12-bit digital value. It also can convert one internal signal into a 12-bit digital value.

The accompanying block diagram shows the overall internal structure of the A/D converter, together with its associated registers.



**A/D Converter Structure**

Note: AN1~AN8 is the external analog input channels, while AN9 is the internal analog input channel.

### A/D Converter Data Registers – ADRL, ADRH

The device, which has an internal 12-bit A/D converter, requires two data registers, a high byte register, known as ADRH, and a low byte register, known as ADRL. After the conversion process takes place, these registers can be directly read by the microcontroller to obtain the digitised conversion value. Only the high byte register, ADRH, utilises its full 8-bit contents. The low byte register utilises only 4 bit of its 8-bit contents as it contains only the lowest bits of the 12-bit converted value.

In the following table, D0~D11 is the A/D conversion data result bits.

Bit	ADRH								ADRL							
	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
Name	D11	D10	D9	D8	D7	D6	D5	D4	D3	D2	D1	D0	—	—	—	—
R/W	R	R	R	R	R	R	R	R	R	R	R	R	—	—	—	—
POR	x	x	x	x	x	x	x	x	x	x	x	x	—	—	—	—

“x”: unknown

“—”: Unimplemented, read as “0”

**D11~D0**: ADC conversion data

### A/D Converter Control Registers – ADCR, ACSR, PCRL, PCRH

To control the function and operation of the A/D converter, several control registers known as ADCR, ACSR, PCRL and PCRH are provided. These 8-bit registers define functions such as the selection of which analog channel is connected to the internal A/D converter, the digitised data format, the A/D clock source as well as controlling the start function and monitoring the A/D converter end of conversion status.

The ACS3~ACS0 bits in the ADCR register are used to determine which external input is selected to be converted. As the device contains only one actual analog to digital converter circuit, each of the individual 9 analog inputs must be routed to the converter. It is the function of the ACS3~ACS0 bits in the ADCR register to determine which analog channel is actually connected to the internal A/D converter.

The PCRH and PCRL control registers also contain the PCR9~PCR1 bits which determine which pins on PA~PB are used as analog inputs for the A/D converter and which pins are to be used as normal I/O pins.

#### • ADCR Register

Bit	7	6	5	4	3	2	1	0
Name	START	EOCB	—	—	ACS3	ACS2	ACS1	ACS0
R/W	R/W	R	—	—	R/W	R/W	R/W	R/W
POR	0	1	—	—	0	0	0	0

- Bit 7      **START:** Start the A/D conversion  
 0→1→0: Start  
 0→1: Reset the A/D converter and set EOCB to “1”
- Bit 6      **EOCB:** End of A/D conversion flag  
 0: A/D conversion ended  
 1: A/D conversion in progress  
 This read only flag is used to indicate when an A/D conversion process has completed. When the conversion process is running, the bit will be high.
- Bit 5~4    Unimplemented, read as “0”
- Bit 3~0    **ACS3~ACS0:** A/D Converter external analog input channel selection  
 0000: Undefined, cannot be used  
 0001: AN1  
 0010: AN2  
 0011: AN3  
 0100: AN4  
 0101: AN5  
 0110: AN6  
 0111: AN7  
 1000: AN8  
 1001: Internal analog input channel AN9  
 1010~1111: Undefined, cannot be used

#### • ACSR Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	ADCS1	ADCS0
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

- Bit 7~2    Unimplemented, read as “0”
- Bit 1~0    **ADCS1~ADCS0:** ADC clock rate selection bit  
 00:  $f_{SYS}/2$   
 01:  $f_{SYS}/8$   
 10:  $f_{SYS}/32$   
 11: Undefined, cannot be used

• **PCRL Register**

Bit	7	6	5	4	3	2	1	0
Name	PCR7	PCR6	PCR5	PCR4	PCR3	PCR2	PCR1	—
R/W	—							
POR	0	0	0	0	0	0	0	—

- Bit 7      **PCR7:** Define PA5 is A/D input or not  
0: Not A/D input  
1: A/D input, AN7
- Bit 6      **PCR6:** Define PA6 is A/D input or not  
0: Not A/D input  
1: A/D input, AN6
- Bit 5      **PCR5:** Define PA7 is A/D input or not  
0: Not A/D input  
1: A/D input, AN5
- Bit 4      **PCR4:** Define PA4 is A/D input or not  
0: Not A/D input  
1: A/D input, AN4
- Bit 3      **PCR3:** Define PB4 is A/D input or not  
0: Not A/D input  
1: A/D input, AN3
- Bit 2      **PCR2:** Define PB3 is A/D input or not  
0: Not A/D input  
1: A/D input, AN2
- Bit 1      **PCR1:** Define PB2 is A/D input or not  
0: Not A/D input  
1: A/D input, AN1
- Bit 0      Unimplemented, read as “0”

• **PCRH Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	PCR9	PCR8
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

- Bit 7~2    Unimplemented, read as “0”
- Bit 1      **PCR9:** Internal analog input channel AN9 enable  
0: Internal analog input channel AN9 disabled  
1: Internal analog input channel AN9 enabled
- Bit 0      **PCR8:** Define PA1 is A/D input or not  
0: Not A/D input  
1: A/D input, AN8

**A/D Operation**

The START bit in the ADCR register is used to start and reset the A/D converter. When the microcontroller sets this bit from low to high and then low again, an analog to digital conversion cycle will be initiated. When the START bit is brought from low to high but not low again, the EOCB bit in the ADCR register will be set high and the analog to digital converter will be reset. It is the START bit that is used to control the overall start operation of the internal analog to digital converter.

The EOCB bit in the ADCR register is used to indicate when the analog to digital conversion process is complete. This bit will be automatically set to zero by the microcontroller after a conversion cycle has ended. In addition, the corresponding A/D interrupt request flag will be set in the interrupt

control register, and if the interrupts are enabled, an appropriate internal interrupt signal will be generated. This A/D internal interrupt signal will direct the program flow to the associated A/D internal interrupt address for processing. If the A/D internal interrupt is disabled, the microcontroller can be used to poll the EOCB bit in the ADCR register to check whether it has been cleared as an alternative method of detecting the end of an A/D conversion cycle.

The clock source for the A/D converter originates from the subdivided version of  $f_{SYS}$ . The division ratio value is determined by the ADCS1~ADCS0 bits in the ACSR register.

Although the A/D clock source is determined by the system clock  $f_{SYS}$ , and by bits ADCS1~ADCS0, there are some limitations on the maximum A/D clock source speed that can be selected. As the minimum value of permissible A/D clock period,  $t_{ADCK}$ , is 0.5  $\mu s$ , care must be taken for system clock frequencies equal to or greater than 4MHz. For example, if the system clock operates at a frequency of 4MHz, the ADCS1~ADCS0 bits should not be set to “00”. Doing so will give A/D clock periods that are less than the minimum A/D clock period which may result in inaccurate A/D conversion values. Refer to the following table for examples, where values marked with an asterisk \* show where, depending upon the device, special care must be taken, as the values may be less than the specified minimum A/D Clock Period.

$f_{SYS}$	A/D Clock Period( $t_{ADCK}$ )			
	ADCS1~ADCS0=00 ( $f_{SYS}/2$ )	ADCS1~ADCS0=01 ( $f_{SYS}/8$ )	ADCS1~ADCS0=10 ( $f_{SYS}/32$ )	ADCS1~ADCS0=11
1MHz	2 $\mu s$	8 $\mu s$	32 $\mu s$	Undefined
2MHz	1 $\mu s$	4 $\mu s$	16 $\mu s$	Undefined
4MHz	500ns	2 $\mu s$	8 $\mu s$	Undefined
8MHz	250ns*	1 $\mu s$	4 $\mu s$	Undefined

**A/D Clock Period Examples**

### A/D Input Pins

All of the A/D analog input pins are pin-shared with the I/O pins as well as other functions. The PCR9~PCR1 bits in the PCRH and PCRL register, determine whether the input pins are setup as A/D converter analog inputs or whether they have other functions. In this way, pins can be changed under program control to change their function between A/D inputs and other functions. All pull-high resistors, which are setup through configuration options, will be automatically disconnected if the pins are setup as A/D inputs. Note that it is not necessary to first setup the A/D pin as an input in the PA~PB port control register to enable the A/D input as when the PCR9~PCR1 bits enable an A/D input, the status of the port control register will be overridden.

### Summary of A/D Conversion Steps

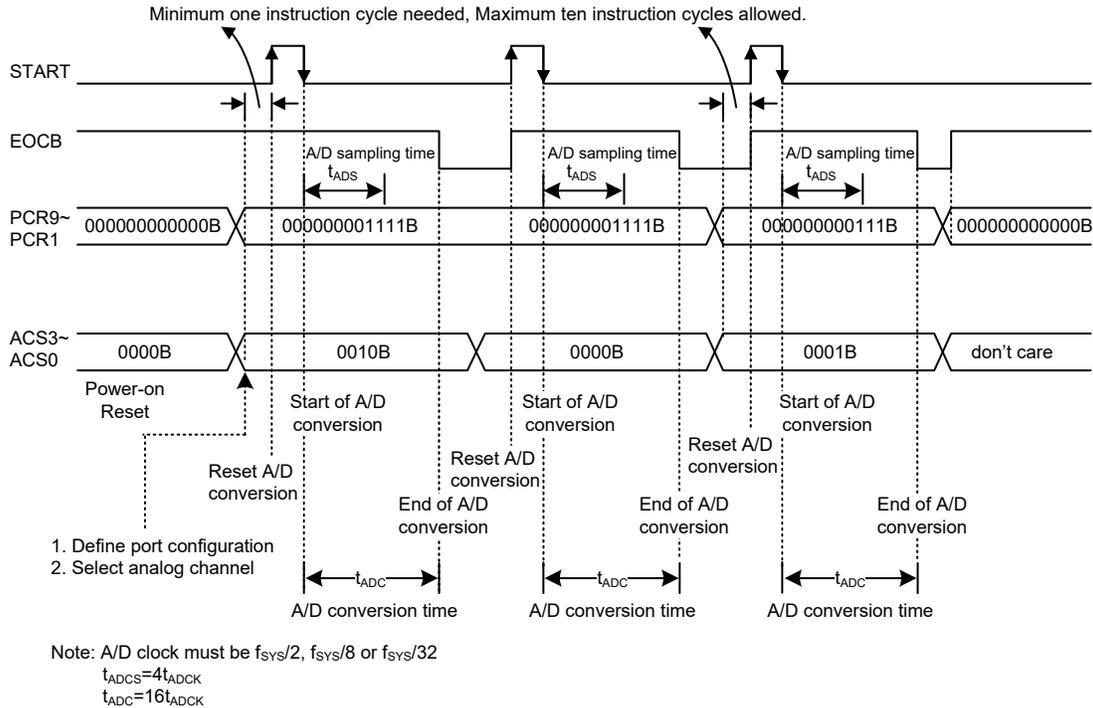
The following summarises the individual steps that should be executed in order to implement an A/D conversion process.

- Step 1  
Select the required A/D conversion clock by correctly programming bits ADCS1~ADCS0 in the ACSR register.
- Step 2  
Select which pins are to be used as A/D inputs and configure them as A/D input pins by correctly programming the PCR9~PCR1 bits in the PCRH and PCRL registers.
- Step 3  
Select which channel is to be connected to the internal A/D converter by correctly programming the ACS3~ACS0 bits in the ADCR register.

- Step 4  
 If the interrupts are to be used, the interrupt control registers must be correctly configured to ensure the A/D converter interrupt function is active. The master interrupt control bit, EMI, the INTC0 interrupt control register must be set to “1”, the A/D converter interrupt bit, ADE, must also be set to “1”.
- Step 5  
 The analog to digital conversion process can now be initialised by setting the START bit in the ADCR register from low to high and then low again. Note that this bit should have been originally cleared to zero.
- Step 6  
 To check when the analog to digital conversion process is complete, the EOCB bit in the ADCR register can be polled. The conversion process is complete when this bit goes low. When this occurs the A/D data registers ADRL and ADRH can be read to obtain the conversion value. As an alternative method, if the interrupts are enabled and the stack is not full, the program can wait for an A/D interrupt to occur.

Note: When checking for the end of the conversion process, if the method of polling the EOCB bit in the ADCR register is used, the interrupt enable step above can be omitted.

The accompanying diagram shows graphically the various stages involved in an analog to digital conversion process and its associated timing. After an A/D conversion process has been initiated by the application program, the microcontroller internal hardware will begin to carry out the conversion, during which time the program can continue with other functions. The time taken for the A/D conversion is  $16 t_{ADCK}$  where  $t_{ADCK}$  is equal to the A/D clock period.



**A/D Conversion Timing**

## Programming Considerations

When programming, the special attention must be given to the PCR9~PCR1 bits in the PCRH and PCRL registers. If these bits are all cleared to zero, no external pins will be selected for use as A/D input pins allowing the pins to be used as normal I/O pins. If the A/D converter is not used, the internal A/D circuitry should be power down. This may be an important consideration in power sensitive applications.

## A/D Transfer Function

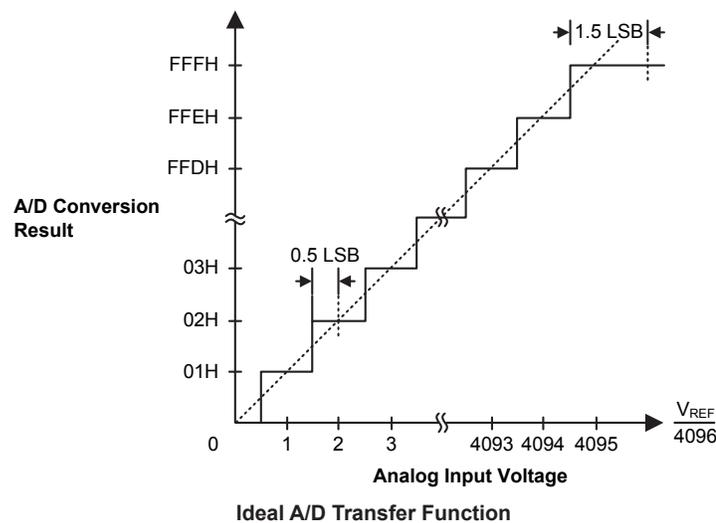
As the device contains a 12-bit A/D converter, its full-scale converted digitised value is equal to 3FFH. Since the full-scale analog input value is equal to the actual A/D converter reference voltage,  $V_{REF}$ , this gives a single bit analog input value of  $V_{REF}$  divided by 4096.

$$1 \text{ LSB} = V_{REF} \div 4096$$

The A/D Converter input voltage value can be calculated using the following equation:

$$\text{A/D input voltage} = \text{A/D output digital value} \times V_{REF} \div 4096$$

The diagram shows the ideal transfer function between the analog input value and the digitised output value for the A/D converter. Except for the digitised zero value, the subsequent digitised values will change at a point 0.5 LSB below where they would change without the offset, and the last full scale digitized value will change at a point 1.5 LSB below the  $V_{REF}$  level.



## A/D Programming Example

The following two programming examples illustrate how to setup and implement an A/D conversion. In the first example, the method of polling the EOCB bit in the ADCR register is used to detect when the conversion cycle is complete, whereas in the second example, the A/D interrupt is used to determine when the conversion is complete.

### Example: using an EOCB polling method to detect the end of conversion

```

clr ADE          ; disable ADC interrupt
mov a, 01H
mov ACSR, a      ; select fsys/8 as A/D clock
mov a, 02H
mov PCRL, a      ; setup PCRL register to configure PB2 as A/D input
mov a, 01h
mov ADCR, a      ; and select AN1 to be connected to the A/D converter
    
```

```
:
:
start_conversion:
clr  START          ; high pulse on start bit to initiate conversion
set  START          ; reset A/D
clr  START          ; start A/D
polling_EOC:
sz   EOCB           ; poll the ADCR register EOCB bit to detect end of A/D conversion
jmp  polling_EOC    ; continue polling
mov  a, ADRL        ; read low byte conversion result value
mov  ADRL_buffer, a ; save result to user defined register
mov  a, ADRH        ; read high byte conversion result value
mov  ADRH_buffer, a ; save result to user defined register
:
:
jmp  start_conversion ; start next A/D conversion
```

**Example: using the interrupt method to detect the end of conversion**

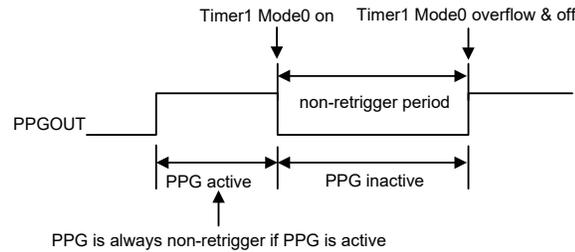
```
clr  ADE            ; disable ADC interrupt
mov  a, 01H
mov  ACSR, a        ; select fsys/8 as A/D clock
mov  a, 02H
mov  PCRL, a        ; setup PCRL register to configure PB2 as A/D input
mov  a, 01h
mov  ADCR, a        ; and select AN1 to be connected to the A/D converter
:
:
Start_conversion:
clr  START          ; high pulse on start bit to initiate conversion
set  START          ; reset A/D
clr  START          ; start A/D
clr  ADF            ; clear ADC interrupt request flag
set  ADE            ; enable ADC interrupt
set  EMI            ; enable global interrupt
:
:
; ADC interrupt service routine
ADC_ISR:
mov  acc_stack, a   ; save ACC to user defined memory
mov  a, STATUS
mov  status_stack, a ; save STATUS to user defined memory
:
:
mov  a, ADRL        ; read low byte conversion result value
mov  adrl_buffer, a ; save result to user defined register
mov  a, ADRH        ; read high byte conversion result value
mov  adrh_buffer, a ; save result to user defined register
:
:
EXIT_ISR:
mov  a, status_stack
mov  STATUS, a      ; restore STATUS from user defined memory
mov  a, acc_stack
mov  acc_stack, a   ; restore ACC from user defined memory
clr  ADF            ; clear ADC interrupt flag
reti
```



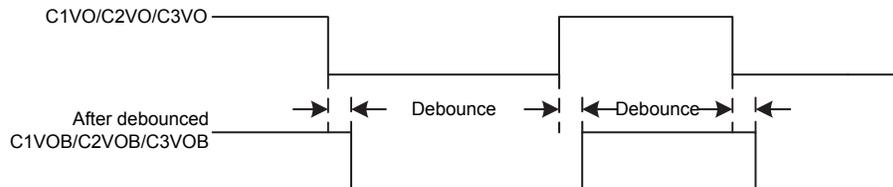
### Non-retriggered function

The PPG unit has non-retriggered function to inhibit further PPG trigger. The PPG will be non-triggered by one of following condition:

1. PPG is active.
2. During the non-retriggered period which starts counting once PPG stopped (Only available by using with mode 0 of Timer/Event Counter 1, the non-trigger period is decided by Timer/Event Counter 1).



- Debounce for C1VO/C2VO/C3VO



Note: The C1VOB/C2VOB/C3VOB signals are the debounced signals of the comparator outputs, C1VO/C2VO/C3VO, respectively.

### Pulse width limit function

The PPG unit has pulse width limit function to stop PPG output. The PPG output will be stopped once the pulse width reaching the limit. This function is implemented by a pulse width limit timer which starts counting once PPG is triggered and stops once overflow or PPG is stopped. The pulse width limit is  $(256-PWLT) / (f_{SYS}/2)$ , where PWLT is pulse width limit timer register.

To start the PPG operation:

- Enable PPG function by configuration option.
- Set the PPG output active level by configuration option.
- Set the PPG timer start counting is synchronized with system clock  $f_{SYS}$  or not by configuration option.
- Set the PPG input mode selection by the PRSEN bit and PSPEN bit of the PPGC register.
- Set the PPG output pulse width. Writing data to PPGTA, PPGTB and PPGTEX registers.
- Decide using C1VOB falling edge to enable the reload function from preload register B or not by setting the C1RLEN bit of the PPGC register and the C1CTL bit in the CMPSC register.
- Decide using the non-retriggered period function or not by using with mode 0 of Timer/Event Counter 1.
- Set pulse width limit timer for pulse width limit function by the PWLT register.
- When PPG input is triggered by INT00 falling edge transition or triggered by a software bit PST being set to "1", the PPG will start counting from the current content of preload register. When PPG input is triggered by a software bit PST being cleared to zero, C2VOB falling edge, PPG timer overflow occurs or a pulse width limit condition occurs, the PPG will stop counting.

• **PPGC Register**

Bit	7	6	5	4	3	2	1	0
Name	PST	PRSEN	PSPEN	RLBF	—	—	TRGMOD	C1RLEN
R/W	R/W	R/W	R/W	R/W	—	—	R/W	R/W
POR	0	0	0	0	—	—	0	0

- Bit 7      **PST**: PPG software trigger bit  
0: Stop PPG  
1: Restart PPG
- Bit 6      **PRSEN**: Restarting the PPG timer using INT00 trigger input enable control  
0: Disable  
1: Enable  
Disable restarting the PPG timer using INT00 trigger input, PPG module output can be restarted by software control bit PST only. Enable restarting the PPG timer using INT00 trigger input, PPG module output can be restarted by INT00 falling edge trigger or software control (PST is set to “1”).
- Bit 5      **PSPEN**: Stopping the PPG timer using the trigger input of C2VOB or C3VOB enable control  
0: Disable  
1: Enable  
The C2VOB or C3VOB signal is the inverted or non-inverted signal from the output of comparator 2 or comparator 3, which is selected by software option  
Disable stopping the PPG timer using C2VOB or C3VOB trigger input, PPG module output can be stopped by software control bit PST only. Enable stopping the PPG timer using C2VOB or C3VOB trigger input, PPG module output can be stopped by C2VOB or C3VOB falling edge trigger or software control (PST is set to “0”).
- Bit 4      **RLBF**: PPG reload control bit  
0: From PPGTA  
1: From PPGTB
- Bit 3~2      Unimplemented, read as “0”
- Bit 1      **TRGMOD**: Select single or double falling edge of INT0 as the input of trigger delay circuit which produce INT00  
0: Single  
1: Double
- Bit 0      **C1RLEN**: C1VOB falling edge to set RLBF for PPG timer reloads from preload register B enable control  
0: Disable  
1: Enable  
This bit is available when C1CTL=0. CP1P and CP1N are connected to programmable internal reference voltage  $V_{RI}$  and PA3/CP0P/TMR0 pin respectively.  
When C1CTL=0 and C1VO=1, if the C1RLEN is changed from 1 to 0, the RLBF bit will be set. So note that when the C1RLEN is changed from 1 to 0, the RLBF bit should be cleared to 0 by software, otherwise the PPG width will be determined by the PPGTB.

Normally, PPG timer is reloaded from preload register A if RLBF=0. If C1RLEN is set and C1CTL is cleared, the C1VOB falling edge caused by comparator output will set RLBF and then PPG timer will be reload from preload register B until RLBF is cleared by software.

The PRSEN is the PPG restarting enable or disable bit using INT00 trigger input. If this bit is enabled, the PPG timer restarting input can be trigger by INT00 falling edge.

The PSPEN is the PPG stopping enable or disable bit using the trigger input of C2VOB or C3VOB. If this bit is enabled, the PPG timer stopping input can be triggered and an IOFF output will active high by the falling edge of C2VOB or C3VOB. The IOFF output is floating during reset, high when

PPG is inactive and low when PPG is active. The PRSEN bit will be cleared by the falling edge of C2VOB or C3VOB, no matter the PPG is in active period or not. This will prevent PPG module output be restarted by INT00 falling edge again, only restarted by software control is permitted until PRSEN is set again by software.

The PST is a software trigger bit, if this bit is set to “1”, the PPG timer will start counting and this bit will be cleared when the PPG timer overflow occurs or PPG timer stop counting. If this bit is cleared to “0”, the PPG timer will stop counting. When the PPG timer is counting and if a falling edge generates from INT00 or a software control bit PST is set, the PPG timer counter is not affected, the trigger from INT00 or PST is not useful. The PST can also be used as a status bit of PPG timer output.

The PPG module output pulse active level is decided by configuration option. Another function is provided, which is the PPG timer start counting is synchronized with clock or not, decided by configuration option.

• **PPGTA Register**

Bit	7	6	5	4	3	2	1	0
Name	PGTA7	PGTA6	PGTA5	PGTA4	PGTA3	PGTA2	PGTA1	PGTA0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

“x”: unknown

• **PPGTB Register**

Bit	7	6	5	4	3	2	1	0
Name	PGTB7	PGTB6	PGTB5	PGTB4	PGTB3	PGTB2	PGTB1	PGTB0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	x	x	x	x	x	x	x	x

“x”: unknown

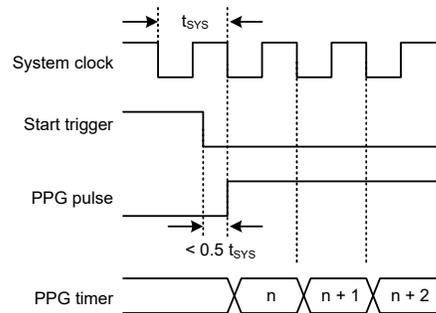
• **PPGTEX Register**

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	PGTB8	—	—	—	PGTA8
R/W	—	—	—	R/W	—	—	—	R/W
POR	—	—	—	x	—	—	—	x

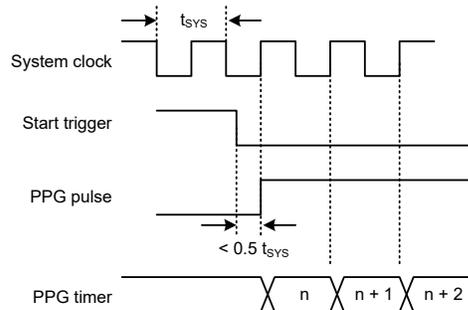
“x”: unknown

To control PPG pulse starting delay time is less than or equal to  $0.5 \times (1/f_{SYS})$  when start synchronized with clock is selected, clock ( $f_{SYS}$ ) edge trigger type (raising or falling), which triggers PPG, varies with next coming clock transition once PPG starts. After PPG starts, the PPG output becomes active and begins to count as soon as first falling or rising transition of system clock comes. After first trigger done, the following clock edge trigger type is decided by the first one. For example, once PPG starts and next coming clock transition is falling edge, the PPG will be trigger by falling edge until PPG stops and vice versa.

EX1: Since the first trigger type is falling edge after PPG starts, the PPG timer is triggered by falling edge until PPG stops.



EX2: Since the first trigger type is raising edge after PPG starts, the PPG timer is triggered by raising edge until PPG stops.



Any action causing PPG stop, such as PPG timer overflow, C2VOB falling edge (if PSPEN=1), software stop (PST=1 → 0) or reaching pulse limit, will cause actions as following:

- PPG timer will be reloaded
- PST is cleared
- PPG is inactive

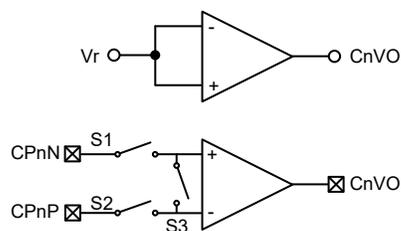
## Comparators

The device includes 4 integrated comparators in PPG module. Either inputs of the comparator 0 can be connected to AN9 selected by C0N2AN9 bit in the CMPSC register. The non-inverting input of the comparator 1 is connected to the programmable internal reference voltage  $V_{R1}$ . The inverting input of comparator 1 can be connected to PA5/CP1N/AN7 or PA3/CP0P/TMR0 by the C1CTL bit. The non-inverting input of the comparator 2 can be connected to PA7/CP2P/AN5 or programmable internal reference voltage  $V_{R2}$  by configuration options. The inverting input of the comparator 2 can be connected to PA6/CP2N/AN6 or programmable internal reference voltage  $V_{R3}$  by configuration options. The non-inverting input of the comparator 3 is connected to programmable internal reference voltage  $V_{R4}$ . The inverting input of comparator 3 can be connected to PA3/CP0P/TMR0 or OPA0 by configuration options

The input offset is adjustable by using a common mode input to calibrate the offset value.

The calibration steps are as the follows:

- (1) Setting CnCOFM=1 to offset cancellation mode (S3 is closed)
- (2) Setting CnCRS to select which input pin as reference voltage (S1 or S2 is closed)
- (3) Adjusting CnCOF0~CnCOF4 until output status is changed
- (4) Setting CnCOFM=0 to normal comparator mode



Comparator (n=0~3)

### Comparator registers

Full control over the internal comparator is provided via the control registers, CMP0C, CMP1C, CMP2C, CMP3C, and CMPSC. These registers are used to select the input path and the offset voltage cancellation function.

Register Name	Bit							
	7	6	5	4	3	2	1	0
CMP0C	C0CMPOP	C0COFM	C0CRS	C0COF4	C0COF3	C0COF2	C0COF1	C0COF0
CMP1C	C1CMPOP	C1COFM	C1CRS	C1COF4	C1COF3	C1COF2	C1COF1	C1COF0
CMP2C	C2CMPOP	C2COFM	C2CRS	C2COF4	C2COF3	C2COF2	C2COF1	C2COF0
CMP3C	C3CMPOP	C3COFM	C3CRS	C3COF4	C3COF3	C3COF2	C3COF1	C3COF0
CMPSC	C3HYSON	C2HYSON	C1HYSON	C0HYSON	—	—	C1CTL	C0N2AN9

Comparators Register List

#### • CMP0C Register

Bit	7	6	5	4	3	2	1	0
Name	C0CMPOP	C0COFM	C0CRS	C0COF4	C0COF3	C0COF2	C0COF1	C0COF0
R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	1	0	0	0	0

- Bit 7 **C0CMPOP**: Comparator 0 positive logic output
- Bit 6 **C0COFM**: Input offset voltage cancellation mode and comparator mode selection  
0: Comparator mode  
1: Input offset voltage cancellation mode
- Bit 5 **C0CRS**: Comparator 0 input offset voltage cancellation reference selection bit  
0: Select CP0N as the reference input  
1: Select CP0P as the reference input
- Bit 4~0 **C0COF4~C0COF0**: Comparator 0 input offset voltage cancellation control bits

#### • CMP1C Register

Bit	7	6	5	4	3	2	1	0
Name	C1CMPOP	C1COFM	C1CRS	C1COF4	C1COF3	C1COF2	C1COF1	C1COF0
R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	1	0	0	0	0

- Bit 7 **C1CMPOP**: Comparator 1 positive logic output
- Bit 6 **C1COFM**: Input offset voltage cancellation mode and comparator mode selection  
0: Comparator mode  
1: Input offset voltage cancellation mode
- Bit 5 **C1CRS**: Comparator 1 input offset voltage cancellation reference selection bit  
0: Select CP1N as the reference input  
1: Select CP1P as the reference input
- Bit 4~0 **C1COF4~C1COF0**: Comparator 1 input offset voltage cancellation control bits

• **CMP2C Register**

Bit	7	6	5	4	3	2	1	0
Name	C2CMPOP	C2COFM	C2CRS	C2COF4	C2COF3	C2COF2	C2COF1	C2COF0
R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	1	0	0	0	0

- Bit 7      **C2CMPOP**: Comparator 2 positive logic output
- Bit 6      **C2COFM**: Input offset voltage cancellation mode and comparator mode selection  
0: Comparator mode  
1: Input offset voltage cancellation mode
- Bit 5      **C2CRS**: Comparator 2 input offset voltage cancellation reference selection bit  
0: Select CP2N as the reference input  
1: Select CP2P as the reference input
- Bit 4~0    **C2COF4~C2COF0**: Comparator 2 input offset voltage cancellation control bits

• **CMP3C Register**

Bit	7	6	5	4	3	2	1	0
Name	C3CMPOP	C3COFM	C3CRS	C3COF4	C3COF3	C3COF2	C3COF1	C3COF0
R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	0	0	0	0	0

- Bit 7      **C3CMPOP**: Comparator 3 positive logic output
- Bit 6      **C3COFM**: Input offset voltage cancellation mode and comparator mode selection  
0: Comparator mode  
1: Input offset voltage cancellation mode
- Bit 5      **C3CRS**: Comparator 3 input offset voltage cancellation reference selection bit  
0: Select CP3N as the reference input  
1: Select CP3P as the reference input
- Bit 4~0    **C3COF4~C3COF0**: Comparator 3 input offset voltage cancellation control bits

• **CMPSC Register**

Bit	7	6	5	4	3	2	1	0
Name	C3HYSON	C2HYSON	C1HYSON	C0HYSON	—	—	C1CTL	C0N2AN9
R/W	R/W	R/W	R/W	R/W	—	—	R/W	R/W
POR	0	0	0	0	—	—	0	0

- Bit 7      **C3HYSON**: Comparator 3 hysteresis control bit  
0: Disable  
1: Enable
- Bit 6      **C2HYSON**: Comparator 2 hysteresis control bit  
0: Disable  
1: Enable
- Bit 5      **C1HYSON**: Comparator 1 hysteresis control bit  
0: Disable  
1: Enable
- Bit 4      **C0HYSON**: Comparator 0 hysteresis control bit  
0: Disable  
1: Enable
- Bit 3~2    Unimplemented, read as “0”
- Bit 1      **C1CTL**: Comparator 1 inputs connection control bit  
0: The inverting input of comparator 1 is connected to PA3/CP0P/TMR0  
1: The inverting input of comparator 1 is connected to PA5/CP1N/AN7
- Bit 0      **C0N2AN9**: Comparator 0 inputs connection  
0: CP0P is connected to AN9  
1: CP0N is connected to AN9

## Operational Amplifier

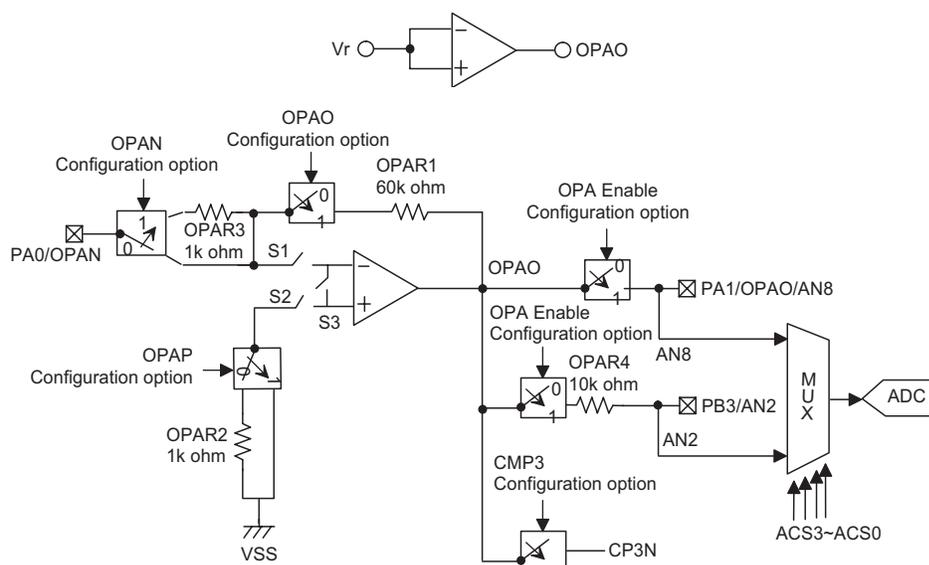
The device includes an integrated operational amplifier which is used to amplify the small analog input signal. It can be enabled or disabled by configuration options, it is only available when HIRC is selected in Oscillator configuration.

The non-inverting input can be connected to  $V_{SS}$  directly or through a  $1k\Omega$  resistor determined by configuration option. The inverting input can be connected to OPAN pin directly or through a  $1k\Omega$  resistor determined by configuration option. The output can be connected back to the inverting input through a  $60k\Omega$  resistor or not determined by configuration option. The output can be connected to PA1/OPAO/AN8 pin or connected to PB3/AN2 pin through a  $10k\Omega$  resistor by configuration options.

The input offset is adjustable by using a common mode input to calibrate the offset value.

The calibration steps are as the follows:

- (1) Setting OPAFM=1 to offset cancellation mode (S3 is closed)
- (2) Setting OPARS to select which input pin as reference voltage (S1 or S2 is closed)
- (3) Adjusting OPAOF0~OPAOF4 until output status is changed.
- (4) Setting OPAOFM=0 to normal operational amplifier mode



## Operational Amplifier register

The overall function is controlled by the OPAC register, it is used to select the operational amplifier input offset voltage cancellation and the input path.

### • OPAC Register

Bit	7	6	5	4	3	2	1	0
Name	OPAOP	OPAOFM	OPARS	OPAOF4	OPAOF3	OPAOF2	OPAOF1	OPAOF0
R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	0	0	1	0	0	0	0

Bit 7 **OPAOP**: Operational amplifier positive logic output

Bit 6 **OPAOFM**: Input offset voltage cancellation mode and operational amplifier mode selection

0: Operational amplifier mode

1: Input offset voltage cancellation mode

- Bit 5      **OPARS**: Operational amplifier input offset voltage cancellation reference selection bit  
             0: Select OPAN as the reference input  
             1: Select OPAP as the reference input
- Bit 4~0    **OPAOF4~OPAOF0**: Operational amplifier input offset voltage cancellation control bits

## Peripheral Clock Output

The Peripheral Clock Output allows the device to supply external hardware with a clock signal synchronised to the microcontroller clock.

### Peripheral Clock Operation

The peripheral clock output pin PCK is pin-shared with the I/O pin PB1. To operate as a peripheral clock output and not as an I/O pin, the configuration option must be set properly and a zero value must also be written to the bit PBC1 in the I/O port control register to ensure that the corresponding peripheral clock output pin is setup as an output. After these two initial steps have been carried out, writing a high value to the PB1 bit in the output data register will enable the PCK output function. Writing a zero value will disable the PCK output function and force the output low. In this way, the Port data output registers can be used as an on/off control for the PCK function. Note that if the configuration options have selected the PCK function, but a high value has been written to its corresponding bit in the PBC control register to configure the pin as an input, then the pin can still function as a normal input line, with pull-high resistor options.

The clock source for the Peripheral Clock Output can originate from the subdivided version of  $f_{SYS}$ . The division ratio value is determined by the PCKPSC bit in the CTRL0 register.

#### • CTRL0 Register

Bit	7	6	5	4	3	2	1	0
Name	C2VOINV	TMR0ECS	—	—	PCKPSC	—	LVDO	LVDC
R/W	R/W	R/W	—	—	R/W	—	R	R/W
POR	0	0	—	—	0	—	0	0

Bit 7      **C2VOINV**: Inverting control of the comparator 2 output signal  
 Described elsewhere

Bit 6      **TMR0ECS**: Select TMR0 external clock source  
 Described elsewhere

Bit 5~4    Unimplemented, read as “0”

Bit 3      **PCKPSC**: Peripheral clock prescaler stage  
 $f_{PCK} =$   
             0:  $f_{SYS}/2048$   
             1:  $f_{SYS}/4096$

Bit 2      Unimplemented, read as “0”

Bit 1      **LVDO**: Low voltage detector output  
 Described elsewhere

Bit 0      **LVDC**: Low voltage detector control  
 Described elsewhere

## Interrupts

Interrupts are an important part of any microcontroller system. When an external event or an internal function such as a Timer/Event Counter or an A/D converter requires microcontroller attention, their corresponding interrupt will enforce a temporary suspension of the main program allowing the microcontroller to direct attention to their respective needs. The device contains one external interrupts and several internal interrupts functions. The external interrupt is generated by the action of the external  $\overline{\text{INT}}$  pin, while the internal interrupts are generated by various internal functions such as Timer/Event Counters, comparator, LVD and A/D converter.

### Interrupt Registers

All the interrupt enable bits and the request flags are controlled by the INTC0, INTC1 and the MFI registers. By controlling the corresponding interrupt enable bits can enable or disable the interrupts. If an interrupt occurs, the request flag will be set and the EMI bit will be automatically cleared to disable other interrupts.

#### • INTC0 Register

Bit	7	6	5	4	3	2	1	0
Name	—	INTF_LVDF	CP1F	MFF	INTE_LVDE	CP1E	MFE	EMI
R/W	—	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	0	0	0	0	0	0	0

- Bit 7      Unimplemented, read as “0”
- Bit 6      **INTF\_LVDF**: External interrupt or LVD interrupt Request Flag  
             0: No request  
             1: Interrupt request  
             Selected by the configuration options
- Bit 5      **CP1F**: Comparator 1 Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 4      **MFF**: Multi-function Interrupt Request Flag  
             0: No request  
             1: Interrupt request
- Bit 3      **INTE\_LVDE**: External interrupt or LVD interrupt Control  
             0: Disable  
             1: Enable  
             Selected by the configuration options
- Bit 2      **CP1E**: Comparator 1 interrupt control  
             0: Disable  
             1: Enable
- Bit 1      **MFE**: Multi-function Interrupt Control  
             0: Disable  
             1: Enable
- Bit 0      **EMI**: Global Interrupt Control  
             0: Disable  
             1: Enable

• **INTC1 Register**

Bit	7	6	5	4	3	2	1	0
Name	—	T2F_ADF	T0F	—	—	T2E_ADE	T0E	—
R/W	—	R/W	R/W	—	—	R/W	R/W	—
POR	—	0	0	—	—	0	0	—

- Bit 7 Unimplemented, read as “0”
- Bit 6 **T2F\_ADF**: Timer/Event Counter 2 or A/D converter Interrupt Request Flag  
0: No request  
1: Interrupt request  
Selected by the configuration options
- Bit 5 **T0F**: Timer/Event Counter 0 interrupt request flag  
0: No request  
1: Interrupt request
- Bit 4~3 Unimplemented, read as “0”
- Bit 2 **T2E\_ADE**: Timer/Event Counter 2 or A/D converter Interrupt Control  
0: Disable  
1: Enable  
Selected by the configuration options
- Bit 1 **T0E**: Timer/Event Counter 0 Interrupt Control  
0: Disable  
1: Enable
- Bit 0 Unimplemented, read as “0”

• **MFI Register**

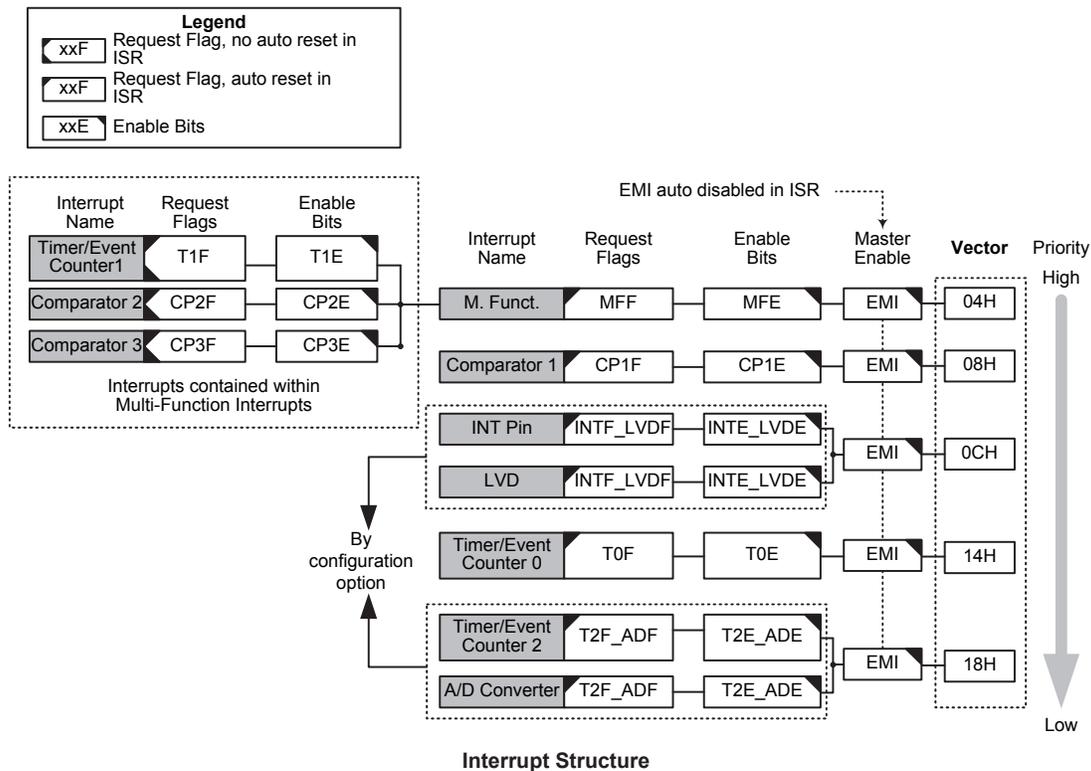
Bit	7	6	5	4	3	2	1	0
Name	—	T1F	CP2F	CP3F	—	T1E	CP2E	CP3E
R/W	—	R/W	R/W	R/W	—	R/W	R/W	R/W
POR	—	0	0	0	—	0	0	0

- Bit 7 Unimplemented, read as “0”
- Bit 6 **T1F**: Timer/Event Counter 1 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 5 **CP2F**: Comparator 2 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 4 **CP3F**: Comparator 3 Interrupt Request Flag  
0: No request  
1: Interrupt request
- Bit 3 Unimplemented, read as “0”
- Bit 2 **T1E**: Timer/Event Counter 1 Interrupt Control  
0: Disable  
1: Enable
- Bit 1 **CP2E**: Comparator 2 Interrupt Control  
0: Disable  
1: Enable
- Bit 0 **CP3E**: Comparator 3 Interrupt Control  
0: Disable  
1: Enable

### Interrupt Operation

When the conditions for an interrupt event occur, such as a Timer/Event Counter overflow, or A/D conversion completion, the relevant interrupt request flag will be set. When an interrupt is generated, the Program Counter, which stores the address of the next instruction to be executed, will be transferred onto the stack. The Program Counter will then be loaded with a new address which will be the value of the corresponding interrupt vector. The microcontroller will then fetch its next instruction from this interrupt vector. The instruction at this vector will usually be a “JMP” which will jump to another section of program which is known as the interrupt service routine. Here is located the code to control the appropriate interrupt. The interrupt service routine must be terminated with a “RETI”, which retrieves the original Program Counter address from the stack and allows the microcontroller to continue with normal execution at the point where the interrupt occurred.

The various interrupt enable bits, together with their associated request flags, are shown in the Accompanying diagrams with their order of priority.



Once an interrupt subroutine is serviced, all the other interrupts will be blocked, as the global interrupt enable bit, EMI bit will be cleared automatically. This will prevent any further interrupt nesting from occurring. However, if other interrupt requests occur during this interval, although the interrupt will not be immediately serviced, the request flag will still be recorded. If an interrupt requires immediate servicing while the program is already in another interrupt service routine, the EMI bit should be set after entering the routine, to allow interrupt nesting. If the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the Stack Pointer is decremented. If immediate service is desired, the stack must be prevented from becoming full.

### Interrupt Priority

Interrupts, occurring in the interval between the rising edges of two consecutive T2 pulses, will be serviced on the latter of the two T2 pulses, if the corresponding interrupts are enabled. In case of simultaneous requests, the following table shows the priority that is applied.

Interrupt sources	Priority	Vector
Multi-function interrupt	1	04H
Comparator 1 output interrupt	2	08H
External interrupt or LVD interrupt (Selected by configuration option)	3	0CH
Timer/Event Counter 0 overflow	4	14H
Timer/Event Counter 2 overflow or A/D converter interrupt (Selected by configuration option)	5	18H

The Timer/Event Counter 1 interrupt, comparator 2 interrupt and the comparator 3 interrupt share the same interrupt vector which is 04H. Each of these interrupts has their own individual interrupt flag but also share the same MFF interrupt flag. The MFF flag will be cleared by hardware once the Multi-function interrupt is serviced, however the individual interrupts that have triggered the Multi-function interrupt need to be cleared by the application program.

### External Interrupt

To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and respective external interrupt enable bit, INTE\_LVDE, must first be set. The external interrupt is triggered by falling edge transition of INT0 (inverted or non-inverted de-bounce signal from  $\overline{\text{INT}}$  or comparator 0 inverting input by software option), an external interrupt request will take place when the external interrupt request flag, INTF\_LVDF, is set. As the external interrupt pin is pin-shared with the PA2 pin, it can only be configured as the external interrupt pin if the external interrupt enable bit in the corresponding interrupt register has been set. The pin must also be setup as an input by setting the corresponding bit in the port control register. When the interrupt is enabled, the stack is not full, a subroutine call to the external interrupt vector, will take place. When the interrupt is serviced, the external interrupt request flag, INTF\_LVDF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts. Note that any pull-high resistor selections on the external interrupt pin will remain valid even if the pin is used as an external interrupt input.

Note that the trigger source of 0CH vector can be LVD interrupt or the external interrupt, which is selected by the configuration options. The external interrupt must be first selected by the configuration options before setting the external interrupt function.

### LVD Interrupt

An LVD Interrupt request will take place when the LVD Interrupt request flag, INTF\_LVDF, is set, which occurs when the Low Voltage Detector function detects a low power supply voltage. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and Low Voltage Interrupt enable bit, INTE\_LVDF, must first be set. When the interrupt is enabled, the stack is not full and a low voltage condition occurs, a subroutine call to the LVD Interrupt vector, will take place. When the Low Voltage Interrupt is serviced, the respective interrupt request flag, INTF\_LVDF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

Note that the trigger source of 0CH vector can be LVD interrupt or the external interrupt, which is selected by the configuration options. The LVD interrupt must be first selected by the configuration options before setting the LVD interrupt function.

### **A/D Converter Interrupt**

The A/D Converter Interrupt is controlled by the termination of an A/D conversion process. An A/D Converter Interrupt request will take place when the A/D Converter Interrupt request flag, T2F\_ADF, is set, which occurs when the A/D conversion process finishes. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and A/D Interrupt enable bit, T2E\_ADE, must first be set. When the interrupt is enabled, the stack is not full and the A/D conversion process has ended, a subroutine call to the A/D Converter Interrupt vector, will take place. When the interrupt is serviced, the A/D Converter Interrupt flag, T2F\_ADF, will be automatically cleared. The EMI bit will also be automatically cleared to disable other interrupts.

Note that the trigger source of 18H vector can be Timer/Event Counter 2 interrupt or the A/D converter interrupt, which is selected by the configuration options. The A/D converter interrupt must be first selected by the configuration options before setting the A/D converter interrupt function.

### **Multi-function Interrupt**

Within the device there is one Multi-function interrupt. Unlike the other independent interrupts, these interrupts have no independent source, but rather are formed from other existing interrupt sources, namely Timer/Event Counter 1 interrupt, comparator 2 interrupt and comparator 3 interrupt.

A Multi-function interrupt request will take place the Multi-function interrupt request flag, MFIF is set. The Multi-function interrupt flag will be set when any of its included functions generate an interrupt request flag. To allow the program to branch to its respective interrupt vector address, when the Multi-function interrupt is enabled and the stack is not full and either one of the interrupts contained within the Multi-function interrupt occurs, a subroutine call to the Multi-function interrupt vector will take place. When the interrupt is serviced, the related Multi-Function request flag will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

However, it must be noted that, although the Multi-function Interrupt flag will be automatically reset when the interrupt is serviced, the request flags from the original source of the Multi-function interrupt, namely Timer/Event Counter 1 interrupt, comparator 2 interrupt and comparator 3 interrupt, will not be automatically reset and must be manually reset by the application program.

### **Timer/Event Counter Interrupts**

There are three Timer/Event Counter interrupts, the Timer/Event Counter interrupt 0/2 is an independent interrupt and the Timer/Event Counter interrupt 1 is contained within the Multi-function Interrupt

For the Timer/Event Counter 0/2 interrupt to occur, the global interrupt enable bit, EMI, and the corresponding timer interrupt enable bit, T0E or T2E\_ADE, must first be set. An actual Timer/Event Counter interrupt will take place when the Timer/Event Counter request flag, T0F or T2F\_ADF, is set, a situation that will occur when the relevant Timer/Event Counter overflows. When the interrupt is enabled, the stack is not full and a Timer/Event Counter 0/2 overflow occurs, a subroutine call to the relevant timer interrupt vector, will take place. When the interrupt is serviced, the timer interrupt request flag, T0F or T2F\_ADF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

Note that the trigger source of 18H vector can be Timer/Event Counter 2 interrupt or the A/D converter interrupt, which is selected by the configuration options. The Timer/Event Counter 2 interrupt must be first selected by the configuration options before setting the Timer/Event Counter 2 interrupt function.

For the Timer/Event Counter 1 interrupt to occur, the global interrupt enable bit, EMI, the relevant Multi-function Interrupt enable bit, MFE and the corresponding timer interrupt enable bit, T1E, must first be set. An actual Timer/Event Counter interrupt will take place when the Timer/Event Counter request flag, T1F, is set, a situation that will occur when the relevant Timer/Event Counter overflows. When the interrupt is enabled, the stack is not full and a Timer/Event Counter 1 overflow occurs, a subroutine call to the relevant timer interrupt vector, will take place. When the interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts. however only the related MFF flag will be automatically cleared. As the Timer/Event Counter 1 interrupt request flag will not be automatically cleared, it has to be cleared by the application program.

### **Comparator Interrupt**

The comparator interrupt is controlled by three internal comparators. The comparator 1 is an independent interrupt and the comparator 2/3 is contained within the Multi-function Interrupt.

The comparator 1 interrupt request will take place when the comparator 1 interrupt request flag, CP1F is set, a situation that will occur when the comparator output bit changes state. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, and comparator interrupt enable bit, CP1E, must first be set. When the interrupt is enabled, the stack is not full and the comparator 1 input generate a comparator output transition, a subroutine call to the comparator interrupt vector, will take place. When the interrupt is serviced, the comparator interrupt request flag, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

The comparator 2/3 interrupt request will take place when the comparator 2/3 interrupt request flags, CP2F or CP3F, are set, a situation that will occur when the comparator output bit changes state. To allow the program to branch to its respective interrupt vector address, the global interrupt enable bit, EMI, the relevant Multi-function Interrupt enable bit, MFE and comparator interrupt enable bits, CP2E and CP3E, must first be set. When the interrupt is enabled, the stack is not full and the comparator inputs generate a comparator output transition, a subroutine call to the comparator interrupt vector, will take place. When the interrupt is serviced, the EMI bit will be automatically cleared to disable other interrupts. however only the related MFF flag will be automatically cleared. As the comparator 2/3 interrupt request flag will not be automatically cleared, it has to be cleared by the application program.

### **Interrupt Wake-up Function**

Each of the interrupt functions has the capability of waking up the microcontroller when in the Power down Mode. A wake-up is generated when an interrupt request flag changes from low to high and is independent of whether the interrupt is enabled or not. Therefore, even though the device is in the Power down Mode and its system oscillator stopped, situations such as external edge transitions on the external interrupt pin, a low power supply voltage may cause their respective interrupt flag to be set high and consequently generate an interrupt. Care must therefore be taken if spurious wake-up situations are to be avoided. If an interrupt wake-up function is to be disabled then the corresponding interrupt request flag should be set high before the device enters the Power down Mode. The interrupt enable bits have no effect on the interrupt wake-up function.

### **Programming Considerations**

By disabling the relevant interrupt enable bits, a requested interrupt can be prevented from being serviced, however, once an interrupt request flag is set, it will remain in this condition in the interrupt register until the corresponding interrupt is serviced or until the request flag is cleared by the application program.

It is recommended that programs do not use the “CALL” instruction within the interrupt service subroutine. Interrupts often occur in an unpredictable manner or need to be serviced immediately. If only one stack is left and the interrupt is not well controlled, the original control sequence will be damaged once a CALL subroutine is executed in the interrupt subroutine.

Every interrupt has the capability of waking up the microcontroller when it is in Power down Mode, the wake up being generated when the interrupt request flag changes from low to high. If it is required to prevent a certain interrupt from waking up the microcontroller then its respective request flag should be first set high before enter Power down Mode.

As only the Program Counter is pushed onto the stack, then when the interrupt is serviced, if the contents of the accumulator, status register or other registers are altered by the interrupt service program, their contents should be saved to the memory at the beginning of the interrupt service routine.

## Low Voltage Detector – LVD

Each device has a Low Voltage Detector function, also known as LVD. This enabled the device to monitor the power supply voltage,  $V_{DD}$ , and provide a warning signal should it fall below a certain level. This function may be especially useful in battery applications where the supply voltage will gradually reduce as the battery ages, as it allows an early warning battery low signal to be generated. The Low Voltage Detector also has the capability of generating an interrupt signal.

### LVD Register

LVD function can be enabled or disabled by configuration option. If the LVD configuration option is enabled, the user can use LVDC bit in system control register 0 to enable/disable the LVD circuit and read the LVD detector status from LVDO bit in system control register 0. The LVD voltage is fixed at 4.4V. The device also provides an interrupt for low voltage detection, see the section of interrupt for the details. In the power down mode, the LVD function will be disabled.

#### • CTRL0 Register

Bit	7	6	5	4	3	2	1	0
Name	C2VOINV	TMR0ECS	—	—	PCKPSC	—	LVDO	LVDC
R/W	R/W	R/W	—	—	R/W	—	R	R/W
POR	0	0	—	—	0	—	0	0

Bit 7      **C2VOINV**: Inverting control of the comparator 2 output signal  
 Described elsewhere

Bit 6      **TMR0ECS**: Select TMR0 external clock source  
 Described elsewhere

Bit 5~4    Unimplemented, read as “0”

Bit 3      **PCKPSC**: Peripheral clock prescaler stage  
 Described elsewhere

Bit 2      Unimplemented, read as “0”

Bit 1      **LVDO**: Low voltage detector output  
 0: Normal voltage  
 1: Low voltage detected

Bit 0      **LVDC**: Low voltage detector control  
 0: Disable  
 1: Enable

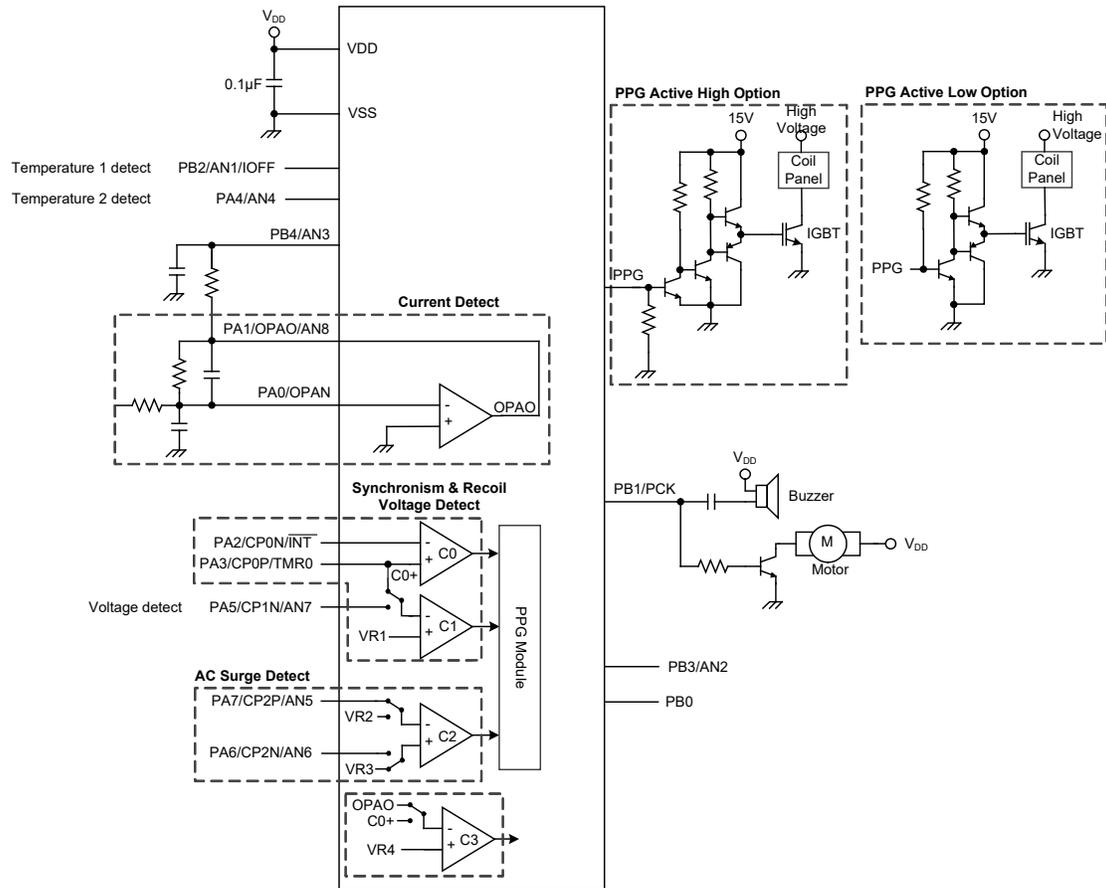
## Configuration Option

Configuration options refer to certain options within the MCU that are programmed into the device during the programming process. During the development process, these options are selected using the HT-IDE software development tools. As these options are programmed into the device using the hardware programming tools, once they are selected they cannot be changed later using the application program. All options must be defined for proper system function, the details of which are shown in the table.

No.	Options
<b>Wake-up Option</b>	
1	PA0~PA7 Wake-up Function Control 1. No wake-up 2. wake-up
<b>Pull-high Resistor Option</b>	
2	PA0~PA7 Pull- high Resistor Control 1. No pull-high 2. Pull-high
3	PB0~PB4 Pull- high Resistor Control 1. No pull-high 2. Pull-high
<b>I/O or other Options</b>	
4	PB2/IOFF selection: 1. PB2 2. IOFF
5	PB1/PCK selection: 1. PB1 2. PCK
<b>OPA Option</b>	
6	OPA Function: 1. Disable 2. Enable with output to PA1/OPAO/AN8 3. Enable with output through OPAR4 to PB3/AN2
7	OPAN Option: 1. Without OPAR3 2. With OPAR3
8	OPAP Option: 1. Directly to VSS 2. With OPAR2 to VSS
9	OPAO Option: 1. No feedback 2. With OPAR1 feedback
<b>PPG Option</b>	
10	PPG polarity control 1. PPG output active high 2. PPG output active low
11	PPG start count synchronize with clock 1. Synchronized with clock 2. Asynchronous with clock
<b>Comparator Option</b>	
12	Comparator 0 Control 1. Disable 2. Enable
13	Comparator 1 Control 1. Disable 2. Enable

No.	Options
14	Comparator 1 de-bounce Control 1. Disable 2. Enable
15	Comparator 2 Control 1. Disable 2. Enable, CP2N=VR3, CP2P=PA7 3. Enable, CP2N=PA6, CP2P=VR2 4. Enable, CP2N=PA6, CP2P=PA7
16	Comparator 2 de-bounce Control 1. Disable 2. Enable
17	Comparator 3 de-bounce Control 1. Disable 2. Enable
<b>Interrupt Option</b>	
18	Interrupt vector 18H source selection 1. Timer/Event Counter 2 interrupt 2. A/D Converter interrupt
19	Interrupt vector 0CH source selection 1. External interrupt 2. LVD interrupt
<b>WDT Option</b>	
20	WDT control 1. Disable 2. Enable
21	WDT time-out Period Selection 1. $2^{16}/f_s$ 2. $2^{15}/f_s$ 3. $2^{14}/f_s$ 4. $2^{13}/f_s$
22	WDT Clock Source Selection 1. LIRC 2. $f_{sys}/4$
23	Clear WDT Instruction 1. 1 instruction 2. 2 instructions
<b>LVD Selection</b>	
24	LVD function 1. Disable 2. Enable

Application Circuits



Note: The resistance and capacitance for reset circuit should be designed in such a way as to ensure that the  $V_{DD}$  is stable .

## **Instruction Set**

### **Introduction**

Central to the successful operation of any microcontroller is its instruction set, which is a set of program instruction codes that directs the microcontroller to perform certain operations. In the case of Holtek microcontroller, a comprehensive and flexible set of over 60 instructions is provided to enable programmers to implement their application with the minimum of programming overheads.

For easier understanding of the various instruction codes, they have been subdivided into several functional groupings.

### **Instruction Timing**

Most instructions are implemented within one instruction cycle. The exceptions to this are branch, call, or table read instructions where two instruction cycles are required. One instruction cycle is equal to 4 system clock cycles, therefore in the case of an 8MHz system oscillator, most instructions would be implemented within 0.5 $\mu$ s and branch or call instructions would be implemented within 1 $\mu$ s. Although instructions which require one more cycle to implement are generally limited to the JMP, CALL, RET, RETI and table read instructions, it is important to realize that any other instructions which involve manipulation of the Program Counter Low register or PCL will also take one more cycle to implement. As instructions which change the contents of the PCL will imply a direct jump to that new address, one more cycle will be required. Examples of such instructions would be “CLR PCL” or “MOV PCL, A”. For the case of skip instructions, it must be noted that if the result of the comparison involves a skip operation then this will also take one more cycle, if no skip is involved then only one cycle is required.

### **Moving and Transferring Data**

The transfer of data within the microcontroller program is one of the most frequently used operations. Making use of several kinds of MOV instructions, data can be transferred from registers to the Accumulator and vice-versa as well as being able to move specific immediate data directly into the Accumulator. One of the most important data transfer applications is to receive data from the input ports and transfer data to the output ports.

### **Arithmetic Operations**

The ability to perform certain arithmetic operations and data manipulation is a necessary feature of most microcontroller applications. Within the Holtek microcontroller instruction set are a range of add and subtract instruction mnemonics to enable the necessary arithmetic to be carried out. Care must be taken to ensure correct handling of carry and borrow data when results exceed 255 for addition and less than 0 for subtraction. The increment and decrement instructions such as INC, INCA, DEC and DECA provide a simple means of increasing or decreasing by a value of one of the values in the destination specified.

## Logical and Rotate Operation

The standard logical operations such as AND, OR, XOR and CPL all have their own instruction within the Holtek microcontroller instruction set. As with the case of most instructions involving data manipulation, data must pass through the Accumulator which may involve additional programming steps. In all logical data operations, the zero flag may be set if the result of the operation is zero. Another form of logical data manipulation comes from the rotate instructions such as RR, RL, RRC and RLC which provide a simple means of rotating one bit right or left. Different rotate instructions exist depending on program requirements. Rotate instructions are useful for serial port programming applications where data can be rotated from an internal register into the Carry bit from where it can be examined and the necessary serial bit set high or low. Another application which rotate data operations are used is to implement multiplication and division calculations.

## Branches and Control Transfer

Program branching takes the form of either jumps to specified locations using the JMP instruction or to a subroutine using the CALL instruction. They differ in the sense that in the case of a subroutine call, the program must return to the instruction immediately when the subroutine has been carried out. This is done by placing a return instruction “RET” in the subroutine which will cause the program to jump back to the address right after the CALL instruction. In the case of a JMP instruction, the program simply jumps to the desired location. There is no requirement to jump back to the original jumping off point as in the case of the CALL instruction. One special and extremely useful set of branch instructions are the conditional branches. Here a decision is first made regarding the condition of a certain data memory or individual bits. Depending upon the conditions, the program will continue with the next instruction or skip over it and jump to the following instruction. These instructions are the key to decision making and branching within the program perhaps determined by the condition of certain input switches or by the condition of internal data bits.

## Bit Operations

The ability to provide single bit operations on Data Memory is an extremely flexible feature of all Holtek microcontrollers. This feature is especially useful for output port bit programming where individual bits or port pins can be directly set high or low using either the “SET [m].i” or “CLR [m].i” instructions respectively. The feature removes the need for programmers to first read the 8-bit output port, manipulate the input data to ensure that other bits are not changed and then output the port with the correct new data. This read-modify-write process is taken care of automatically when these bit operation instructions are used.

## Table Read Operations

Data storage is normally implemented by using registers. However, when working with large amounts of fixed data, the volume involved often makes it inconvenient to store the fixed data in the Data Memory. To overcome this problem, Holtek microcontrollers allow an area of Program Memory to be setup as a table where data can be directly stored. A set of easy to use instructions provides the means by which this fixed data can be referenced and retrieved from the Program Memory.

## Other Operations

In addition to the above functional instructions, a range of other instructions also exist such as the “HALT” instruction for Power-down operations and instructions to control the operation of the Watchdog Timer for reliable program operations under extreme electric or electromagnetic environments. For their relevant operations, refer to the functional related sections.

## Instruction Set Summary

The following table depicts a summary of the instruction set categorised according to function and can be consulted as a basic instruction reference using the following listed conventions.

### Table Conventions

x: Bits immediate data  
m: Data Memory address  
A: Accumulator  
i: 0~7 number of bits  
addr: Program memory address

Mnemonic	Description	Cycles	Flag Affected
<b>Arithmetic</b>			
ADD A,[m]	Add Data Memory to ACC	1	Z, C, AC, OV
ADDM A,[m]	Add ACC to Data Memory	1 <sup>Note</sup>	Z, C, AC, OV
ADD A,x	Add immediate data to ACC	1	Z, C, AC, OV
ADC A,[m]	Add Data Memory to ACC with Carry	1	Z, C, AC, OV
ADCM A,[m]	Add ACC to Data memory with Carry	1 <sup>Note</sup>	Z, C, AC, OV
SUB A,x	Subtract immediate data from the ACC	1	Z, C, AC, OV
SUB A,[m]	Subtract Data Memory from ACC	1	Z, C, AC, OV
SUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV
SBC A,[m]	Subtract Data Memory from ACC with Carry	1	Z, C, AC, OV
SBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	1 <sup>Note</sup>	Z, C, AC, OV
DAA [m]	Decimal adjust ACC for Addition with result in Data Memory	1 <sup>Note</sup>	C
<b>Logic Operation</b>			
AND A,[m]	Logical AND Data Memory to ACC	1	Z
OR A,[m]	Logical OR Data Memory to ACC	1	Z
XOR A,[m]	Logical XOR Data Memory to ACC	1	Z
ANDM A,[m]	Logical AND ACC to Data Memory	1 <sup>Note</sup>	Z
ORM A,[m]	Logical OR ACC to Data Memory	1 <sup>Note</sup>	Z
XORM A,[m]	Logical XOR ACC to Data Memory	1 <sup>Note</sup>	Z
AND A,x	Logical AND immediate Data to ACC	1	Z
OR A,x	Logical OR immediate Data to ACC	1	Z
XOR A,x	Logical XOR immediate Data to ACC	1	Z
CPL [m]	Complement Data Memory	1 <sup>Note</sup>	Z
CPLA [m]	Complement Data Memory with result in ACC	1	Z
<b>Increment &amp; Decrement</b>			
INCA [m]	Increment Data Memory with result in ACC	1	Z
INC [m]	Increment Data Memory	1 <sup>Note</sup>	Z
DECA [m]	Decrement Data Memory with result in ACC	1	Z
DEC [m]	Decrement Data Memory	1 <sup>Note</sup>	Z
<b>Rotate</b>			
RRA [m]	Rotate Data Memory right with result in ACC	1	None
RR [m]	Rotate Data Memory right	1 <sup>Note</sup>	None
RRCA [m]	Rotate Data Memory right through Carry with result in ACC	1	C
RRC [m]	Rotate Data Memory right through Carry	1 <sup>Note</sup>	C
RLA [m]	Rotate Data Memory left with result in ACC	1	None
RL [m]	Rotate Data Memory left	1 <sup>Note</sup>	None
RLCA [m]	Rotate Data Memory left through Carry with result in ACC	1	C
RLC [m]	Rotate Data Memory left through Carry	1 <sup>Note</sup>	C

Mnemonic	Description	Cycles	Flag Affected
<b>Data Move</b>			
MOV A,[m]	Move Data Memory to ACC	1	None
MOV [m],A	Move ACC to Data Memory	1 <sup>Note</sup>	None
MOV A,x	Move immediate data to ACC	1	None
<b>Bit Operation</b>			
CLR [m].i	Clear bit of Data Memory	1 <sup>Note</sup>	None
SET [m].i	Set bit of Data Memory	1 <sup>Note</sup>	None
<b>Branch Operation</b>			
JMP addr	Jump unconditionally	2	None
SZ [m]	Skip if Data Memory is zero	1 <sup>Note</sup>	None
SZA [m]	Skip if Data Memory is zero with data movement to ACC	1 <sup>Note</sup>	None
SZ [m].i	Skip if bit i of Data Memory is zero	1 <sup>Note</sup>	None
SNZ [m].i	Skip if bit i of Data Memory is not zero	1 <sup>Note</sup>	None
SIZ [m]	Skip if increment Data Memory is zero	1 <sup>Note</sup>	None
SDZ [m]	Skip if decrement Data Memory is zero	1 <sup>Note</sup>	None
SIZA [m]	Skip if increment Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
SDZA [m]	Skip if decrement Data Memory is zero with result in ACC	1 <sup>Note</sup>	None
CALL addr	Subroutine call	2	None
RET	Return from subroutine	2	None
RET A,x	Return from subroutine and load immediate data to ACC	2	None
RETI	Return from interrupt	2	None
<b>Table Read Operation</b>			
TABRD [m]	Read table (specific page) to TBLH and Data Memory	2 <sup>Note</sup>	None
TABRDC [m]	Read table (current page) to TBLH and Data Memory	2 <sup>Note</sup>	None
TABRDL [m]	Read table (last page) to TBLH and Data Memory	2 <sup>Note</sup>	None
<b>Miscellaneous</b>			
NOP	No operation	1	None
CLR [m]	Clear Data Memory	1 <sup>Note</sup>	None
SET [m]	Set Data Memory	1 <sup>Note</sup>	None
CLR WDT	Clear Watchdog Timer	1	TO, PDF
CLR WDT1	Pre-clear Watchdog Timer	1	TO, PDF
CLR WDT2	Pre-clear Watchdog Timer	1	TO, PDF
SWAP [m]	Swap nibbles of Data Memory	1 <sup>Note</sup>	None
SWAPA [m]	Swap nibbles of Data Memory with result in ACC	1	None
HALT	Enter power down mode	1	TO, PDF

Note: 1. For skip instructions, if the result of the comparison involves a skip then two cycles are required, if no skip takes place only one cycle is required.

2. Any instruction which changes the contents of the PCL will also require 2 cycles for execution.
3. For the “CLR WDT1” and “CLR WDT2” instructions the TO and PDF flags may be affected by the execution status. The TO and PDF flags are cleared after both “CLR WDT1” and “CLR WDT2” instructions are consecutively executed. Otherwise the TO and PDF flags remain unchanged.

## Instruction Definition

<b>ADC A,[m]</b>	Add Data Memory to ACC with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C
<b>ADCM A,[m]</b>	Add ACC to Data Memory with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C
<b>ADD A,[m]</b>	Add Data Memory to ACC
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C
<b>ADD A,x</b>	Add immediate data to ACC
Description	The contents of the Accumulator and the specified immediate data are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + x$
Affected flag(s)	OV, Z, AC, C
<b>ADDM A,[m]</b>	Add ACC to Data Memory
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C
<b>AND A,[m]</b>	Logical AND Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
<b>AND A,x</b>	Logical AND immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bit wise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } x$
Affected flag(s)	Z
<b>ANDM A,[m]</b>	Logical AND ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z

<b>CALL addr</b>	Subroutine call
Description	Unconditionally calls a subroutine at the specified address. The Program Counter then increments by 1 to obtain the address of the next instruction which is then pushed onto the stack. The specified address is then loaded and the program continues execution from this new address. As this instruction requires an additional operation, it is a two cycle instruction.
Operation	Stack $\leftarrow$ Program Counter + 1 Program Counter $\leftarrow$ addr
Affected flag(s)	None
<b>CLR [m]</b>	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	[m] $\leftarrow$ 00H
Affected flag(s)	None
<b>CLR [m].i</b>	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	[m].i $\leftarrow$ 0
Affected flag(s)	None
<b>CLR WDT</b>	Clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared.
Operation	WDT cleared TO $\leftarrow$ 0 PDF $\leftarrow$ 0
Affected flag(s)	TO, PDF
<b>CLR WDT1</b>	Pre-clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared. Note that this instruction works in conjunction with CLR WDT2 and must be executed alternately with CLR WDT2 to have effect. Repetitively executing this instruction without alternately executing CLR WDT2 will have no effect.
Operation	WDT cleared TO $\leftarrow$ 0 PDF $\leftarrow$ 0
Affected flag(s)	TO, PDF
<b>CLR WDT2</b>	Pre-clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared. Note that this instruction works in conjunction with CLR WDT1 and must be executed alternately with CLR WDT1 to have effect. Repetitively executing this instruction without alternately executing CLR WDT1 will have no effect.
Operation	WDT cleared TO $\leftarrow$ 0 PDF $\leftarrow$ 0
Affected flag(s)	TO, PDF
<b>CPL [m]</b>	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	[m] $\leftarrow$ $\overline{[m]}$
Affected flag(s)	Z

<b>CPLA [m]</b>	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow \overline{[m]}$
Affected flag(s)	Z
<b>DAA [m]</b>	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	$[m] \leftarrow ACC + 00H$ or $[m] \leftarrow ACC + 06H$ or $[m] \leftarrow ACC + 60H$ or $[m] \leftarrow ACC + 66H$
Affected flag(s)	C
<b>DEC [m]</b>	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	$[m] \leftarrow [m] - 1$
Affected flag(s)	Z
<b>DECA [m]</b>	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Z
<b>HALT</b>	Enter power down mode
Description	This instruction stops the program execution and turns off the system clock. The contents of the Data Memory and registers are retained. The WDT and prescaler are cleared. The power down flag PDF is set and the WDT time-out flag TO is cleared.
Operation	TO $\leftarrow$ 0 PDF $\leftarrow$ 1
Affected flag(s)	TO, PDF
<b>INC [m]</b>	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	$[m] \leftarrow [m] + 1$
Affected flag(s)	Z
<b>INCA [m]</b>	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] + 1$
Affected flag(s)	Z

<b>JMP addr</b>	Jump unconditionally
Description	The contents of the Program Counter are replaced with the specified address. Program execution then continues from this new address. As this requires the insertion of a dummy instruction while the new address is loaded, it is a two cycle instruction.
Operation	Program Counter $\leftarrow$ addr
Affected flag(s)	None
<b>MOV A,[m]</b>	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	ACC $\leftarrow$ [m]
Affected flag(s)	None
<b>MOV A,x</b>	Move immediate data to ACC
Description	The immediate data specified is loaded into the Accumulator.
Operation	ACC $\leftarrow$ x
Affected flag(s)	None
<b>MOV [m],A</b>	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	[m] $\leftarrow$ ACC
Affected flag(s)	None
<b>NOP</b>	No operation
Description	No operation is performed. Execution continues with the next instruction.
Operation	No operation
Affected flag(s)	None
<b>OR A,[m]</b>	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC $\leftarrow$ ACC "OR" [m]
Affected flag(s)	Z
<b>OR A,x</b>	Logical OR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC $\leftarrow$ ACC "OR" x
Affected flag(s)	Z
<b>ORM A,[m]</b>	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	[m] $\leftarrow$ ACC "OR" [m]
Affected flag(s)	Z
<b>RET</b>	Return from subroutine
Description	The Program Counter is restored from the stack. Program execution continues at the restored address.
Operation	Program Counter $\leftarrow$ Stack
Affected flag(s)	None

<b>RET A,x</b>	Return from subroutine and load immediate data to ACC
Description	The Program Counter is restored from the stack and the Accumulator loaded with the specified immediate data. Program execution continues at the restored address.
Operation	Program Counter ← Stack ACC ← x
Affected flag(s)	None
<b>RETI</b>	Return from interrupt
Description	The Program Counter is restored from the stack and the interrupts are re-enabled by setting the EMI bit. EMI is the master interrupt global enable bit. If an interrupt was pending when the RETI instruction is executed, the pending Interrupt routine will be processed before returning to the main program.
Operation	Program Counter ← Stack EMI ← 1
Affected flag(s)	None
<b>RL [m]</b>	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	[m].(i+1) ← [m].i; (i=0~6) [m].0 ← [m].7
Affected flag(s)	None
<b>RLA [m]</b>	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.(i+1) ← [m].i; (i=0~6) ACC.0 ← [m].7
Affected flag(s)	None
<b>RLC [m]</b>	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	[m].(i+1) ← [m].i; (i=0~6) [m].0 ← C C ← [m].7
Affected flag(s)	C
<b>RLCA [m]</b>	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.(i+1) ← [m].i; (i=0~6) ACC.0 ← C C ← [m].7
Affected flag(s)	C
<b>RR [m]</b>	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	[m].i ← [m].(i+1); (i=0~6) [m].7 ← [m].0
Affected flag(s)	None

<b>RRA [m]</b>	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory is rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim6)$ $ACC.7 \leftarrow [m].0$
Affected flag(s)	None
<b>RRC [m]</b>	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	$[m].i \leftarrow [m].(i+1); (i=0\sim6)$ $[m].7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>RRCA [m]</b>	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i=0\sim6)$ $ACC.7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
<b>SBC A,[m]</b>	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - \bar{C}$
Affected flag(s)	OV, Z, AC, C
<b>SBCM A,[m]</b>	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m] - \bar{C}$
Affected flag(s)	OV, Z, AC, C
<b>SDZ [m]</b>	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] - 1$ Skip if $[m]=0$
Affected flag(s)	None

<b>SDZA [m]</b>	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] - 1$ Skip if ACC=0
Affected flag(s)	None
<b>SET [m]</b>	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
<b>SET [m].i</b>	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None
<b>SIZ [m]</b>	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if [m]=0
Affected flag(s)	None
<b>SIZA [m]</b>	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if ACC=0
Affected flag(s)	None
<b>SNZ [m].i</b>	Skip if bit i of Data Memory is not 0
Description	If bit i of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None
<b>SUB A,[m]</b>	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C

<b>SUBM A,[m]</b>	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C
<b>SUB A,x</b>	Subtract immediate data from ACC
Description	The immediate data specified by the code is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - x$
Affected flag(s)	OV, Z, AC, C
<b>SWAP [m]</b>	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	$[m].3 \sim [m].0 \leftrightarrow [m].7 \sim [m].4$
Affected flag(s)	None
<b>SWAPA [m]</b>	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC.3 \sim ACC.0 \leftarrow [m].7 \sim [m].4$ $ACC.7 \sim ACC.4 \leftarrow [m].3 \sim [m].0$
Affected flag(s)	None
<b>SZ [m]</b>	Skip if Data Memory is 0
Description	If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if $[m]=0$
Affected flag(s)	None
<b>SZA [m]</b>	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m]$ Skip if $[m]=0$
Affected flag(s)	None
<b>SZ [m].i</b>	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if $[m].i=0$
Affected flag(s)	None

<b>TABRD [m]</b>	Read table (specific page) to TBLH and Data Memory
Description	The low byte of the program code (specific page) addressed by the table pointer pair (TBHP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>TABRDC [m]</b>	Read table (current page) to TBLH and Data Memory
Description	The low byte of the program code (current page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>TABRDL [m]</b>	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
<b>XOR A,[m]</b>	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" [m]
Affected flag(s)	Z
<b>XORM A,[m]</b>	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "XOR" [m]
Affected flag(s)	Z
<b>XOR A,x</b>	Logical XOR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" x
Affected flag(s)	Z

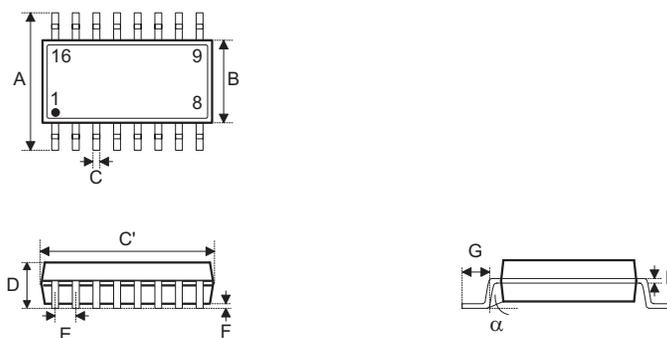
## Package Information

Note that the package information provided here is for consultation purposes only. As this information may be updated at regular intervals users are reminded to consult the [Holtek website](#) for the latest version of the [Package/Carton Information](#).

Additional supplementary information with regard to packaging is listed below. Click on the relevant section to be transferred to the relevant website page.

- Package Information (include Outline Dimensions, Product Tape and Reel Specifications)
- The Operation Instruction of Packing Materials
- Carton information

16-pin NSOP (150mil) Outline Dimensions



Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	—	0.236 BSC	—
B	—	0.154 BSC	—
C	0.012	—	0.020
C'	—	0.390 BSC	—
D	—	—	0.069
E	—	0.050 BSC	—
F	0.004	—	0.010
G	0.016	—	0.050
H	0.004	—	0.010
$\alpha$	0°	—	8°

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	—	6.00 BSC	—
B	—	3.90 BSC	—
C	0.31	—	0.51
C'	—	9.90 BSC	—
D	—	—	1.75
E	—	1.27 BSC	—
F	0.10	—	0.25
G	0.40	—	1.27
H	0.10	—	0.25
$\alpha$	0°	—	8°

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